## **CHAPTER 6: Basic SQL**

#### **Answers to Selected Exercises**

**6.5** - Consider the database shown in Figure 1.2, whose schema is shown in Figure 2.1. What are the referential integrity constraints that should hold on the schema? Write appropriate SQL DDL statements to define the database.

#### Answer:

The following referential integrity constraints should hold (we use the notation: R.(A1, ..., An) --> S.(B1, ..., Bn)

to represent a foreign key from the attributes A1, ..., An of R (the referencing relation)

to S (the referenced relation)):

PREREQUISITE.(CourseNumber) --> COURSE.(CourseNumber)

PREREQUISITE.(PrerequisiteNumber) --> COURSE.(CourseNumber)

SECTION.(CourseNumber) --> COURSE.(CourseNumber)

GRADE REPORT.(StudentNumber) --> STUDENT.(StudentNumber)

GRADE REPORT.(SectionIdentifier) --> SECTION.(SectionIdentifier)

One possible set of CREATE TABLE statements to define the database is given below.

CREATE TABLE STUDENT (Name VARCHAR(30) NOT NULL,

StudentNumber INTEGER NOT NULL,

Class CHAR NOT NULL,

Major CHAR(4),

PRIMARY KEY (StudentNumber) );

CREATE TABLE COURSE ( CourseName VARCHAR(30) NOT NULL,

CourseNumber CHAR(8) NOT NULL,

CreditHours INTEGER,

Department CHAR(4),

PRIMARY KEY (CourseNumber),

UNIQUE (CourseName));

CREATE TABLE PREREQUISITE (CourseNumber CHAR(8) NOT NULL,

PrerequisiteNumber CHAR(8) NOT NULL,

PRIMARY KEY (CourseNumber, PrerequisiteNumber),

FOREIGN KEY (CourseNumber) REFERENCES

COURSE (CourseNumber),

FOREIGN KEY (PrerequisiteNumber) REFERENCES

COURSE (CourseNumber) );

CREATE TABLE SECTION (SectionIdentifier INTEGER NOT NULL,

CourseNumber CHAR(8) NOT NULL,

Semester VARCHAR(6) NOT NULL,

Year CHAR(4) NOT NULL,

Instructor VARCHAR(15),

PRIMARY KEY (SectionIdentifier),

FOREIGN KEY (CourseNumber) REFERENCES

COURSE (CourseNumber) );

CREATE TABLE GRADE\_REPORT ( StudentNumber INTEGER NOT NULL,

SectionIdentifier INTEGER NOT NULL,

Grade CHAR,

PRIMARY KEY (StudentNumber, SectionIdentifier),

FOREIGN KEY (StudentNumber) REFERENCES

STUDENT (StudentNumber),

FOREIGN KEY (SectionIdentifier) REFERENCES

SECTION (SectionIdentifier) );

# **6.6** - Repeat Exercise 4.5, but use the AIRLINE schema of Figure 3.8.

## Answer:

```
The following referential integrity constraints should hold:
FLIGHT LEG.(FLIGHT NUMBER) --> FLIGHT.(NUMBER)
FLIGHT LEG.(DEPARTURE AIRPORT CODE) --> AIRPORT.(AIRPORT CODE)
FLIGHT LEG.(ARRIVAL AIRPORT CODE) --> AIRPORT.(AIRPORT CODE)
LEG_INSTANCE.(FLIGHT_NUMBER, LEG_NUMBER) -->
FLIGHT LEG.(FLIGHT NUMBER, LEG NUMBER)
LEG INSTANCE.(AIRPLANE ID) --> AIRPLANE.(AIRPLANE ID)
LEG_INSTANCE.(DEPARTURE_AIRPORT_CODE) --> AIRPORT.(AIRPORT_CODE)
LEG_INSTANCE.(ARRIVAL_AIRPORT_CODE) --> AIRPORT.(AIRPORT_CODE)
FARES.(FLIGHT_NUMBER) --> FLIGHT.(NUMBER)
CAN_LAND.(AIRPLANE_TYPE_NAME) --> AIRPLANE_TYPE.(TYPE_NAME)
CAN_LAND.(AIRPORT_CODE) --> AIRPORT.(AIRPORT_CODE)
AIRPLANE.(AIRPLANE TYPE) --> AIRPLANE TYPE.(TYPE NAME)
SEAT_RESERVATION.(FLIGHT_NUMBER, LEG_NUMBER, DATE) -->
LEG_INSTANCE.(FLIGHT_NUMBER, LEG_NUMBER, DATE)
One possible set of CREATE TABLE statements to define the database is given below.
CREATE TABLE AIRPORT (AIRPORT CODE CHAR(3) NOT NULL,
NAME VARCHAR(30) NOT NULL,
CITY VARCHAR(30) NOT NULL,
STATE VARCHAR(30),
PRIMARY KEY (AIRPORT CODE) );
CREATE TABLE FLIGHT ( NUMBER VARCHAR(6) NOT NULL,
AIRLINE VARCHAR(20) NOT NULL,
WEEKDAYS VARCHAR(10) NOT NULL,
PRIMARY KEY (NUMBER) );
CREATE TABLE FLIGHT LEG (FLIGHT NUMBER VARCHAR(6) NOT NULL,
LEG_NUMBER INTEGER NOT NULL,
DEPARTURE_AIRPORT_CODE CHAR(3) NOT NULL,
SCHEDULED DEPARTURE TIME TIMESTAMP WITH TIME ZONE,
ARRIVAL_AIRPORT_CODE CHAR(3) NOT NULL,
SCHEDULED_ARRIVAL_TIME TIMESTAMP WITH TIME ZONE,
PRIMARY KEY (FLIGHT_NUMBER, LEG_NUMBER),
FOREIGN KEY (FLIGHT NUMBER) REFERENCES FLIGHT (NUMBER),
FOREIGN KEY (DEPARTURE_AIRPORT_CODE) REFERENCES
AIRPORT (AIRPORT CODE),
FOREIGN KEY (ARRIVAL_AIRPORT_CODE) REFERENCES
AIRPORT (AIRPORT_CODE) );
CREATE TABLE LEG_INSTANCE (FLIGHT_NUMBER VARCHAR(6) NOT NULL,
LEG_NUMBER INTEGER NOT NULL,
LEG DATE DATE NOT NULL,
NO OF_AVAILABLE_SEATS INTEGER,
AIRPLANE_ID INTEGER,
DEPARTURE_AIRPORT_CODE CHAR(3),
DEPARTURE_TIME TIMESTAMP WITH TIME ZONE,
ARRIVAL_AIRPORT_CODE CHAR(3),
ARRIVAL_TIME TIMESTAMP WITH TIME ZONE,
PRIMARY KEY (FLIGHT NUMBER, LEG NUMBER, LEG DATE),
FOREIGN KEY (FLIGHT NUMBER, LEG NUMBER) REFERENCES
FLIGHT LEG (FLIGHT NUMBER, LEG NUMBER),
FOREIGN KEY (AIRPLANE ID) REFERENCES
AIRPLANE (AIRPLANE ID),
FOREIGN KEY (DEPARTURE_AIRPORT_CODE) REFERENCES
AIRPORT (AIRPORT_CODE),
FOREIGN KEY (ARRIVAL AIRPORT CODE) REFERENCES
AIRPORT (AIRPORT_CODE) );
CREATE TABLE FARES (FLIGHT NUMBER VARCHAR(6) NOT NULL,
FARE_CODE VARCHAR(10) NOT NULL,
```

```
AMOUNT DECIMAL(8,2) NOT NULL,
RESTRICTIONS VARCHAR(200),
PRIMARY KEY (FLIGHT_NUMBÉR, FARE_CODE),
FOREIGN KEY (FLIGHT NUMBER) REFERENCES FLIGHT (NUMBER) );
CREATE TABLE AIRPLANE_TYPE (TYPE_NAME VARCHAR(20) NOT NULL,
MAX SEATS INTEGER NOT NULL,
COMPANY VARCHAR(15) NOT NULL,
PRIMARY KEY (TYPE_NAME));
CREATE TABLE CAN_LAND ( AIRPLANE_TYPE_NAME VARCHAR(20) NOT NULL,
AIRPORT_CODE CHAR(3) NOT NULL,
PRIMARY KEY (AIRPLANE TYPE NAME, AIRPORT CODE),
FOREIGN KEY (AIRPLANE TYPE NAME) REFERENCES
AIRPLANE TYPE (TYPE NAME),
FOREIGN KEY (AIRPORT CODE) REFERENCES
AIRPORT (AIRPORT_CODE) );
CREATE TABLE AIRPLANE (AIRPLANE ID INTEGER NOT NULL,
TOTAL_NUMBER_OF_SEATS INTEGER NOT NULL,
AIRPLANE TYPE VARCHAR(20) NOT NULL,
PRIMARY KEY (AIRPLANE_ID),
FOREIGN KEY (AIRPLANE_TYPE) REFERENCES AIRPLANE_TYPE (TYPE_NAME) );
CREATE TABLE SEAT_RESERVATION (FLIGHT_NUMBER VARCHAR(6) NOT NULL,
LEG NUMBER INTEGER NOT NULL,
LEG DATE DATE NOT NULL,
SEAT_NUMBER VARCHAR(4),
CUSTOMER NAME VARCHAR(30) NOT NULL,
CUSTOMER PHONE CHAR(12),
PRIMARY KEY (FLIGHT NUMBER, LEG NUMBER, LEG DATE, SEAT NUMBER),
FOREIGN KEY (FLIGHT NUMBER, LEG NUMBER, LEG DATE) REFERENCES
LEG INSTANCE (FLIGHT NUMBER, LEG NUMBER, LEG DATE) );
```

**6.7** - Consider the LIBRARY relational database schema of Figure 4.6. Choose the appropriate action (reject, cascade, set to null, set to default) for each referential integrity constraint, both for the *deletion* of a referenced tuple, and for the *update* of a primary key attribute value in a referenced tuple. Justify your choices.

#### Answer:

Below are possible choices. In general, if it is not clear which action to choose, REJECT should be chosen, since it will not permit automatic changes to happen (by update propagation) that may be unintended.

BOOK\_AUTHORS.(Bookld) --> BOOK.(Bookld)

CASCADE on both DELETE or UPDATE (since this corresponds to a multi-valued attribute of BOOK (see the solution to Exercise 6.27); hence, if a BOOK is deleted, or the value of its Bookld is updated (changed), the deletion or change is automatically propagated to the referencing BOOK\_AUTHORS tuples)

BOOK.(PublisherName) --> PUBLISHER.(Name)

REJECT on DELETE (we should not delete a PUBLISHER tuple which has existing BOOK tuples that reference the PUBLISHER)

CASCADE on UPDATE (if a PUBLISHER's Name is updated, the change should be propagated automatically to all referencing BOOK tuples)

BOOK LOANS.(Bookld) --> BOOK.(Bookld)

CASCADE on both DELETE or UPDATE (if a BOOK is deleted, or the value of its Bookld is updated (changed), the deletion or change is automatically propagated to the referencing BOOK LOANS tuples) (Note: One could also choose REJECT on DELETE)

BOOK COPIES.(Bookld) --> BOOK.(Bookld)

CASCADE on both DELETE or UPDATE (if a BOOK is deleted, or the value of its Bookld is updated (changed), the deletion or change is automatically propagated to the referencing

**BOOK COPIES tuples)** 

BOOK\_LOANS.(CardNo) --> BORROWER.(CardNo)

CASCADE on both DELETE or UPDATE (if a BORROWER tuple is deleted, or the value of its CardNo is updated (changed), the deletion or change is automatically propagated to the referencing BOOK\_LOANS tuples) (Note: One could also choose REJECT on DELETE, with the idea that if a BORROWER is deleted, it is necessary first to make a printout of all BOOK\_LOANS outstanding before deleting the BORROWER; in this case, the tuples in BOOK\_LOANS that reference the BORROWER being deleted would first be explicitly deleted after making the printout, and before the BORROWER is deleted)

BOOK COPIES.(Branchld) --> LIBRARY BRANCH.(Branchld)

CASCADE on both DELETE or UPDATE (if a LIBRARY\_BRANCH is deleted, or the value of its Branchld is updated (changed), the deletion or change is automatically propagated to the referencing BOOK\_COPIES tuples) (Note: One could also choose REJECT on DELETE) BOOK LOANS.(Branchld) --> LIBRARY BRANCH.(Branchld)

CASCADE on both DELETE or UPDATE (if a LIBRARY\_BRANCH is deleted, or the value of its BranchId is updated (changed), the deletion or change is automatically propagated to the referencing BOOK\_LOANS tuples) (Note: One could also choose REJECT on DELETE)

**6.8** - Write appropriate SQL DDL statements for declaring the LIBRARY relational database schema of Figure 4.6. Specify the keys and referential triggered actions.

## Answer:

One possible set of CREATE TABLE statements is given below:

CREATE TABLE BOOK (Bookld CHAR(20) NOT NULL,

Title VARCHAR(30) NOT NULL,

PublisherName VARCHAR(20),

PRIMARY KEY (Bookld),

FOREIGN KEY (PublisherName) REFERENCES PUBLISHER (Name) ON UPDATE CASCADE );

CREATE TABLE BOOK AUTHORS (Bookld CHAR(20) NOT NULL,

AuthorName VARCHAR(30) NOT NULL,

PRIMARY KEY (Bookld, AuthorName),

FOREIGN KEY (BookId) REFERENCES BOOK (BookId)

ON DELETE CASCADE ON UPDATE CASCADE );

CREATE TABLE PUBLISHER (Name VARCHAR(20) NOT NULL,

Address VARCHAR(40) NOT NULL,

Phone CHAR(12),

PRIMARY KEY (Name) );

CREATE TABLE BOOK COPIES (Bookld CHAR(20) NOT NULL,

Branchid INTEGER NOT NULL,

No Of Copies INTEGER NOT NULL,

PRIMARY KEY (Bookld, Branchld),

FOREIGN KEY (BookId) REFERENCES BOOK (BookId)

ON DELETE CASCADE ON UPDATE CASCADE,

FOREIGN KEY (Branchid) REFERENCES BRANCH (Branchid)

ON DELETE CASCADE ON UPDATE CASCADE );

CREATE TABLE BORROWER ( CardNo INTEGER NOT NULL,

Name VARCHAR(30) NOT NULL,

Address VARCHAR(40) NOT NULL,

Phone CHAR(12),

PRIMARY KEY (CardNo));

CREATE TABLE BOOK LOANS (CardNo INTEGER NOT NULL,

Bookld CHAR(20) NOT NULL,

Branchid INTEGER NOT NULL,
DateOut DATE NOT NULL,
DueDate DATE NOT NULL,
PRIMARY KEY (CardNo, Bookld, Branchid),
FOREIGN KEY (CardNo) REFERENCES BORROWER (CardNo)
ON DELETE CASCADE ON UPDATE CASCADE,
FOREIGN KEY (Branchid) REFERENCES LIBRARY\_BRANCH (Branchid)
ON DELETE CASCADE ON UPDATE CASCADE,
FOREIGN KEY (Bookld) REFERENCES BOOK (Bookld)
ON DELETE CASCADE ON UPDATE CASCADE);
CREATE TABLE LIBRARY\_BRANCH (Branchid INTEGER NOT NULL,
BranchName VARCHAR(20) NOT NULL,
Address VARCHAR(40) NOT NULL,
PRIMARY KEY (Branchid));

**6.9** - How can the key and foreign key constraints be enforced by the DBMS? Is the enforcement technique you suggest difficult to implement? Can the constraint checks be executed in an efficient manner when updates are applied to the database?

#### Answer:

One possible technique that is often used to check efficiently for the key constraint is to create an index on the combination of attributes that form each key (primary or secondary). Before inserting a new record (tuple), each index is searched to check that no value currently exists in the index that matches the key value in the new record. If this is the case, the record is inserted successfully. For checking the foreign key constraint, an index on the primary key of each referenced relation will make this check relatively efficient. Whenever a new record is inserted in a referencing relation, its foreign key value is used to search the index for the primary key of the referenced relation, and if the referenced record exists, then the new record can be successfully inserted in the referencing relation. For deletion of a referenced record, it is useful to have an index on the foreign key of each referencing relation so as to be able to determine efficiently whether any records reference the record being deleted. If the indexes described above do not exist, and no alternative access structure (for example, hashing) is used in their place, then it is necessary to do linear searches to check for any of the above constraints, making the checks quite inefficient.

- **6.10** No solution provided.
- **6.11** Specify the updates of Exercise 3.11 using the SQL update commands.

# Answers:

Below, we show how each of the updates may be specified in SQL. Notice that some of these updates violate integrity constraints as discussed in the solution to Exercise 5.10, and hence should be rejected if executed on the database of Figure 5.6.

```
(a) Insert < 'Robert', 'F', 'Scott', '943775543', '21-JUN-42', '2365 Newcastle Rd, Bellaire, TX', M, 58000, '888665555', 1 > into EMPLOYEE.
INSERT INTO EMPLOYEE
VALUES ('Robert', 'F', 'Scott', '943775543', '21-JUN-42', '2365 Newcastle Rd, Bellaire, TX', M, 58000, '888665555', 1)
```

(b) Insert < 'ProductA', 4, 'Bellaire', 2 > into PROJECT. INSERT INTO PROJECT VALUES ('ProductA', 4, 'Bellaire', 2)

- (c) Insert < 'Production', 4, '943775543', '01-OCT-88' > into DEPARTMENT. INSERT INTO DEPARTMENT VALUES ('Production', 4, '943775543', '01-OCT-88')
- (d) Insert < '677678989', null, '40.0' > into WORKS\_ON. INSERT INTO WORKS\_ON VALUES ('677678989', NULL, '40.0')
- (e) Insert < '453453453', 'John', M, '12-DEC-60', 'SPOUSE' > into DEPENDENT. INSERT INTO DEPENDENT VALUES ('453453453', 'John', M, '12-DEC-60', 'SPOUSE')
- (f) Delete the WORKS\_ON tuples with ESSN= '333445555'. DELETE FROM WORKS\_ON WHERE ESSN= '333445555'
- (g) Delete the EMPLOYEE tuple with SSN= '987654321'. DELETE FROM EMPLOYEE WHERE SSN= '987654321'
- (h) Delete the PROJECT tuple with PNAME= 'ProductX'. DELETE FROM PROJECT WHERE PNAME= 'ProductX'
- (i) Modify the MGRSSN and MGRSTARTDATE of the DEPARTMENT tuple with DNUMBER=
  5 to '123456789' and '01-OCT-88', respectively.
  UPDATE DEPARTMENT
  SET MGRSSN = '123456789', MGRSTARTDATE = '01-OCT-88'
  WHERE DNUMBER= 5
- (j) Modify the SUPERSSN attribute of the EMPLOYEE tuple with SSN= '999887777' to '943775543'.

  UPDATE EMPLOYEE
  SET SUPERSSN = '943775543'

  WHERE SSN= '999887777'
- (k) Modify the HOURS attribute of the WORKS\_ON tuple with ESSN= '999887777' and PNO= 10 to '5.0'.

  UPDATE WORKS\_ON
  SET HOURS = '5.0'

  WHERE ESSN= '999887777' AND PNO= 10
- **6.12** Specify the following queries in SQL on the database schema of Figure 1.2.
- (a) Retrieve the names of all senior students majoring in 'COSC' (computer science).
- (b) Retrieve the names of all courses taught by professor King in 85 and 86.
- (c) For each section taught by professor King, retrieve the course number, semester, year, and number of students who took the section.
- (d) Retrieve the name and transcript of each senior student (Class=5) majoring in COSC. Transcript includes course name, course number, credit hours, semester, year, and grade for each course completed by the student.

- (e) Retrieve the names and major departments of all straight A students (students who have a grade of A in all their courses).
- (f) Retrieve the names and major departments of all students who do not have any grade of A in any of their courses.

#### Answers:

(a) SELECT Name FROM STUDENT WHERE Major='COSC'

(b) SELECT CourseName
FROM COURSE, SECTION
WHERE COURSE.CourseNumber=SECTION.CourseNumber AND Instructor='King'
AND (Year='85' OR Year='86')

Another possible SQL query uses nesting as follows:

SELECT CourseName FROM COURSE

WHERE CourseNumber IN ( SELECT CourseNumber

FROM SECTION

WHERE Instructor='King' AND (Year='85' OR Year='86'))

(c) SELECT CourseNumber, Semester, Year, COUNT(\*)

FROM SECTION, GRADE REPORT

WHERE Instructor='King' AND SECTION.SectionIdentifier=GRADE\_REPORT.SectionIdentifier GROUP BY CourseNumber, Semester, Year

- (d) SELECT Name, CourseName, C.CourseNumber, CreditHours, Semester, Year, Grade FROM STUDENT ST, COURSE C, SECTION S, GRADE\_REPORT G WHERE Class=5 AND Major='COSC' AND ST.StudentNumber=G.StudentNumber AND G.SectionIdentifier=S.SectionIdentifier AND S.CourseNumber=C.CourseNumber
- (e) SELECT Name, Major
  FROM STUDENT
  WHERE NOT EXISTS ( SELECT \*
  FROM GRADE\_REPORT
  WHERE StudentNumber= STUDENT.StudentNumber AND NOT(Grade='A'))
- (f) SELECT Name, Major
  FROM STUDENT
  WHERE NOT EXISTS ( SELECT \*
  FROM GRADE\_REPORT
  WHERE StudentNumber= STUDENT.StudentNumber AND Grade='A')
- **6.13** Write SQL update statements to do the following on the database schema shown in Figure 1.2.
- (a) Insert a new student <'Johnson', 25, 1, 'MATH'> in the database.
- (b) Change the class of student 'Smith' to 2.
- (c) Insert a new course <'Knowledge Engineering','COSC4390', 3,'COSC'>.
- (d) Delete the record for the student whose name is 'Smith' and student number is 17.

## Answers:

(a) INSERT INTO STUDENT VALUES ('Johnson', 25, 1, 'MATH')

(b) UPDATE STUDENT SET CLASS = 2 WHERE Name='Smith'

(c) INSERT INTO COURSE VALUES ('Knowledge Engineering', 'COSC4390', 3, 'COSC')

(d) DELETE FROM STUDENT WHERE Name='Smith' AND StudentNumber=17

## **6.14** – No solution provided

**6.15** - Consider the EMPLOYEE table's constraint EMPSUPERFK as specified in Figure 4.2 is changed to read as follows:

CONSTRAINT EMPSUPERFK

FOREIGN KEY (SUPERSSN) REFERNCES EMPLOYEE(SSN)

ON DELETE CASCADE ON UPDATE CASCADE.

Answer the following questions:

a. What happens when the following command is run on the database state shown in Figure 5.6?

DELETE EMPLOYEE WHERE LNAME = 'Borg'

b. Is it better to CASCADE or SET NULL in case of EMPSUPERFK constraint ON DELETE?

#### Answers:

- a) The James E. Borg entry is deleted from the table, and each employee with him as a supervisor is also (and their supervisees, and so on). In total, 8 rows are deleted and the table is empty.
- b) It is better to SET NULL, since an employee is not fired (DELETED) when their supervisor is deleted. Instead, their SUPERSSN should be SET NULL so that they can later get a new supervisor.
- **6.16** Write SQL statements to create a table EMPLOYEE\_BACKUP backup of EMPLOYEE table shown in Figure 3.6.

## Answer:

INSERT INTO EMPLOYEE\_BACKUP VALUES ( SELECT \* FROM EMPLOYEE )