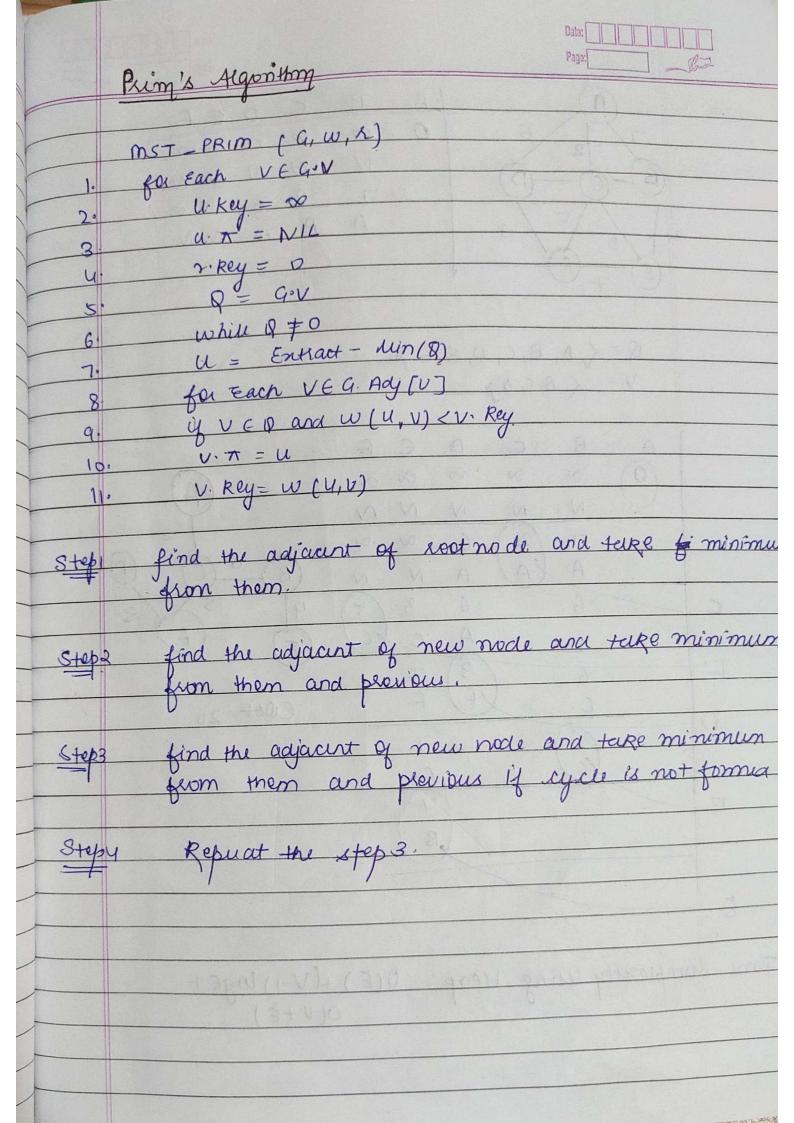
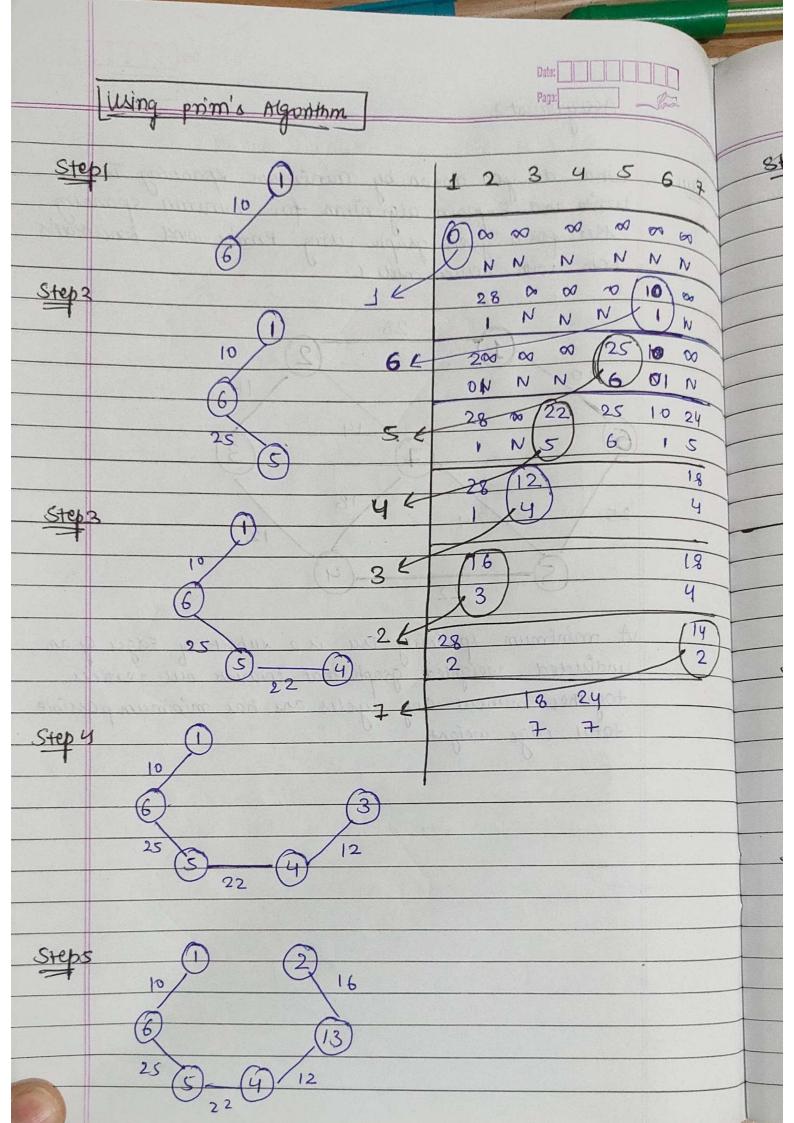
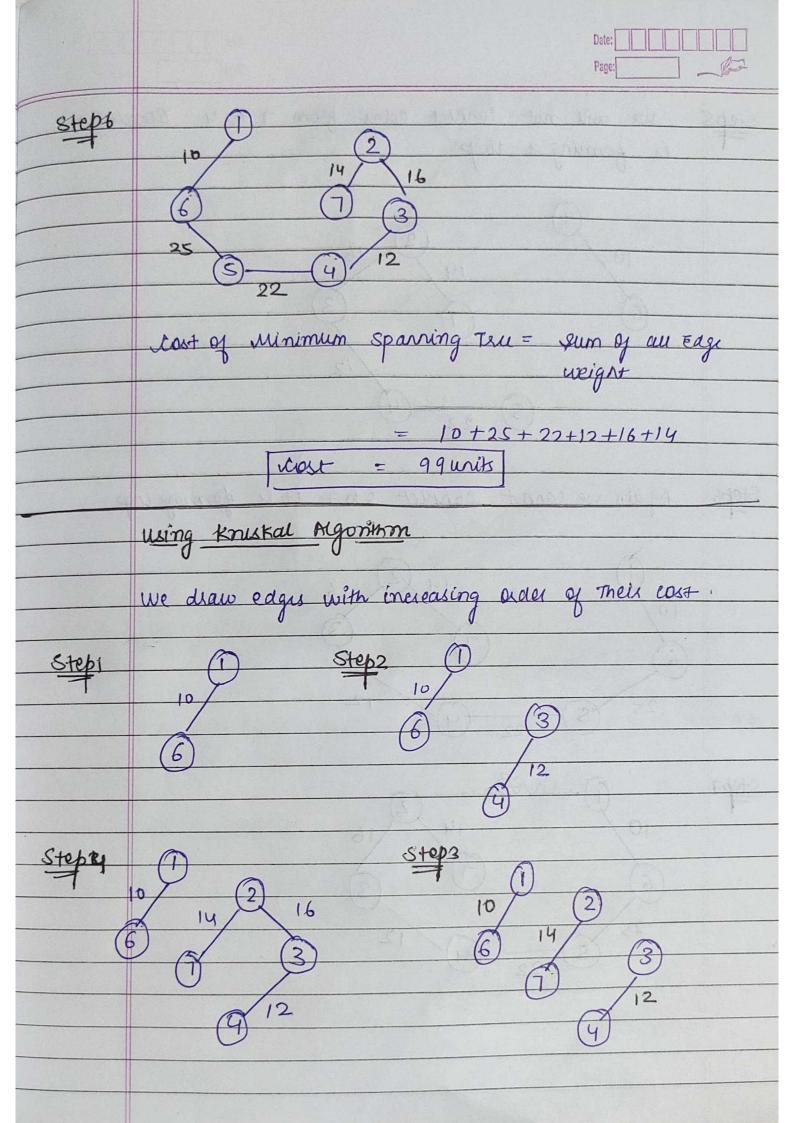
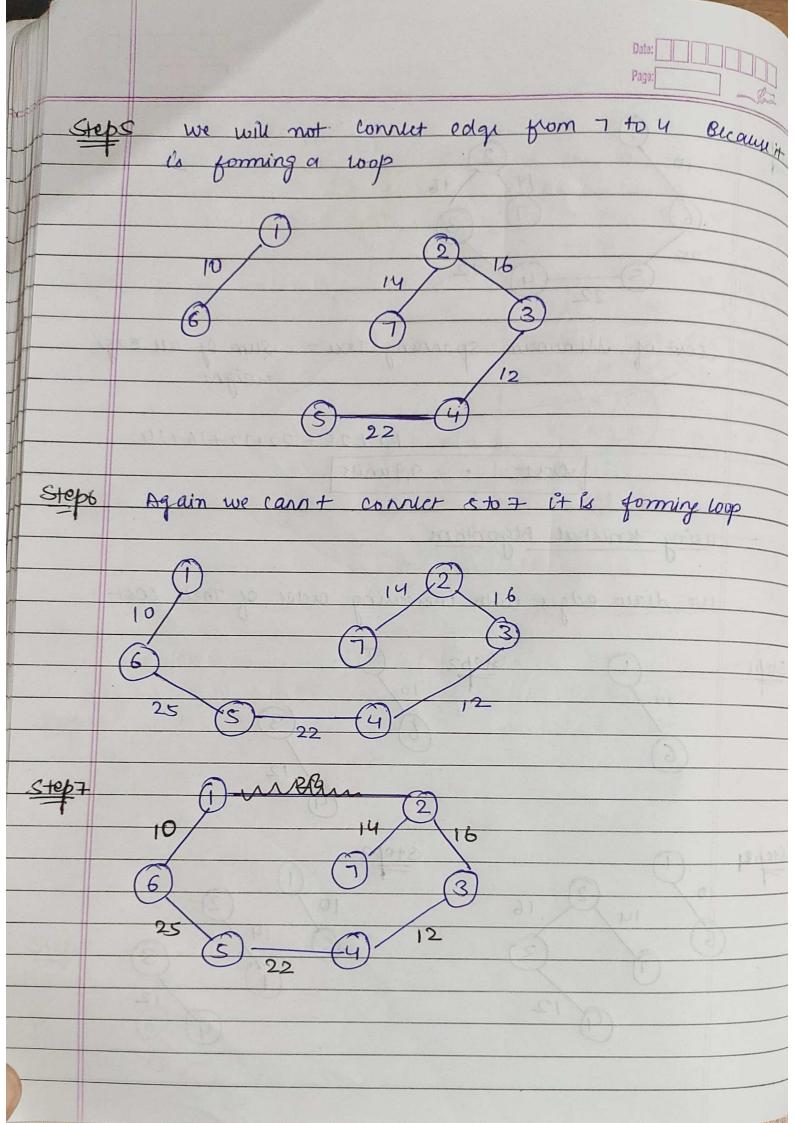


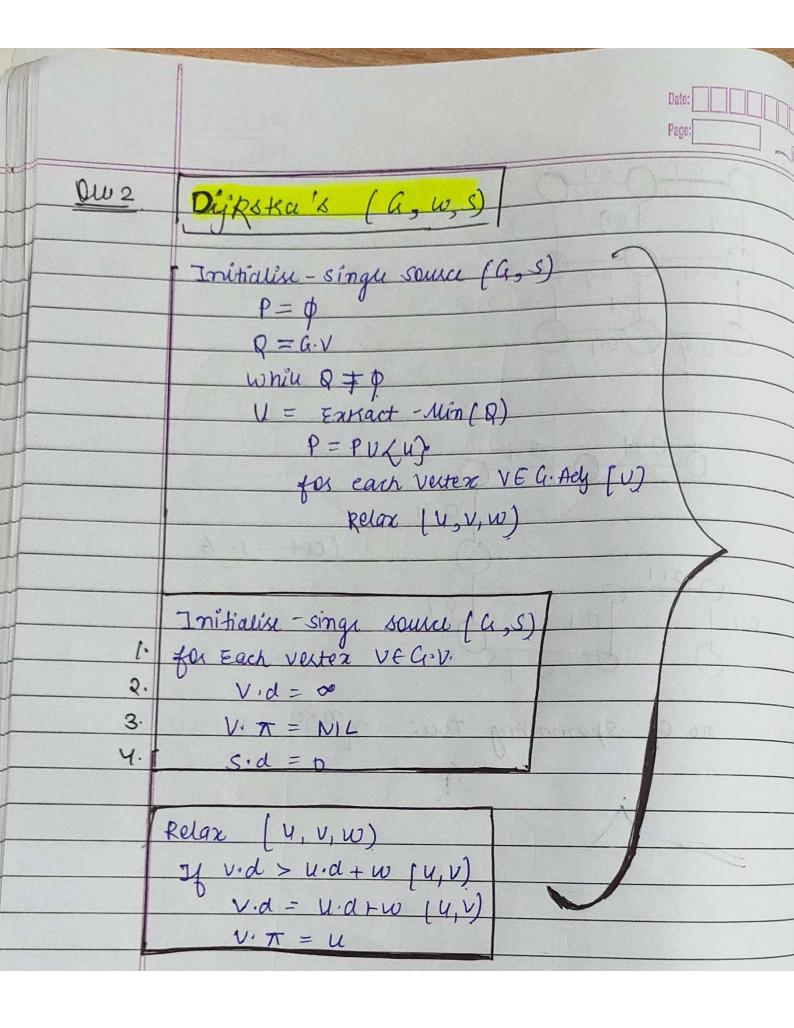
	Date:
	Kruskal's Algorithm
Stepi	Sort are the Edges of a graph in non-decreasing
-	Sort au the Edges of a graph in non-decreasing order by weight
step2	Take the minimum weight Edge and add 2 MST
TANK	Take the minimum weight Edge and add 2 MST [Minimum sparning Tsue).
Step3	Take the minimum weight Edge and add 2 ms 7
Step4	Take the next minimum weight Edge and and 2 MST & cycle is not formed
一	2 MST & cycle le not formed
Stops	Replat Step 4.
	Implimintation Algorithm
1.	MST_ KRVSKAL (G, W) Lg=graph
ļ·	MST_ KRVSKAL (G, w) $\mathcal{L}g = graph$ $A = \phi \qquad \qquad w = weightage$
1. 2. 3	MST_ KRVSKAL (G, w) $\mathcal{L}g = graph$ $A = \phi$ $w = weightage$ for each vertex $v \in G \cdot v$. $A = st$ Make- $st \in (v)$
1· 2· 3	MST KRVSKAL (G, w) $\mathcal{L}g = graph$ $A = \phi$ $w = weightage$ for each vertex $v \in G \cdot v$. $A = sit$ Make-sit (v) Sout the Edges of $G \cdot E$ into non-decreasing order
1· 2· 3	MST_ KRVSKAL (G, w) $\mathcal{L}g = graph$ $A = \phi$ $w = weightage$ for each vertex $v \in G \cdot v$. $A = sut$ Make-sut (v) Sort the Edges of $G \cdot E$ into non-decreasing order by weight w .
1· 2· 3 4	MST_ KRVSKAL (G, w) $\mathcal{L}_{g=graph}$ $A = \phi$ $w = weightage$ for each vertex $v \in G \cdot v$. $A = sut$ Make-sut (v) Sort the Edges of $G \cdot E$ into non-decreasing order by weight $w \cdot v$ for Each Edge $(u, v) \in G \cdot E$, in non-decreasing order
1· 2· 3 4	MST_ KRVSKAL (G, w) $\mathcal{L}_{g=graph}$ $A = \phi$ $w = weightage$ for each vertex $v \in G \cdot V \cdot A = sit$ Make-sit (V) Sort the Edges of $G \cdot E$ into non-decreasing order by weight $w \cdot V \cdot E \cdot E \cdot V \cdot E \cdot E \cdot E \cdot E \cdot E \cdot E$
1: 2: 3 4 5:	MST_ KRVSKAL [G, w)
	MST_ KRVSKAL (G, w) $\mathcal{L}_{g=graph}$ $A = \phi$ $w = weightage$ for each vertex $v \in G \cdot V \cdot A = sit$ Make-sit (V) Sort the Edges of $G \cdot E$ into non-decreasing order by weight $w \cdot V \cdot E \cdot E \cdot V \cdot E \cdot E \cdot E \cdot E \cdot E \cdot E$
7	MST_ KRVSKAL (G, w) $\mathcal{L}_{G} = graph$ $A = \phi$ $w = weightage$ for each vertex $v \in G \cdot v$. $A = sut$ Make-sut (v) Sort the Edges of $G \cdot E$ into non-decreasing order by weight w . for Each Edge $(u, v) \in G \cdot E$, in mon-decreasing order By weight $V = v \cdot v$
7.	MST_ KRVSKAL (G, W)
7.	MST_ KRVSKAL (G, W)

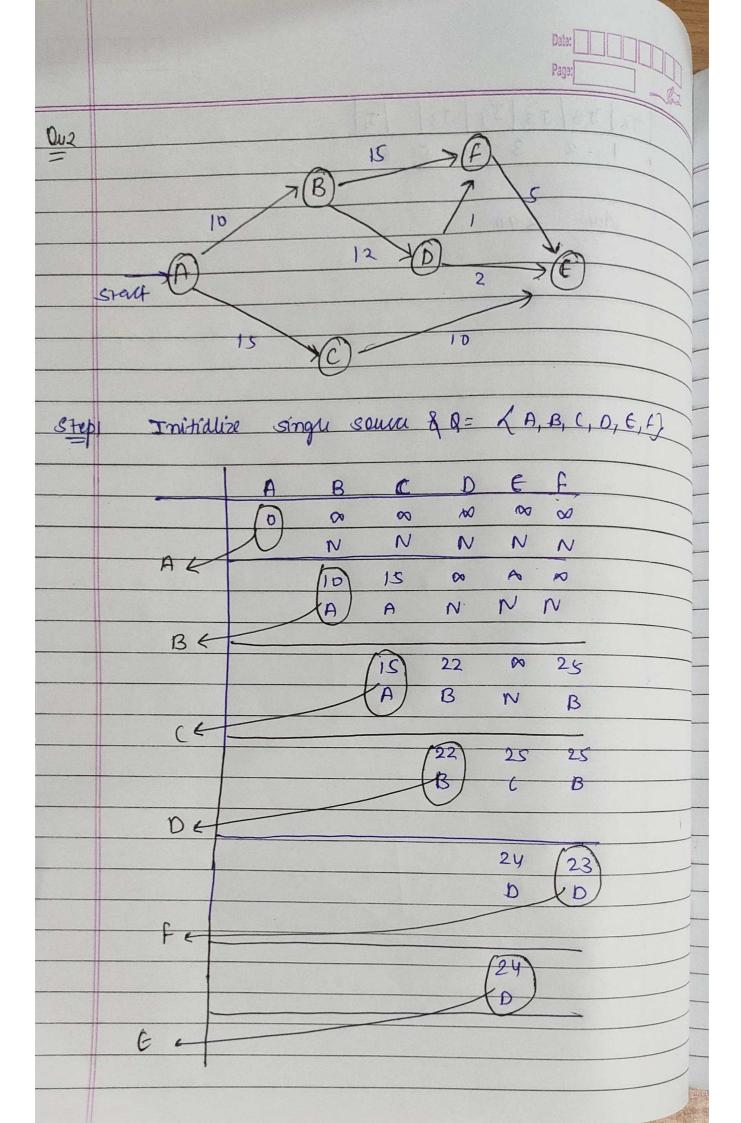












	Date:
	Page:
A)	$A \rightarrow B \rightarrow C \rightarrow D \rightarrow F \rightarrow E$
O	2
B)	24
C)	$A \rightarrow B \rightarrow D \rightarrow E$

- Quiz	Bulman-Josa Algorithm	Date: Page:
1· 2· 3· 4· 5· 6·	BF (G, W, S) Initialise - single source (G, S) for i=1 to G. V -1 for Each Edge (U, V) & G. & Relav U, V, W) for Each Edge [U, V) & G. & Y V. d > U. a + W[4, V)	
7. 8.	Relax $[4, v, v_0]$ $v \cdot d = u \cdot d + w(u, v)$	
	Initidin Single Source (G,S) for Each Vertex VEG.V. V.a = D S-a = 0 V-T = NIL	2

