

Discrete Mathematics

BCSC 0010

Module 2

Normal Forms

Normal forms

- If the number of variables , involved in a given statement are more, then the construction of truth tables may not be practical, therefore, we consider other method known as reduction to **normal form**.
- In this method, we use the word
- “**product**” in place of “**conjunction**”
- “**sum**” in place of “**disjunction**”

Some Basic terms Related to Normal Form

- **Elementary Product:** A product of the variables and their negations is called an elementary product.
Eg: Let p and q be any two variables, then p , q , $p \wedge q$, $\sim p \wedge \sim q$, $p \wedge \sim q$, $\sim p \wedge q$ etc. are elementary products.
- **Elementary Sum:** A Sum of the variables and their negations is called an elementary sum.
Eg: p , q , $p \vee q$, $p \vee \sim q$, $\sim p \vee \sim q$, $\sim p \vee q$ etc. are elementary sums.

Types of Normal Form

- Disjunctive Normal Form(DNF)
- Conjunctive Normal Form(CNF)

Disjunctive Normal Form(DNF)

- A statement which consists of a **sum of elementary products** of propositional variables and is equivalent to the given compound statement, is called a **disjunctive normal form** of the given statement.
- This form is **not unique** for the given statement.

Method to obtain disjunctive normal form

- **Step 1.** replace \rightarrow and \leftrightarrow by \wedge , \vee , and \sim in the statement.
- **Step 2.** Manipulate to get an equivalent form which is the **sum of elementary product terms**.

Example

- Obtain DNF of $(p \sqcup q)^{\wedge \sim q}$
- **Solution:**
- $(p \sqcup q)^{\wedge \sim q}$
- $(\sim p \vee q)^{\wedge \sim q}$ By logical equivalence
- $(\sim p^{\wedge \sim q}) \vee (q^{\wedge \sim q})$ By Distributive law
- This is the required DNF

Conjunctive Normal Form(CNF)

- A statement which consists of a **product of elementary sums** of propositional variables and is equivalent to the given compound statement, is called a **Conjunctive normal form** of the given statement.
- This form is **not unique** for the given statement.

Method to obtain conjunctive normal form

- **Step 1.** replace \rightarrow and \leftrightarrow by \wedge , \vee , and \sim in the statement.
- **Step 2.** Manipulate to get an equivalent form which is the **product of elementary sum** terms.

Example

- Obtain CNF of $(p \sqcup q) \wedge \sim q$
- **Solution:**
- $(p \sqcup q) \wedge \sim q$
- $(\sim p \vee q) \wedge \sim q$ By logical equivalence
- This is the required CNF

Example

• Obtain DNF of $p \vee (\sim p \wedge (q \vee (q \wedge \sim r)))$

• $p \vee (\sim p \wedge (q \vee (q \wedge \sim r)))$

• $p \vee (\sim p \wedge (q \vee (\sim q \vee \sim r)))$

By logical equivalence

• $p \vee (\sim p \wedge (q \vee \sim q) \vee \sim r)$

Rearranged

• $p \vee p \vee q \vee \sim q \vee \sim r$

By logical equivalence

This is the required DNF

Minterms

- Let p and q be two propositional variables.
- All possible formulas which consist of **product of p or its negation** and **product of q or its negation**, but should not contain both the variable and its negation in any one of the formula are called **minterms** of p and q .
- Eg: For two variables p and q , there are $2^2 = 4$ minterms, namely, $p \wedge q$, $p \wedge \sim q$, $\sim p \wedge q$, $\sim p \wedge \sim q$
- **If there are n variables then number of minterms = 2^n**

Maxterms

- Let p and q be two propositional variables.
- All possible formulas which consist of **sum of p or its negation** and **sum of q or its negation**, but should not contain both the variable and its negation in any one of the formula are called **maxterm** of p and q .
- Eg: For two variables p and q , there are $2^2 = 4$ maxterms, namely, $p \vee q$, $p \vee \sim q$, $\sim p \vee q$, $\sim p \vee \sim q$.
- **If there are n variables then number of maxterms = 2^n**

Principal disjunctive normal form(PDNF)

- For a given formula, an equivalent formula consisting of **disjunction of minterms** only is known as its Principal disjunctive normal form (PDNF) or Sum of products canonical form.
- Each minterm has **truth value T** for exactly one combination of truth values
- A formula which is **tautology** will have all **minterms**
- A formula which is **contradiction** will have no **minterms**

Method to find PDNF

- 1. Construct truth table for the given formula.
- 2. Identify the rows in which the formula has true – **as truth value.**
- 3. Construct **minterms** from each such rows by taking
 - (i) the variable with **true – as variable itself** and
 - (ii) the variable with **false – as negated variable.**
- 4. **Sum of these minterms will be the required PDNF.**

Example

p	q	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

p	q	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

PDNF for $p \square q$ will be $(p \wedge q) \vee (\neg p \wedge q) \vee (\neg p \wedge \neg q)$

PDNF for $p \leftrightarrow q$ will be $(p \wedge q) \vee (\neg p \wedge \neg q)$

Obtain PDNF of $(p \wedge q) \vee (\sim p \wedge r) \vee (q \wedge r)$

p	q	r	$\sim p$	$p \wedge q$	$\sim p \wedge r$	$q \wedge r$	$(p \wedge q) \vee (\sim p \wedge r)$	$(p \wedge q) \vee (\sim p \wedge r) \vee (q \wedge r)$
T	T	T	F	T	F	T	T	T
T	T	F	F	T	F	F	T	T
T	F	T	F	F	F	F	F	F
T	F	F	F	F	F	F	F	F
F	T	T	T	F	T	T	T	T
F	T	F	T	F	F	F	F	F
F	F	T	T	F	T	F	T	T
F	F	F	T	F	F	F	F	F

PDNF: $(p \wedge q \wedge r) \vee (p \wedge q \wedge \sim r) \vee (\sim p \wedge q \wedge r) \vee (\sim p \wedge \sim q \wedge r)$

Principal conjunctive normal form(PCNF)

- For a given formula, an equivalent formula consisting of **conjunction of maxterms** only is known as its Principal disjunctive normal form(PCNF) or product of sums canonical form.
- Each maxterm has **truth value F** for exactly one combination of truth values
- A formula which is **tautology** will have no **maxterms**
- A formula which is **contradiction** will have all **maxterms**

Method to find PCNF

- 1. Construct truth table for the given formula.
- 2. Identify the rows in which the formula has **false – as truth value.**
- 3. Construct **maxterm** from each such row by taking
 - (i) the variable with **true -as negated variable** and
 - (ii) the variable with **false- as the variable itself.**
- **4. Product of these maxterms will be the required PCNF.**

Example

p	q	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

p	q	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

PCNF for $p \rightarrow q$ will be $(\neg p \vee q)$

PCNF for $p \leftrightarrow q$ will be $(\neg p \vee q) \wedge (p \vee \neg q)$

Obtain PCNF of $(p \wedge q) \vee (\sim p \wedge r) \vee (q \wedge r)$

p	q	r	$\sim p$	$p \wedge q$	$\sim p \wedge r$	$q \wedge r$	$(p \wedge q) \vee (\sim p \wedge r)$	$(p \wedge q) \vee (\sim p \wedge r) \vee (q \wedge r)$
T	T	T	F	T	F	T	T	T
T	T	F	F	T	F	F	T	T
T	F	T	F	F	F	F	F	F
T	F	F	F	F	F	F	F	F
F	T	T	T	F	T	T	T	T
F	T	F	T	F	F	F	F	F
F	F	T	T	F	T	F	T	T
F	F	F	T	F	F	F	F	F

PCNF: $(\sim p \vee q \vee \sim r) \wedge (\sim p \vee q \vee r) \wedge (p \vee \sim q \vee r) \wedge (p \vee q \vee r)$

Example

- Obtain PCNF of $(p \vee q) \wedge \sim p \square \sim q$
- Obtain CNF of $\sim(p \vee q) \leftrightarrow (p \wedge q)$
- Obtain DNF of $\sim(p \vee q) \leftrightarrow (p \wedge q)$

Next Topic

- **Predicate and Quantifiers**