

## Gauss Elimination

② 2 steps in Gauss elimination

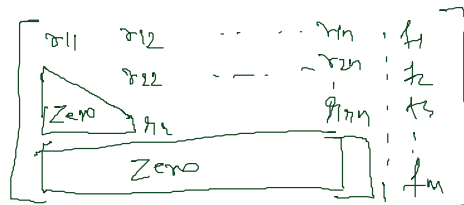
- ① Forward elimination
- ② Backward Substitution

→ Reduce  $AX = b$  to upper triangular matrix (or REF)  
 $\therefore TX = b'$  } Reduced

→ solve  $TX = b'$  for  $x$  Same  $x$  is the solution for  $A$  { given that it is unique }

Before the back substitution the augmented matrix (Reduced) will look like this:

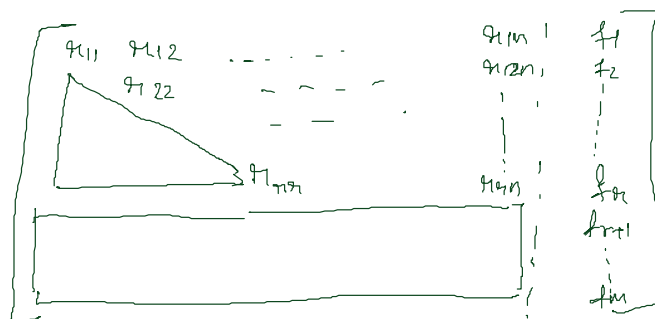
- Reduced matrix = upper triangular
- first  $r$  - non zero rows
- R.H.S may have last  $r$  rows = zero



$m$  rows  
 $n$  columns

up till  $r$   
 we have non zero rows

$r = \text{rank of matrix}$



- ★ if any of the last  $r$  rows are Non zero: then
- ★ if  $m = \text{no of rows}$  then Complexity =  $O(m^3)$

System of Equations is inconsistent

① Inconsistency

Eg:

Reduced matrix  $\rightarrow \begin{bmatrix} 1 & 1 & 1 & 3 \\ 0 & -2 & 1 & -1 \\ 0 & 0 & 0 & -4 \end{bmatrix}$

$\therefore r < m$   
 $r = \text{no of Equations}$   
 $m = \text{no of rows}$

$r = 2$   
 $m = 3$

$\therefore Rx = f$  is inconsistent

$\therefore Ax = B$  is inconsistent and  $\Rightarrow$  No Solution

② Consistency with  $f_r = 0$

## ② Consistency with $fr = 0$

Consider an example of Reduced matrix

Here also  $r < m$ ;  $r = 2, m = 3$

but  $fr = 0$   $\therefore$  (underdetermined)

System is Consistent

and have infinite many solutions.

$$\begin{bmatrix} 1 & 1 & 1 & 3 \\ 0 & -2 & 1 & -1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \leftarrow fr = 0$$

In order to solve this:

Assume  $x_3 = c$

Backward Substitution:

$$\begin{aligned} x_1 + x_2 + x_3 &= 3 \Rightarrow x_1 + x_2 = 3 - c \\ -2x_2 + x_3 &= -1 \Rightarrow -2x_2 = -1 - c \\ x_2 &= \frac{1+c}{2} \end{aligned}$$

$$\Rightarrow x_1 + \left(\frac{1+c}{2}\right) = 3 - c$$

$$x_1 = 3 - c - \left(\frac{1+c}{2}\right)$$

$$\Rightarrow \frac{6 - 2c - 1 - c}{2}$$

$$x_1 \Rightarrow \frac{5 - 3c}{2}$$

$$\therefore x = \begin{bmatrix} c \\ (1+c)/2 \\ (5-3c)/2 \end{bmatrix}$$

$\therefore$  Consistent = at least 1 solution  
Inconsistent = No solution

Overdetermined  
determined  
underdetermined

$\begin{cases} \text{if } m > n \leftarrow \\ \text{if } m = n \leftarrow \text{unique solution} \\ \text{if } m < n \leftarrow \text{infinite solution} \end{cases}$

$\begin{cases} \text{if } fr = 0 \rightarrow \text{infinite solution} \\ \text{rank} = \text{rank (columns)} \rightarrow \text{unique solution} \end{cases}$

## Pitfalls of Gauss elimination method

① Division by zero - when pivot element is 0.  
even if we exchange rows there may

② Rounding off error -

Computers carry limited significant figures so it might round off for decimal points. Resulting in inaccurate answer.

be case where we

few divide by zero.

[Be Careful]

$$\begin{bmatrix} 0 & 2 & 8 \\ 4 & 6 & 7 \\ 2 & 1 & 6 \end{bmatrix}$$

solution: Pivoting

\* Pivoting: Replacing the 0 pivot element row with the row with large number

\* Scaling - used to reduce the round off error accuracy

\* Partial pivoting: Search largest element then replace it with pivot element

\* Scaling - Use the round off error and improve accuracy

\* Partial pivoting: Search largest element then replace it with pivot element.

\* Use 5 significant figure to reduce rounding off error

\* Counting no of operations: To judge the quality of numerical method:

- Space complexity
- Time complexity

Eg, Consider  $2 \times 2$  matrix

$$\begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \end{bmatrix}$$

Augmented  $[A|b]$

$$\Rightarrow \begin{bmatrix} a_{11} & a_{12} & b_1 \\ 0 & 0 & b_2 \end{bmatrix}$$

Consider worst case  
No zero value initially

$$R_2 \leftarrow R_2 - \frac{a_{21}}{a_{11}} R_1$$

↑  
div

first element of  $R_2 \rightarrow 0$

5 multiply  
5 add  
no  $a_{21}$  is 0  
 $0 \cdot 4 = 0$   
 $4 \cdot 4 = 16$

\* Doing for all rows:

$$R_2 \leftarrow R_2 - \frac{a_{21}}{a_{11}} R_1$$

$$R_3 \leftarrow R_3 - \frac{a_{31}}{a_{11}} R_1$$

total 3 division

4 multiply  
4 add

$\therefore n-2$  : divisions

$(n-1)(n-2)$  = multiplications  
 $(n-1)(n-2)$  = additions

total no of ops required is

$$2 \sum_{k=1}^{n-1} (n-k)(n-k+1) = O(n^3)$$

Total no of divisions =  $O(n^2)$

$$= \sum_{k=1}^{n-1} (n-k) = O(n^2)$$

In back substitution:

$$\text{No of operations: } \left[ 2 \sum_{k=1}^n (n-k) \right] + n = O(n^2)$$

Consider:  $n = 1000$ ,  $n = 100000$