

18.404/6.840 Lecture 19

Last time:

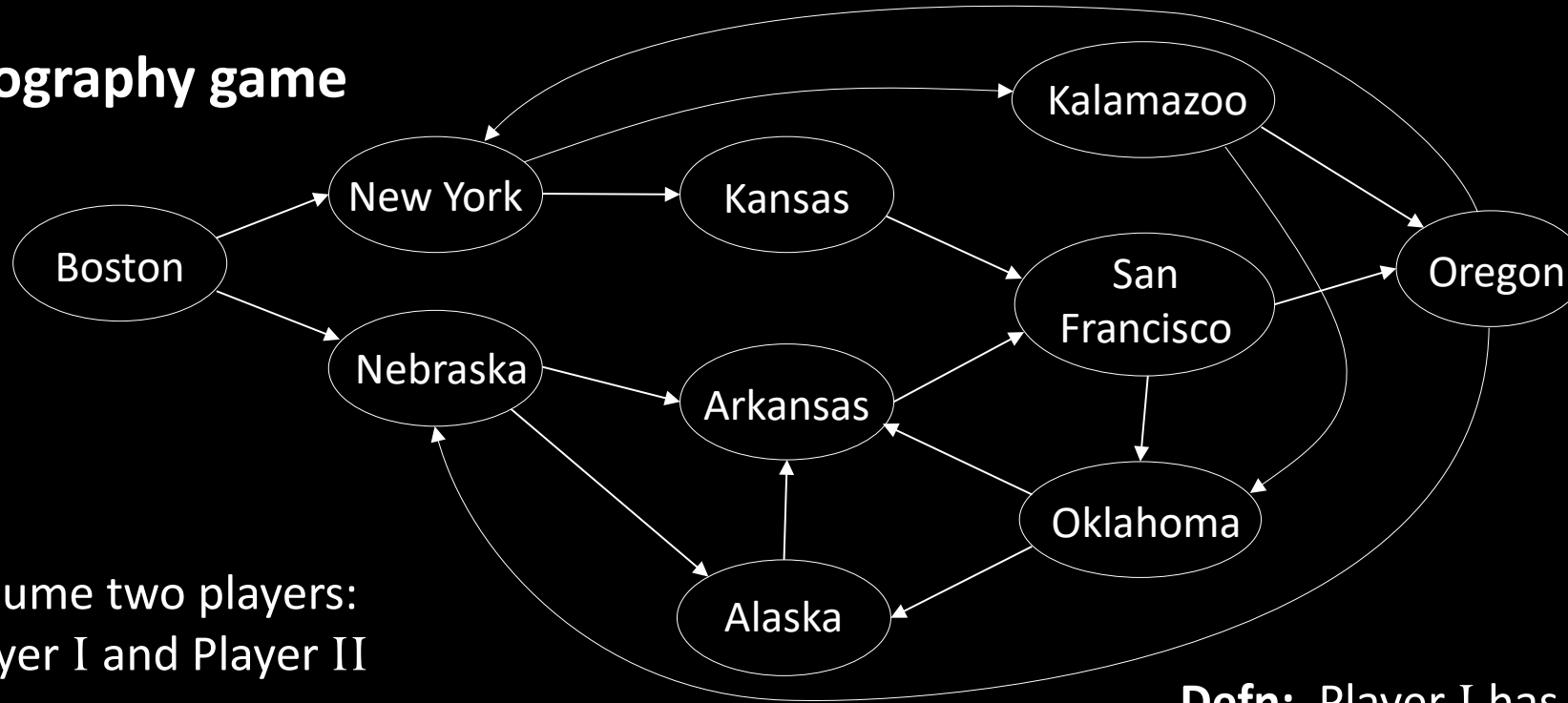
- Review $PSPACE$
- Savitch's Theorem: $NSPACESPACE$
- is $PSPACE$ -complete

Today:

- Games and Quantifiers
- The Formula Game
- Generalized Geography is $PSPACE$ -complete
- Logspace: L and NL

Games and Complexity

Geography game



Assume two players:
Player I and Player II

Players take turns picking places that start with the letter which ended the previous place. No repeats allowed. The first player stuck (= cannot move) loses.

Generalized Geography Game

Played on any directed graph.
Players take turns picking nodes that form a simple path.
The first player stuck loses.

Defn: Player I has a forced win in Generalized Geography on graph starting at node .

“forced win” also called a “winning strategy” means that the player will win if both players play optimally.

Theorem: is PSPACE-complete

Games and Quantifiers

The Formula Game

Given QBF

There are two Players \exists and \forall .

Player \exists assigns values to the \exists -quantified variables.

Player \forall assigns values to the \forall -quantified variables.

The players choose the values according to the order of the quantifiers in ϕ .

After all variables have been assigned values, we determine the winner:

Player \exists wins if the assignment satisfies ϕ .

Player \forall wins if not.

Claim: Player \exists has a forced win in the formula game on ϕ iff ϕ is TRUE.

Therefore Player \exists has a forced win on ϕ .

Next: show.

is PSPACE-complete

Theorem: is PSPACE-complete

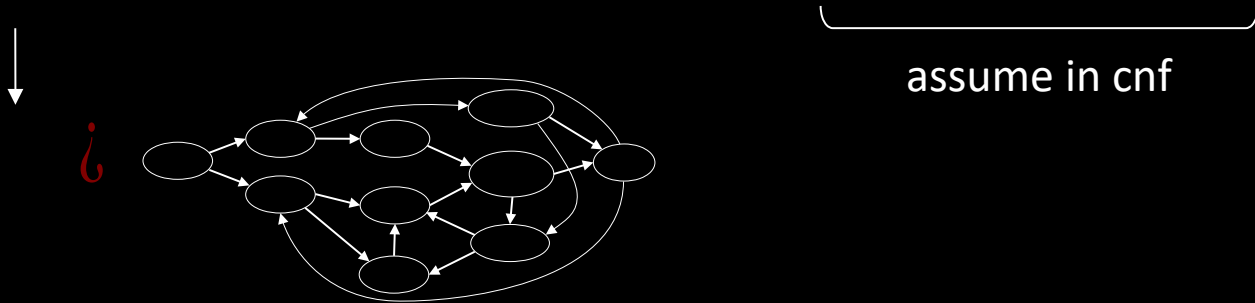
Proof: 1) PSPACE (recursive algorithm, exercise)
2)

Give reduction from to .

Construct to mimic the formula game on .

has Players I and II

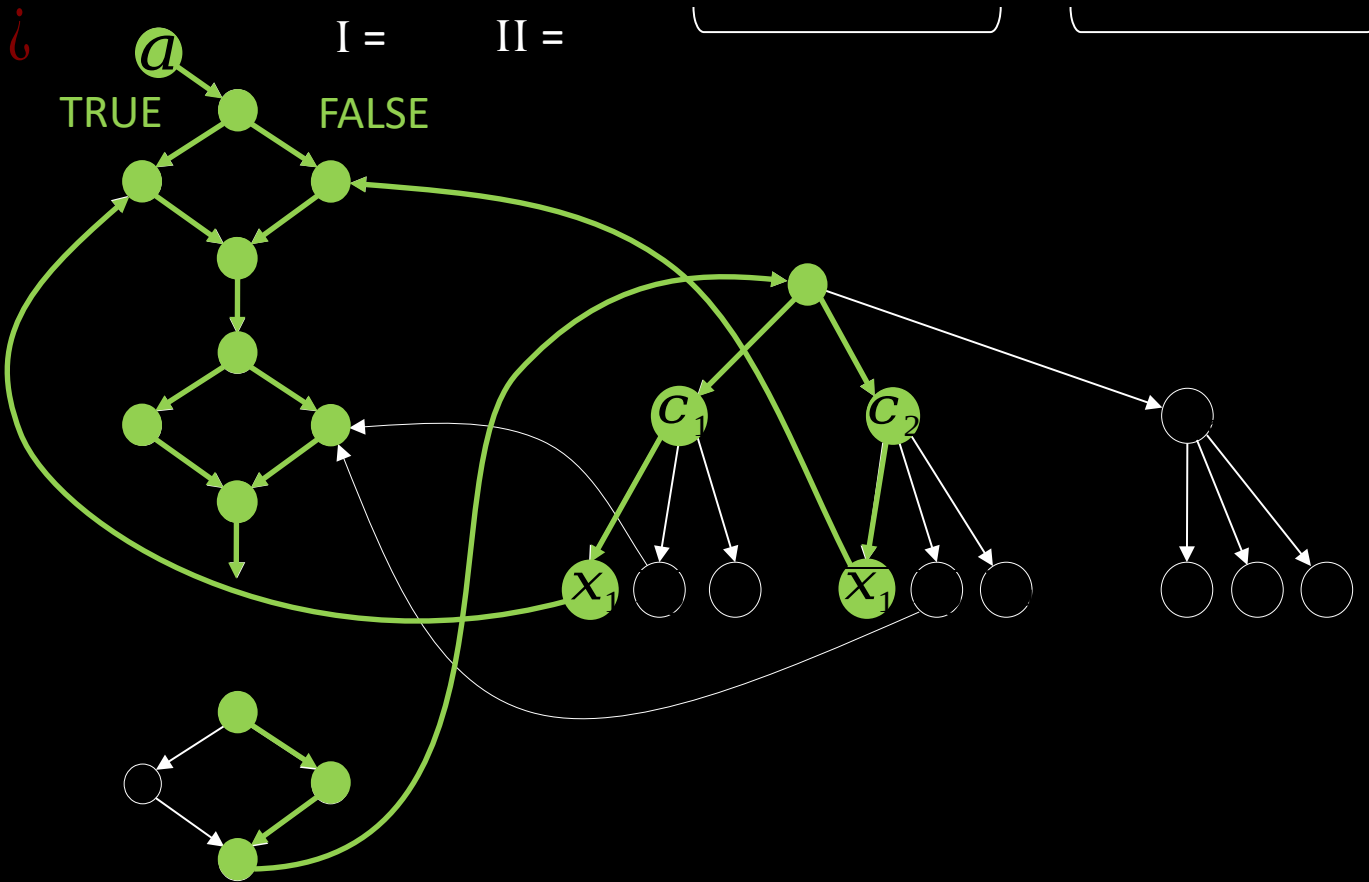
Player I plays role of \neg -Player in . Ditto for Player II and the \neg -Player.



Constructing the graph

Illustrate construction by example

Say



Endgame

should win if assignment satisfied all clauses
should win if some unsatisfied clause

Implementation

picks clause node claimed unsatisfied
picks literal node claimed to satisfy the clause
liar will be stuck



Log space

To define sublinear space computation, do not count input as part of space used.
Use 2-tape TM model with read-only input tape.

Defn: $L = \text{SPACE}$

$NL = \text{NSPACE}$

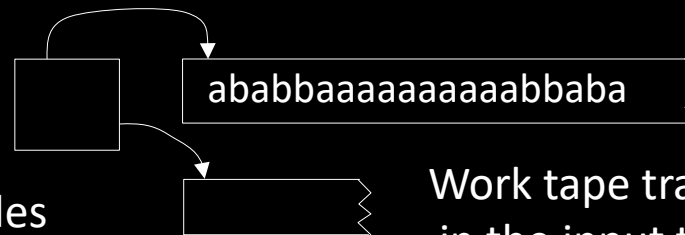
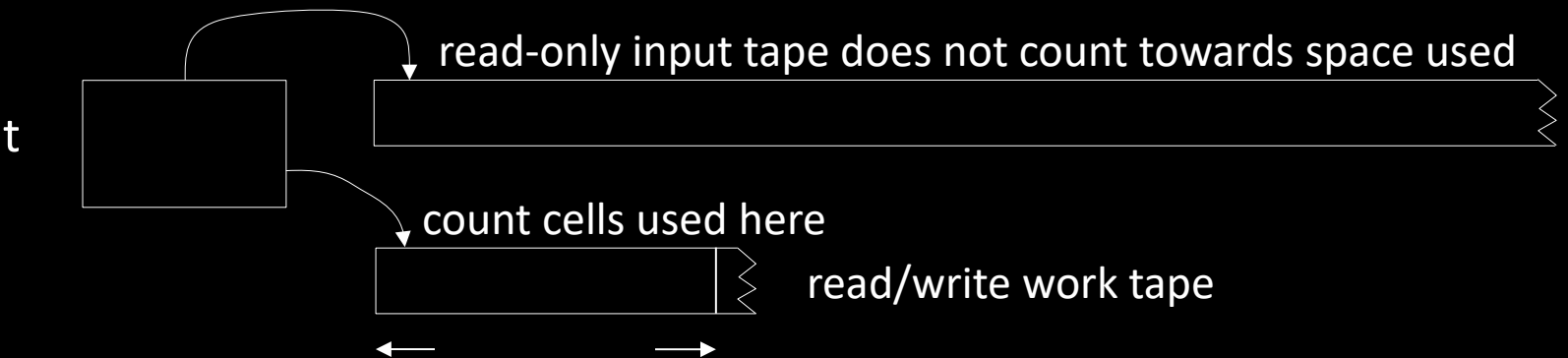
Log space can represent a constant number of pointers into the input.

Examples

1. L

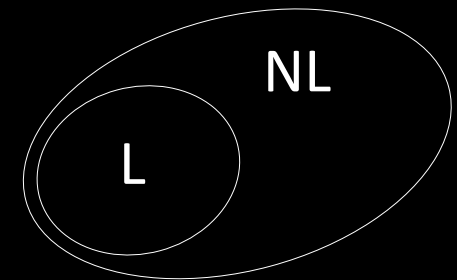
2. NL

Nondeterministically select the nodes
of a path connecting to .



Work tape tracks corresponding locations
in the input tape.

$L = NL?$ Unsolved



Log space properties

Theorem: $L \subseteq P$

Proof: Say M decides L in space $f(n)$.

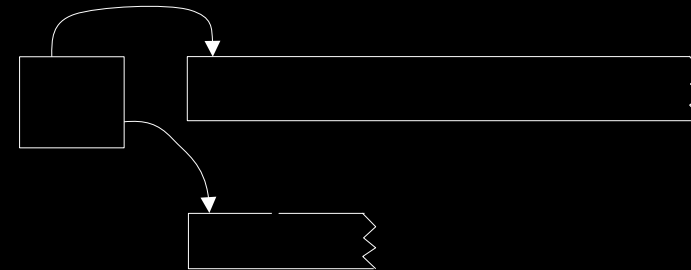
Defn: A configuration for M on x is where q is a state, i and j are the tape head positions, and γ is the tape contents. The number of such configurations is $2^{O(f(n))}$ for some c .

Therefore M runs in polynomial time.

Conclusion: $L \subseteq P$

Theorem: $NL = coNL$

Proof: Savitch's theorem works for log space



NL properties

Theorem: $NL = P$

Proof: Say NTM decides in space $f(n)$.

Defn: The configuration graph for M on x has

nodes: all configurations for M on x

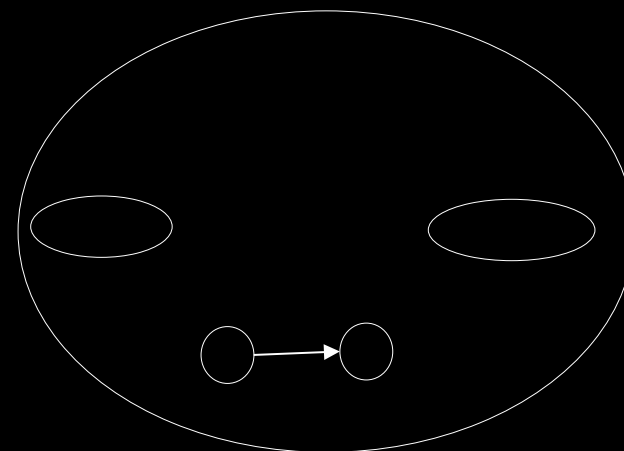
edges: edge from C_1 to C_2 if C_1 can yield C_2 in 1 step.

Claim: M accepts x iff the configuration graph has a path from C_{start} to C_{accept} .

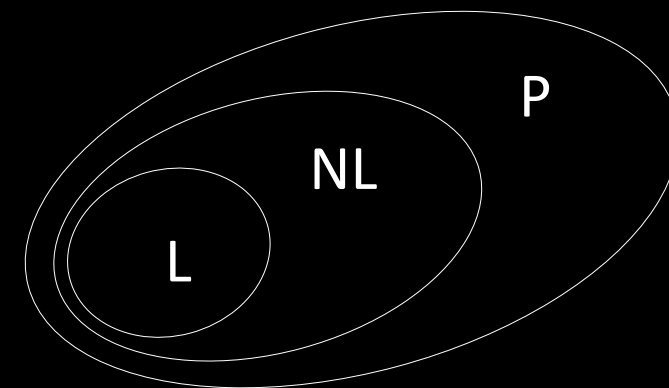
Polynomial time algorithm for $NL = P$:

“On input x ”

1. Construct the configuration graph G_x .
2. *Accept* if there is a path from C_{start} to C_{accept} .
Reject if not.”



$L = P$? Unsolved



Quick review of today

1. The Formula Game
2. Generalized Geography is PSPACE-complete
3. Log space: L and NL
4. Configuration graph
5. NL \subseteq P