18.404/6.840 Intro to the Theory of Computation

Instructor: Mike Sipser

Office Hours 4:00 – 5:30 Tuesdays

TAs: Office Hours TBD

- Fadi Atieh, Damian Barabonkov,
- Alex Dimitrakakis, Thomas Xiong,
- Abbas Zeitoun, and Emily Liu

Recitations start Friday

- Optional unless you need them!
- Hourly 10-2pm, online. On Sept 11, noon and 2pm → in-person

Homework, Exams, Quizzes

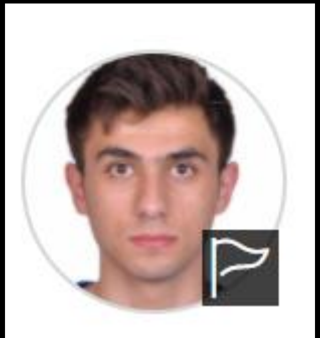
- See Course Information on homepage math.mit.edu/18.404
- First Pset due Sept 10. Posted on homepage

Our TAs

Alex

Thomas

Emily



Fadi

Abbas

Damian

18.404 Course Outline

Computability Theory 1930s – 1950s

- What is computable… or not?
- Examples:
  program verification, mathematical truth
- Models of Computation:
  Finite automata, Turing machines, …

Complexity Theory 1960s – present

- What is computable <u>in practice</u>?
- Example: factoring problem
- P versus NP problem
- Measures of complexity: Time and Space
- Models: Probabilistic and Interactive computation

Course Mechanics

Zoom Lectures

- Live and Interactive via Chat
- Live lectures are recorded for later viewing

Zoom Recitations starting this Friday

- Not recorded; notes will be posted
- Two convert to in-person on Sept 11
- Review concepts and more examples
- Optional unless you are having difficulty
  Participation can raise low grades
- Attend any recitation

Homework bi-weekly – 35%

- More information to follow

Midterm (15%) and Final exam (25%)

- Open book and notes

Check-in quizzes for credit – 25%

- Distinct Live and Recorded versions
- Complete either one for credit within 48 hours
- Initially ungraded; full credit for participation

Check-in 1.1

Course Expectations

Prerequisites

Prior substantial experience and comfort with mathematical concepts, theorems, and proofs. <u>Creativity will be needed for psets and exams.</u>

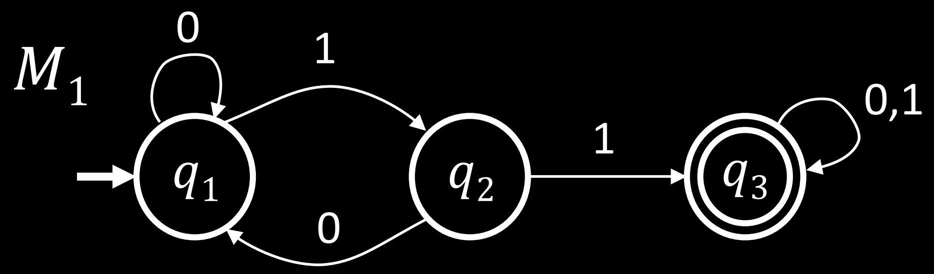
Collaboration policy on homework

- Allowed. But try problems yourself first.
- <u>Write up your own solutions.</u>
- <u>No bibles or online materials.</u>

Role of Theory in Computer Science

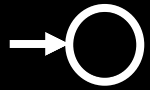
1. **Applications**
2. **Basic Research**
3. **Connections to other fields**
4. **What is the nature of computation?**

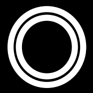
Let's begin: Finite Automata



States: $q_1\ q_2\ q_3$

Transitions: $\xrightarrow{1}$

Start state: 

Accept states: 

Input: finite string
Output: <u>Accept</u> or <u>Reject</u>

Computation process: Begin at start state, read input symbols, follow corresponding transitions, <u>Accept</u> if end with accept state, <u>Reject</u> if not.

Examples: 01101 → Accept
00101 → Reject

M_1 accepts exactly those strings in A where
$A = \{w \mid w$ contains substring $11\}$.

Say that A is the language of M_1 and that M_1 recognizes A and that $A = L(M_1)$.

Check-in 1.2

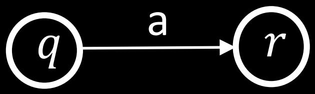
Finite Automata – Formal Definition

Defn: A <u>finite automaton</u> M is a 5-tuple $(Q, \Sigma, \delta, q_0, F)$

Q finite set of states

Σ finite set of alphabet symbols

δ transition function $\delta\colon Q \times \Sigma \to Q$

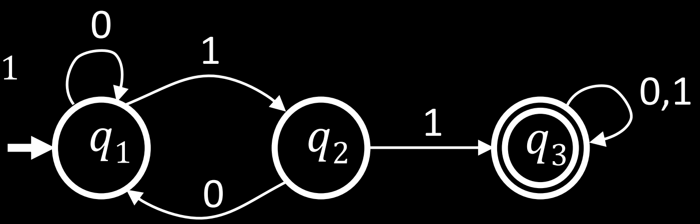
$\delta(q, a) = r$ means 

q_0 start state

F set of accept states

Example:

M_1



$M_1 = (Q, \Sigma, \delta, q_1, F)$

$Q = \{q_1, q_2, q_3\}$

$\Sigma = \{0, 1\}$

$F = \{q_3\}$

$\delta =$

	0	1
q_1	q_1	q_2
q_2	q_1	q_3
q_3	q_3	q_3

Finite Automata – Computation

Strings and languages

- A <u>string</u> is a finite sequence of symbols in Σ
- A <u>language</u> is a set of strings (finite or infinite)
- The <u>empty string</u> ε is the string of length 0
- The <u>empty language</u> ø is the set with no strings

Defn: M <u>accepts string</u> $w = w_1 w_2 \ldots w_n$ each $w_i \epsilon \Sigma$
if there is a sequence of states $r_0, r_1, r_2, , \ldots, r_n \ \epsilon \ Q$
where:

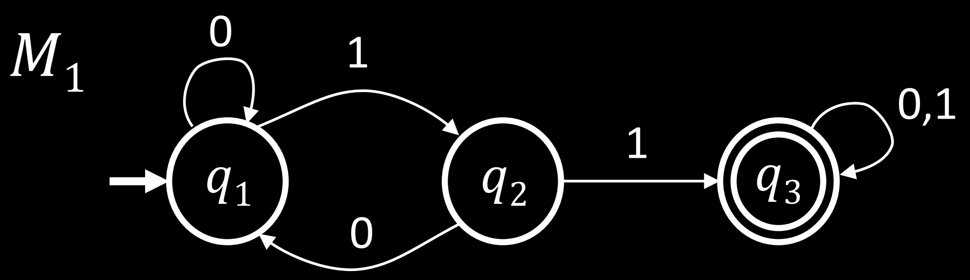
- $r_0 = q_0$
- $r_i = \delta(r_{i-1}, w_i)$ for $1 \le i \le n$
- $r_n \epsilon F$

Recognizing languages

- $L(M) = \{w | \, M \text{ accepts } w\}$
- $L(M)$ is <u>the language of</u> M
- M <u>recognizes</u> $L(M)$

Defn: A language is <u>regular</u> if some finite automaton recognizes it.

Regular Languages – Examples



$L(M_1) = \{w \mid w \text{ contains substring } 11\} = A$

Therefore A is regular

More examples:

Let $B = \{w \mid w$ has an even number of 1s$\}$
B is regular (make automaton for practice).

Let $C = \{w \mid w$ has equal numbers of 0s and 1s$\}$
C is not regular (we will prove).

Goal: Understand the regular languages

Regular Expressions

Regular operations. Let A, B be languages:

- Union: $A \cup B = \{w \mid w \in A \text{ or } w \in B\}$
- Concatenation: $A \circ B = \{xy \mid x \in A \text{ and } y \in B\} = AB$
- Star: $A^* = \{x_1 \dots x_k \mid \text{each } x_i \in A \text{ for } k \geq 0\}$

  Note: $\varepsilon \in A^*$ always

Example. Let $A = \{\text{good, bad}\}$ and $B = \{\text{boy, girl}\}$.

- $A \cup B = \{\text{good, bad, boy, girl}\}$
- $A \circ B = AB = \{\text{goodboy, goodgirl, badboy, badgirl}\}$
- $A^* = \{\varepsilon, \text{good, bad, goodgood, goodbad, badgood, badbad, goodgoodgood, goodgoodbad}, \dots\}$

Regular expressions

- Built from Σ, members $\Sigma, \emptyset, \varepsilon$ [Atomic]
- By using $\cup, \circ, *$ [Composite]

Examples:

- $(0 \cup 1)^* = \Sigma^*$ gives all strings over Σ
- $\Sigma^* 1$ gives all strings that end with 1
- $\Sigma^* 11 \Sigma^* =$ all strings that contain $11 = L(M_1)$

Goal: Show finite automata equivalent to regular expressions

Closure Properties for Regular Languages

Theorem: If A_1, A_2 are regular languages, so is $A_1 \cup A_2$ (closure under ∪)

Proof: Let $M_1 = (Q_1, \Sigma, \delta_1, q_1, F_1)$ recognize A_1
$M_2 = (Q_2, \Sigma, \delta_2, q_2, F_2)$ recognize A_2

Construct $M = (Q, \Sigma, \delta, q_0, F)$ recognizing $A_1 \cup A_2$

M should accept input w if either M_1 or M_2 accept w.

Mini-quiz 3

In the proof, if M_1 and M_2 are finite automata where M_1 has k_1 states and M_2 has k_2 states Then how many states does M have?

(a) $k_1 + k_2$
(b) $(k_1)^2 + (k_2)^2$
(c) $k_1 \times k_2$

Components of M:

$Q = Q_1 \times Q_2$
$= \{(q_1, q_2) | q_1 \in Q_1 \text{ and } q_2 \in Q_2\}$

$q_0 = (q_1, q_2)$

$\delta((q, r), a) = (\delta_1(q, a), \delta_2(r, a))$

$F = \cancel{F_1 \times F_2}$ NO! [gives intersection]

$F = (F_1 \times Q_2) \cup (Q_1 \times F_2)$

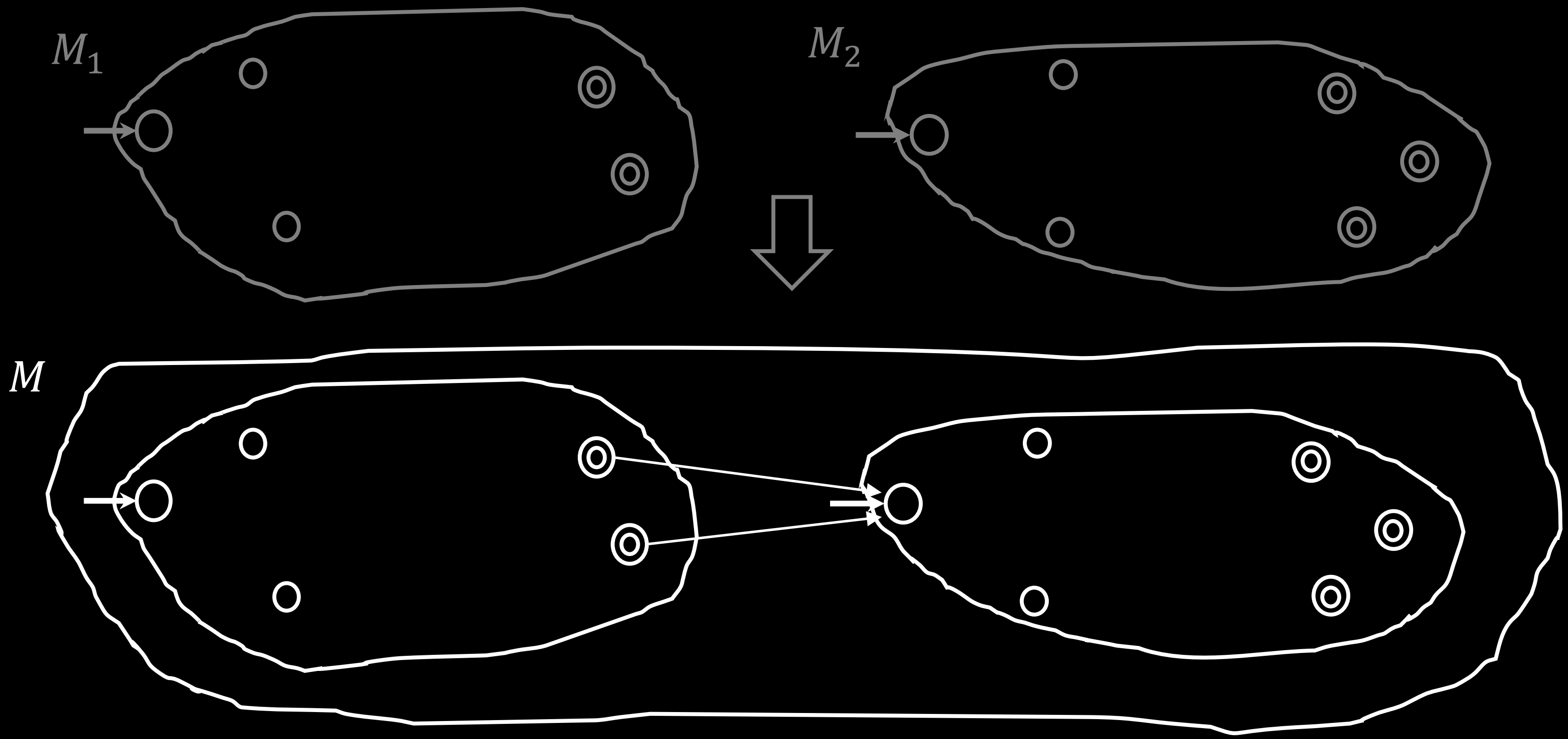
Check-in 1.3

Closure Properties continued

Theorem: If A_1, A_2 are regular languages, so is A_1A_2 (closure under $\circ$)

Proof: Let $M_1 = (Q_1, \Sigma, \delta_1, q_1, F_1)$ recognize A_1
$M_2 = (Q_2, \Sigma, \delta_2, q_2, F_2)$ recognize A_2

Construct $M = (Q, \Sigma, \delta, q_0, F)$ recognizing A_1A_2



M should accept input w
if $w = xy$ where
M_1 accepts x and M_2 accepts y.

Doesn't work: Where to split w?

Quick review of today

1. Introduction, outline, mechanics, expectations
2. Finite Automata, formal definition, regular languages
3. Regular Operations and Regular Expressions
4. Proved: Class of regular languages is closed under $\cup$
5. Started: Closure under $\circ$, to be continued…

Check-in 1.4