

18.404/6.840 Lecture 2

Last time:

- Finite automata, regular languages
- Regular operations
- Regular expressions
- Closure under

Today:

- Nondeterminism
- Closure under and
- Regular expressions → finite automata

Goal: Show finite automata equivalent to regular expressions

- This week's check-ins will not be counted
- TA office hours will be posted tomorrow
- Chat is restricted to TAs only.

Problem Sets

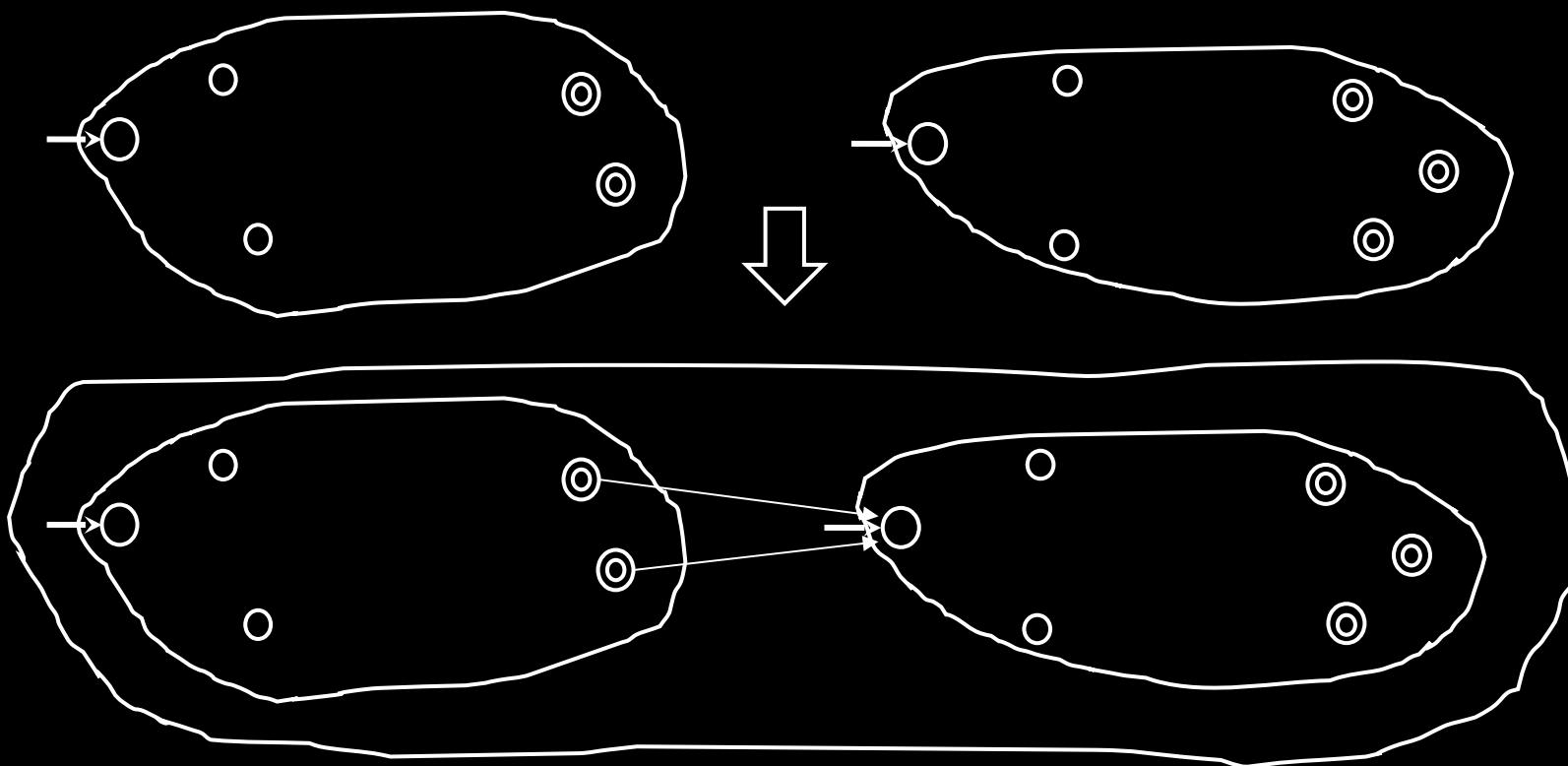
- 35% of overall grade
- Problems are hard! Leave time to think about them.
- Writeups need to be clear and understandable, handwritten ok.
Level of detail in proofs comparable to lecture: focus on main ideas.
Don't need to include minor details.
- Submit via gradescope (see Canvas) by 2:30pm Cambridge time.
Late submission accepted (on gradescope) until 11:59pm following day:
1 point (out of 10 points) per late problem penalty.
After that solutions are posted so not accepted without S3 excuse.
- Optional problems:
Don't count towards grade except for A+.
Value to you (besides the challenge):
Recommendations, employment (future grading, TA, UROP)
- Problem Set 1 is due in one week.

Closure Properties for Regular Languages

Theorem: If L_1 and L_2 are regular languages, so is $L_1 \cup L_2$ (closure under union)

Recall proof attempt: Let M_1 recognize L_1 and M_2 recognize L_2 .

Construct M recognizing $L_1 \cup L_2$:



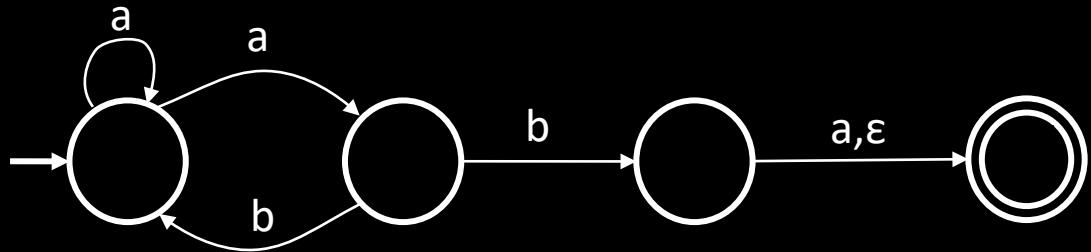
should accept input
if M_1 accepts
and M_2 accepts .

Doesn't work: Where to split ?

Hold off. Need new concept.



Nondeterministic Finite Automata



New features of nondeterminism:

- multiple paths possible (0, 1 or many at each step)
- ϵ -transition is a “free” move without reading input
- Accept input if some path leads to accept

Example inputs:

- ab
- aa
- aba
- abb

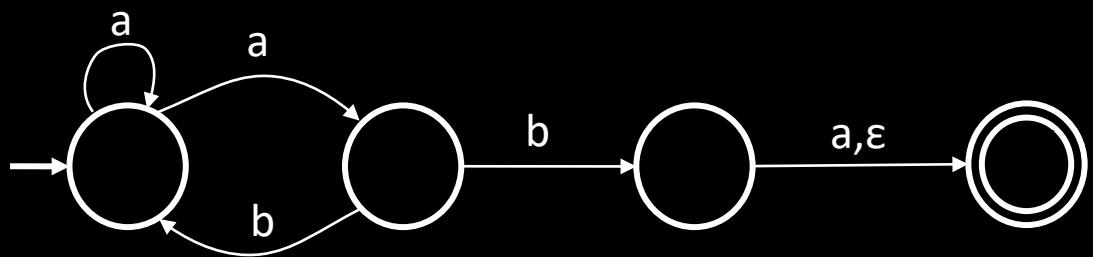
Check-in 2.1

What does do on input aab ?

- (a) Accept
- (b) Reject
- (c) Both Accept and Reject

Check-in 2.1

NFA – Formal Definition



Defn: A nondeterministic finite automaton (NFA) is a 5-tuple

states alphabet transition function
start state accept states

- all same as before except
-  power set
- In the example:

Ways to think about nondeterminism:

Computational: Fork new parallel thread and accept if any thread leads to an accept state.

Mathematical: Tree with branches.
Accept if any branch leads to an accept state.

Magical: Guess at each nondeterministic step which way to go. Machine always makes the right guess that leads to accepting, if possible.

Converting NFAs to DFAs

Theorem: If an NFA recognizes L then L is regular

Proof: Let \mathcal{N} be an NFA recognizing L .

Construct DFA \mathcal{D} recognizing L .

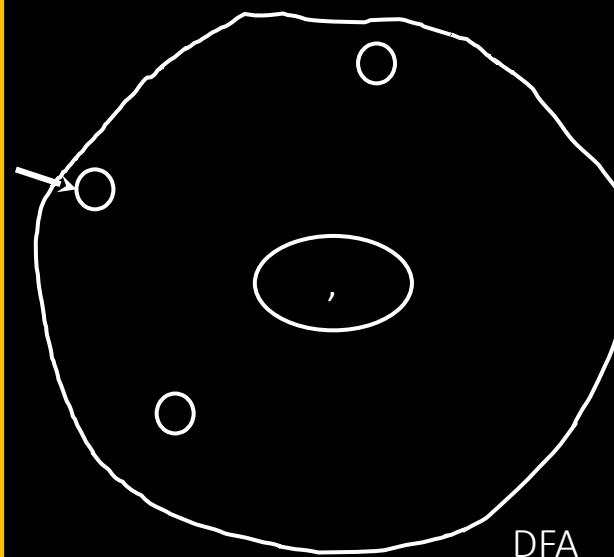
(Ignore the ϵ -transitions, can easily modify to handle them)

IDEA: DFA keeps track of the subset of possible states in NFA.

Check-in 2.2

If \mathcal{N} has n states, how many states does \mathcal{D} have by this construction?

- (a)
- (b)
- (c)



Construction of :

for some

intersects

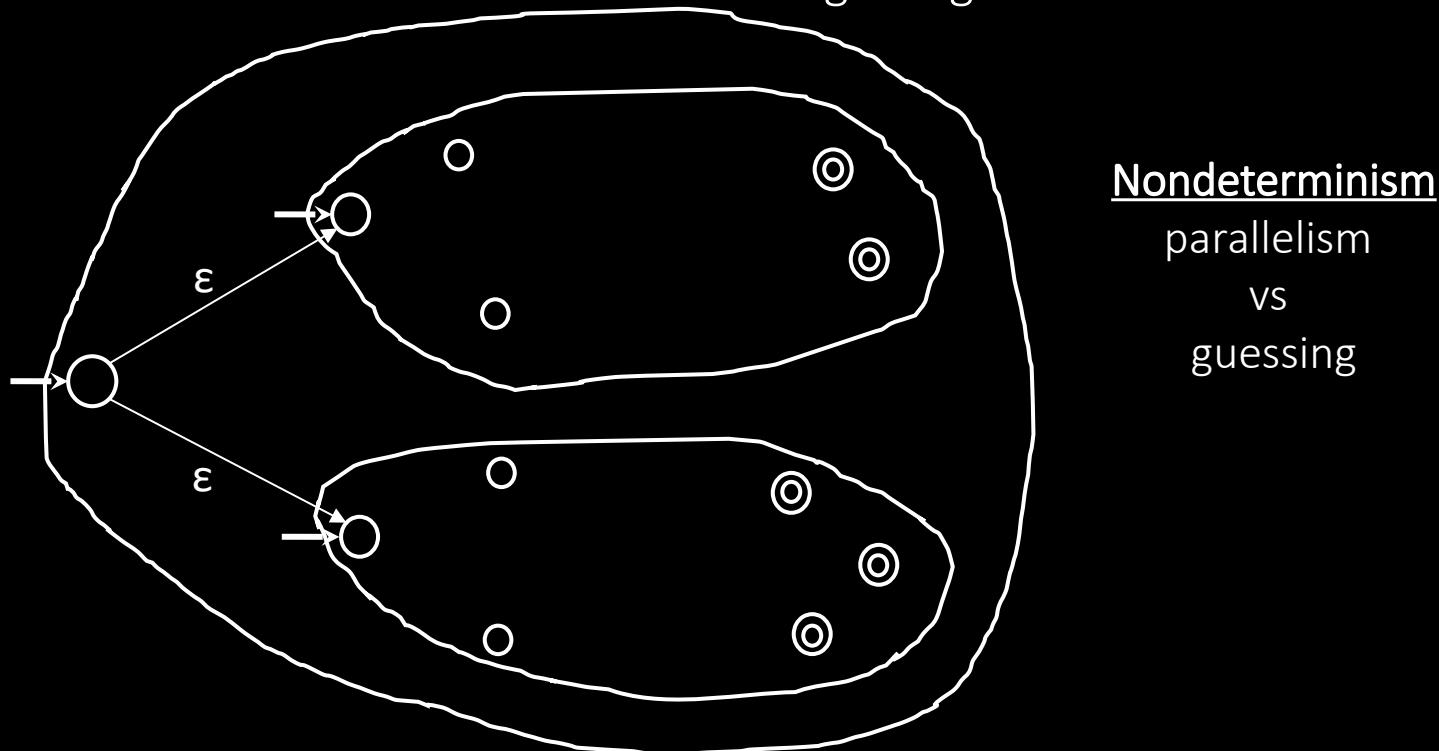
Check-in 2.2



Return to Closure Properties

Recall Theorem: If L_1 and L_2 are regular languages, so is $L_1 \cup L_2$.
(The class of regular languages is closed under union)

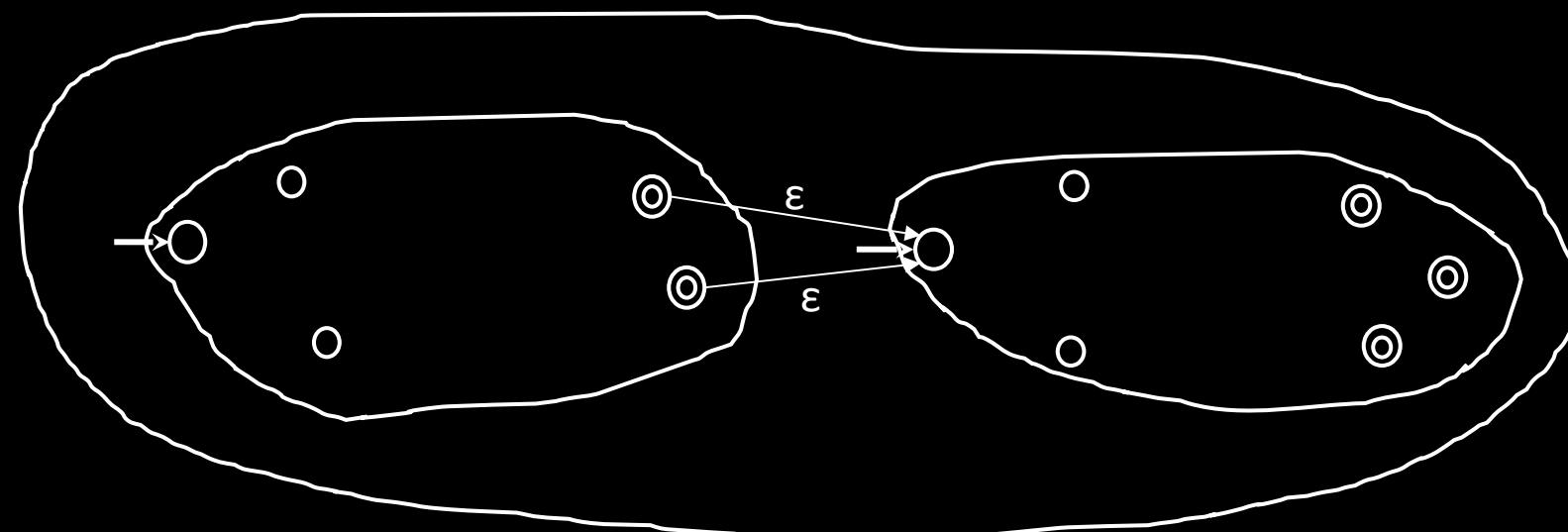
New Proof (sketch): Given DFAs M_1 and M_2 recognizing L_1 and L_2 respectively.
Construct NFA M recognizing $L_1 \cup L_2$.



Closure under (concatenation)

Theorem: If L_1 and L_2 are regular languages, so is $L_1 L_2$.

Proof sketch: Given DFAs M_1 and M_2 recognizing L_1 and L_2 respectively.
Construct NFA recognizing $L_1 L_2$.



should accept input
if where
accepts and accepts .

i ————— |

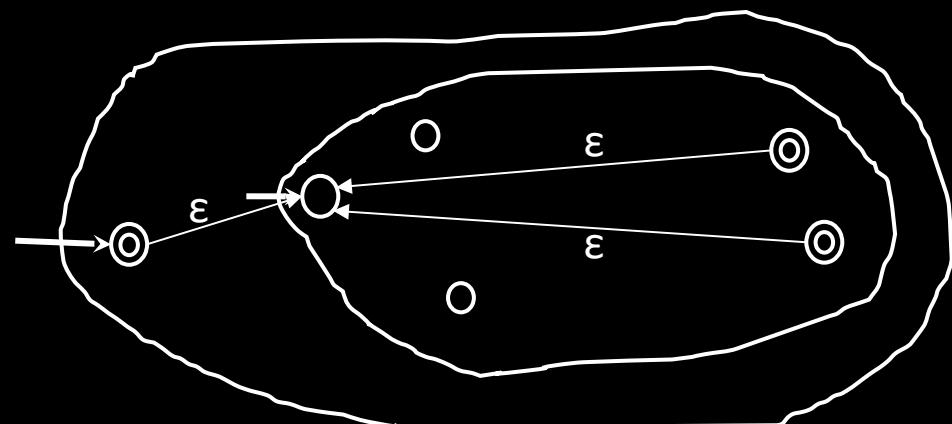
Nondeterministic has the option to
jump to when accepts.



Closure under \star (star)

Theorem: If L is a regular language, so is L^\star

Proof sketch: Given DFA M recognizing L
Construct NFA M' recognizing L^\star



Make sure M accepts ϵ

Check-in 2.3

If M has n states, how many states does M' have by this construction?

- (a)
- (b)
- (c)

Check-in 2.3

Regular Expressions NFA

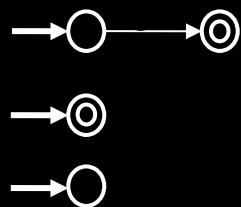
Theorem: If α is a regular expr and β then $\alpha\beta$ is regular

Proof: Convert α to equivalent NFA :

If α is atomic:

for

Equivalent is:



If α is composite:



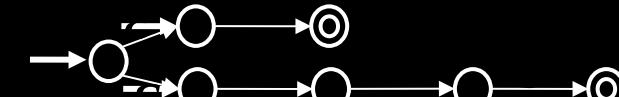
Use closure constructions

Example:

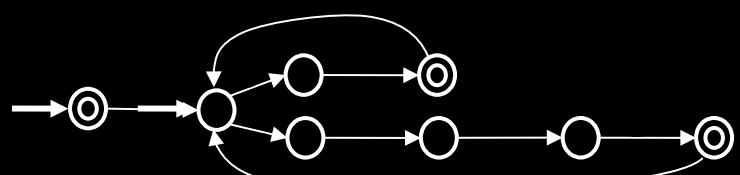
Convert α^* to equivalent NFA



:



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Quick review of today

1. Nondeterministic finite automata (NFA)
2. Proved: NFA and DFA are equivalent in power
3. Proved: Class of regular languages is closed under
4. Conversion of regular expressions to NFA

Check-in 2.4

Recitations start tomorrow online (same link as for lectures).

They are optional, unless you need more help.

You may attend any recitation(s).

Which do you think you'll attend? (you may check several)

- (a) 10:00 (b) 11:00 (c) 12:00
(d) 1:00 (e) 2:00 (f) I prefer a different time (please post on piazza, but no promises)

Check-in 2.4