

# Practical – 5

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**Batch:** B4

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**Aim:** Implement a dynamic algorithm for Longest Common Subsequence (LCS) to find the length and LCS for DNA sequences.

## Problem Statement:

(i) DNA sequences can be viewed as strings of A, C, G, and T characters, which represent nucleotides. Finding the similarities between two DNA sequences are an important computation performed in bioinformatics.

[Note that a subsequence might not include consecutive elements of the original sequence.]

**TASK 1:** Find the similarity between the given X and Y sequence.

X=AGCCCTAAGGGCTACCTAGCTT

Y= GACAGCCTACAAGCGTTAGCTTG

**Output:** Cost matrix with all costs and direction, final cost of LCS and the LCS.

Length of LCS=16

CODE:

```
def LCS(x, y):
    m, n = len(x), len(y)

    # Initialize Cost Matrix (c) and direction matrices (b)
    c = [[0] * (n + 1) for _ in range(m + 1)]
    b = [["]" * (n + 1) for _ in range(m + 1)]
```

```

for i in range(1, m + 1):
    for j in range(1, n + 1):
        if x[i - 1] == y[j - 1]:
            c[i][j] = c[i - 1][j - 1] + 1
            b[i][j] = "D"
        elif c[i - 1][j] >= c[i][j - 1]:
            c[i][j] = c[i - 1][j]
            b[i][j] = "U"
        else:
            c[i][j] = c[i][j - 1]
            b[i][j] = "L"
return c, b

def get_LCS(b, x, i, j):
    """Backtrack to find the LCS string"""
    if i == 0 or j == 0:
        return ""
    if b[i][j] == "D":
        return get_LCS(b, x, i - 1, j - 1) + x[i - 1]
    elif b[i][j] == "U":
        return get_LCS(b, x, i - 1, j)
    else:
        return get_LCS(b, x, i, j - 1)

# Input sequences
X = "AGCCCTAACGGCTACCTAGCTT"
Y = "GACAGCCTACAAGCGTTAGCTTG"

c, b = LCS(X, Y)
lcs_string = get_LCS(b, X, len(X), len(Y))

print("Final Cost of LCS:", c[len(X)][len(Y)])
print("LCS String:", lcs_string)

# Optional: print cost matrix
print("\nCost Matrix:")
for row in c:
    print(row)

# Printing Direction Matrix
print("\nDirection Matrix:")
for row in b:
    print(row)

```

## Code Screenshot:

```
Task-1.py > LCS
1  def LCS(x, y):
2      m, n = len(x), len(y)
3
4      # Initialize Cost Matrix (c) and direction matrices (b)
5      c = [[0] * (n + 1) for _ in range(m + 1)]
6      b = [["] * (n + 1) for _ in range(m + 1)]
7
8      for i in range(1, m + 1):
9          for j in range(1, n + 1):
10             if x[i - 1] == y[j - 1]:
11                 c[i][j] = c[i - 1][j - 1] + 1
12                 b[i][j] = "D"
13             elif c[i - 1][j] >= c[i][j - 1]:
14                 c[i][j] = c[i - 1][j]
15                 b[i][j] = "U"
16             else:
17                 c[i][j] = c[i][j - 1]
18                 b[i][j] = "L"
19
20
21
22     def get_LCS(b, x, i, j):
23         """Backtrack to find the LCS string"""
24         if i == 0 or j == 0:
25             return ""
26         if b[i][j] == "D":
27             return get_LCS(b, x, i - 1, j - 1) + x[i - 1]
28         elif b[i][j] == "U":
29             return get_LCS(b, x, i - 1, j)
30         else:
31             return get_LCS(b, x, i, j - 1)
```

```
33
34     # Input sequences
35     X = "AGCCCTAACGGCTACCTAGCTT"
36     Y = "GACAGCCTACAAGCGTTAGCTTG"
37
38     c, b = LCS(X, Y)
39     lcs_string = get_LCS(b, X, len(X), len(Y))
40
41     print("Final Cost of LCS:", c[len(X)][len(Y)])
42     print("LCS String:", lcs_string)
43
44     # Optional: print cost matrix
45     print("\nCost Matrix:")
46     for row in c:
47         print(row)
48
49     # Printing Direction Matrix
50     print("\nDirection Matrix:")
51     for row in b:
52         print(row)
53
```

## Code Output:

**TASK-2:** Find the longest repeating subsequence (LRS). Consider it as a variation of the longest common subsequence (LCS) problem. Let the given string be S. You need to find the LRS within S. To use the LCS framework, you effectively compare S with itself. So, consider string<sub>1</sub> = S and string<sub>2</sub> = S.

Example:

AABCBDC

LRS= ABC or ABD

CODE:

```
def LRS(s):
    n = len(s)
    # Step 1: Initialize DP table
    c = [[0] * (n + 1) for _ in range(n + 1)]

    # Step 2: Fill DP table
    for i in range(1, n + 1):
        for j in range(1, n + 1):
            if s[i - 1] == s[j - 1] and i != j:
                c[i][j] = 1 + c[i - 1][j - 1] # diagonal
            else:
                c[i][j] = max(c[i - 1][j], c[i][j - 1]) # top or left

    # Step 3: Backtracking to reconstruct LRS
    i, j = n, n
    lrs_seq = []
    while i > 0 and j > 0:
        if s[i - 1] == s[j - 1] and i != j:
            lrs_seq.append(s[i - 1])
            i -= 1
            j -= 1
        elif c[i - 1][j] >= c[i][j - 1]:
            i -= 1
        else:
            j -= 1

    return c, c[n][n], "".join(reversed(lrs_seq))

# Example Run
S = "AABEBCDD"
dp, length, lrs_str = LRS(S)

print("LRS Length:", length)
print("LRS:", lrs_str)
```

```

print("\nDP Matrix:")
for row in dp:
    print(row)

```

Code Screenshot:

```

Task-2.py > ...
1  def LRS(s):
2      n = len(s)
3      # Step 1: Initialize DP table
4      c = [[0] * (n + 1) for _ in range(n + 1)]
5
6      # Step 2: Fill DP table
7      for i in range(1, n + 1):
8          for j in range(1, n + 1):
9              if s[i - 1] == s[j - 1] and i != j:
10                  c[i][j] = 1 + c[i - 1][j - 1] # diagonal
11              else:
12                  c[i][j] = max(c[i - 1][j], c[i][j - 1]) # top or left
13
14      # Step 3: Backtracking to reconstruct LRS
15      i, j = n, n
16      lrs_seq = []
17      while i > 0 and j > 0:
18          if s[i - 1] == s[j - 1] and i != j:
19              lrs_seq.append(s[i - 1])
20              i -= 1
21              j -= 1
22          elif c[i - 1][j] >= c[i][j - 1]:
23              i -= 1
24          else:
25              j -= 1
26
27      return c, c[n][n], "".join(reversed(lrs_seq))
28
29
30  # Example Run
31  S = "AABEBCDD"
32  dp, length, lrs_str = LRS(S)
33
34  print("LRS Length:", length)
35  print("LRS:", lrs_str)
36
37  print("\nDP Matrix:")
38  for row in dp:
39      print(row)
40

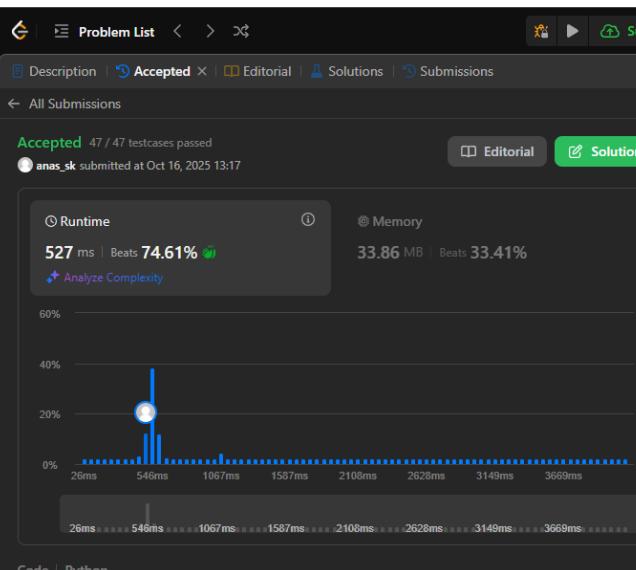
```

## Code Output:

```
[Running] python -u "c:\Users\Krish\OneDrive\Desktop\RBU\RBU-Sem-3\LAB  
LRS Length: 3  
LRS: ABD  
  
DP Matrix:  
[0, 0, 0, 0, 0, 0, 0, 0, 0]  
[0, 0, 1, 1, 1, 1, 1, 1, 1]  
[0, 1, 1, 1, 1, 1, 1, 1, 1]  
[0, 1, 1, 1, 1, 2, 2, 2, 2]  
[0, 1, 1, 1, 1, 2, 2, 2, 2]  
[0, 1, 1, 2, 2, 2, 2, 2, 2]  
[0, 1, 1, 2, 2, 2, 2, 2, 2]  
[0, 1, 1, 2, 2, 2, 2, 2, 3]  
[0, 1, 1, 2, 2, 2, 2, 3, 3]  
  
[Done] exited with code=0 in 0.099 seconds
```

## Leetcode Question:

link : <https://leetcode.com/problems/longest-common-subsequence/submissions/1803160931/>



The screenshot shows a Leetcode submission page for a Python solution. The submission was accepted with 47/47 testcases passed, submitted by [anas\\_sk](#) at Oct 16, 2025 13:17. The runtime is 527 ms (Beats 74.61%) and the memory usage is 33.86 MB (Beats 33.41%). The code is a dynamic programming solution for finding the longest common subsequence:

```
1 class Solution(object):  
2     def longestCommonSubsequence(self, text1, text2):  
3         m, n = len(text1), len(text2)  
4         dp = [[0] * (n + 1) for _ in range(m + 1)]  
5  
6         for i in range(1, m + 1):  
7             for j in range(1, n + 1):  
8                 if text1[i - 1] == text2[j - 1]:  
9                     dp[i][j] = dp[i - 1][j - 1] + 1  
10                else:  
11                    dp[i][j] = max(dp[i - 1][j], dp[i][j - 1])  
12  
13         return dp[m][n]
```

The page also includes a complexity analysis chart and a sidebar with saved testcases.

The screenshot shows a programming environment with a dark theme. On the left, the 'Problem List' tab is selected. A specific problem, '1143. Longest Common Subsequence', is highlighted. It has a 'Solved' status with a green circular icon. Below the title, there are buttons for 'Medium', 'Topics', 'Companies', and 'Hint'. The problem description states: "Given two strings `text1` and `text2`, return the length of their longest common subsequence. If there is no common subsequence, return 0." It defines a 'subsequence' as a new string generated from the original string with some characters deleted without changing the relative order of the remaining characters. An example is provided: "ace" is a subsequence of "abcde". A note specifies that a 'common subsequence' is a subsequence that is common to both strings.

**Example 1:**

**Input:** text1 = "abcde", text2 = "ace"  
**Output:** 3  
**Explanation:** The longest common subsequence is "ace" and its length is 3.

**Example 2:**

**Input:** text1 = "abc", text2 = "abc"

On the right, the 'Code' tab is active, showing Python code for solving the problem. The code uses dynamic programming with a 2D array `dp`. It initializes `dp[0][0] = 0` and `dp[i][0] = dp[0][j] = 0` for  $i > 0$  and  $j > 0$ . It then iterates through the strings to calculate the length of the longest common subsequence. The code is as follows:

```
1 class Solution(object):
2     def longestCommonSubsequence(self, text1, text2):
3         m, n = len(text1), len(text2)
4         dp = [[0] * (n + 1) for _ in range(m + 1)]
5
6         for i in range(1, m + 1):
7             for j in range(1, n + 1):
8                 if text1[i - 1] == text2[j - 1]:
9                     dp[i][j] = dp[i - 1][j - 1] + 1
10                else:
11                    dp[i][j] = max(dp[i - 1][j], dp[i][j - 1])
12
13         return dp[m][n]
```

The interface also shows statistics at the bottom: 14.7K views, 242 likes, 128 online submissions, and a 'Testcase' button.

