Sat Solver Optimisations

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This document contains research into optimisations of CNF-SAT problems.

1 Definitions

BOOL *Boolean Function.* A boolean function $f : V \subseteq Variables \to \{0, 1\}$ with truthy sets $P \subseteq \mathcal{P}(Variables)$ yields 1 if $V \in P$, else yields 0.

SAT *Boolean Satisfyability Equation.* We define a boolean formula by the following context-free grammar with start variable S, a given variable set $Variable = \{x_i : i \in [n-1]\}$ of n variables and a set of binary operations $Operation = \{\cdot, +, \oplus, \dots\}$ and an assigned boolean function $f_{Operation} : \mathcal{P} \subseteq \{L, R\} \to \{0, 1\}$, where L and R is the truthiness of left and right respectively.

$$\begin{array}{rcc} S & \to & (S) \\ S & \to & S \; Operation \; S \\ S & \to & \neg S \\ S & \to & Variable \end{array}$$

CNF *Unmixed conjunctive normal form.* We define a boolean formula in conjunctive normal form by the following context- free grammar with start variable S and a given variable set $Variable = \{x_i : i \in [n-1]\}$ of n variables.

$$
\begin{array}{ccc}
S & \rightarrow & (Disjunction) & \text{Short: } \mathbf{S}:\{D\} \\
S & \rightarrow & S \cdot (Disjunction) & \\
Disjunction & \rightarrow & Literal & \text{Short: } \mathbf{D}:\{X\} \\
Disjunction & \rightarrow & Disjunction + Literal & \\
Literal & \rightarrow & Variable & \text{Short: } \mathbf{L}:x_i \\
Literal & \rightarrow & \neg Variable &
\end{array}
$$

The notation $a \cdot \overline{b} + b \cdot \overline{c} \cdot \overline{d} + \overline{a} \cdot \overline{c}$ will be used as a set of literal sets $\{\{a, \neg b\}, \{b, \neg c, \neg d\}, \{\neg a, \neg b\}\}$ in this paper.

2 Transformation Techniques

U-SAT *Unmixed conjunctive normal form.* Given a *CNF-SAT* formula, we call it *unmixed*, if every disjunction consists of either positive or negated literals.

A polynomial time reduction $SAT \leq_p U-SAT$ for an equation $S_E \rightarrow U_E$ would be:

$$\forall D \in S:$$

TODO De Morgan

3-CNF-SAT *Three-literals conjunctive normal form.* Given a *CNF-SAT* formula, we call it $n - SAT$, if every disjunction consists of exactly n-literals.

TODO Proof

U3-SAT *Three-literals unmixed conjunctive normal form.* Given a *CNF-SAT* formula, we call it $U3-SAT$, if it is *unmixed* and every disjunction consists of exactly 3-literals.

TODO

3 Optimisation Techniques

Contradiction. A formula with contradictional disjunctions is not satisfiable.

$$\forall A : X \cdot \neg X \cdot A(X) = \bot$$

Common Literals. A formula with common literals is not satisfiable.

$$\forall A : X \cdot \neg X \cdot A(X) = \bot$$