

CNF · SAT

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TODO A CNF-SAT is an equation ...

Conjunctive normalform A satisfiability equation in conjunctive normalform is a set of literal sets. We denote it with $+$ as the logical operation *or* and \cdot (can be omitted) as *and*.

$$\{\{a, \bar{b}\}, \{b, \bar{c}, \bar{d}\}, \{\bar{a}, \bar{c}\}\}$$
$$(a + \bar{b}) \cdot (b + \bar{c} + \bar{d}) \cdot (\bar{a} + \bar{c})$$

Disjunctive normalform

$$ab + \bar{b}cd + ac$$

Formulas





Zugzwang. If a disjunction consists of one literal, the value of this literal is fixed.

$$\{X\} \in CNF \Rightarrow X = \top$$

Absorption (CNF). If a disjunction is a superset of another disjunction, it can be ignored.

$$\forall X \subseteq Y : X \cdot Y \equiv_{\text{CNF}} X$$

Proof: $(X \Rightarrow Y) \vee \neg X$ makes Y irrelevant.

	X	
Y		
		

The Absorption Rule allows to freely create new disjunctions, as long it's consisting of an already existing disjunction.









Absorption (DNF).

$$\forall X \subseteq Y : X + Y \equiv_{\text{DNF}} X$$

Proof: $(X \Rightarrow \top) \wedge (\neg X \Rightarrow \neg Y)$









Common Contradiction.









$$(A + X) \cdot (B + \overline{X}) \equiv_{\text{CNF}} (BX + A\overline{X})$$

	X			
A				
				
	B			

Resolution.

$$(A + X)(B + \overline{X}) \equiv_{\text{CNF}} (A + X)(B + \overline{X})(A + B)$$

	X			
A				
				
	B			

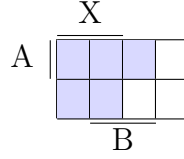
	X			
A				
				
	B			

Proof:

$$(A + X)(B + \overline{X}) \text{ is stronger than } (A + B)$$

Common Part.

$$(A + X) \cdot (B + X) \equiv_{\text{CNF}} (X + AB)$$



Semi-Common Part.

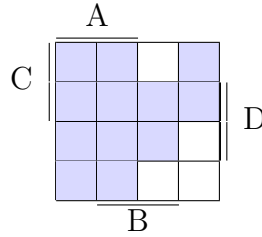
$$(AX + B)(A + C) \equiv_{\text{CNF}} A(X + B) + BC$$

Proof:

$$\begin{aligned} & (AX + B)(A + C) \\ \equiv_{\text{CNF}} & AX + AXC + AB + BC \\ \equiv_{\text{CNF}}^{\text{Absorption}} & AX + AB + BC \end{aligned}$$

Common and Contradicting Part

$$(A + B + C) \cdot (A + \neg B + D) \equiv_{\text{CNF}} (A + BD + C\overline{B})$$



Proof:

$$\begin{aligned} & (A + B + C) \cdot (A + \neg B + D) \\ \equiv_{\text{CNF}} & (A + (B + C)(\neg B + D)) \\ \equiv_{\text{CNF}} & (A + BD + C\overline{B}) \end{aligned}$$

Algorithms

bf0 · Brute Force

1. Try out every combination of variable interpretations.

da0 · Disjunctive Absorption

1. Use trivial optimisations.
2. Count the number of occurrences of every literals (positive and negated) and find the literal with most occurrences.
3. Branch: Assume the variable and accordingly clean up the equation using trivial transformations.
4. Absorption: Absorb as many disjunctions to the two- literals clauses.
5. If the equation is in 2-CNF-SAT, proceed with the next step, otherwise jump to the first step in a recursive step.
6. Solve 2-SAT the equation in polynomial time.