

CS 5/7320  
Artificial Intelligence

# Solving problems by searching

AIMA Chapter 3

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Slides by Michael Hahsler  
based on slides by Svetlana Lazepnik  
with figures from the AIMA textbook.

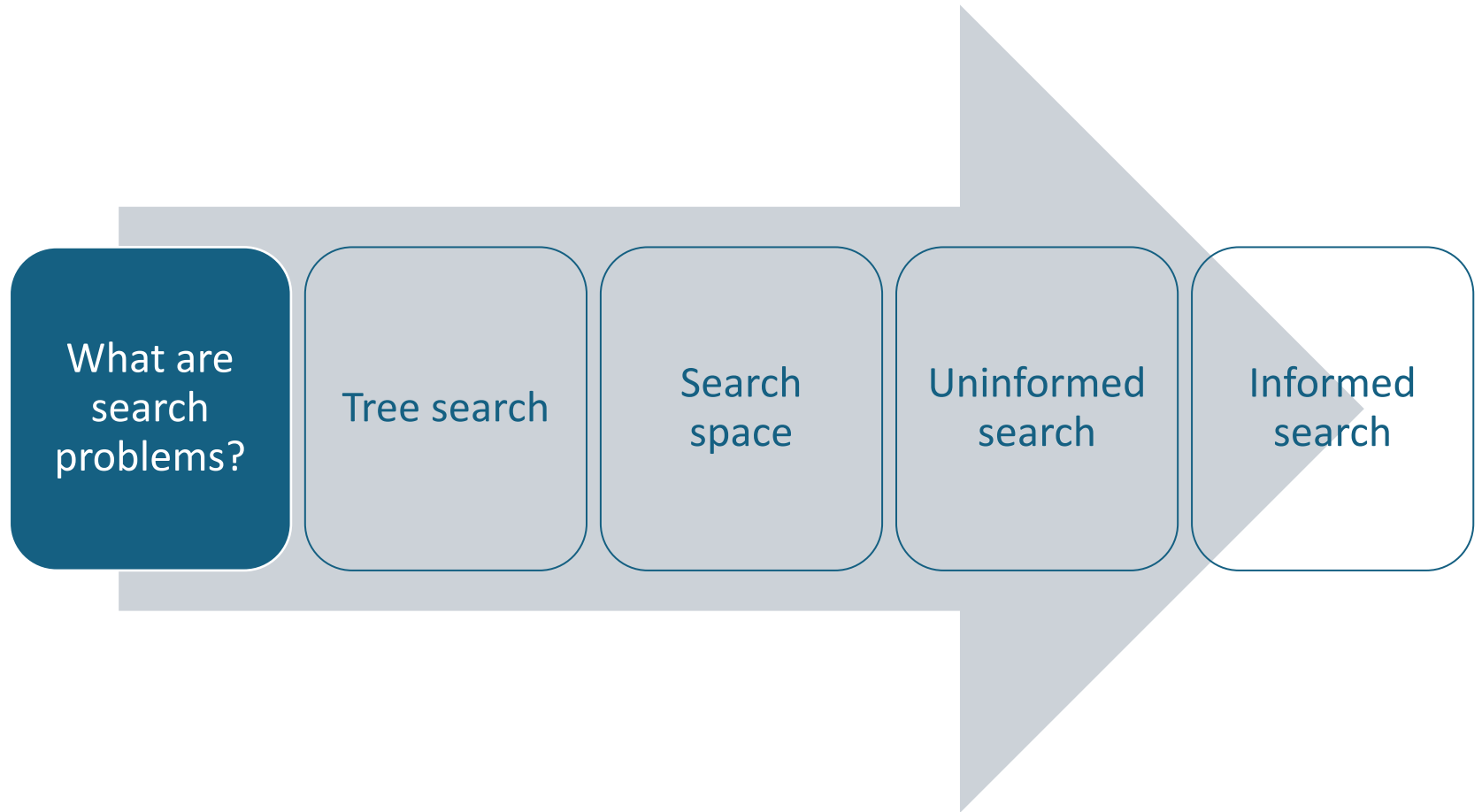


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Online Material

# Contents



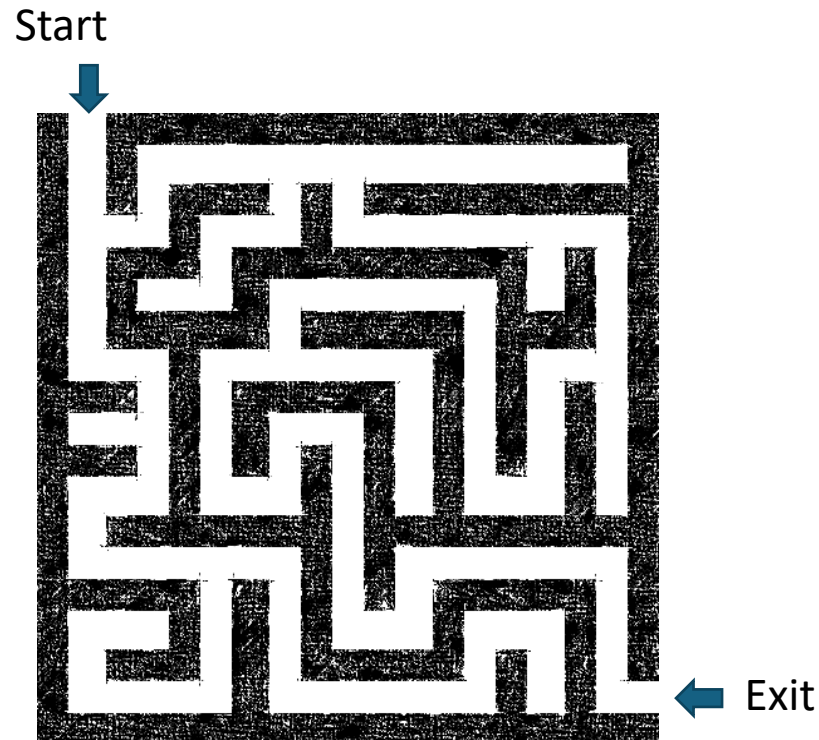


# Search Problems

How do we define a search problem?

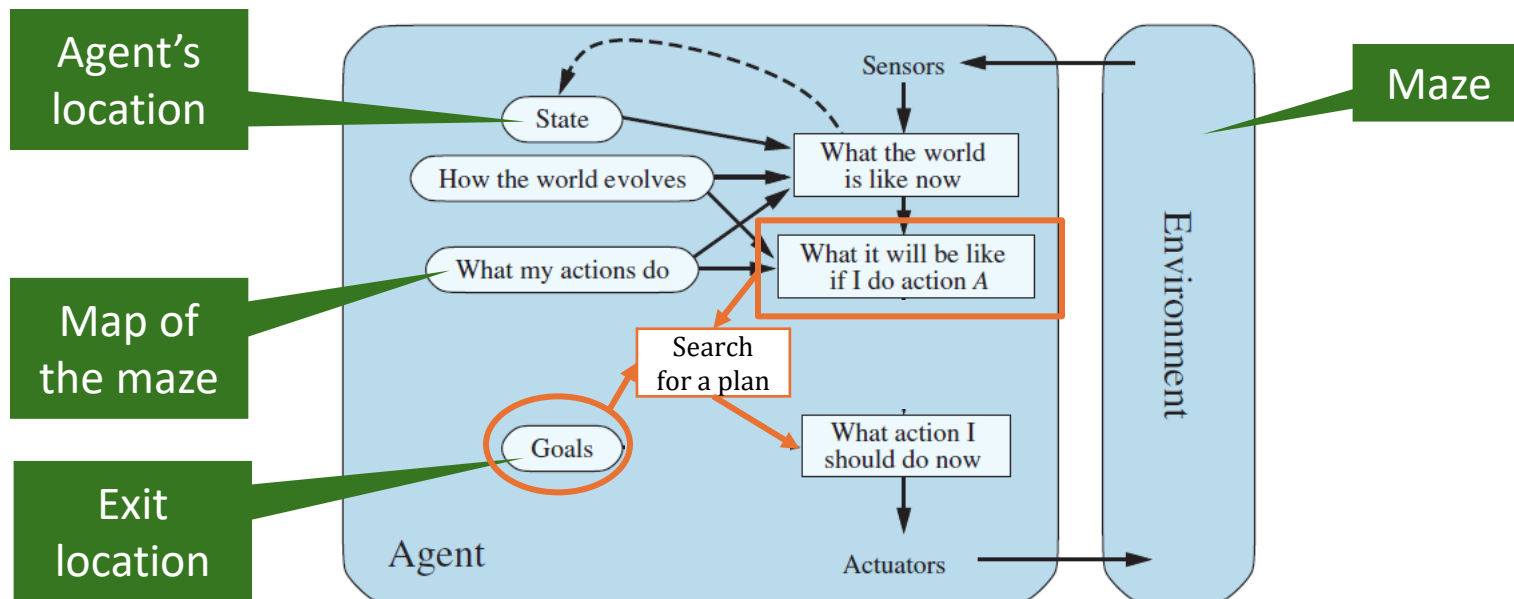
# What are Search Problems?

- We will consider the problem of designing **goal-based agents** in **known, fully observable, and deterministic** environments.
- Example environment:



# Remember: Planning Agent (Goal-based)

- The agent has the task of reaching a defined **goal state**.
- The performance measure is typically the cost to reach the goal.
- We will discuss a special type of **goal-based agents** called **planning agents**, which use **search algorithms** to plan a sequence of actions that lead to the goal.

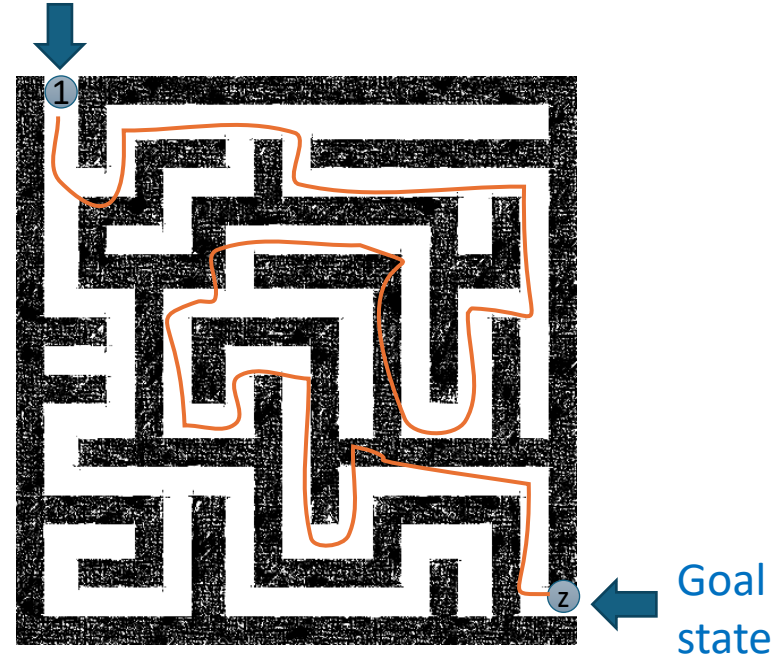


$$a_i = \operatorname{argmin}_{a_i \in A} \left[ \sum_{t=i}^T c_t \mid s_T \in S^{\text{goal}} \right]$$

# Planning for Search Problems

- For now, we consider only a discrete environment using an **atomic state representation** (states are just labeled 1, 2, ...).
- The **state space** is the set of all possible states of the environment. Some states are marked as **goal states**.
- The **optimal solution** is the sequence of actions (or equivalently, a sequence of states) that gives the lowest path cost for reaching the goal.

Initial state

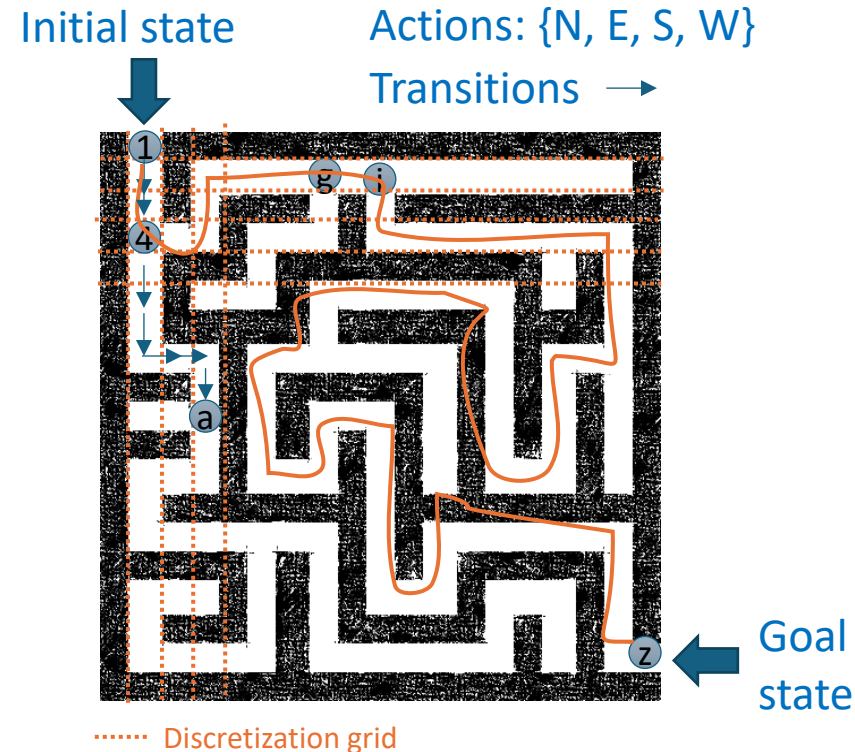


## Phases:

- 1) **Search/Planning**: the process of looking for the **sequence of actions** that reaches a goal state. Requires that the agent knows what happens when it moves!
- 2) **Execution**: Once the agent begins executing the search solution in a deterministic, known environment, it can ignore its percepts (**open-loop system**).

# Definition of a Search Problem

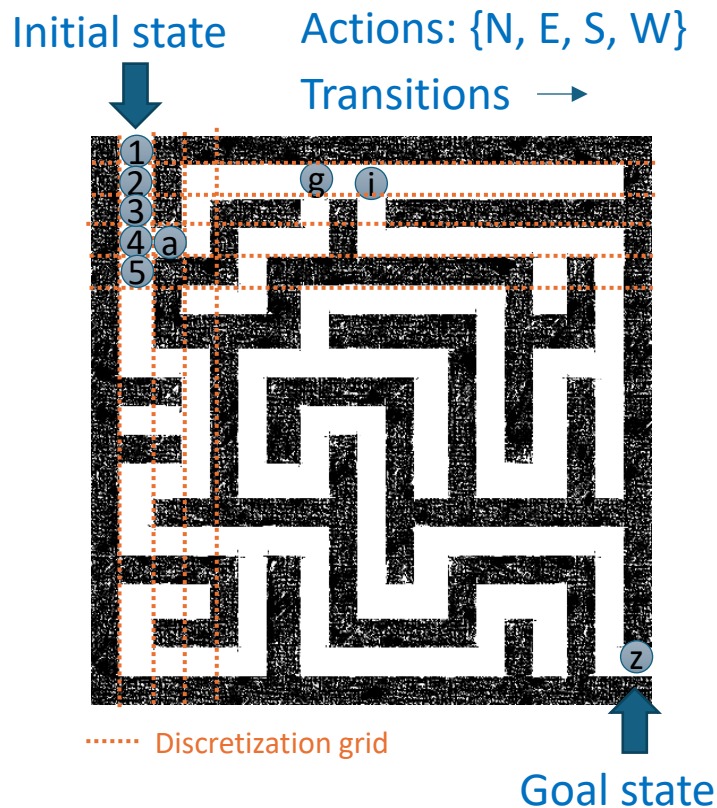
- **Initial state:** state description
- **Actions:** set of possible actions  $A$
- **Transition model:** a function that defines the new state resulting from performing an action in the current state
- **Goal state:** state description
- **Path cost:** the sum of *step costs*



**Important:** The **state space** is typically too large to be enumerated, or it is continuous. Therefore, the problem is defined by initial state, actions and the transition model and not the set of all possible states.

# Transition Function and Available Actions

## Original Description



- Definition as an **action schema**:  
 $Action(go(dir))$   
 PRECOND: no wall in direction  $dir$   
 EFFECT: change the agent's location according to  $dir$
- Definition as a **function**:  
 $f: S \times A \rightarrow S \text{ or } s' = result(s, a)$
- A **graph** with states as vertices and actions as edges.
- Function implemented as a **table** representing the state space as a graph.
- Available actions** in a state come from the transition function. E.g.,  
 $actions(4) = \{E, S, N\}$

$s$	$a$	$s'$
1	S	2
2	N	1
2	S	3
...	...	...
4	E	a
4	S	5
4	N	3
...	...	...

**Note:** Known and deterministic is a property of the transition function!



## Original Description

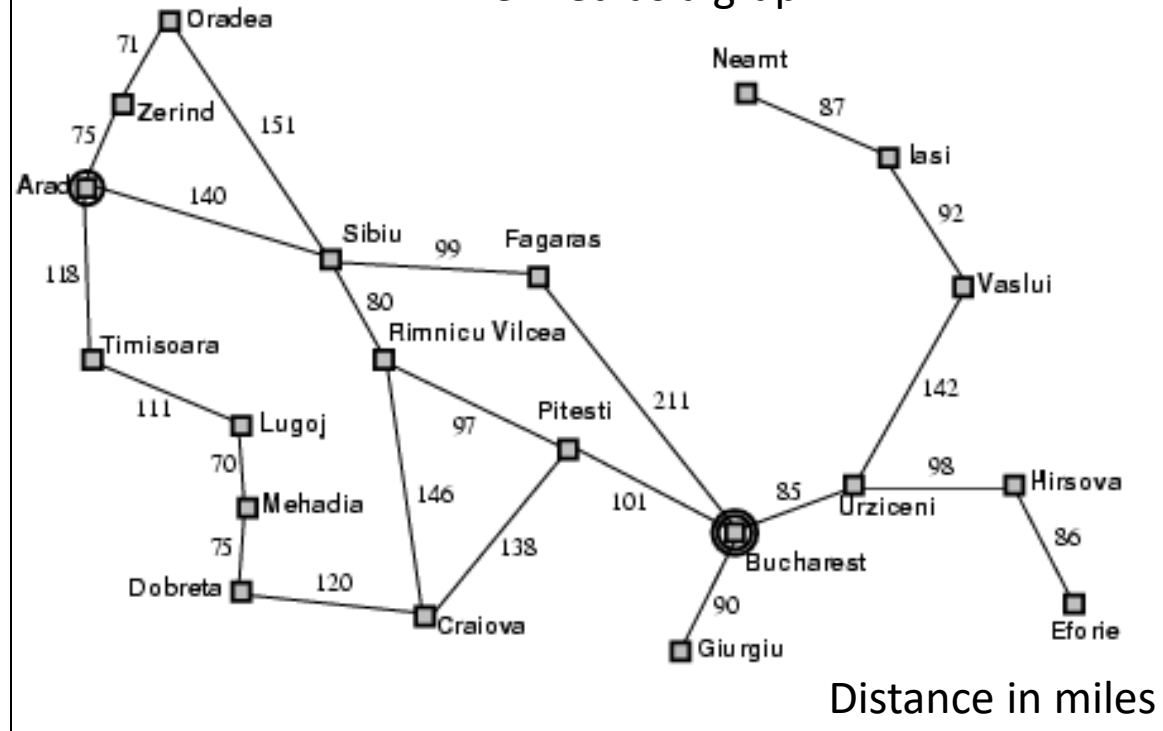


## Example: Romania Vacation

- On vacation in Romania; currently in Arad
- Flight leaves tomorrow from Bucharest

- **Initial state:** Arad
- **Actions:** Drive from one city to another.
- **Transition model and states:** If you go from city A to city B, you end up in city B.
- **Goal state:** Bucharest
- **Path cost:** Sum of edge costs.

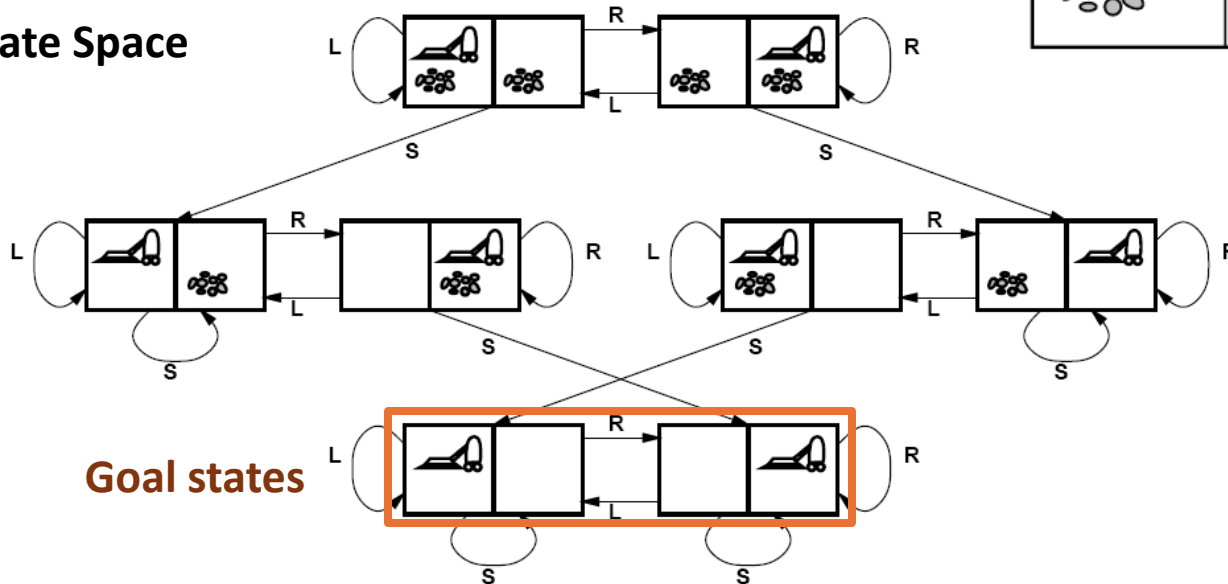
## State Space/Transition model Defined as a graph



Distance in miles

# Example: Vacuum world

State Space



- **Initial State:** Defined by agent location and dirt location.
- **Actions:** Left, right, suck
- **Transition model:** Clean a location or move.
- **Goal state:** All locations are clean.
- **Path cost:** E.g., number of actions

There are 8 possible atomic states of the system.  
Why is the number of states for  $n$  possible locations  $n(2^n)$ ?

# Example: Sliding-tile Puzzle

- **Initial State:** A given configuration.
- **Actions:** Move blank left, right, up, down
- **Transition model:** Move a tile
- **Goal state:** Tiles are arranged empty and 1-8 in order
- **Path cost:** 1 per tile move.

7	2	4
5		6
8	3	1

Start State

	1	2
3	4	5
6	7	8

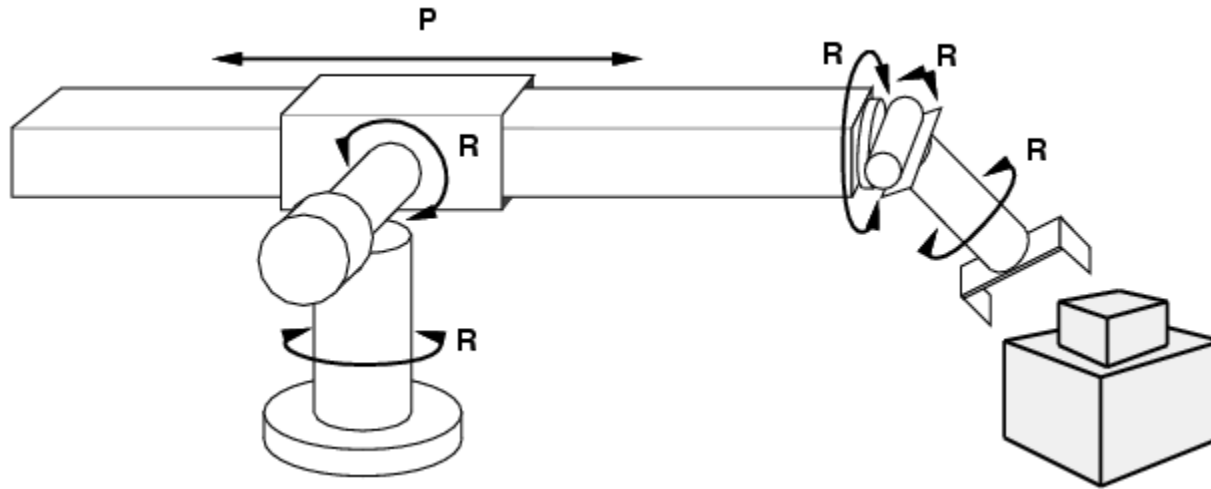
Goal State

## State space size

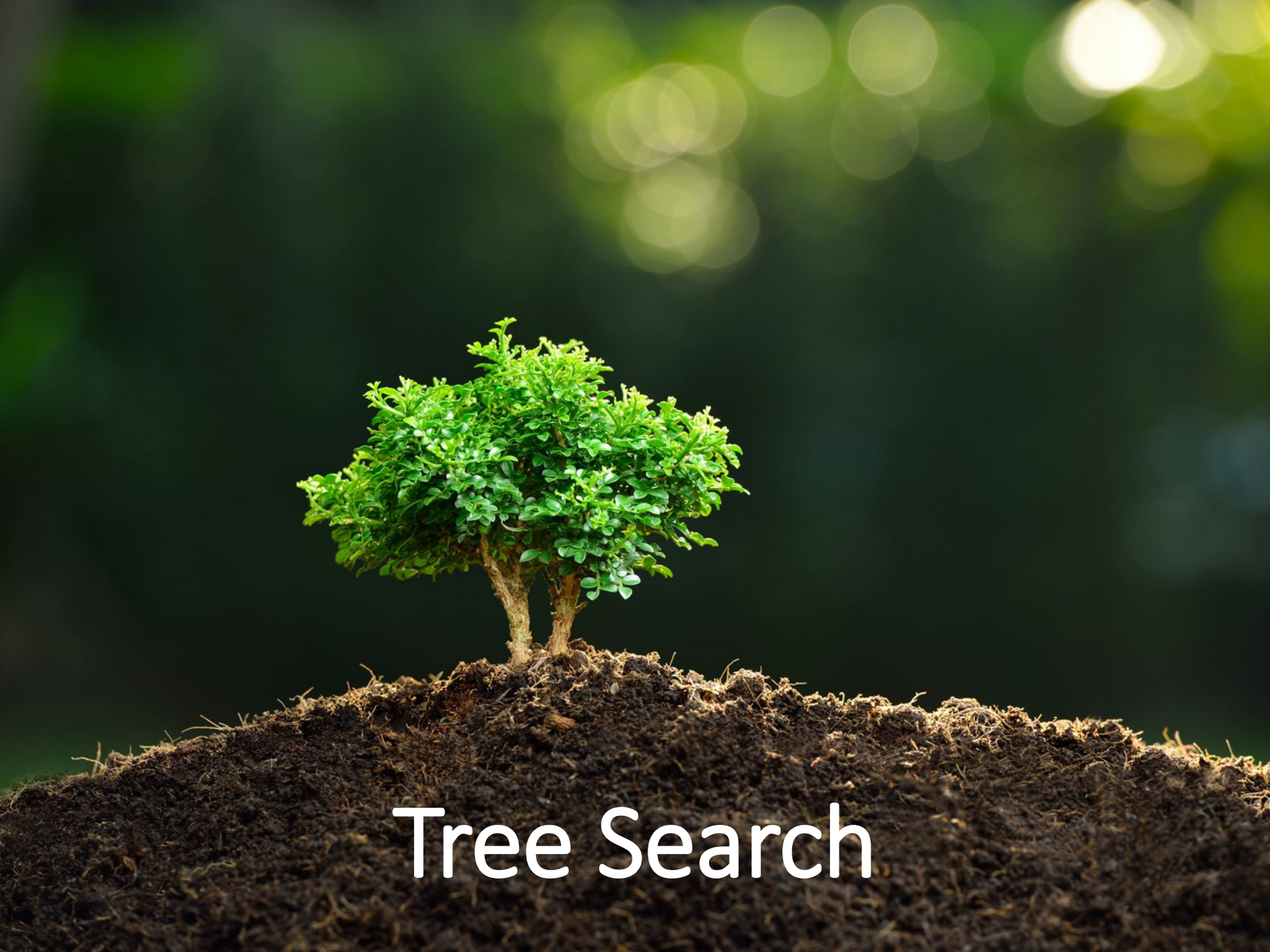
Each state describes the location of each tile (including the empty one).  $\frac{1}{2}$  of the permutations are unreachable.

- 8-puzzle:  $9!/2 = 181,440$  states
- 15-puzzle:  $16!/2 \approx 10^{13}$  states
- 24-puzzle:  $25!/2 \approx 10^{25}$  states

# Example: Robot Motion Planning



- **Initial State:** Current arm position with real-valued coordinates of robot joint angles.
- **Actions:** Continuous motions of robot joints.
- **Transition model:** Movement.
- **Goal state:** Desired final configuration (e.g., object is grasped).
- **Path cost:** Time to execute, smoothness of path, etc.



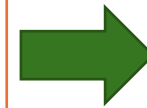
# Tree Search

# Solving Search Problems

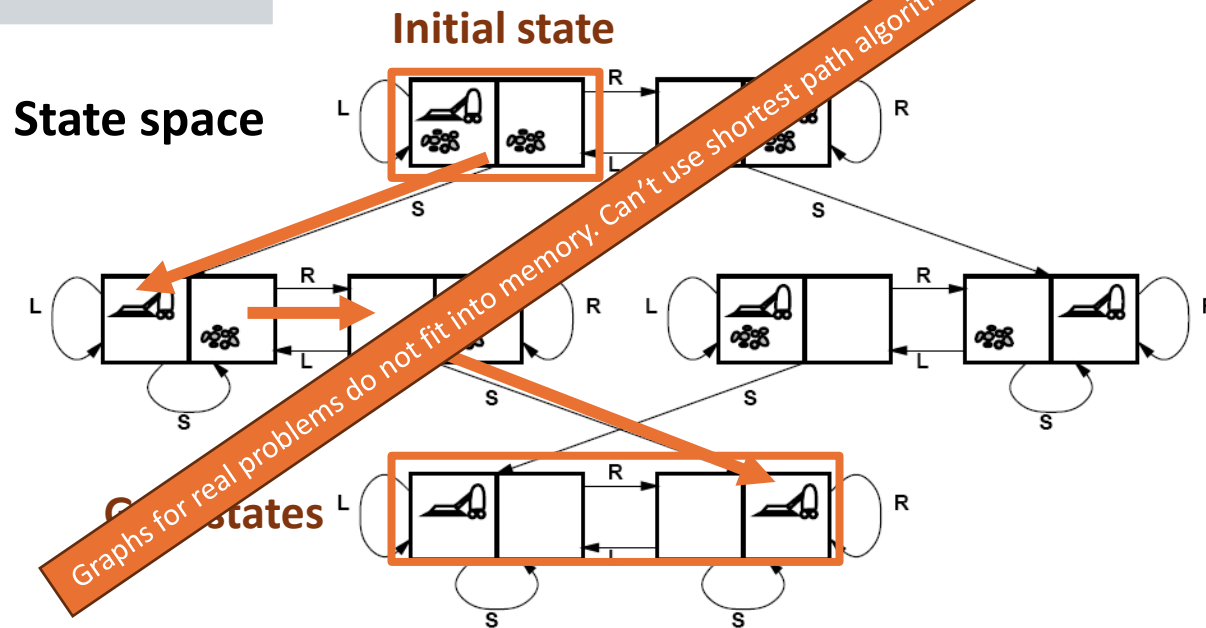
Given a search problem definition

- Initial state
- Actions
- Transition model
- Goal state
- Path cost

How do we find the optimal solution (sequence of actions/states) when shortest path algorithms for graphs are too expensive?



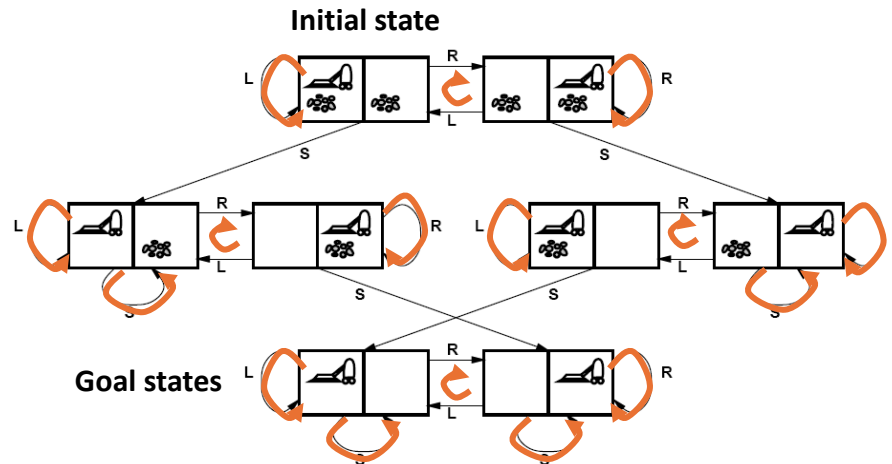
Construct a **search tree** for the state space graph so we can use much cheaper **tree search**!



# Issue: Transition Model is a Graph and Not a Tree!

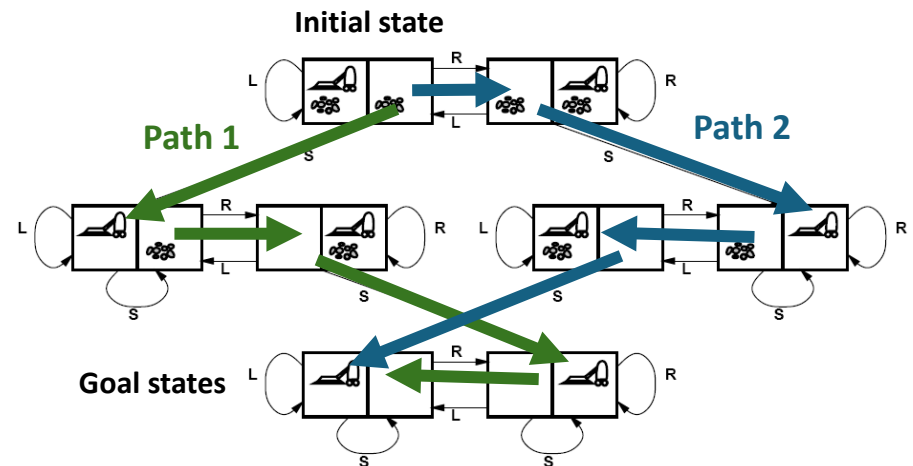
## Cycles

Return to the same state



## Non-cycle redundant paths

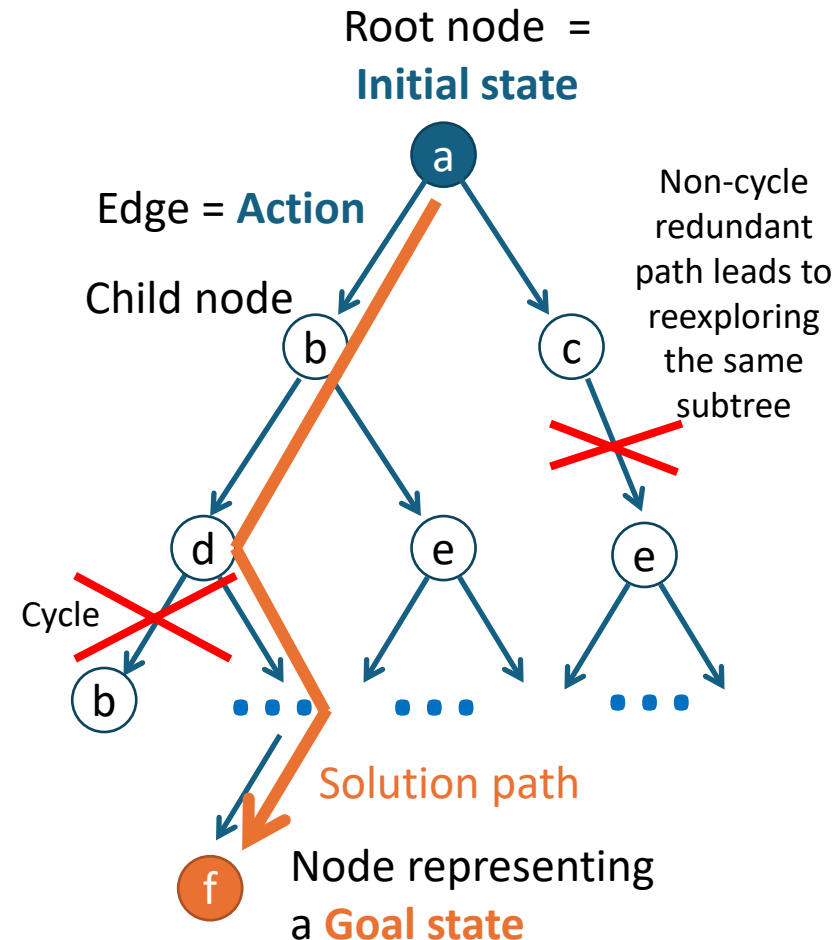
Multiple paths to get to the same state





# Creating a Search Tree

- Superimpose a “what if” tree of possible actions and outcomes (states) on the state space graph.
- The **Root node** represents the initial state.
- An action child node is reached by an **edge** representing an action. The corresponding state is defined by the transition model.
- Trees cannot have **cycles (loops)**. Cycles in the search space must be broken to prevent infinite loops.
- Trees cannot have **multiple paths to the same state**. These are called redundant paths. Removing suboptimal redundant paths improves search efficiency.
- A **path** through the tree corresponds to a sequence of actions (states).
- A **solution** is a path ending in a node representing a goal state.
- **Nodes vs. states**: Each tree node represents a state of the system. If redundant path cannot be prevented then state can be represented by multiple nodes in the tree.





# Differences Between Typical Tree Search and AI Search

## Typical tree search

- Assumes a given tree that fits in memory.
- Trees have by construction no cycles or redundant paths.

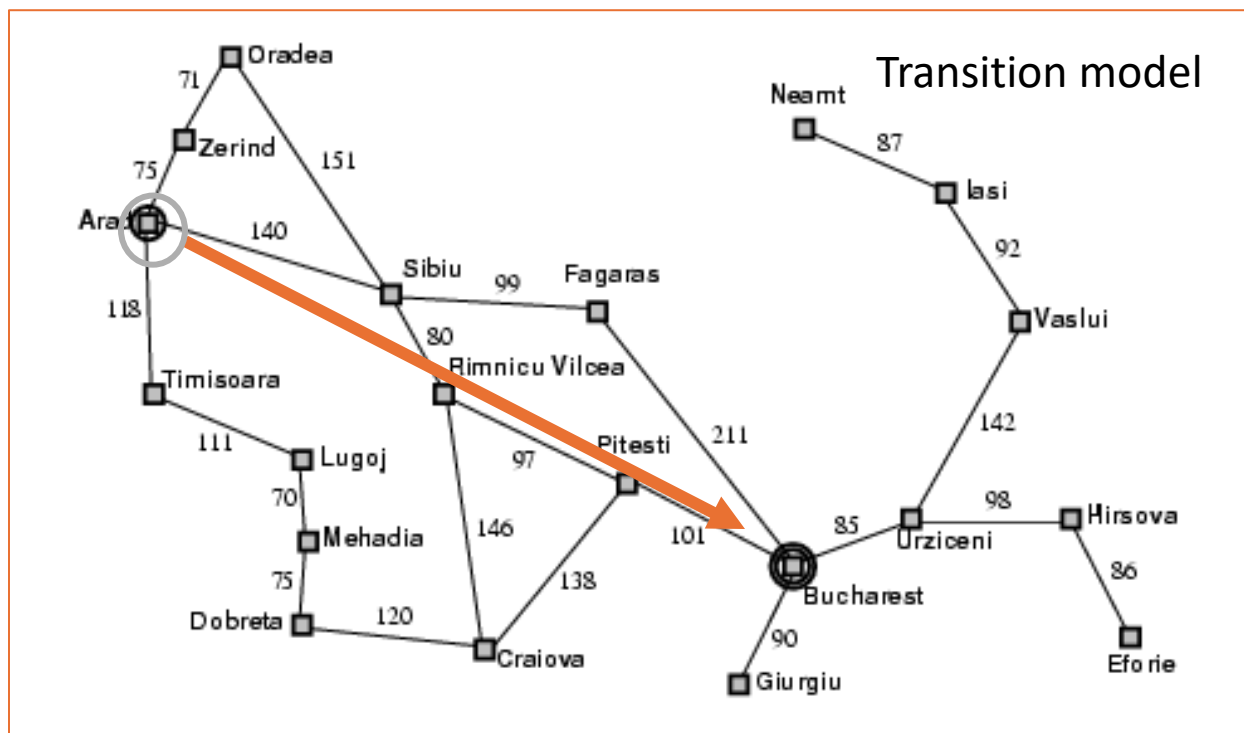
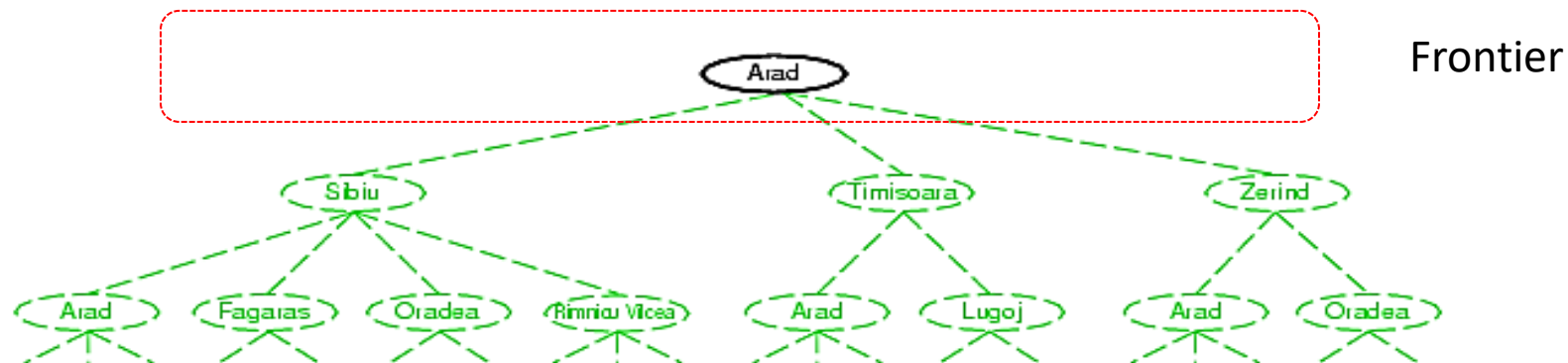
## AI tree/graph search

- The search tree is too large to fit into **memory**.
  - a. **Builds parts of the tree** from the initial state using the transition function representing the graph.
  - b. **Memory management** is very important.
- The search space is typically a very large and complicated graph. Memory-efficient **cycle checking** is very important to avoid infinite loops or minimize searching parts of the search space multiple times.
- Checking redundant paths often requires too much memory and we accept searching the same part multiple times.

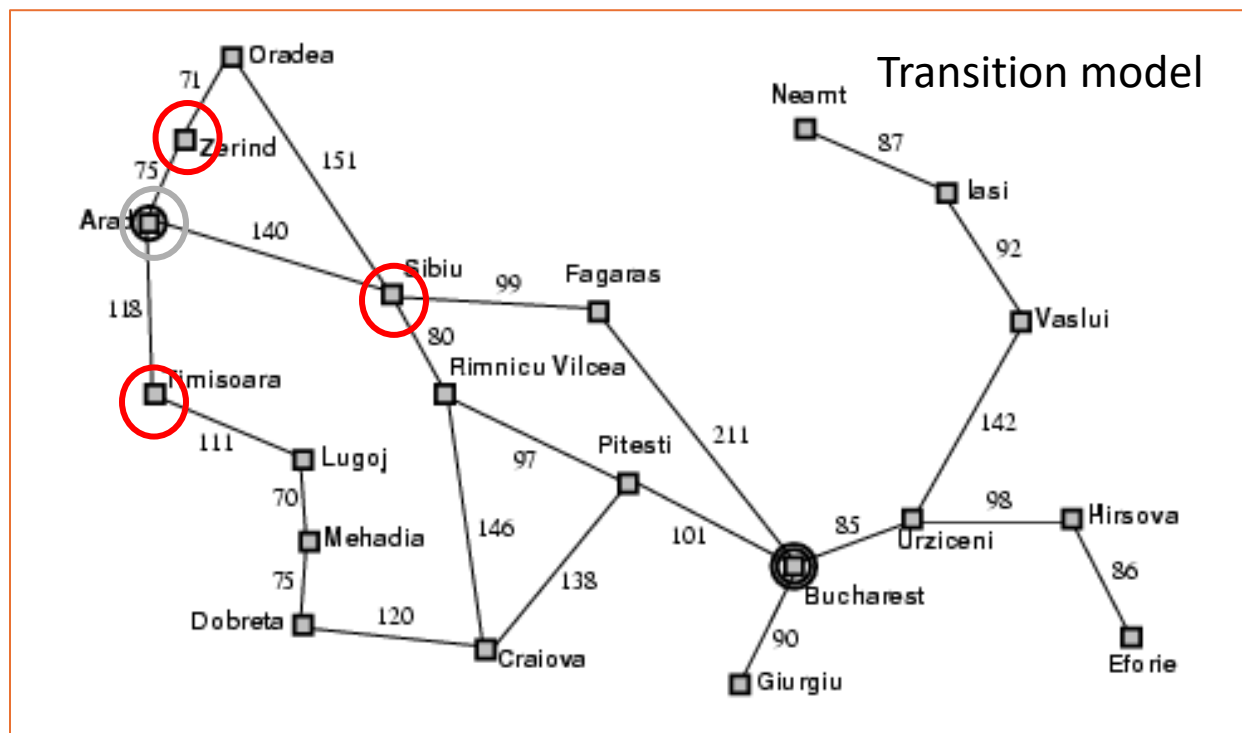
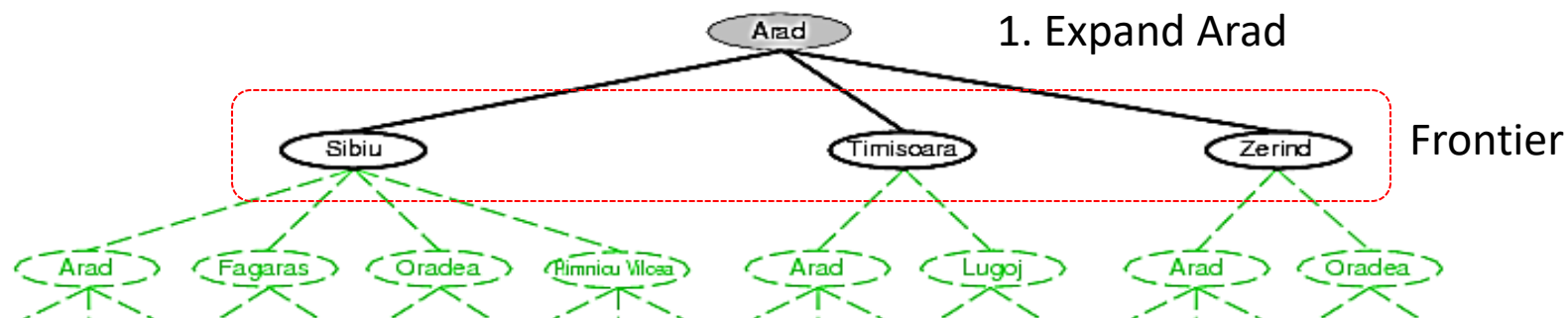
# Tree Search Algorithm Outline

1. Initialize the **frontier** (set of unexplored known nodes) using the **starting state/root node**.
2. While the frontier is not empty:
  - a) Choose the next frontier node to expand according to the **search strategy**.
  - b) If the node represents a **goal state**, return it as the solution.
  - c) Else **expand** the node (i.e., apply all possible actions to the transition model) and add its children nodes representing the newly reached states to the frontier.

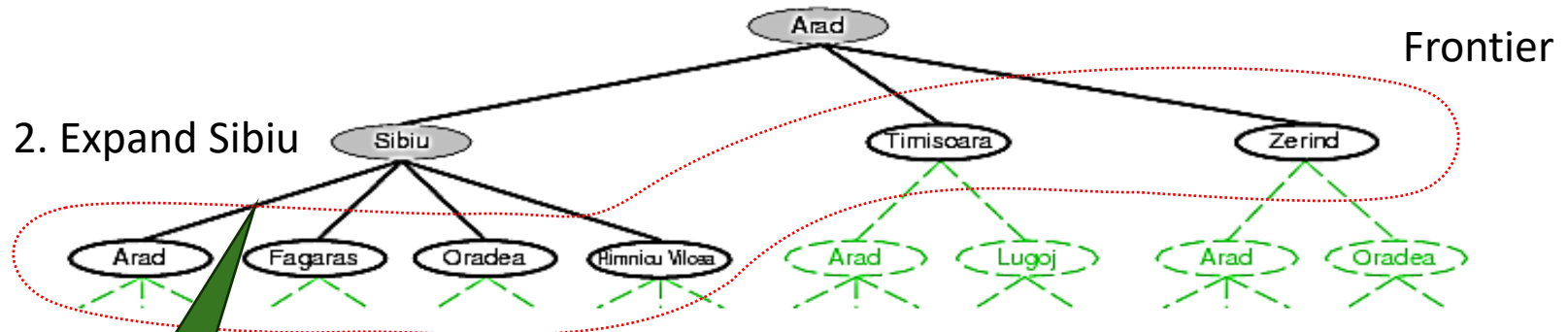
# Tree Search Example



# Tree Search Example

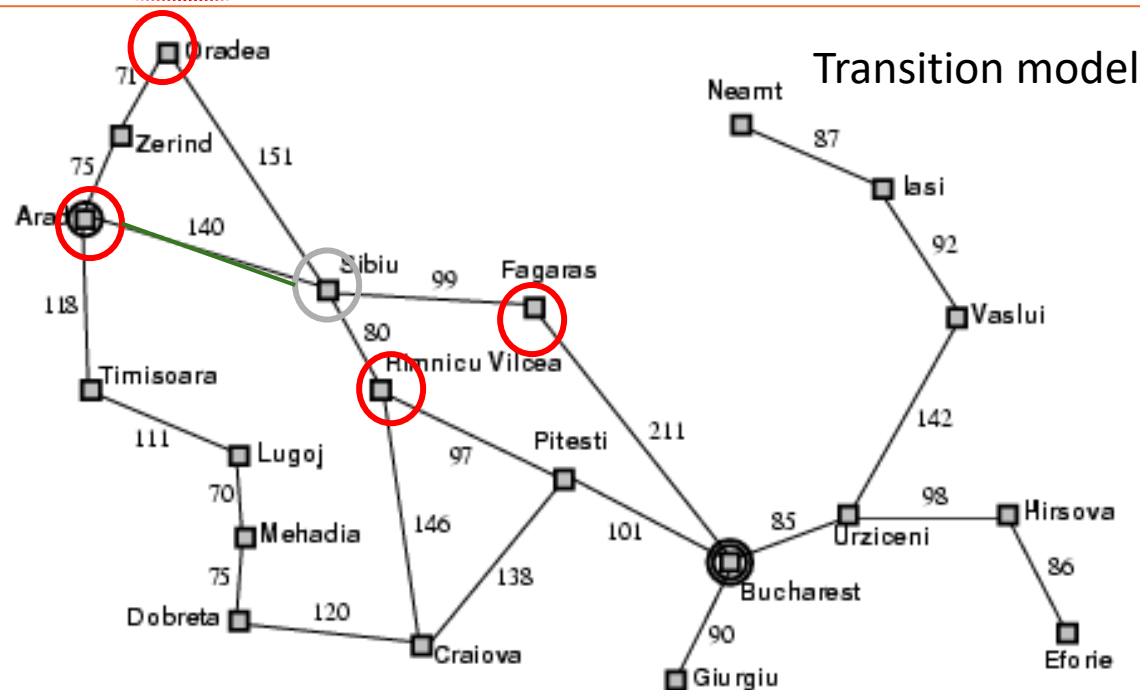


# Tree Search Example



## Example of a cycle

We could have  
also expanded  
Timisoara or  
Zerind!



# Search Strategies: Properties

- A **search strategy** is defined by picking the **order of node expansion**.
- Strategies are evaluated along the following dimensions:
  - **Completeness**: does it always find a solution if one exists?
  - **Optimality**: does it always find a least-cost solution?
  - **Time complexity**: how long does it take?
  - **Space complexity**: how much memory does it need?
- We will discuss different search strategies and use these properties to compare them.



**Space and Time Complexity**

**State Space vs. Search Tree Size**



# State Space vs. Search Tree Size

- Space and time complexity depend on the **number of tree nodes** searched (created and visited) till a goal node is found. For a tree with  $n$  nodes we have:

$$O(n)$$

- **Remember:** For perfect cycle checking and redundant path elimination, we have a 1:1 mapping between nodes and states:

Nodes in the search tree = states in the search space

Otherwise, we may have multiple nodes representing a state.

- We have the following options to estimate  $n$  for a search problem:
  - a. **Estimate the reachable state space size.**
  - b. **Estimate the number of searched tree nodes.**
- Estimating the complexity is important to judge:
  - How difficult is the problem?
  - What algorithm will fit in memory?
  - Can we find a solution fast enough?
  - Can we find the optimal solution, or do we need to use a heuristic?



# State Space Size Estimation

## State Space

- Number of different states the agent and environment can be in.
- **Reachable states** are defined by the initial state and the transition model. Only reachable states are important for search.

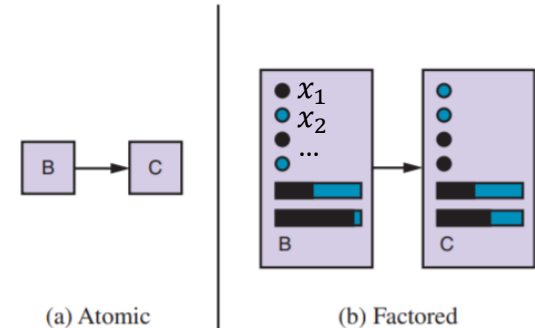
## Estimation

- Even if the used algorithm represents the state space using atomic states, we may know the internal (factored) representation. It can be used to estimate the problem size.
- The basic rule is to estimate the state space size for factored state representation with  $l$  fluents (variables) as:

$$n = |X_1| \times |X_2| \times \dots \times |X_l|$$

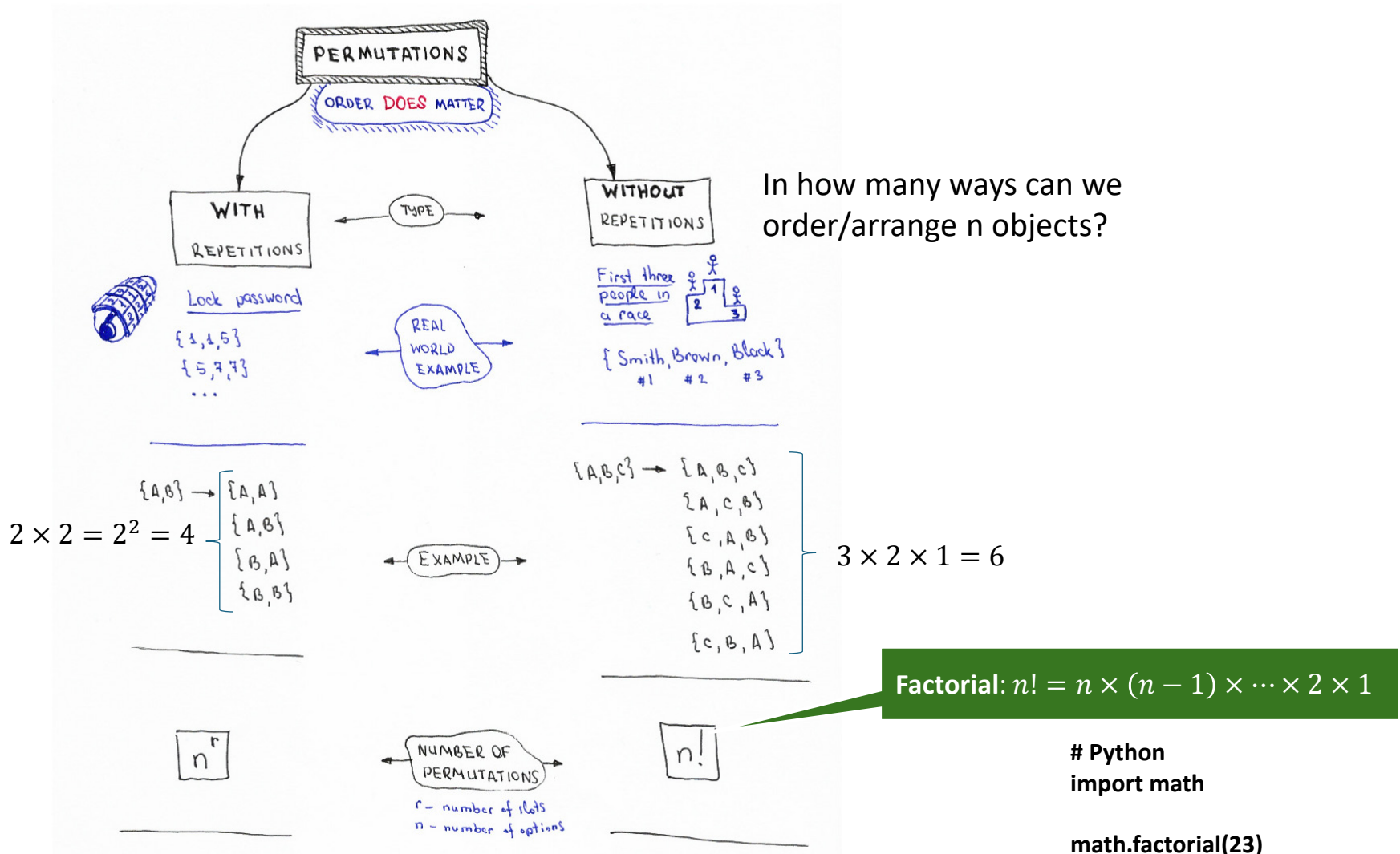
where  $|\cdot|$  is the number of possible values.

## State representation

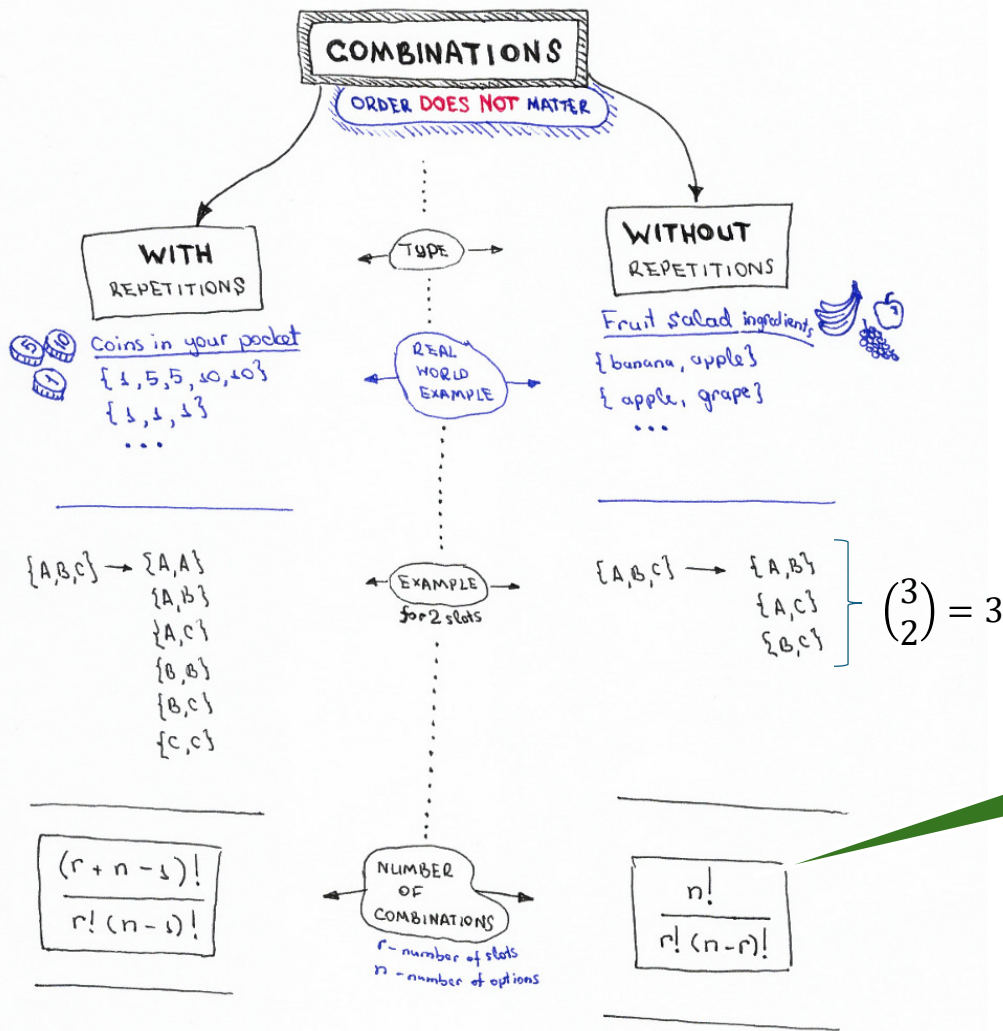


The factored state consists of variables called fluents that represent conditions that can change over time.

# Reminder: Combinatorics - Permutations



# Reminder: Combinatorics - Combinations

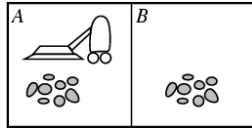


**Binomial Coefficient:**  $\binom{n}{r} = C(n, r) = {}_nC_r$   
 Read as "n choose r" because it is the number of ways can we choose r out of n objects?  
 Special case for  $r = 2$ :  $\binom{n}{2} = \frac{n(n-1)}{2}$

```
# Python
import scipy.special
```

```
# the two give the same results
scipy.special.binom(10, 5)
scipy.special.comb(10, 5)
```

# Example: What is the State Space Size?



## Dirt

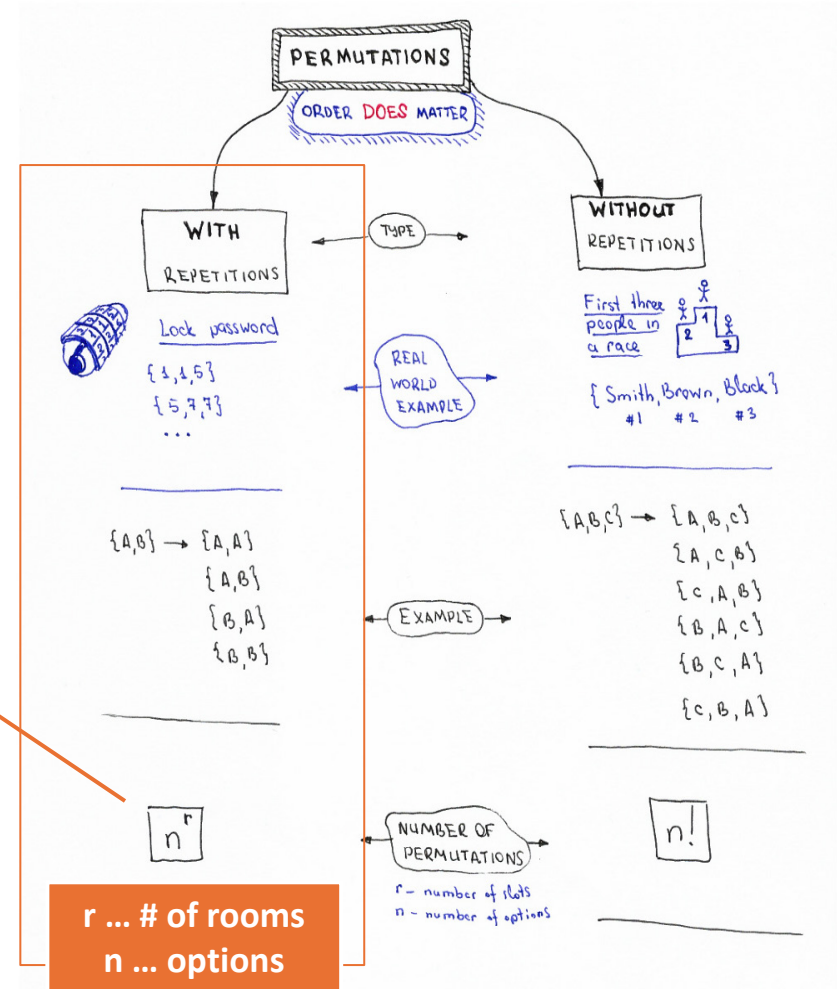
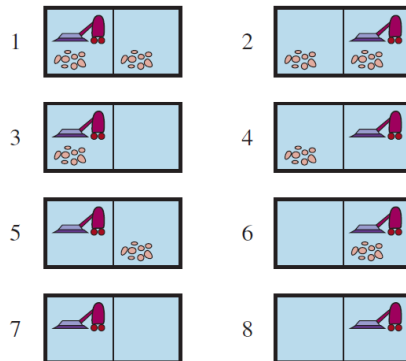
- **Permutation:** A and B are different rooms, order does matter!
- **With repetition:** Dirt can be in both rooms.
- There are 2 options (clean/dirty)

→  $2^2$

## Robot location

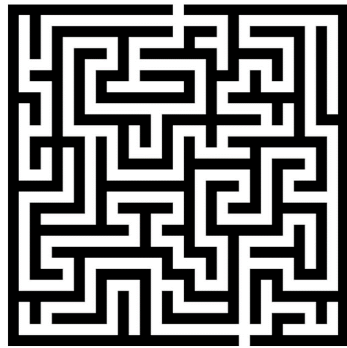
- Can be in 1 out of 2 rooms.  
→ 2

Total:  $n = 2 \times 2^2 = 2^3 = 8$

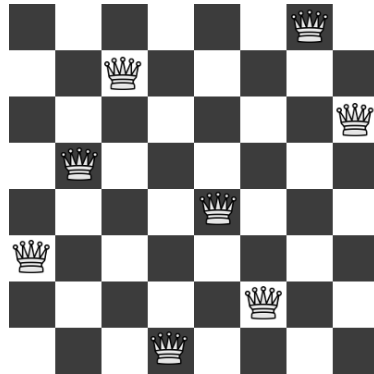


# Examples: What is the State Space Size?

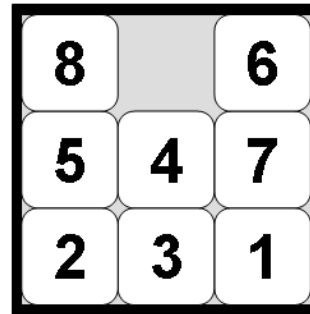
Often a rough upper limit is sufficient to determine how hard the search problem is.



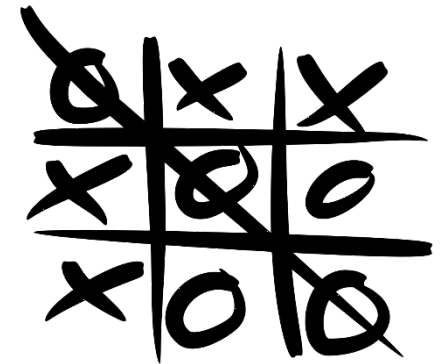
Maze



8-queens problem



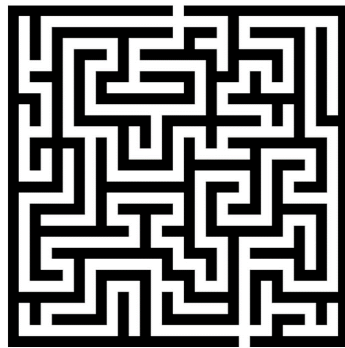
8-puzzle problem



Tic-tac-toe

# Examples: What is the State Space Size?

Often a rough upper limit is sufficient to determine how hard the search problem is.

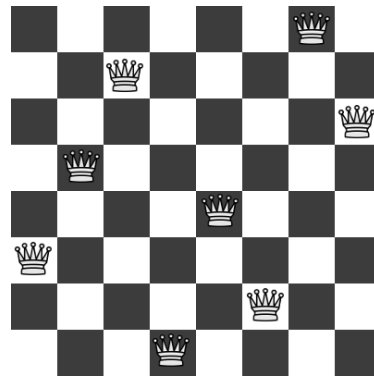


Maze

State: Positions the agent can be in.

$$n \leq |x| \times |y|$$

$n$  = # white squares.



8-queens problem

State: Board

Action: Move one queen at a time

Each square can have a queen:

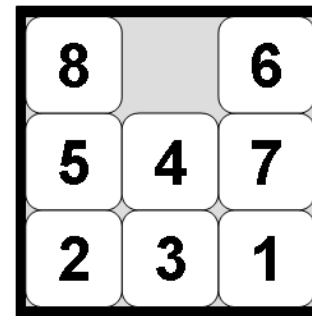
$$n < 2^{64} \approx 1.8 \times 10^{19}$$

Each queen chooses a square:

$$n < 64^8 \approx 2.9 \times 10^{14}$$

Choose 8 squares:

$$n = \binom{64}{8} \approx 4.4 \times 10^9$$



8-puzzle problem

State: Tile

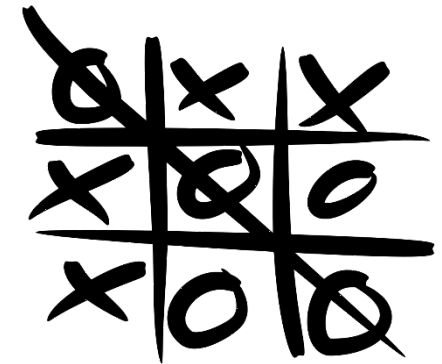
arrangement

All arrangements of 9 elements.

$$n \leq 9!$$

Half is unreachable:

$$n = \frac{9!}{2} = 181,440$$



Tic-tac-toe

State: Board

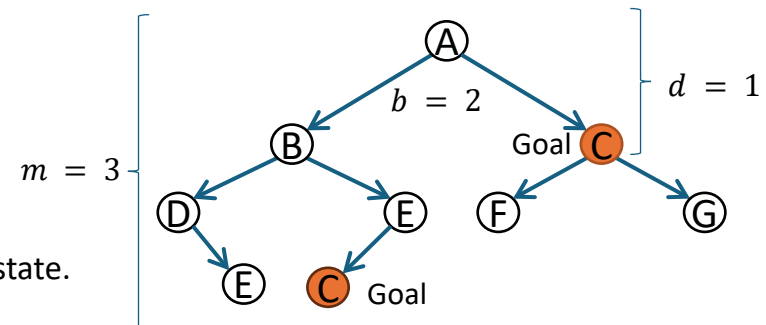
$$n < 3^9 = 19,683$$

Many boards are not legal (e.g., all x's)

The actual number can be obtained by a depth-first traversal of the game tree.

# Estimating the Search Tree Size

- Instead of estimating the state space size, it is often more useful to estimate the number of searched nodes in the search tree.
- This is especially important when **redundant paths are not eliminated**, where one state can be represented by multiple nodes.
- We can base the estimation on the search problem description:
  - initial state
  - Actions
  - transition function.
- Used metrics are:
  - $b$ : maximum branching factor of the search tree  
max. number of available actions.
  - $m$ : maximal tree depth  
length of the longest path with loops removed.
  - $d$ : depth of the optimal solution  
min. length of the path from the initial state to a solution state.
- The number of searched nodes is then a function of  $b, m$  and  $d$ .



$$n = f(b, m, d) \Rightarrow O(f(b, m, d))$$

# Example: What is the Search Complexity?

- $b$ : maximum branching factor  
= max. number of available actions?

3

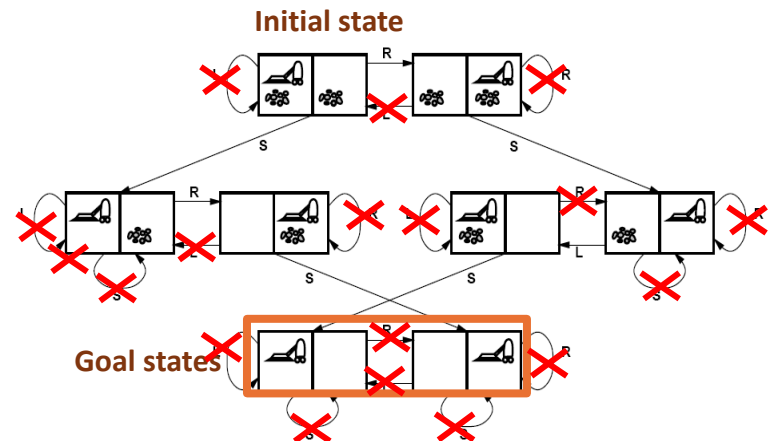
- $m$ : the number of actions in longest path? Without loops!

4

- $d$ : min. depth of the optimal solution?

3

## State Space with Transition Model



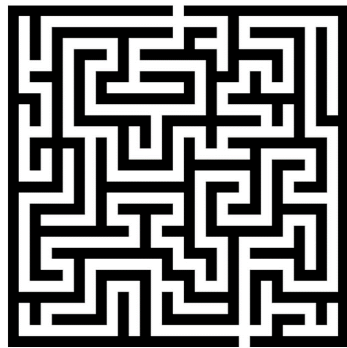
✗ Make sure it is a tree!



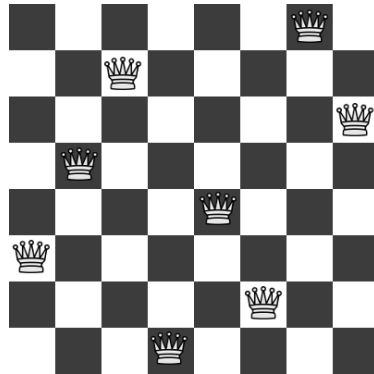
# Examples: What is the Search Complexity?

$b$ : maximum branching factor  
 $m$ : max. depth of tree  
 $d$ : depth of the optimal solution

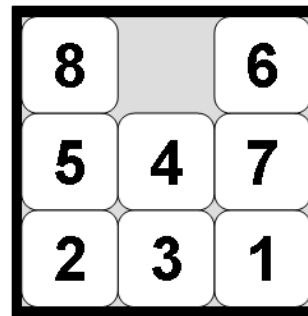
Often a rough upper limit is sufficient to determine how hard the search problem is.



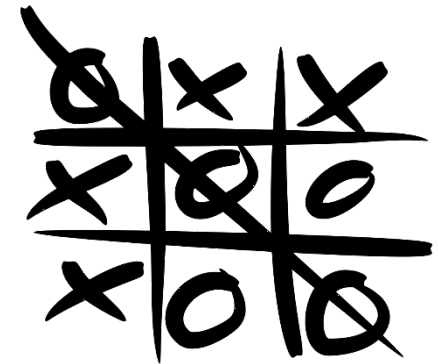
Maze



8-queens problem



8-puzzle problem



Tic-tac-toe

$b =$

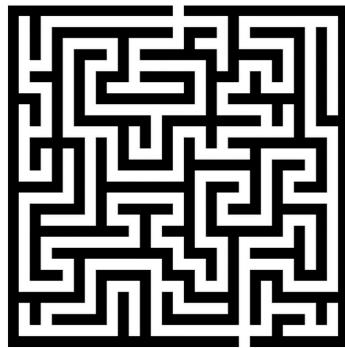
$m =$

$d =$

# Examples: What is the Search Complexity?

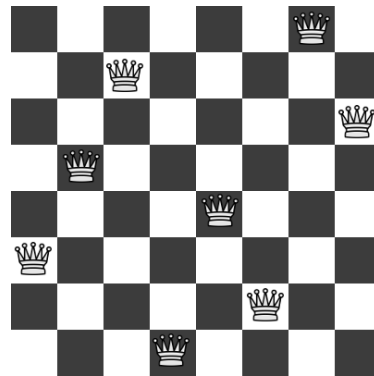
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Often a rough upper limit is sufficient to determine how hard the search problem is.



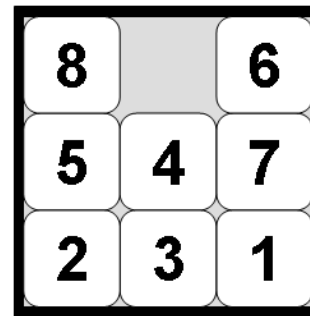
Maze

$b = 4$  actions  
 $m$  = longest path to the goal or a dead end (bounded by  $x \times y$ )  
 $d$  = shortest path to the goal (bounded by  $x \times y$ )



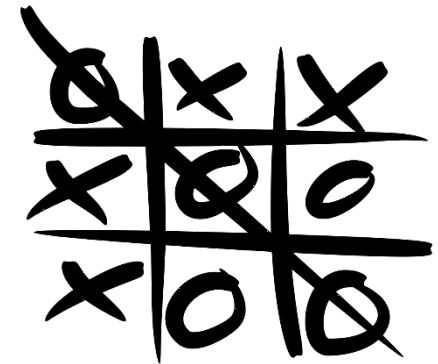
8-queens problem

Action: Move one queen at a time  
 $b = 8 \times (64 - 8) = 448$   
 $m$  = We may have to try all:  $\binom{64}{8} \approx 4.4 \times 10^9$   
 $d$  = move each queen in the right spot = 8



8-puzzle problem

$b = 4$  actions to move the empty tile.  
 $m$  = Try all  $O(9!)$   
 $d$  = ??? We need to solve the problem to know.



Tic-tac-toe

$b = 9$  actions for the first move.  
 $m = 9$   
 $d = 9$  (if both play optimal)



# Things to Remember

- Time and space complexity of search algorithms determine if we can implement a tree search solution!
- We can estimate the complexity by the following methods:
  1. Estimate the **state space size** using a factored state representation
  2. Estimate the **search tree size** using branching factor and tree depth.
- If each node represents exactly one state then both estimates will be equivalent. We will learn soon when this is or is not the case.
- We typically calculate an estimate of the actual size, or we use the Big-O notation if we are interested in how the problem scales with size.

# Uninformed Search



Breadth-First Search

# Uninformed Search Strategies

The search algorithm/planning agent is **not provided with information about how close a state is to the goal state**.

It can only use

- the labels of the **atomic states** and
- the transition function.

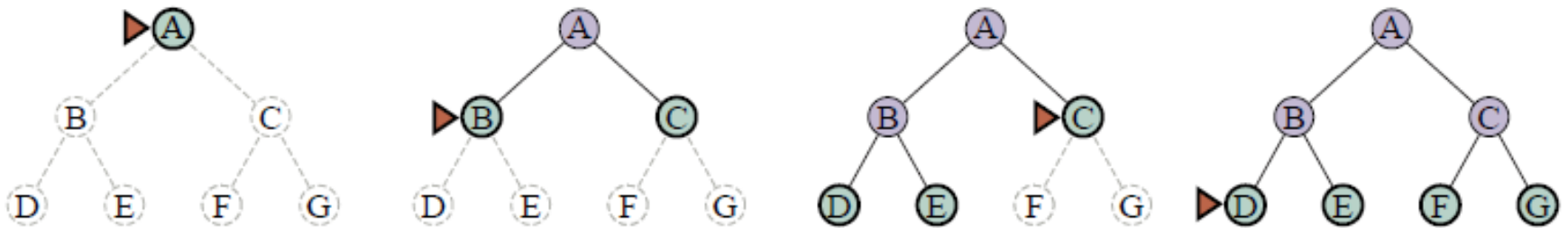
**Idea:** blindly search, following a simple strategy, until the goal state is reached.

**Search strategies:**

- **Breadth-first search strategy:** BFS and uniform-cost search
- **Depth-first search strategy:** DFS and Iterative deepening search

# Breadth-First Search (BFS)

**Expansion rule:** Expand shallowest unexpanded node in the frontier (=FIFO).



**Figure 3.8** Breadth-first search on a simple binary tree. At each stage, the node to be expanded next is indicated by the triangular marker.

## Data Structures

- **Frontier** data structure: holds references to the green nodes (green) and is implemented as a FIFO **queue**.
- **Reached** data structure: holds references to all visited nodes (gray and green) and is used to prevent visiting nodes more than once (cycle and redundant path checking).
- Builds a **complete tree** with links between parent and child.

# Implementation: Breadth-First Search

```
function BREADTH-FIRST-SEARCH(problem) returns a solution node or failure  
  node  $\leftarrow$  NODE(problem.INITIAL)  
  if problem.IS-GOAL(node.STATE) then return node  
  frontier  $\leftarrow$  a FIFO queue, with node as an element  
  reached  $\leftarrow$  {problem.INITIAL}  
  while not IS-EMPTY(frontier) do  
    node  $\leftarrow$  POP(frontier)  
    for each child in EXPAND(problem, node) do  
      s  $\leftarrow$  child.STATE  
      if problem.IS-GOAL(s) then return child  
      if s is not in reached then  
        add s to reached  
        add child to frontier  
  return failure
```

**Expand** adds the next level below node to the frontier.

**reached** makes sure we do not visit nodes twice (e.g., in a cycle or other redundant path). Fast lookup is important.

# Implementation: Expanding the Search Tree

- AI tree search creates the search tree while searching.
- The EXPAND function tries all available actions in the current node using the **transition function** (RESULTS).
- It returns a **list of child nodes** for the frontier.

```
function EXPAND(problem, node) yields nodes  
  s ← node.STATE  
  for each action in problem.ACTIONS(s) do  
    s' ← problem.RESULT(s, action)  
    cost ← node.PATH-COST + problem.ACTION-COST(s, action, s')  
    yield NODE(STATE=s', PARENT=node, ACTION=action, PATH-COST=cost)
```

Transition  
function

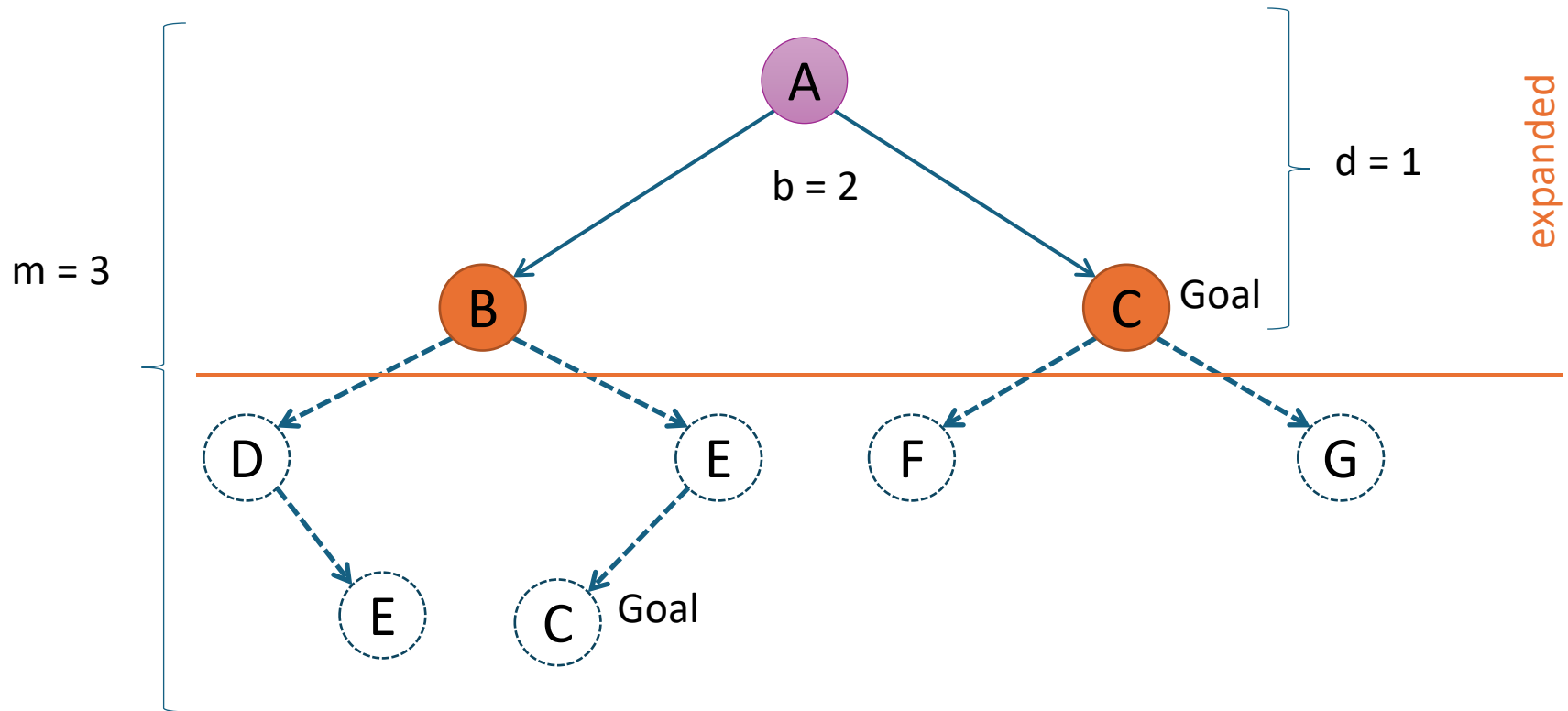
Yield (generator function) can also be  
implemented by returning a list of nodes.



# Time and Space Complexity

## Breadth-First Search

$d$ : depth of the optimal solution  
 $m$ : max. depth of tree  
 $b$ : maximum branching factor



All paths to the depth of the goal are expanded. The search tree size is

$$1 + b + b^2 + \dots + b^d \Rightarrow O(b^d)$$

# Properties of Breadth-First Search

- **Complete?**

Yes

*d*: depth of the optimal solution  
*m*: max. depth of tree  
*b*: maximum branching factor

- **Optimal?**

Yes – if cost is the same per step (action). Otherwise: Use uniform-cost search.

- **Time?**

Number of nodes created:  $O(b^d)$

- **Space?**

Stored nodes:  $O(b^d)$

**Note:**

- In AI, the large space complexity is usually a bigger problem than time!

# Uniform-cost Search

## (= Dijkstra's Shortest Path Algorithm)

- **Expansion rule:** Expand node in the frontier with **the least path cost** from the initial state.
- Implementation: **best-first search** where the frontier is a **priority queue** ordered by lower  $f(n) = \text{path cost}$  (cost of all actions starting from the initial state).
- Breadth-first search is a special case when all step costs being equal, i.e., each action costs the same!

- **Complete?**

Yes, if all step cost is greater than some small positive constant  $\epsilon > 0$

$d$ : depth of the optimal solution  
 $m$ : max. depth of tree  
 $b$ : maximum branching factor

- **Optimal?**

Yes – nodes expanded in increasing order of path cost

- **Time?**

Expands all nodes with path cost  $c \leq C^*$  (cost of optimal solution) leading to  $O(b^{1+C^*/\epsilon})$  for the number of nodes.

Note: This can be greater than BFS's  $O(b^d)$ : the search can explore long paths consisting of small steps before exploring shorter paths consisting of larger steps.

- **Space?**

$O(b^{1+C^*/\epsilon})$

See [Dijkstra's algorithm on Wikipedia](#)

# Implementation: Best-First Search Strategy

**Note:** This generalizes Breadth-First-Search

```
function UNIFORM-COST-SEARCH(problem) returns a solution node, or failure  
return BEST-FIRST-SEARCH(problem, PATH-COST)
```

Path cost with a cost of 1 per action makes the priority queue behave like a FIFO queue!

```
function BEST-FIRST-SEARCH(problem, f) returns a solution node or failure  
  node  $\leftarrow$  NODE(STATE=problem.INITIAL)  
  frontier  $\leftarrow$  a priority queue ordered by f, with node as an element  
  reached  $\leftarrow$  a lookup table, with one entry with key problem.INITIAL and value node  
  while not IS-EMPTY(frontier) do  
    node  $\leftarrow$  POP(frontier)  
    if problem.IS-GOAL(node.STATE) then return node  
    for each child in EXPAND(problem, node) do  
      s  $\leftarrow$  child.STATE  
      if s is not in reached or child.PATH-COST < reached[s].PATH-COST then  
        reached[s]  $\leftarrow$  child  
        add child to frontier  
  return failure
```

The order for expanding the frontier is determined by  $f(n)$  = path cost from the initial state to node  $n$ .

This check is added to BFS! It visits a node again if it can be reached by a better (cheaper) path.

See BFS for function EXPAND.

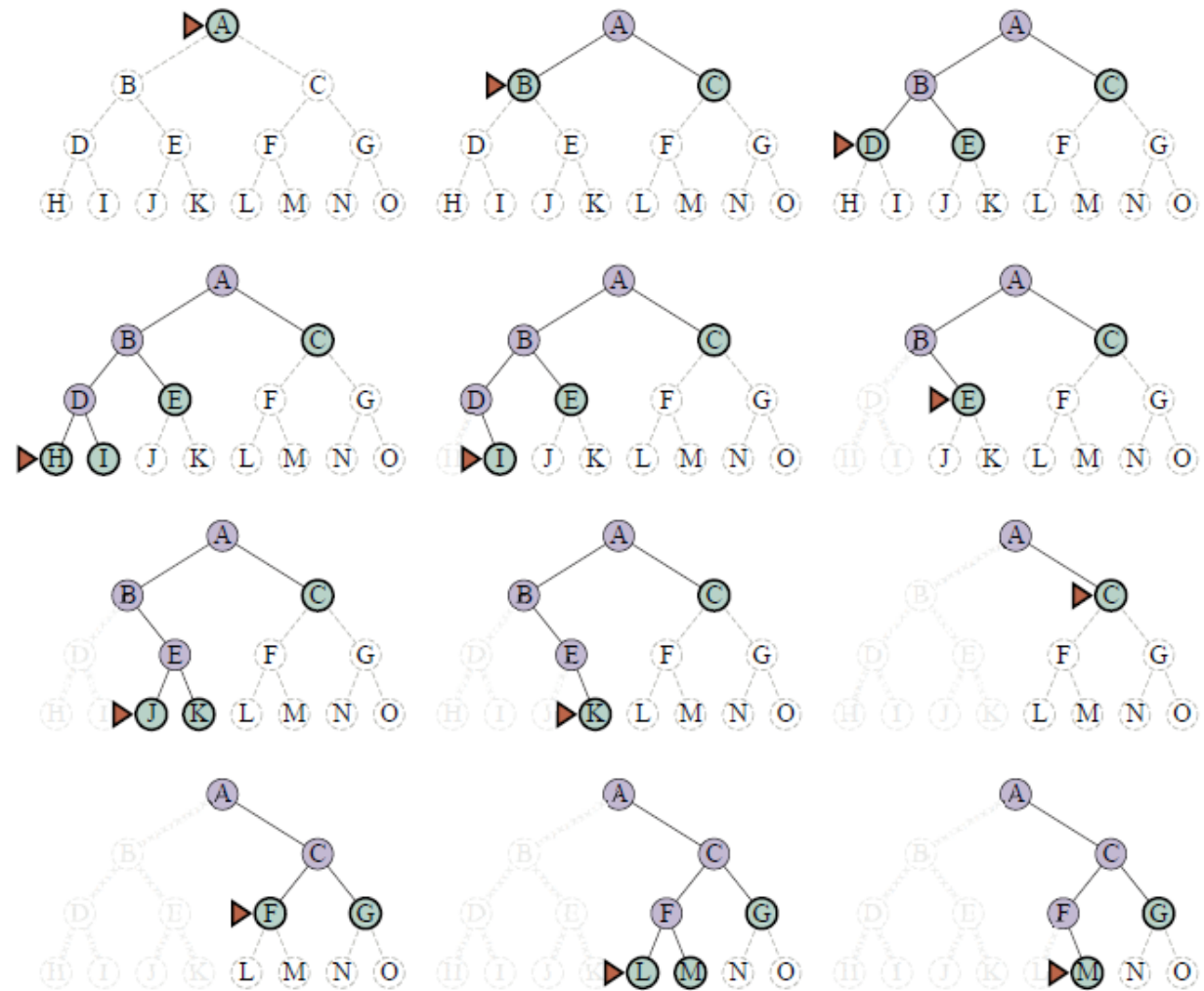
# Uninformed Search



Depth-First Search

# Depth-First Search (DFS)

- **Expansion rule:** Expand deepest unexpanded node in the frontier (last added).
- **Frontier: stack (LIFO)**
- **No reached data structure:** forgets completely explored subtrees.
- **Needs Cycle checking:** don't expand nodes that are already in the current path to the root node.
- **Cannot avoid redundant paths:** Leads to multiple nodes representing the same state and replicated work.



**Figure 3.11** A dozen steps (left to right, top to bottom) in the progress of a depth-first search on a binary tree from start state A to goal M. The frontier is in green, with a triangle marking the node to be expanded next. Previously expanded nodes are lavender, and potential future nodes have faint dashed lines. Expanded nodes with no descendants in the frontier (very faint lines) can be discarded.

# Implementation: DFS

- DFS could be implemented like BFS/Best-first search, just taking the last element from the frontier (LIFO). However, to reduce the space complexity to  $O(bm)$ , **no reached data structure can be used!**
- Options:
  - **Iterative implementation:** Build the tree, and abandoned branches are removed from memory. Cycle checking is only done against the current path. This is similar to Backtracking search.
  - Recursive implementation: Cycle checking is an issue because the current path is stored in the function call stack, which is not accessible to the function. An additional data structure that contains the nodes in the current path can be used.

DFS uses  $\ell = \infty$

```
function DEPTH-LIMITED-SEARCH(problem,  $\ell$ ) returns a node or failure or cutoff
  frontier  $\leftarrow$  a LIFO queue (stack) with NODE(problem.INITIAL) as an element
  result  $\leftarrow$  failure
  while not IS-EMPTY(frontier) do
    node  $\leftarrow$  POP(frontier)
    if problem.IS-GOAL(node.STATE) then return node
    if DEPTH(node) >  $\ell$  then
      result  $\leftarrow$  cutoff
    else if not IS-CYCLE(node) do
      for each child in EXPAND(problem, node) do
        add child to frontier
  return result
```

Memory management: remove nodes for abandoned branches here!

**Cycles:** Prevent cycles by checking against the current path. We also need to ensure that the frontier does not contain the same state more than once.

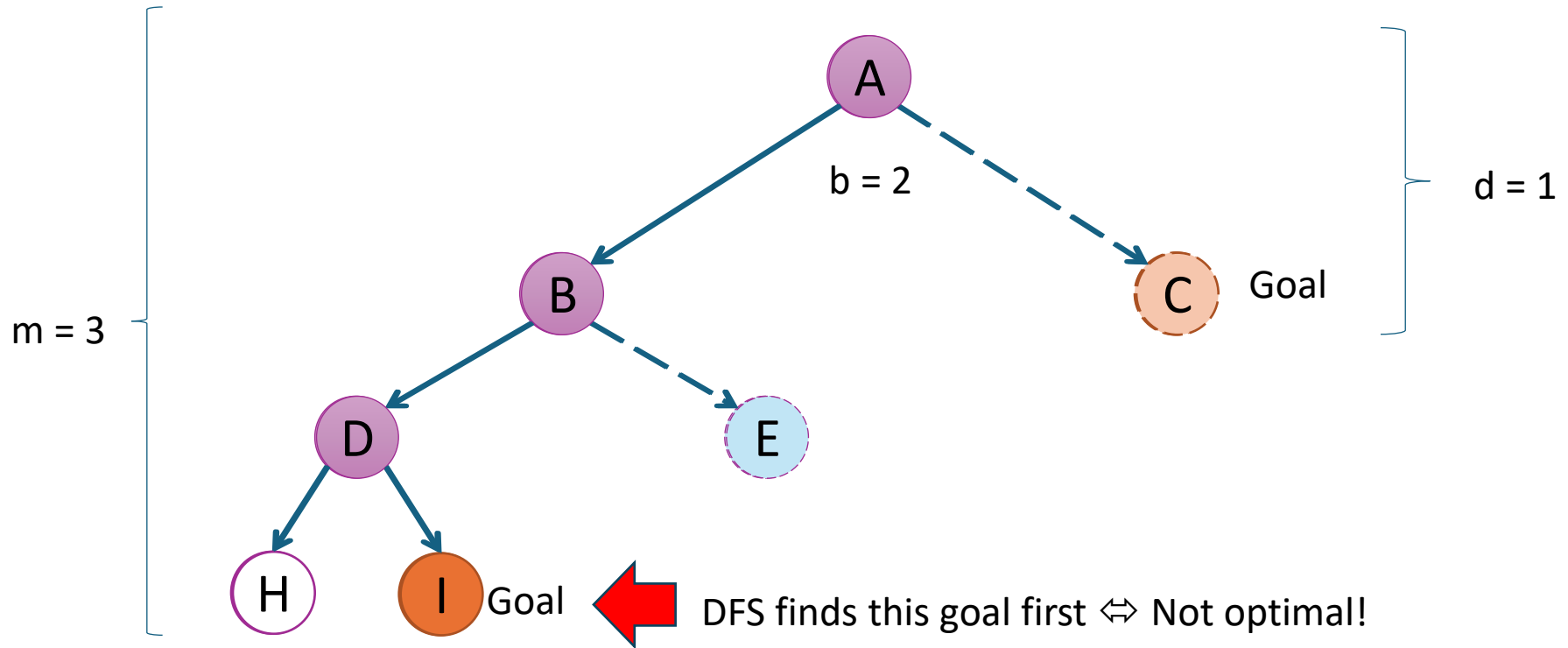
**Redundant paths:** We cannot prevent other redundant paths.

See BFS for function EXPAND.

# Time and Space Complexity

## Depth-First Search

$d$ : depth of the optimal solution  
 $m$ : max. depth of tree  
 $b$ : maximum branching factor



- Time:  $O(b^m)$  – worst case is expanding all paths.
- Space:  $O(bm)$  - if it only stores the frontier nodes and the current path.



# Properties of Depth-First Search

- **Complete?**

- In finite search spaces, cycles are avoided by checking for repeated states along the path.
- Incomplete in infinite search spaces.

- **Optimal?**

No – returns the first solution it finds.

$d$ : depth of the optimal solution  
 $m$ : max. depth of tree  
 $b$ : maximum branching factor

- **Time?**

The worst case is to reach a solution at maximum depth  $m$  in the last path:

$$O(b^m)$$

Terrible compared to BFS if  $m \gg d$ .

- **Space?**

$O(bm)$  is **linear in max. tree depth  $m$**  which is very good but only achieved if **no reached data structure and memory management** is used!

Cycles can be broken but redundant paths cannot be checked.

# Iterative Deepening Search (IDS)

Can we

- **get DFS's good memory footprint,**
- **avoid infinite cycles, and**
- **preserve BFS's optimality guaranty?**

Use depth-restricted DFS and gradually increase the depth.

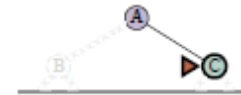
1. Check if the root node is the goal.
2. Do a DFS searching for a path of length 1
3. If goal not found, do a DFS searching for a path of length 2
4. If goal not found, do a DFS searching for a path of length 3
5. ...

# Iterative Deepening Search (IDS)

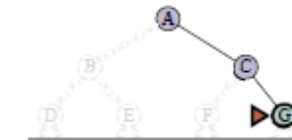
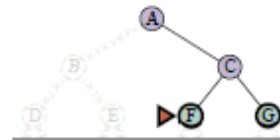
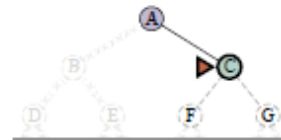
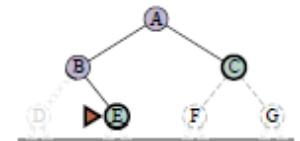
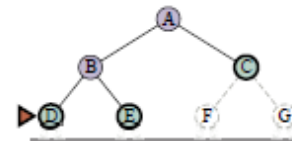
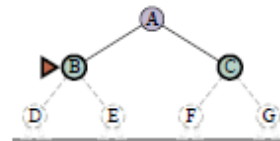
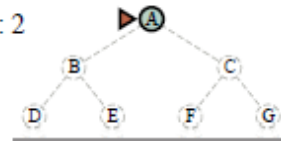
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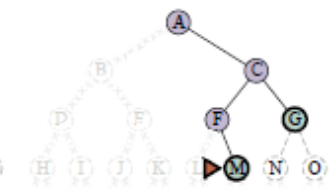
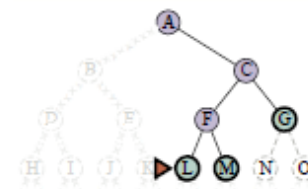
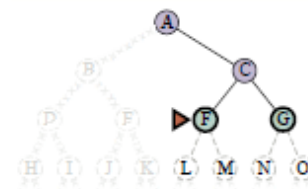
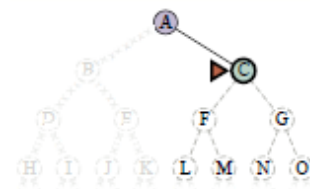
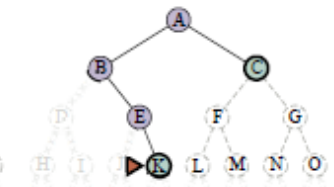
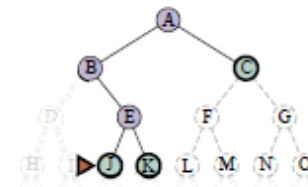
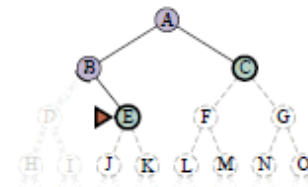
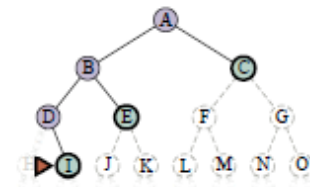
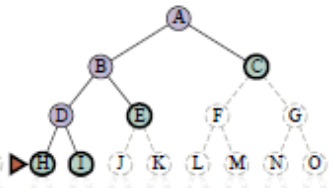
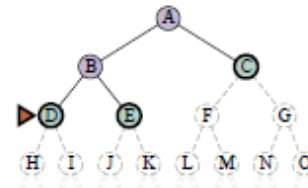
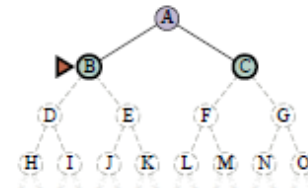
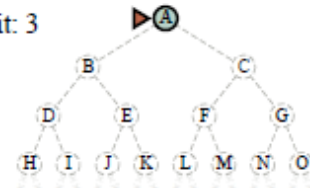
limit: 1



limit: 2



limit: 3



# Implementation: IDS

```
function ITERATIVE-DEEPENING-SEARCH(problem) returns a solution node or failure  
  for depth = 0 to  $\infty$  do  
    result  $\leftarrow$  DEPTH-LIMITED-SEARCH(problem, depth)  
    if result  $\neq$  cutoff then return result
```

```
function DEPTH-LIMITED-SEARCH(problem,  $\ell$ ) returns a node or failure or cutoff  
  frontier  $\leftarrow$  a LIFO queue (stack) with NODE(problem.INITIAL) as an element  
  result  $\leftarrow$  failure  
  while not IS-EMPTY(frontier) do  
    node  $\leftarrow$  POP(frontier)  
    if problem.IS-GOAL(node.STATE) then return node  
    if DEPTH(node) >  $\ell$  then  
      result  $\leftarrow$  cutoff  
    else if not IS-CYCLE(node) do  
      for each child in EXPAND(problem, node) do  
        add child to frontier  
  return result
```

See BFS for function EXPAND.

# Properties of Iterative Deepening Search

- **Complete?**

Yes

$d$ : depth of the optimal solution  
 $m$ : max. depth of tree  
 $b$ : maximum branching factor

- **Optimal?**

Yes, if step cost = 1 (like BFS)

- **Time?**

Consists of rebuilding trees up to  $d$  times

$db + (d-1)b^2 + \dots + 1b^d = O(b^d) \Leftrightarrow$  Slower than BFS, but the same complexity class!

- **Space?**

$O(bd) \Leftrightarrow$  linear space. Even less than DFS since  $m \leq d$ . Cycles need to be handled by the depth-limited DFS implementation.

**Note:** IDS produces the same result as BFS but trades **much better space complexity** for worse run time.

This makes IDS/DFS the  
workhorse of AI.



# Informed Search

## Introduction

# Informed Search

AI search problems typically have a very large search space. We would like to improve efficiency by **expanding as few nodes as possible**.

**Idea:** The agent can use **additional information** in the form of “hints” about what promising states are to explore first. These hints are derived from

- **information** the agent has (e.g., a map with the goal location marked) or
- **percepts** coming from a sensor (e.g., a GPS sensor and coordinates of the goal).

**Method:** The agent uses a **heuristic function  $h(n)$**

- to rank nodes in the frontier based on the additional information, and
- to select the most promising node in the frontier for expansion using the **best-first search** strategy.

**Discussed algorithms:**

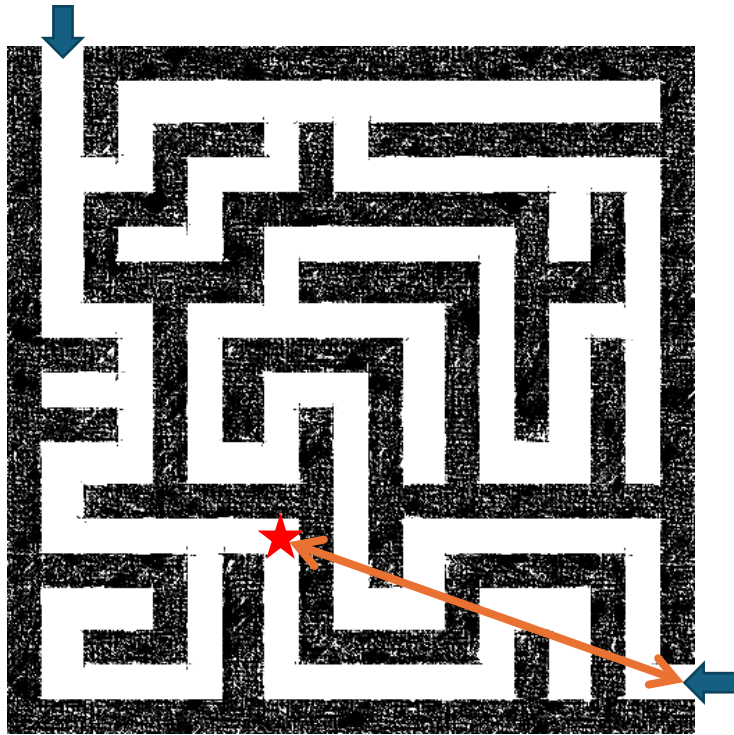
- Greedy best-first search
- A\* search

# Heuristic Function

- **Heuristic function**  $h(n)$  estimates the cost of reaching a node representing the goal state from the currently considered node  $n$ .
- Examples:

**Euclidean distance**

Start state

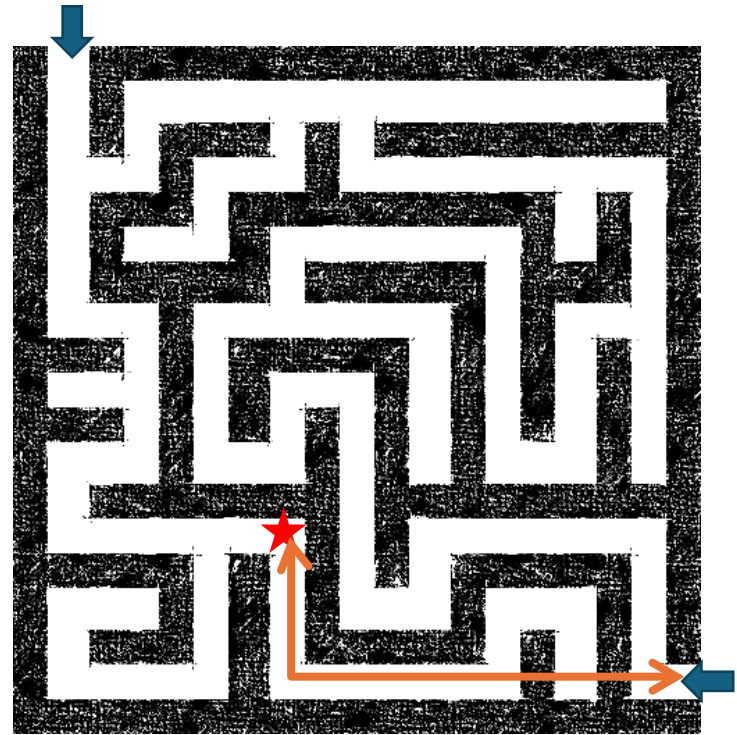


★ State for currently considered node

Goal state

**Manhattan distance**

Start state

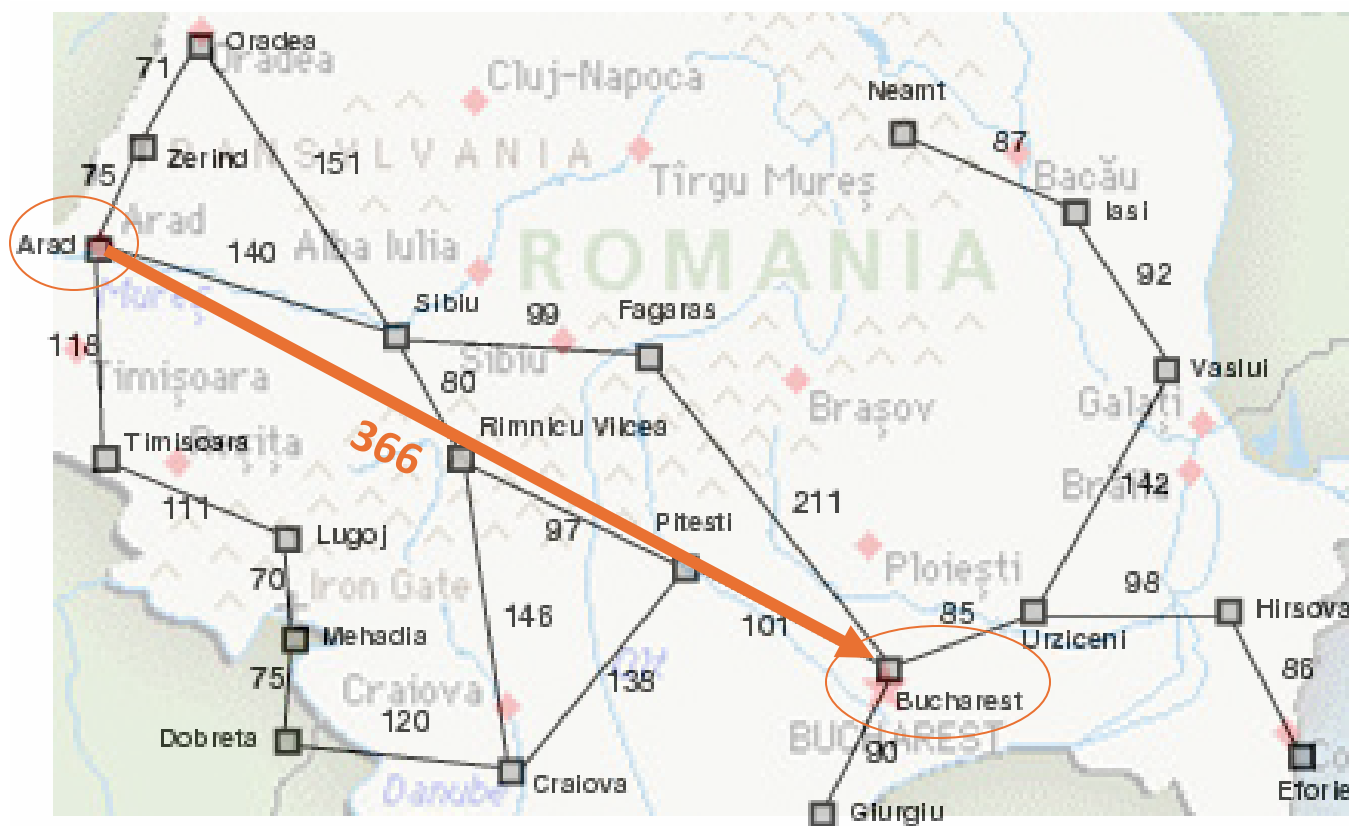


Goal state



# Heuristic for the Romania Problem

Use the map for hints: Estimate the driving distance from Arad to Bucharest using a straight-line distance on the map.

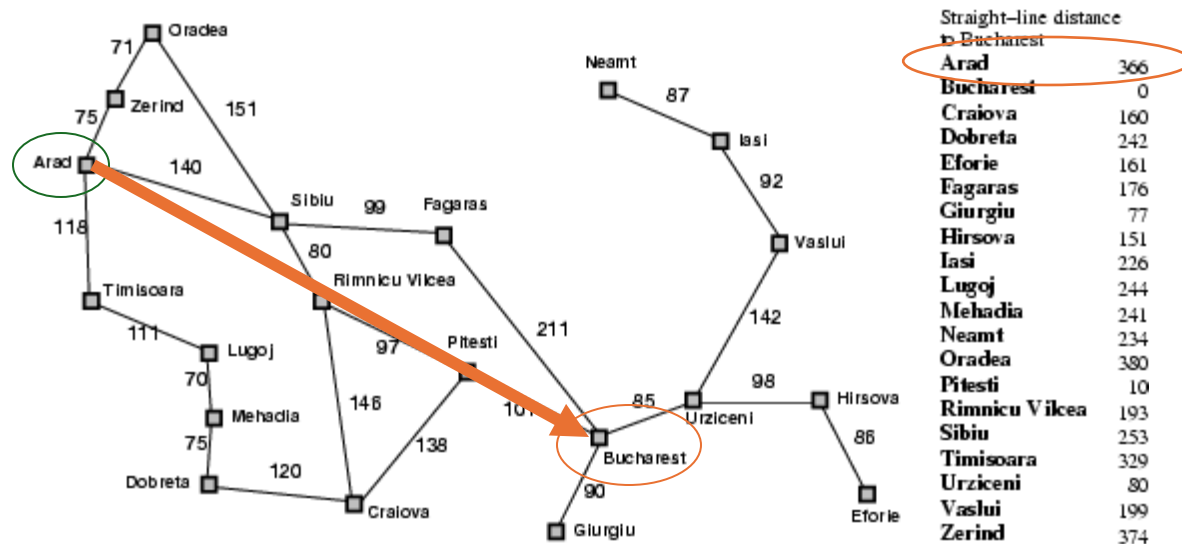


Straight-line distance to Bucharest $h(n)$	
Arad	366
Bucharest	0
Craiova	160
Dobreta	242
Eforie	161
Fagaras	176
Giurgiu	77
Hirsova	151
Iasi	226
Lugoj	244
Mehadia	241
Neamt	234
Oradea	380
Pitesti	10
Rimnicu Vilcea	193
Sibiu	253
Timisoara	329
Urziceni	80
Vaslui	199
Zerind	374

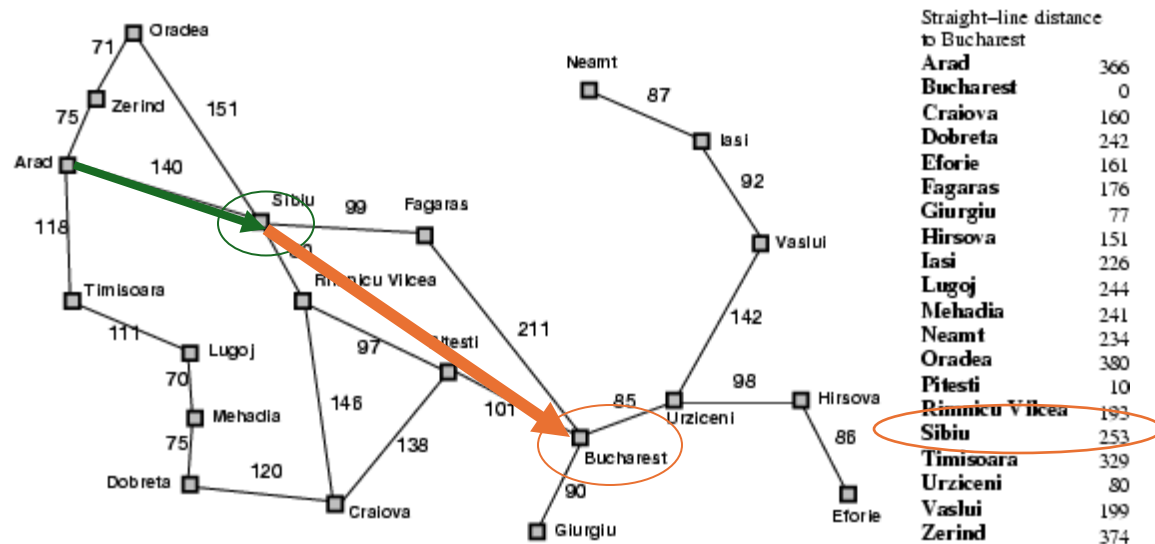
# Greedy Best-First Search Example

Expansion rule: Expand the node that has the lowest value of the heuristic function  $h(n)$

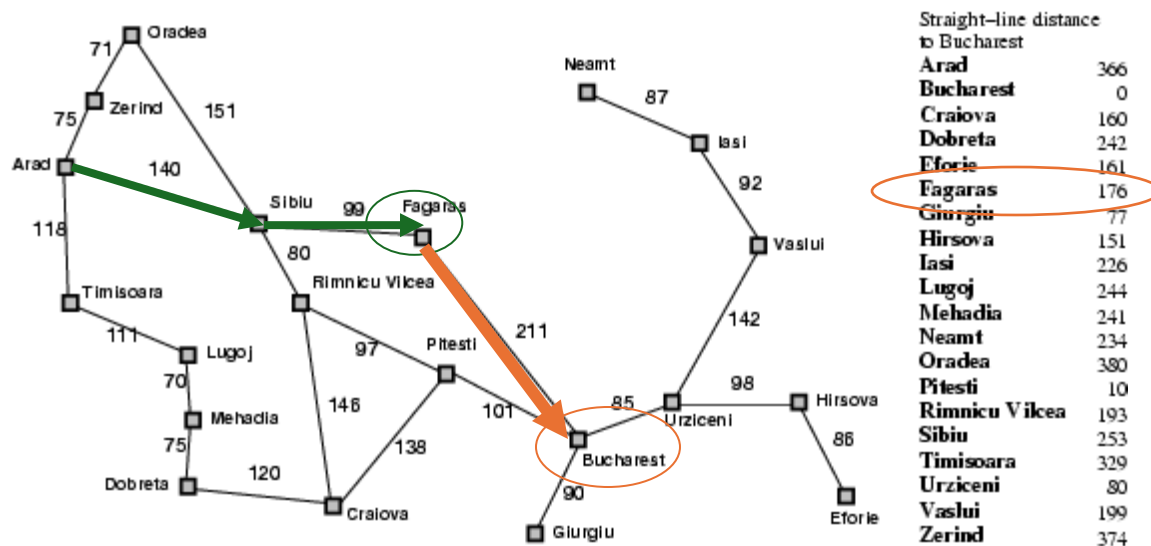
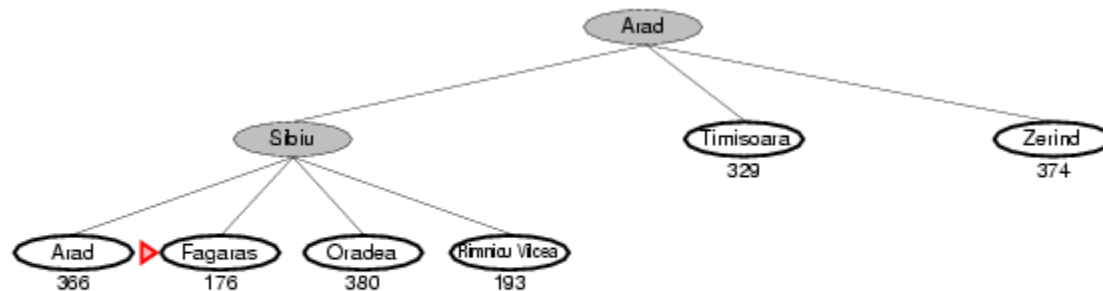
$$h(n) = \text{Arad} = 366$$



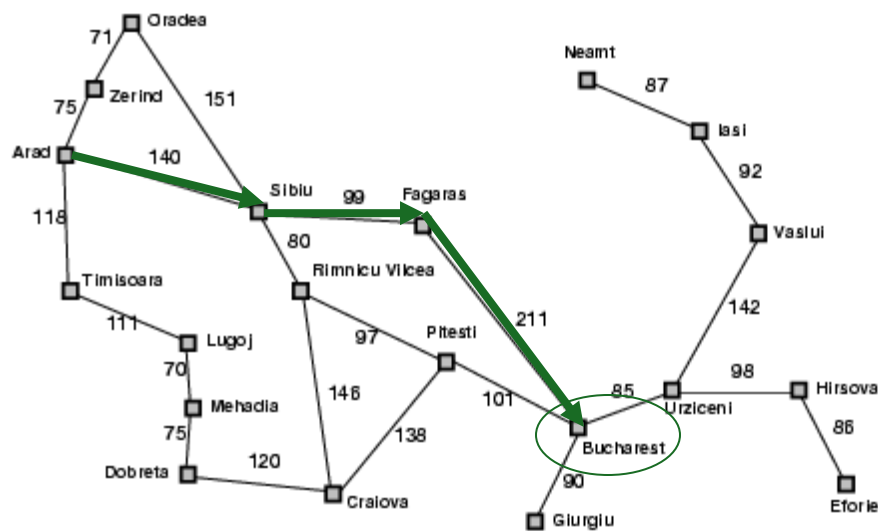
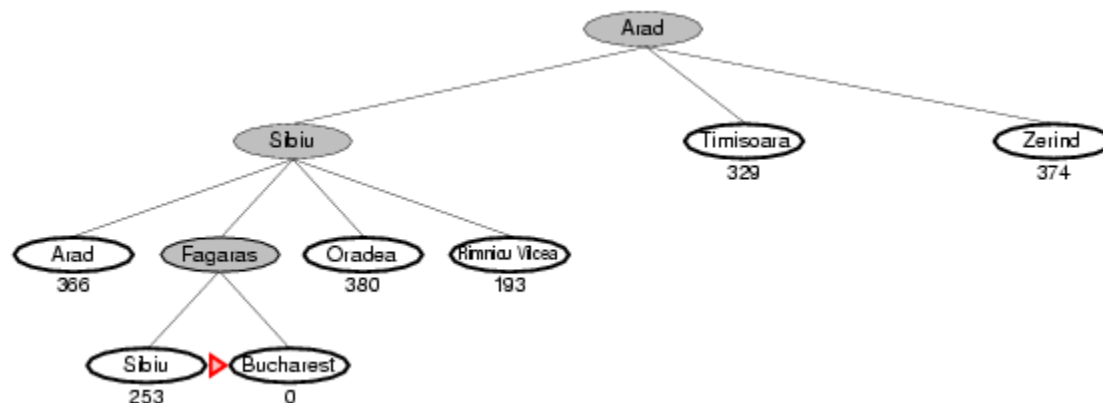
# Greedy Best-First Search Example



# Greedy Best-First Search Example



# Greedy Best-First Search Example



Straight-line distance to Bucharest

Arad	366
Bucharest	0
Craiova	160
Dobreta	242
Eforie	161
Fagaras	176
Giurgiu	77
Hirsova	151
Iasi	226
Lugoj	244
Mehadia	241
Neamt	234
Oradea	380
Pitesti	10
Rimnicu Vilcea	193
Sibiu	253
Timisoara	329
Urziceni	80
Vaslui	199
Zerind	374

Total:

$$140 + 99 + 211 = 450 \text{ miles}$$

# Properties of Greedy Best-First Search

- **Complete?**

Yes – Best-first search is complete in finite spaces.

- **Optimal?**

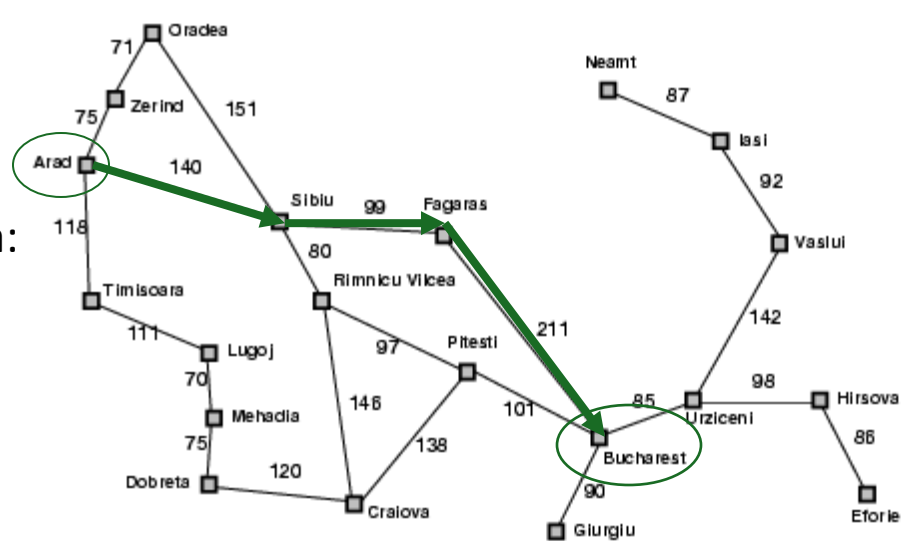
No

Total:

$$140 + 99 + 211 = 450 \text{ miles}$$

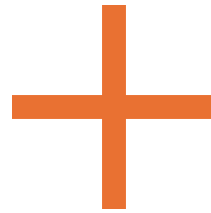
Alternative through Rimnicu Vilcea:

$$140 + 80 + 97 + 101 = 418 \text{ miles}$$



# Implementation of Greedy Best-First search

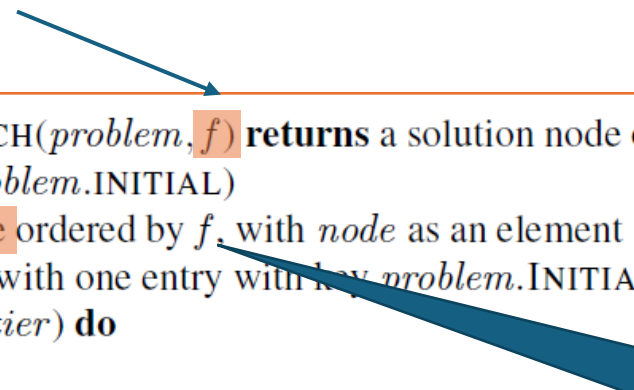
Best-First  
Search



Expand the frontier  
using  
 $f(n) = h(n)$

# Implementation of Greedy Best-First Search

Heuristic  $h(n)$  so we expand the node with the lowest estimated cost



```
function BEST-FIRST-SEARCH(problem, f) returns a solution node or failure  
  node  $\leftarrow$  NODE(STATE=problem.INITIAL)  
  frontier  $\leftarrow$  a priority queue ordered by f, with node as an element  
  reached  $\leftarrow$  a lookup table, with one entry with key problem.INITIAL and value node  
  while not IS-EMPTY(frontier) do  
    node  $\leftarrow$  POP(frontier)  
    if problem.IS-GOAL(node.STATE) then return node  
    for each child in EXPAND(problem, node) do  
      s  $\leftarrow$  child.STATE  
      if s is not in reached or child.PATH-COST < reached[s].PATH-COST then  
        reached[s]  $\leftarrow$  child  
        add child to frontier  
  return failure
```

The order for expanding the frontier is determined by  $f(n)$

See BFS for function EXPAND.



# Properties of Greedy Best-First Search

- **Complete?**

Yes – Best-first search is complete in finite spaces.

- **Optimal?**

No

$d$ : depth of the optimal solution  
 $m$ : max. depth of tree  
 $b$ : maximum branching factor

- **Time?**

Worst case:  $O(b^m) \Leftrightarrow$  like DFS

Best case:  $O(bm)$  – If  $h(n)$  is 100% accurate we only expand a single path.

- **Space?**

Same as time complexity.

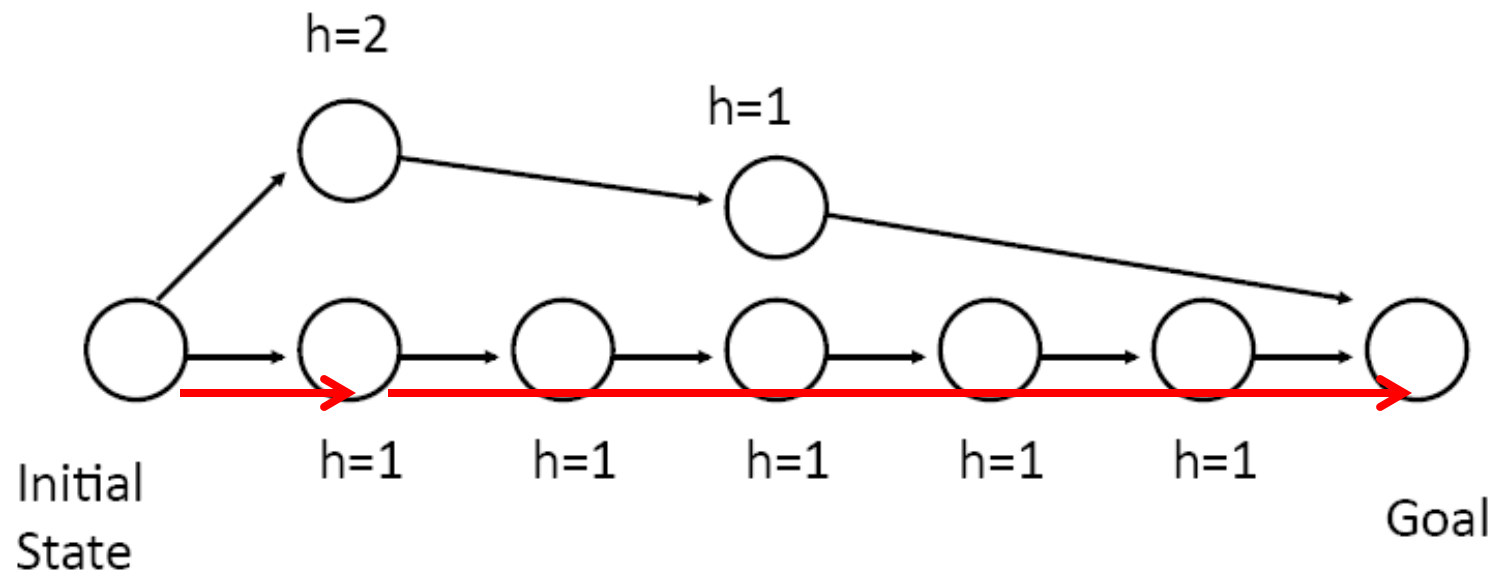


# Informed Search

## A\* Search

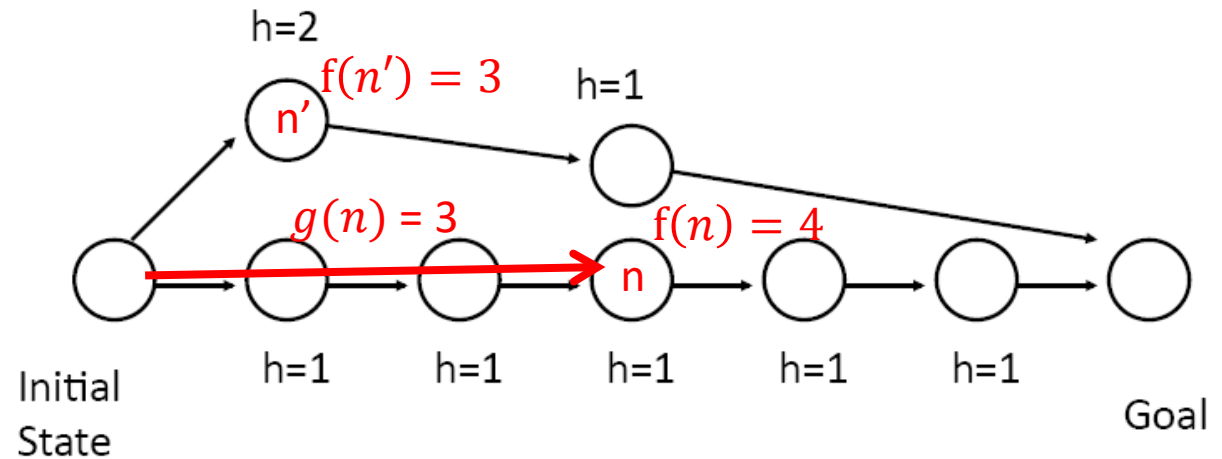
# The Optimality Problem of Greedy Best-First search

Greedy best-first search only considers the estimated cost to the goal.



$h = 1$  is better than  $h = 2$ .  
Greedy best-first will go this way  
and never reconsider!

# A\* Search



- **Idea:** Take the cost of the path to  $n$  called  $g(n)$  into account to avoid expanding paths that are already very expensive.
- The evaluation function  $f(n)$  is the estimated total cost of the path through node  $n$  to the goal:

$$f(n) = g(n) + h(n)$$

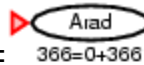
$g(n)$ : cost so far to reach  $n$  (path cost)

$h(n)$ : estimated cost from  $n$  to goal (heuristic)

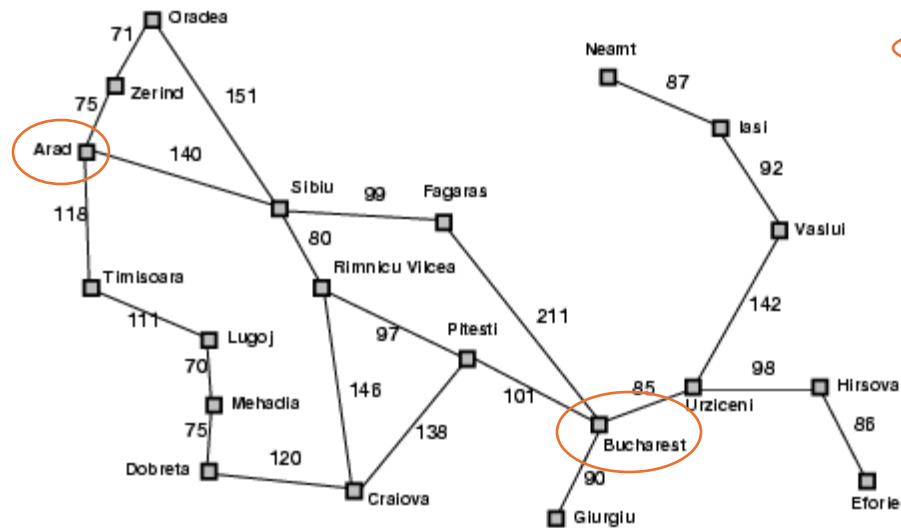
- The agent in the example above will stop at  $n$  with  $f(n) = 3 + 1 = 4$  and chose the path up with a better  $f(n') = 1 + 2 = 3$ .

**Note:** For greedy best-first search we just used  $f(n) = h(n)$ .

# A\* Search Example

Expansion rule:  $f(n) = g(n) + h(n) =$    $366=0+366$

Expand the node with  
the smallest  $f(n)$



$h(n)$

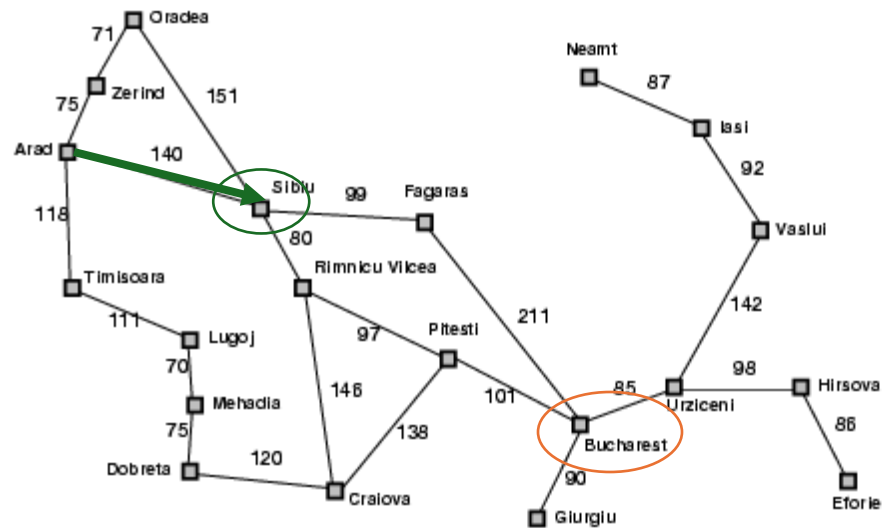
Straight-line distance  
to Bucharest

Arad	366
Bucharest	0
Craiova	160
Dobreta	242
Eforie	161
Fagaras	176
Giurgiu	77
Hirsova	151
Iasi	226
Lugoj	244
Mehadia	241
Neamt	234
Oradea	380
Pitesti	10
Rimnicu Vilcea	193
Sibiu	253
Timisoara	329
Urziceni	80
Vaslui	199
Zerind	374

# A\* Search Example



$$f(n) = g(n) + h(n)$$

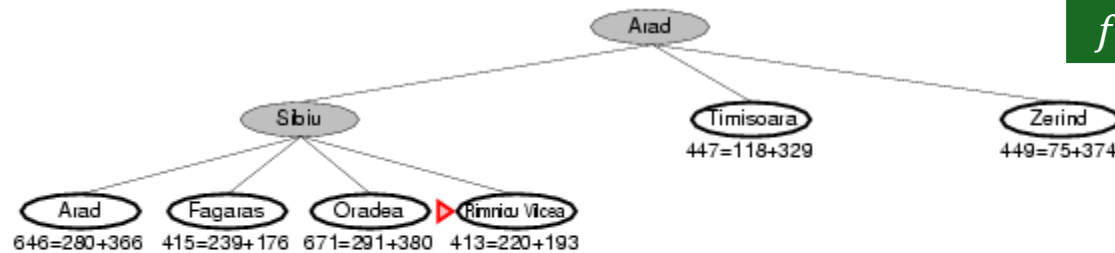


$h(n)$

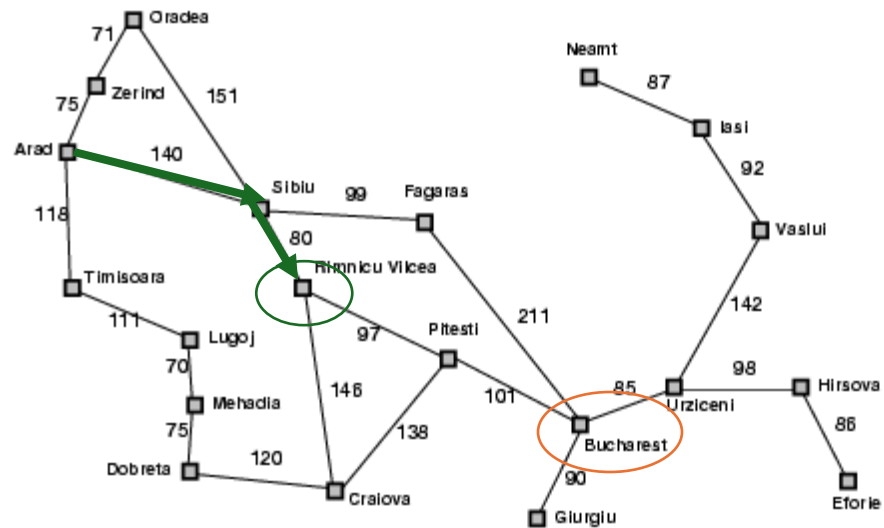
Straight-line distance  
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# A\* Search Example



$$f(n) = g(n) + h(n)$$



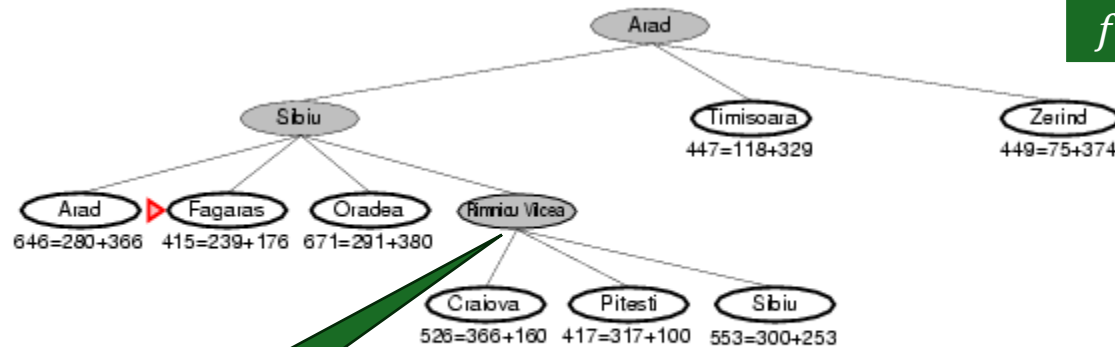
$h(n)$

Straight-line distance to Bucharest

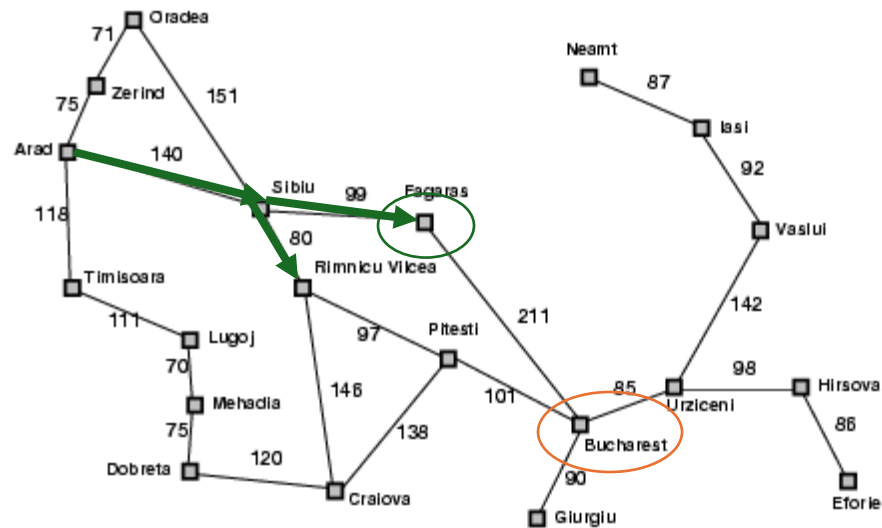
Arad	366
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# A\* Search Example

$$f(n) = g(n) + h(n)$$



Reconsiders Rimnicu Vilcea because Fagaras may have a shorter total cost to Bucharest.

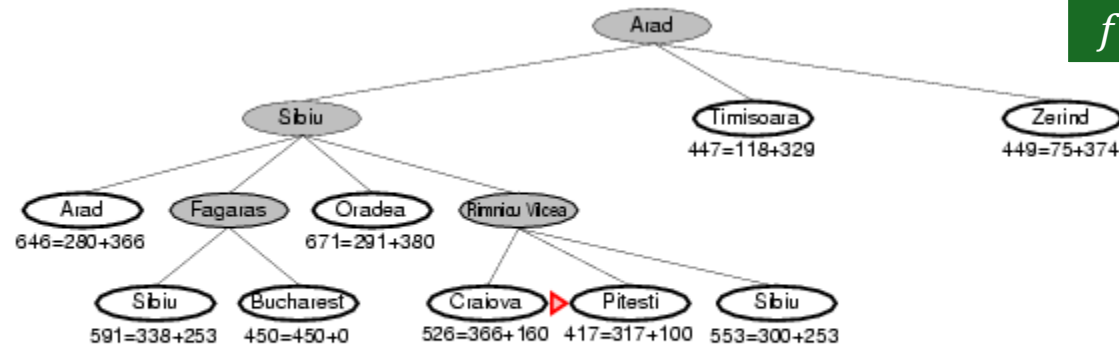


$h(n)$

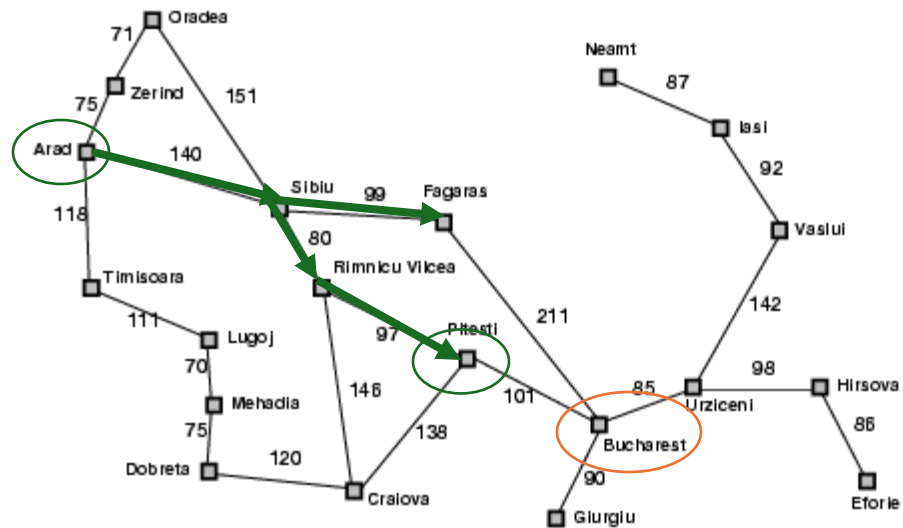
Straight-line distance to Bucharest	
Arad	366
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# A\* Search Example



$$f(n) = g(n) + h(n)$$

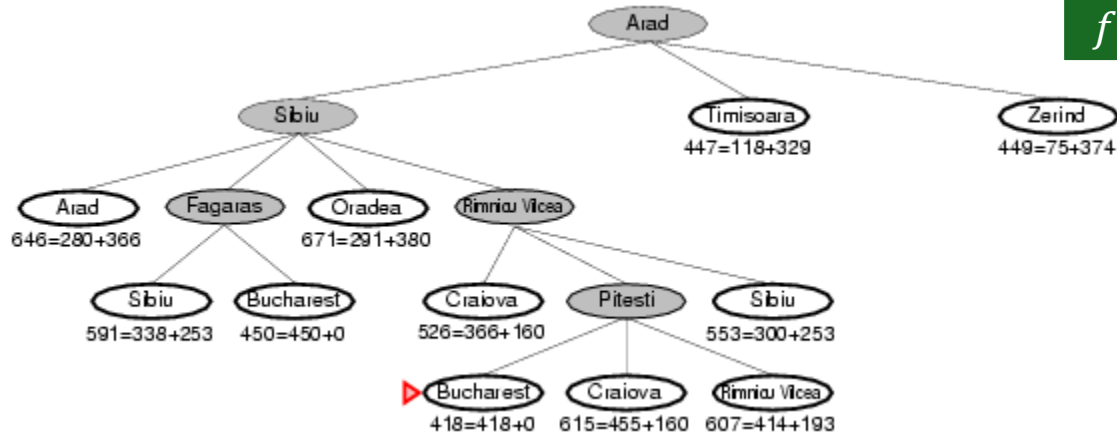


$h(n)$

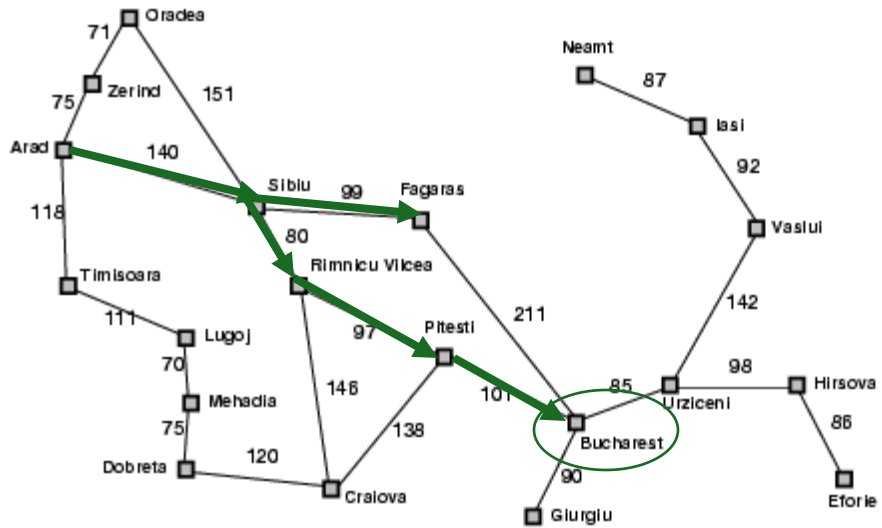
Straight-line distance  
to Bucharest

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# A\* Search Example



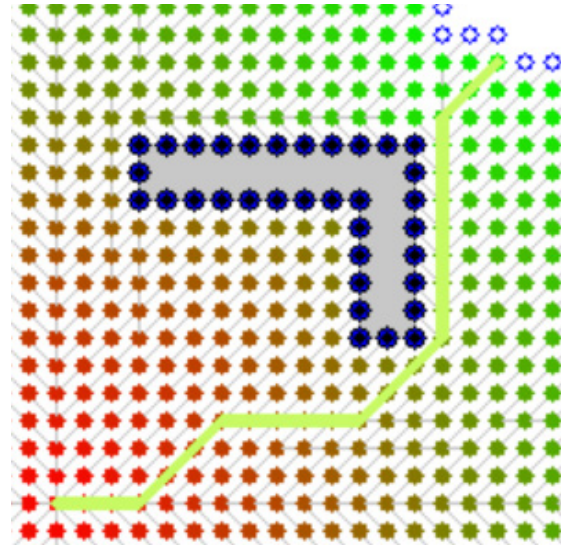
$$f(n) = g(n) + h(n)$$


$$h(n)$$
Straight-line distance  
to Bucharest

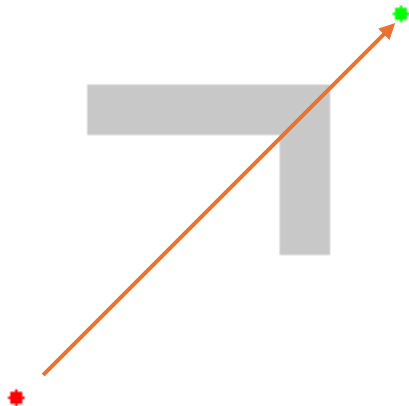
Arad	366
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# BFS vs. A\* Search

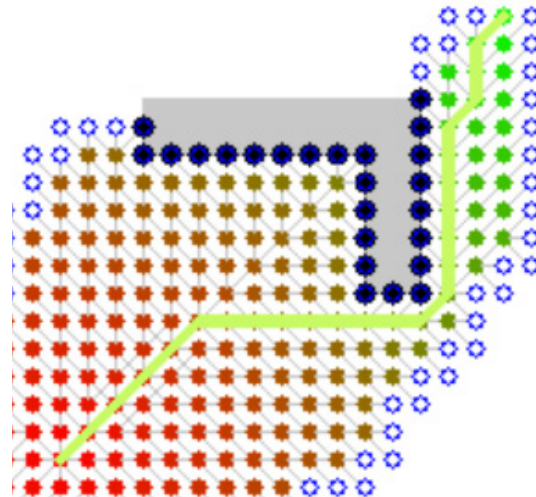
BFS



A\*



A\*

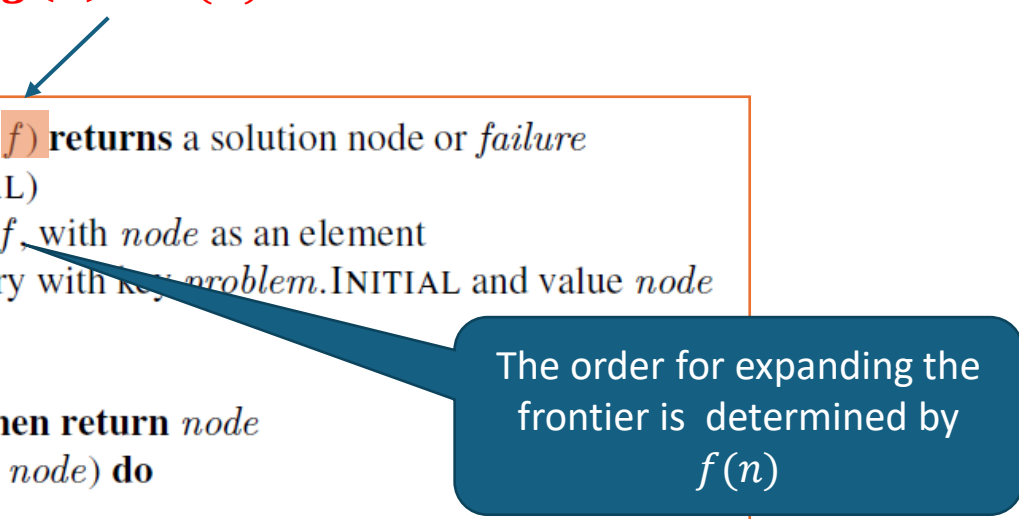


Source: [Wikipedia](https://en.wikipedia.org/wiki/Heuristic_search)

A\* Search expands  
fewer nodes than BFS!

# Implementation of A\* Search

Path cost to  $n$  + heuristic from  $n$  to goal = estimate of the total cost  
 $g(n) + h(n)$



```
function BEST-FIRST-SEARCH(problem, f) returns a solution node or failure  
  node  $\leftarrow$  NODE(STATE=problem.INITIAL)  
  frontier  $\leftarrow$  a priority queue ordered by f, with node as an element  
  reached  $\leftarrow$  a lookup table, with one entry with key problem.INITIAL and value node  
  while not IS-EMPTY(frontier) do  
    node  $\leftarrow$  POP(frontier)  
    if problem.IS-GOAL(node.STATE) then return node  
    for each child in EXPAND(problem, node) do  
      s  $\leftarrow$  child.STATE  
      if s is not in reached or child.PATH-COST < reached[s].PATH-COST then  
        reached[s]  $\leftarrow$  child  
        add child to frontier  
  return failure
```

The order for expanding the frontier is determined by  $f(n)$

See BFS for function EXPAND.

# Optimality: Admissible Heuristics

**Definition:** A heuristic  $h$  is **admissible** if for every node  $n$ ,  $h(n) \leq h^*(n)$ , where  $h^*(n)$  is the true cost to reach the goal state from  $n$ .

I.e., an admissible heuristic is a **lower bound** and never overestimates the true cost to reach the goal.

**Example:** Straight line distance never overestimates the actual road distance.

**Theorem:** If  $h$  is admissible,  $A^*$  is optimal.

# Guarantees of A\* Search

A\* is **optimally efficient**

No other tree-based search algorithm that employs the same heuristic can expand fewer nodes and still guarantee the optimal solution.

**Proof:** Any algorithm that does not expand all nodes with  $f(n) < C^*$  (the lowest cost of going to a goal node) cannot be optimal. It risks missing the optimal solution.

# Properties of A\*Search

- **Complete?**

Yes

- **Optimal?**

Yes

- **Time?**

Number of nodes for which  $f(n) \leq C^*$  in the worst case  $O(b^d)$  like BFS.

- **Space?**

Same as time complexity. This is often too high unless a very good heuristic is known.

# Iterative-Deepening A\* Search – IDA\*

- **Idea:** A\* search without a reached data structure.
- **Remember:** Regular IDA is uninformed and increases the cutoff by one after each iteration.
- IDA\* uses the cost  $f = g + h$  of a node as the cutoff. In each iteration, the cost cutoff increases slightly. It is optimal if steps are small and  $h$  is admissible.
- **Issues:**
  - By how much to increase the cutoff in each iteration.
  - Rebuilds the tree many times.
- **Other memory-bounded variants of A\* search:**
  - Recursive best-first search (RBFS) adds a  $f$ -limit to the depth-first search behavior of best-first search.
  - Simplified memory-bounded A\* (SMA\*) performs A\* till the memory is full and then drops the worst (highest  $f$ ) leaf node from memory. It can rebuild the node later if needed.



A hand holding a compass in a snowy forest. The hand is wearing a dark jacket. The compass is a round, silver-colored one with a black face and white markings. The background is a blurred, snowy forest with white branches and snow-covered ground.

# Informed Search Designing Heuristics

# Designing Heuristic Functions

7	2	4
5		6
8	3	1

Start State

	1	2
3	4	5
6	7	8

Goal State

Example heuristics for the 8-puzzle:

- $h_1(n)$  = number of misplaced tiles
- $h_1(start) = 8$

- $h_2(n)$  = total Manhattan distance (number of squares from the desired location of each tile)
- $h_2(start) = 3 + 1 + 2 + 2 + 2 + 3 + 3 + 2 = 18$

Are  $h_1$  and  $h_2$  admissible?

1 needs to move 3 positions

# Heuristics from Relaxed Problems

- A problem with fewer restrictions on the actions is called a relaxed problem.
- **The cost of an optimal solution to a relaxed problem is an admissible heuristic for the original problem. I.e., the true cost is never smaller.**
- What relaxation is used by  $h_1$  and  $h_2$ ?
  - $h_1$ : If the rules of the 8-puzzle are relaxed so that a tile can move **anywhere**, then  $h_1(n)$  gives the shortest solution.
  - $h_2$ : If the rules are relaxed so that a tile can move to **any adjacent square**, then  $h_2(n)$  gives the shortest solution.

7	2	4
5		6
8	3	1

Start State

	1	2
3	4	5
6	7	8

Goal State

$$h_1(start) = 8$$

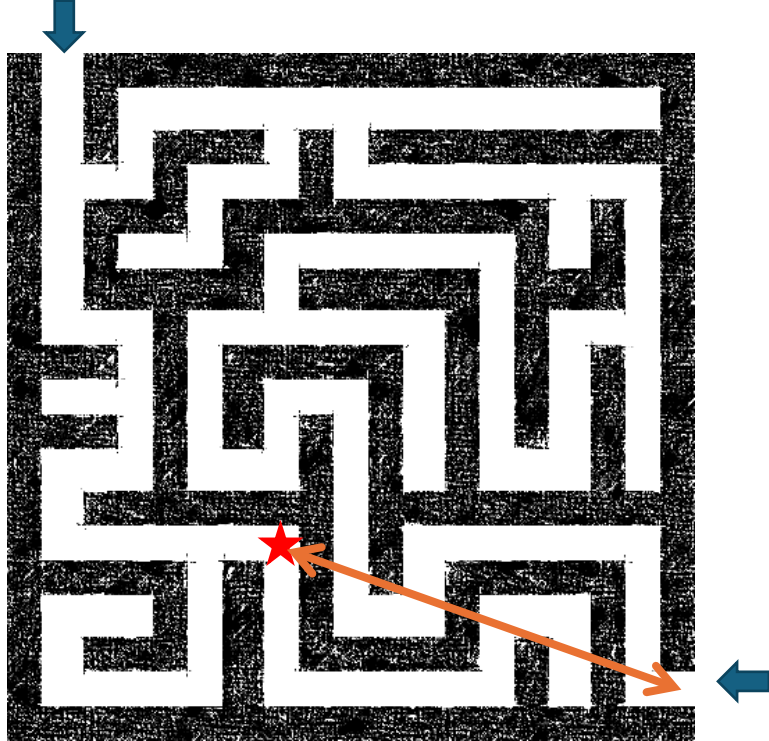
$$\begin{aligned} h_2(start) &= 3 + 1 + 2 + 2 + 2 + 3 + 3 + 2 \\ &= 18 \end{aligned}$$

# Heuristics from Relaxed Problems

**What relaxations are used in these two cases?**

**Euclidean distance**

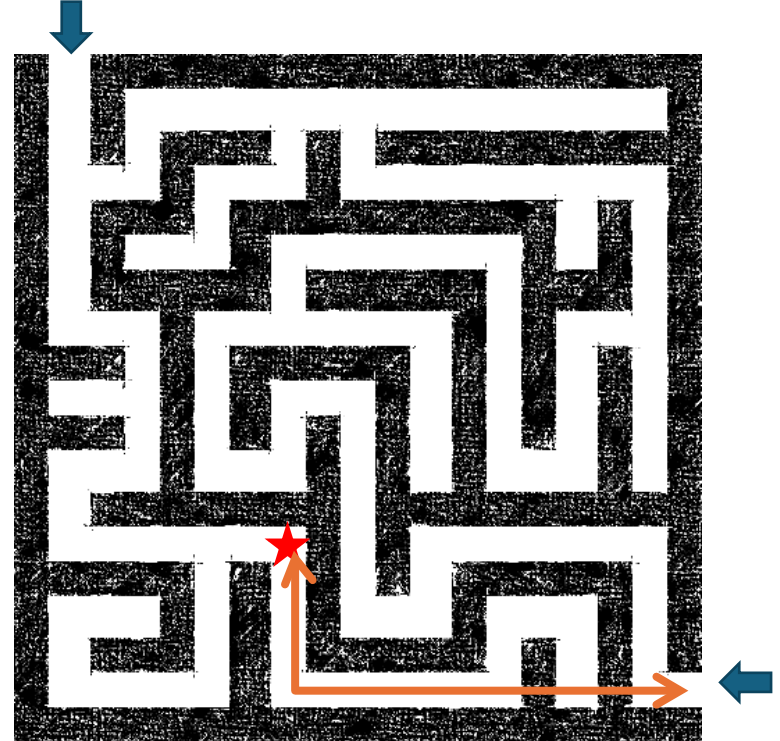
Start state



Goal state

**Manhattan distance**

Start state



Goal state

# Heuristics from Subproblems

- Let  $h_3(n)$  be the cost of getting a subset of tiles (say, 1,2,3,4) into their correct positions. The final order of the \* tiles does not matter.
- Solutions for subproblems are an admissible heuristic.
- Calculation:
  - Very small subproblems are often easy to solve.
  - We can precompute and save the exact solution cost for every or many possible subproblem instances – ***pattern database***.

*	2	4
*		*
*	3	1

Start State

	1	2
3	4	*
*	*	*

Goal State

## Dominance: What Heuristic is Better?

**Definition:** If  $h_1$  and  $h_2$  are both admissible heuristics and  $h_2(n) \geq h_1(n)$  for all  $n$ , then  $h_2$  dominates  $h_1$

Is  $h_1$  or  $h_2$  better for A\* search?

- A\* search expands every node with  $f(n) < C^* \Leftrightarrow h(n) < C^* - g(n)$
- $h_2$  is never smaller than  $h_1$ . A\* search with  $h_2$  will expand less nodes and is therefore better.

# Combining Heuristics

- Suppose we have a collection of admissible heuristics  $h_1, h_2, \dots, h_m$ , but none of them dominates the others.
- Combining them is easy:

$$h(n) = \max\{h_1(n), h_2(n), \dots, h_m(n)\}$$

- That is, always pick for each node the heuristic that is closest to the real cost to the goal  $h^*(n)$ .

# Example: Effect of Information in Search

Typical search costs for the 8-puzzle

- State space:  $\frac{9!}{2} = 1,811,440$  states
- Problem with solution at depth  $d = 12$ 
  - IDS = 3,644,035 nodes
  - $A^*(h_1) = 227$  nodes
  - $A^*(h_2) = 73$  nodes
- Solution at depth  $d = 24$ 
  - IDS  $\approx 54,000,000,000$  nodes
  - $A^*(h_1) = 39,135$  nodes
  - $A^*(h_2) = 1,641$  nodes

7	2	4
5		6
8	3	1

$h_1(n)$  = number of misplaced tiles  
 $h_2(n)$  = total Manhattan distance

Contains many  
redundant paths which  
IDS cannot break!



# Satisficing Search: Weighted A\* Search

- Often it is sufficient to find a **“good enough” solution** if it can be found very quickly or with way less computational resources. I.e., **expanding fewer nodes**.
- We could use inadmissible heuristics in A\* search (e.g., by multiplying  $h(n)$  with a factor  $W$ ) that sometimes overestimate the optimal cost to the goal slightly.
  1. It potentially reduces the number of expanded nodes significantly.
  2. **This will break the algorithm’s optimality guaranty!**

$$f(n) = g(n) + W \times h(n)$$

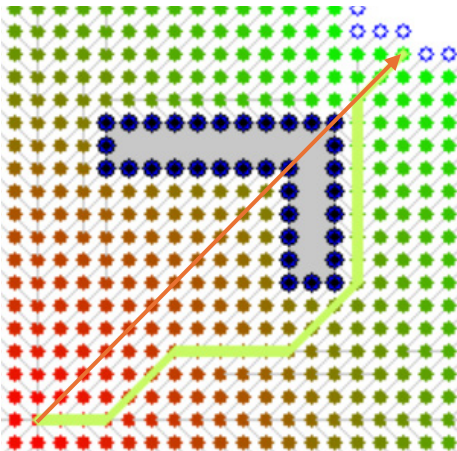
<b>Weighted A* search:</b>	$g(n) + W \times h(n)$	$(1 < W < \infty)$
----------------------------	------------------------	--------------------

The presented algorithms are special cases:

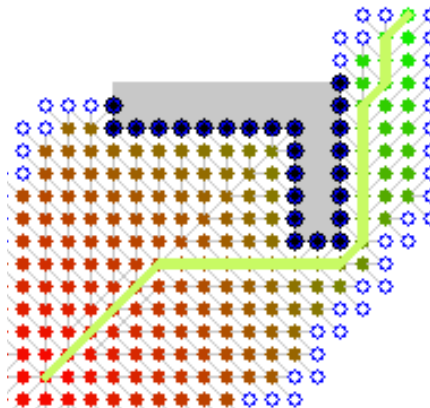
A* search:	$g(n) + h(n)$	$(W = 1)$
Uniform cost search/BFS:	$g(n)$	$(W = 0)$
Greedy best-first search:	$h(n)$	$(W = \infty)$

# Example of Weighted A\* Search

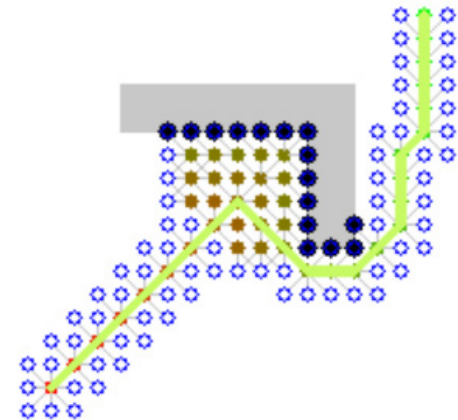
Reduction in the number of expanded nodes



Breadth-first Search (BFS)  
 $f(n) = \# \text{ actions to reach } n$



Exact A\* Search  
 $f(n) = g(n) + h_{Eucl}(n)$



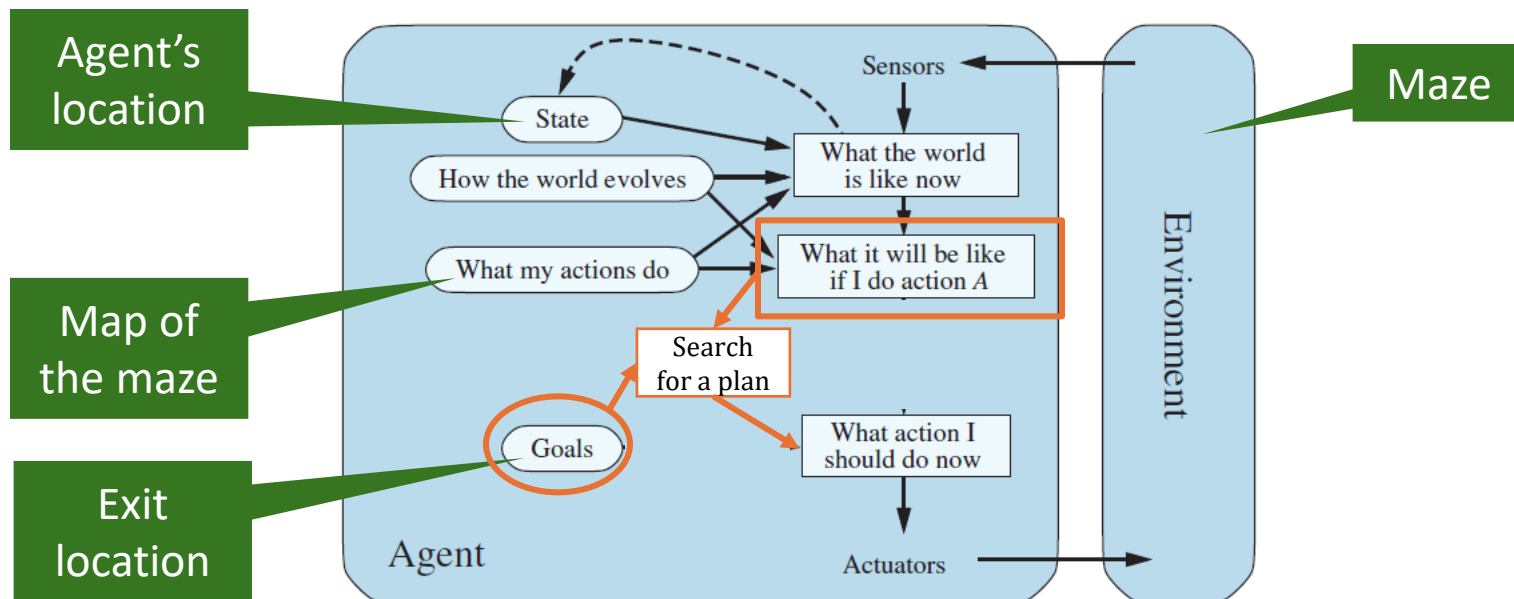
Weighted A\* Search  
 $f(n) = g(n) + 5 h_{Eucl}(n)$



# Summary Planning Agents

# Remember: Planning Agent (Goal-based)

- The agent has the task of reaching a defined **goal state**.
- The performance measure is typically the cost to reach the goal.
- We will discuss a special type of **goal-based agents** called **planning agents**, which use **search algorithms** to plan a sequence of actions that lead to the goal.

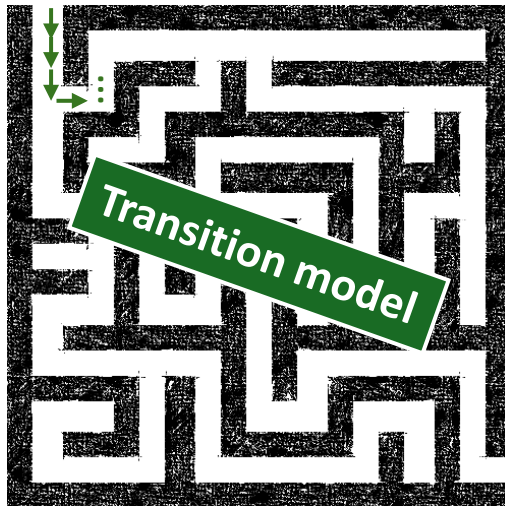


$$a_i = \operatorname{argmin}_{a_i \in A} \left[ \sum_{t=i}^T c_t \mid s_T \in S^{\text{goal}} \right]$$

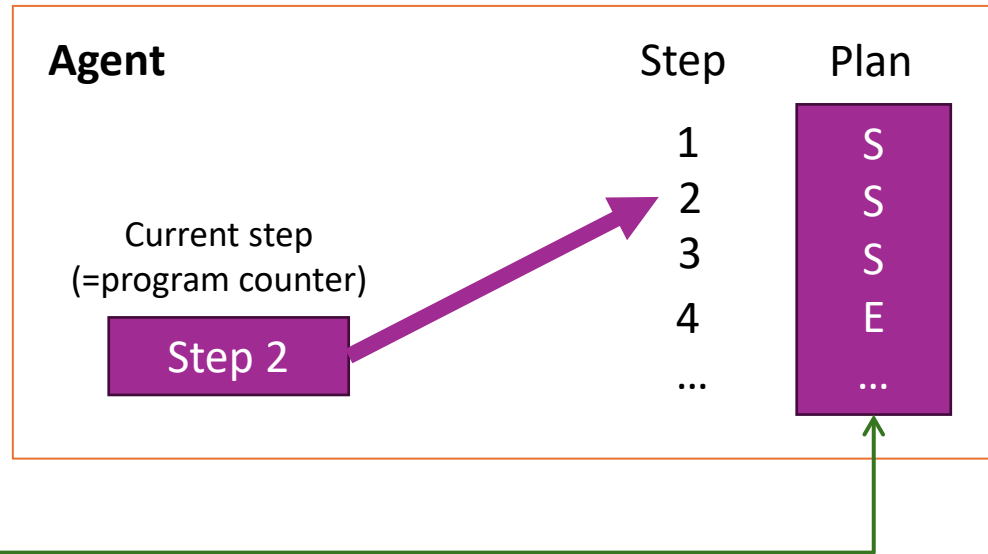
# A Planning Agent: Planning vs. Execution Phase

1. Planning is done by a **planning function** using search. The result is a **plan**.
2. The plan is executed by the agent function, which returns the planned actions from the plan step-by-step.

## Planning function



## Execution of the plan at step 2 returns action S

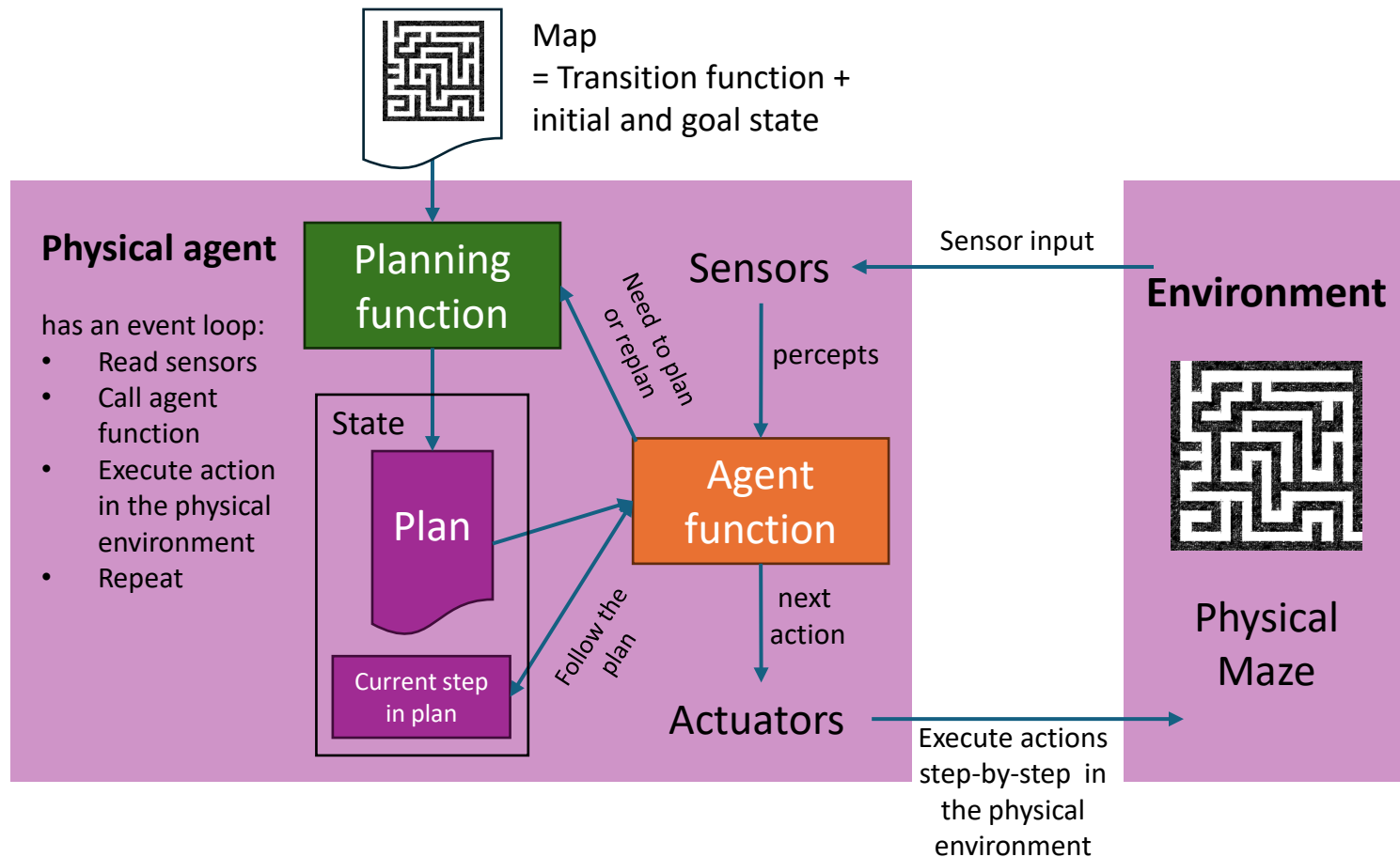


**Note:** The execution does not use percepts or the transition function. It blindly follows the plan.

**Caution:** This only works in an environment with **deterministic transitions**.



# Example: Complete Planning Agent to Solve a Maze



- The event loop calls the agent function for the next action.
- The agent function **follows** the stored plan or calls the planning function if there is no **plan** yet or if it thinks the current plan does not work based on the percepts (**replanning**).

# Summary: All Search Strategies

b: maximum branching factor of the search tree  
 d: depth of the optimal solution  
 m: maximum length of any path in the state space  
 C\*: cost of optimal solution

Needs cycle checking  
Cannot avoid redundant paths

Algorithm	Complete?	Optimal?	Time complexity	Space complexity
<b>BFS (Breadth-first search)</b>	Yes	If all step costs are equal	$O(b^d)$	$O(b^d)$
<b>Uniform-cost Search</b>	Yes	Yes	Number of nodes with $g(n) \leq C^*$	
<b>DFS</b>	In finite spaces (cycles checking)	<b>No</b>	$O(b^m)$	$O(bm)$
<b>IDS</b>	Yes	If all step costs are equal	$O(b^d)$	$O(bd)$
<b>Greedy best-first Search</b>	In finite spaces (cycles checking)	No	Depends on heuristic Best case: $O(bd)$ Worst case: $O(b^m)$	
<b>A* Search</b>	Yes	Yes	Number of nodes with $g(n) + h(n) \leq C^*$ With a good heuristic	

# Implementation as Best-First Search

- All discussed search strategies can be implemented using Best-first search.
- Best-first search expands always the **node with the minimum value of an evaluation function  $f(n)$** .

Search Strategy	Evaluation function $f(n)$
BFS (Breadth-first search)	$g(n)$ (=uniform path cost)
Uniform-cost Search	$g(n)$ (=path cost)
<del>DFS/IDS (see note below!)</del>	<del><math>-g(n)</math></del>
Greedy Best-first Search	$h(n)$
(weighted) A* Search	$g(n) + W \times h(n)$

- **Important note:** Do not implement DFS/IDS using Best-first search! You will get poor space complexity from BFS and the disadvantages of DFS (not optimal and worse time complexity).





# What You Should Know

- Tree search can be used for planning actions for **goal-based agents** in known, fully observable, and deterministic environments.
- Issues are:
  - The large search space typically does not fit into memory. AI uses a transition function as a compact representation of the **transition model**.
  - The search tree is built on the fly, and we must deal with **cycles, redundant paths, and memory management**.
- **Uninformed search:**
  - **BFS** is not usable for large state spaces!
  - AI typically uses **DFS** implemented in a memory-efficient way. It requires cycle checking, is not optimal, and cannot avoid redundant paths.
  - **Iterative Deepening Search (IDS)** is an optimal version of DFS that rebuilds trees level-wise. It is very slow.
- **Informed search** uses heuristics based on knowledge or percepts to improve search performance
  - A\* requires an admissible heuristic and expands fewer nodes than BFS.