

COMP3251

## Lecture 2: Merge Sort (Chapter 2.3)

# Sorting

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**Output:** The same set of  $n$  integers in ascending order.

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# How good is insertion sort?

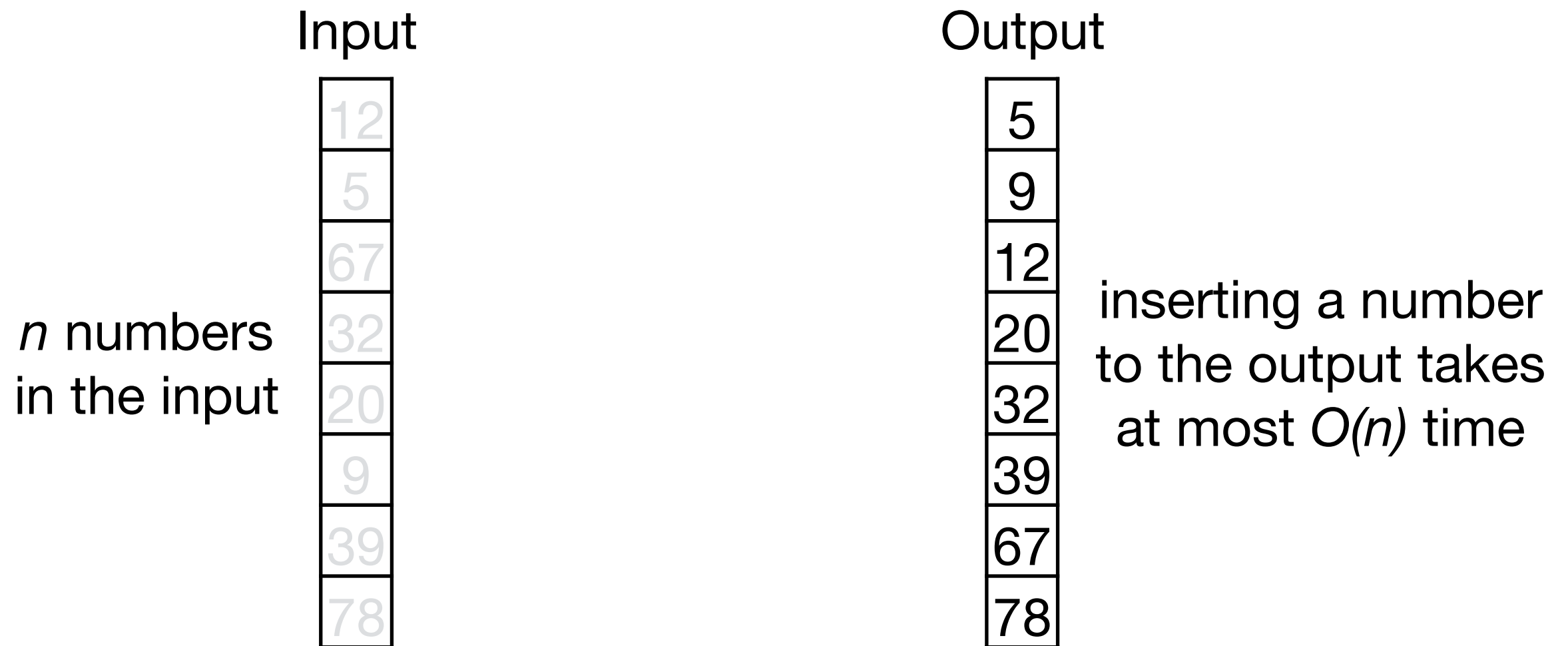
- **Worst-case analysis**
  - For any input length  $n$ , consider the maximum number of steps the algorithm needs on the **worst input** of length  $n$ .
  - Measuring performance instance by instance does not lead to reasonable algorithms (too specialized).
  - Average performance is a reasonable alternative, but **not** the focus of this course.
- **Asymptotic analysis**
  - Big-O notation
  - Focus on the performance when  $n$  gets really big
  - **Example:**  $5n^2+2n+10 = O(n^2)$

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**Warm-up:** The insertion sort algorithm takes  $O(n^2)$  time.



# Divide and Conquer (Ch. 2)

The divide-and-conquer algorithm design paradigm solves a problem as follows:

- 1) **Divide:** Breaking the problem into subproblems that are themselves smaller instances of the same type of problem;
- 2) **Recurse:** Recursively solving these subproblems;
- 3) **Combine:** Appropriately combining their answers to get an answer of the original problem.

**Note:** If the size of a subproblem is small enough, we will stop using the divide-and-conquer strategy; instead, we may solve the subproblem by brute-force.

# Merge Sort

**Input:** A set of  $n$  integers —  $x_1, x_2, \dots, x_n$ .

**Output:** The same set of  $n$  integers **in ascending order**.

**Divide:** Divide the input integers into two subsets  
 $x_1, \dots, x_{n/2}$  and  $x_{n/2+1}, \dots, x_n$ .

**Recurse:** Sort the two subsets ( $n/2$ -size subproblems).  
Let  $y_1, \dots, y_{n/2}$  and  $y_{n/2+1}, \dots, y_n$  be the output.

**Combine:** Merge two sorted list  $y_1, \dots, y_{n/2}$  and  $y_{n/2+1}, \dots, y_n$   
into a single sorted lists of integers.

**Running time:**  $T(n)$ , total time for sorting a set of  $n$  integers.

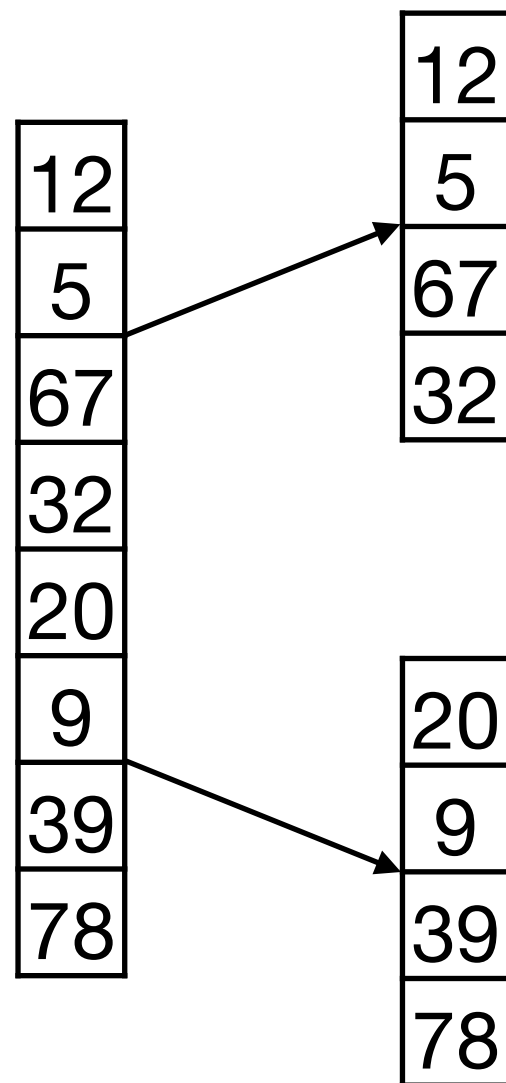
$$T(n) = 2T(n/2) + (\text{running time of merging two } n/2\text{-size sorted lists})$$

# An Example of Merge Sort

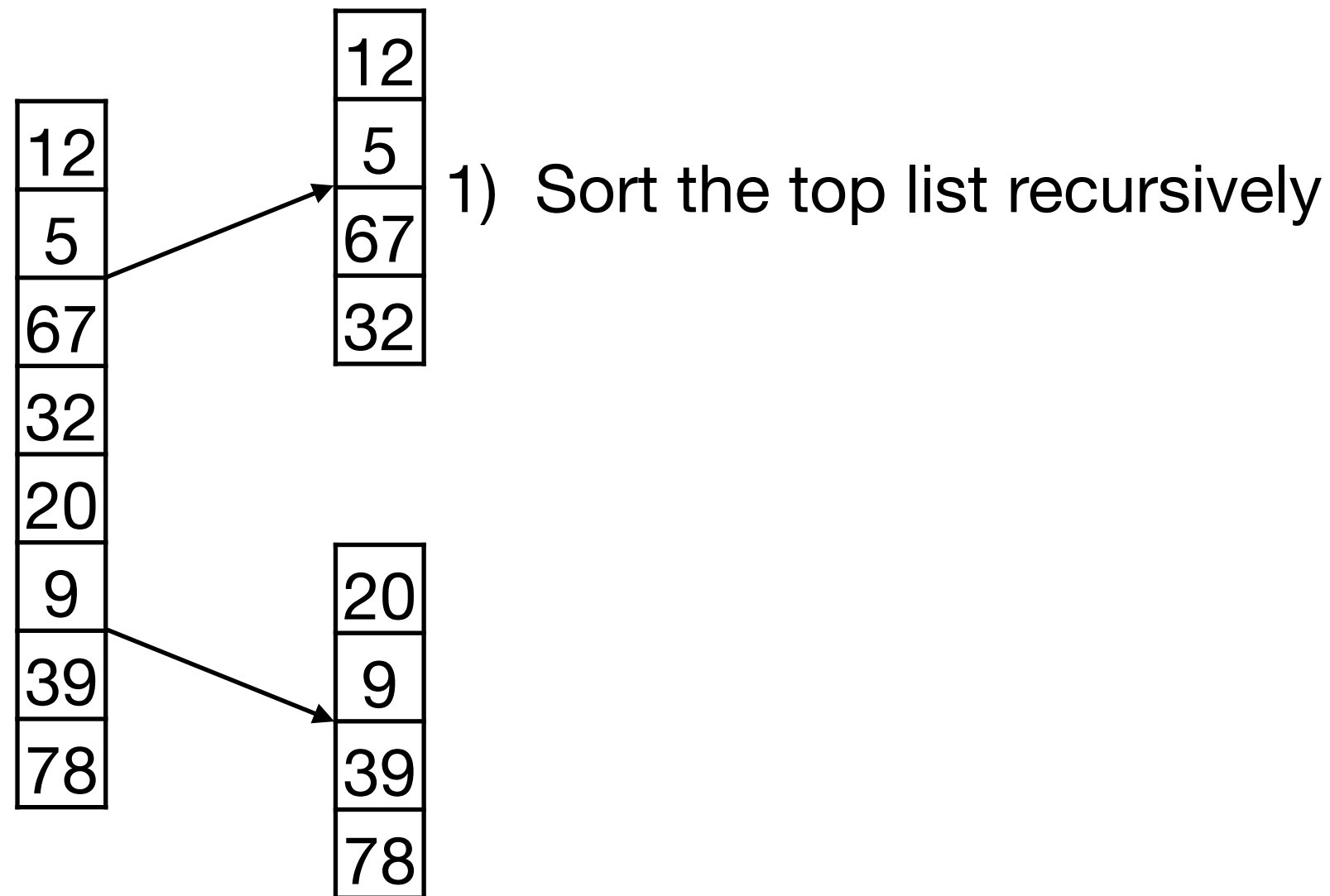
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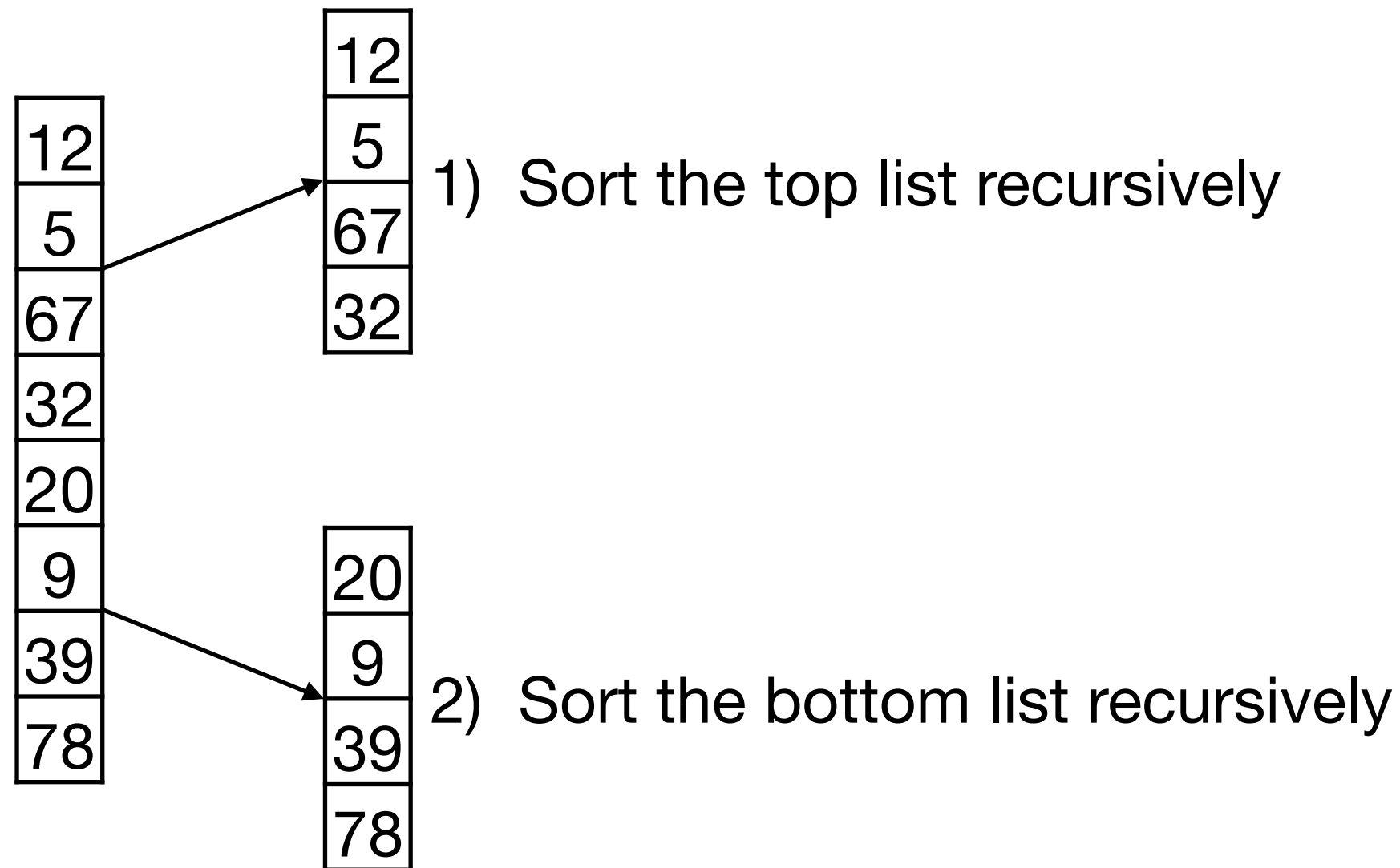
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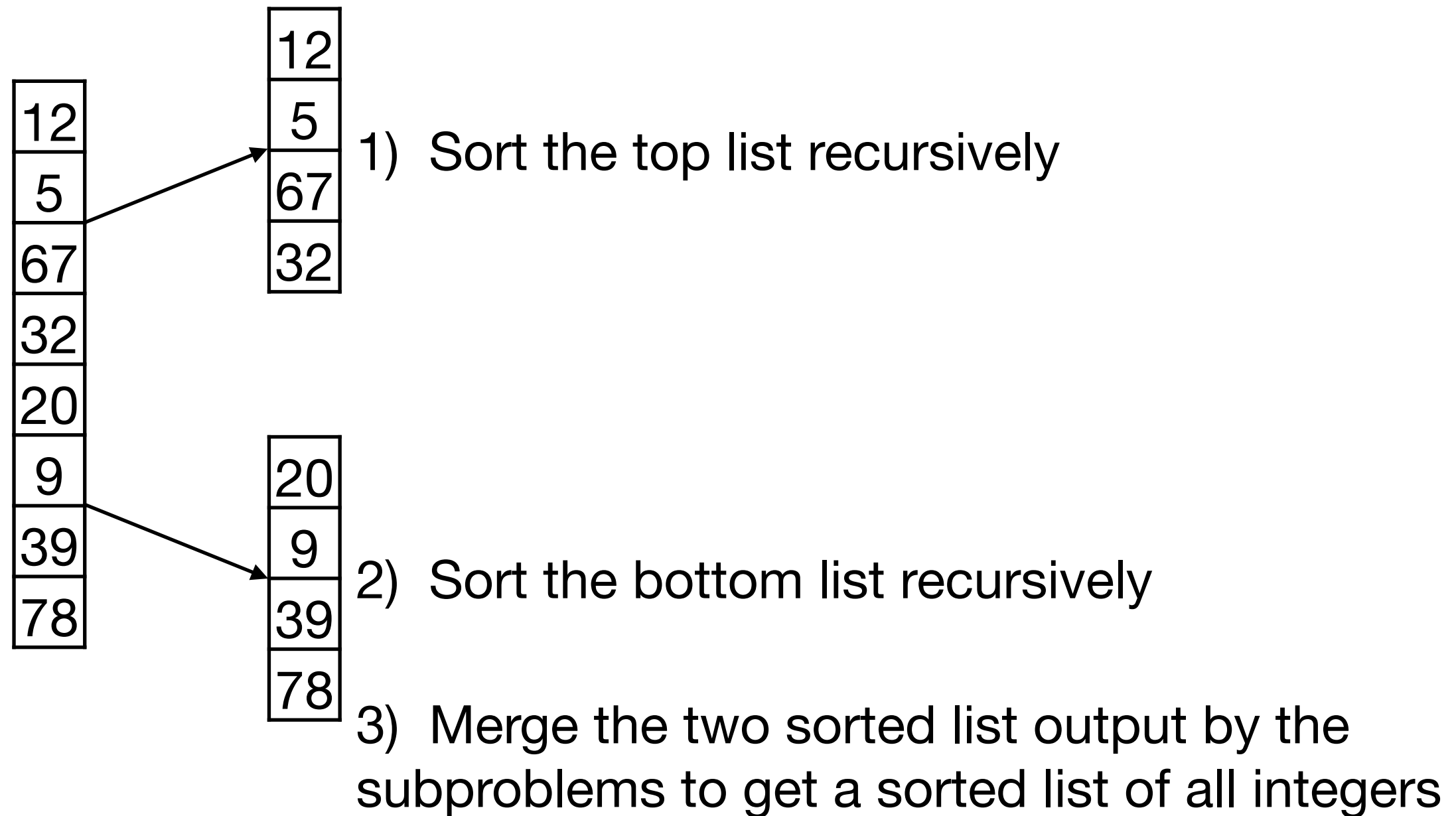
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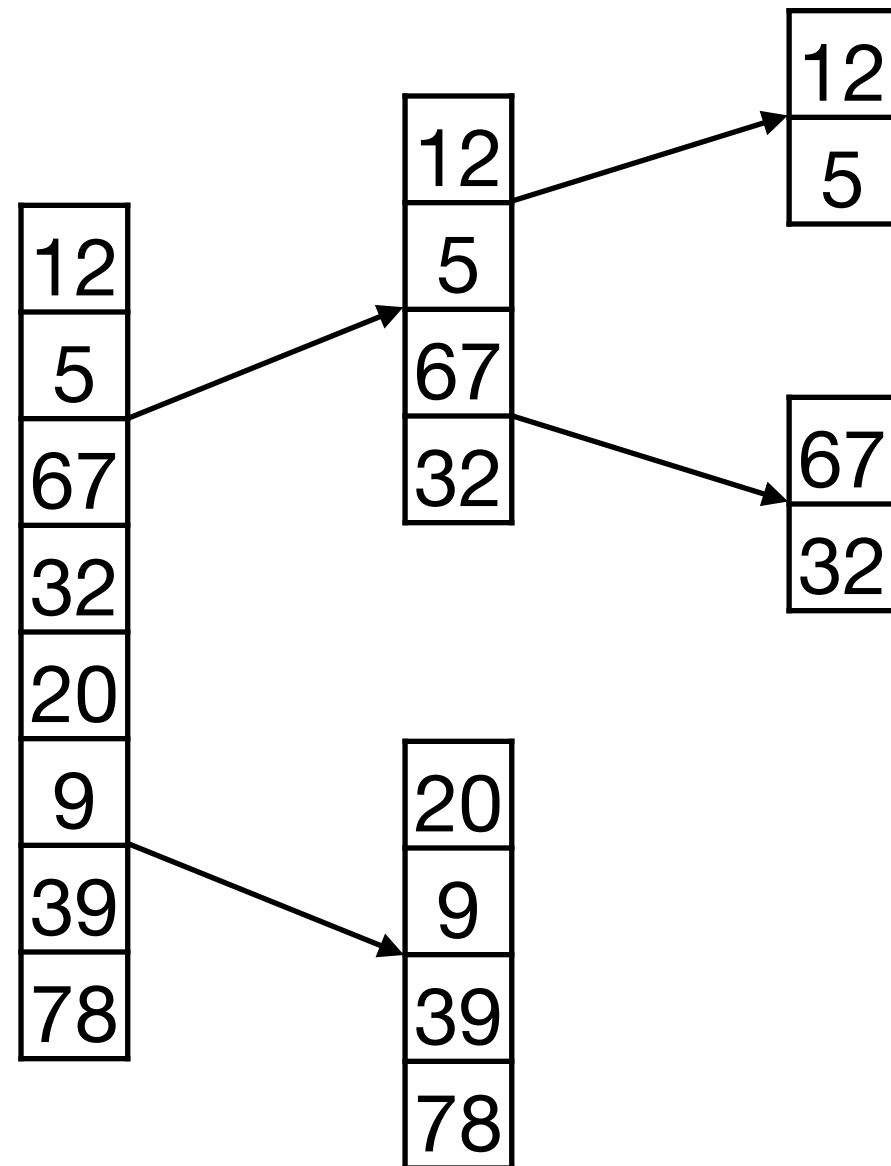
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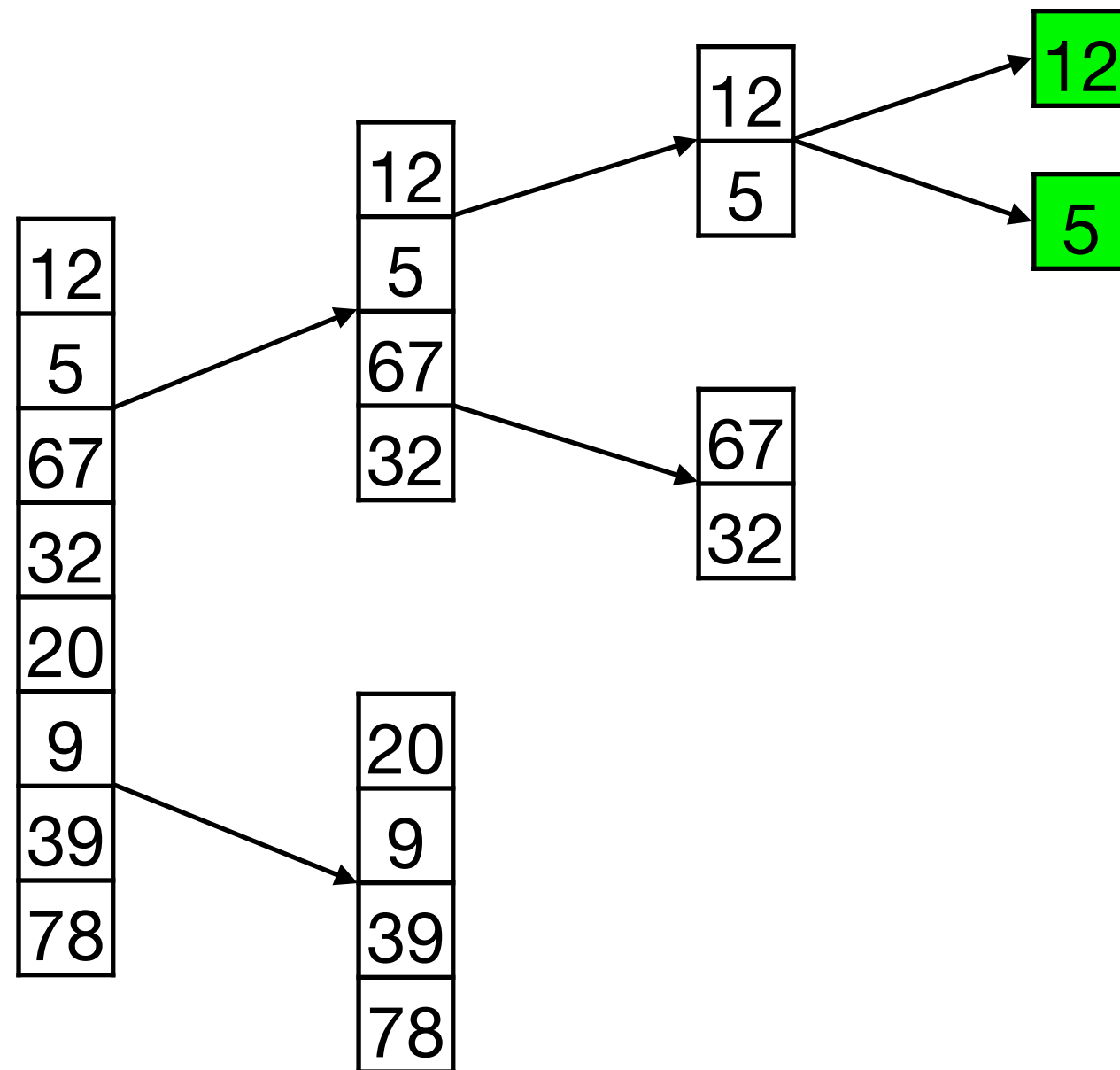
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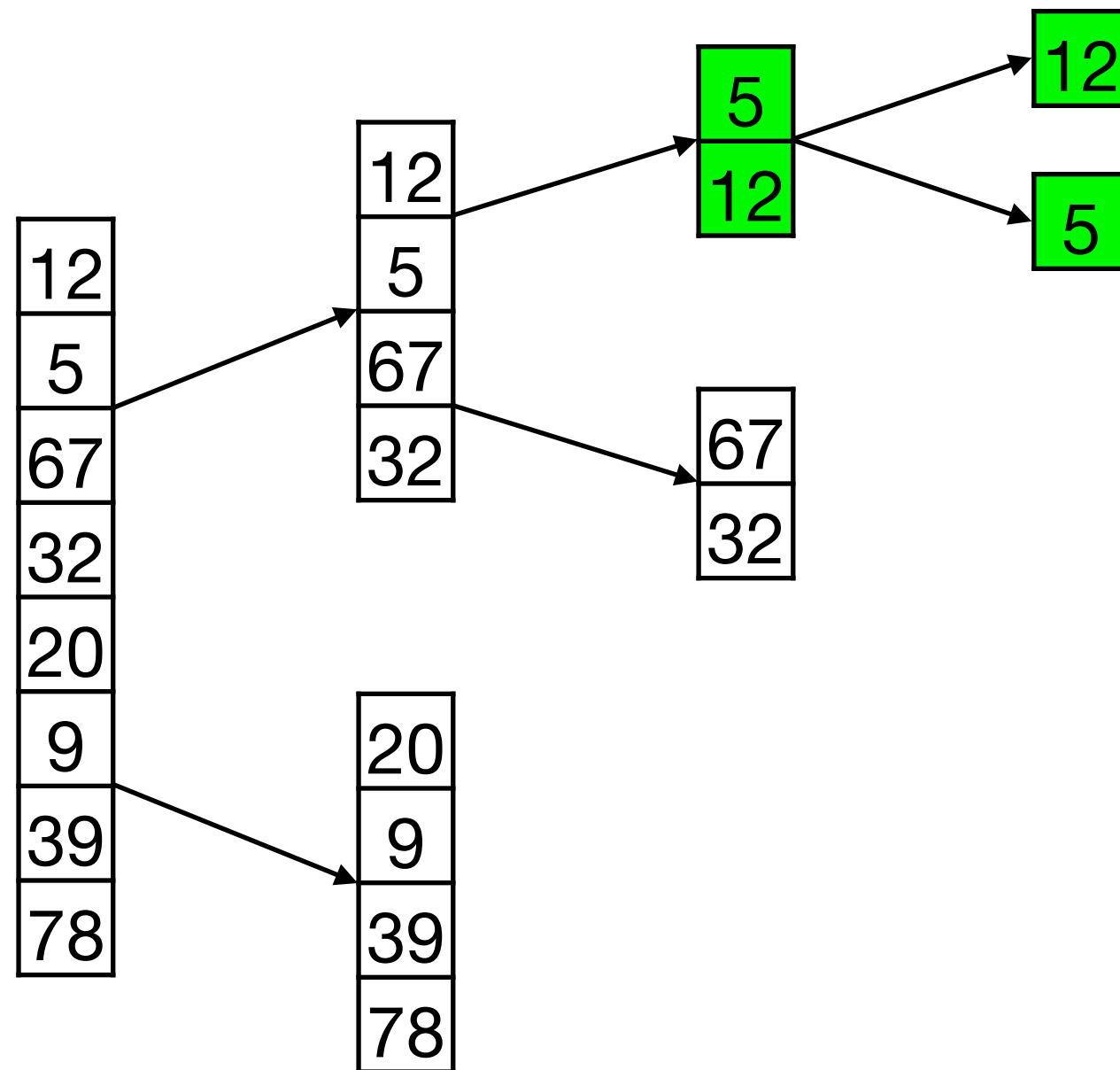
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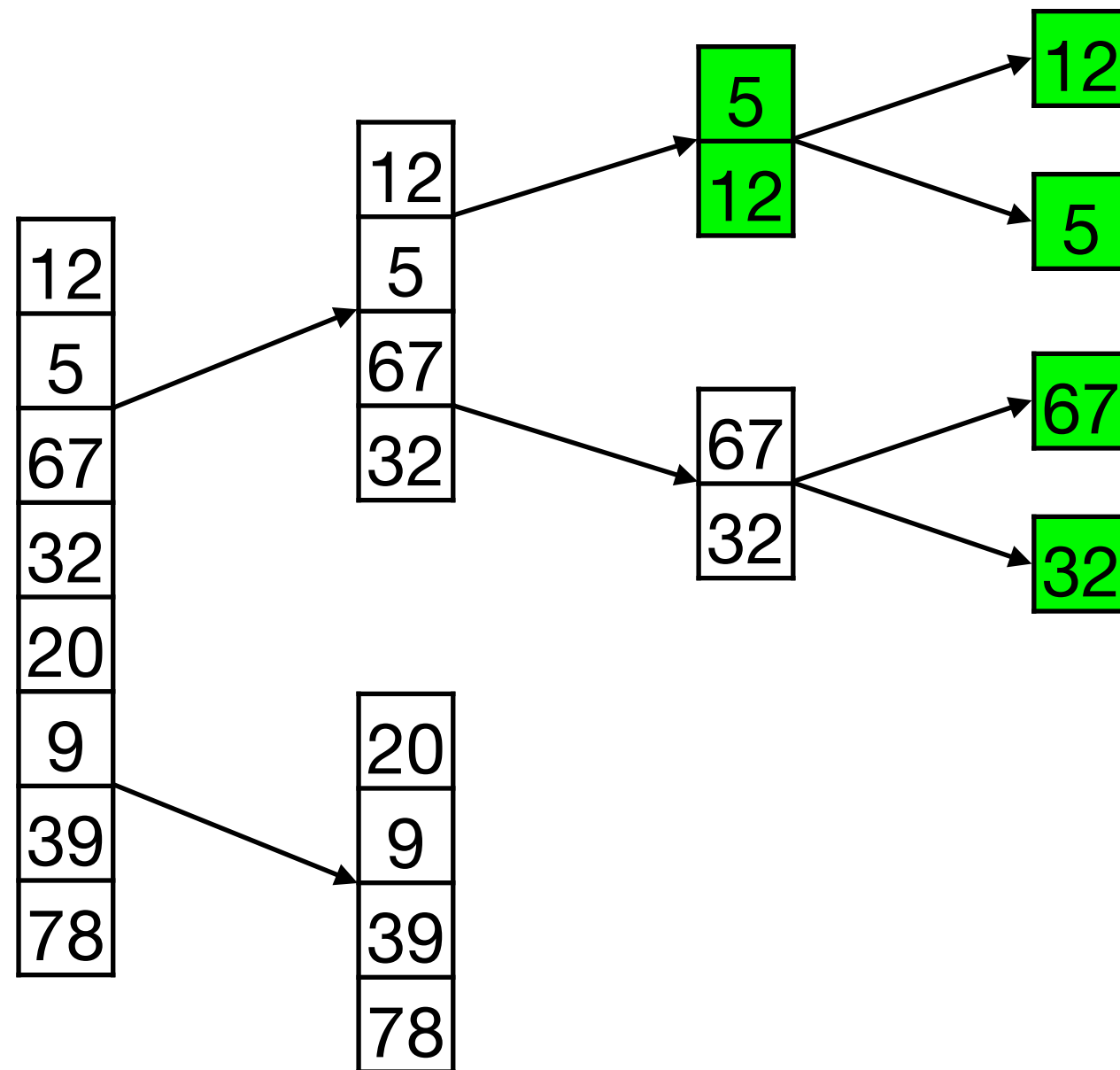
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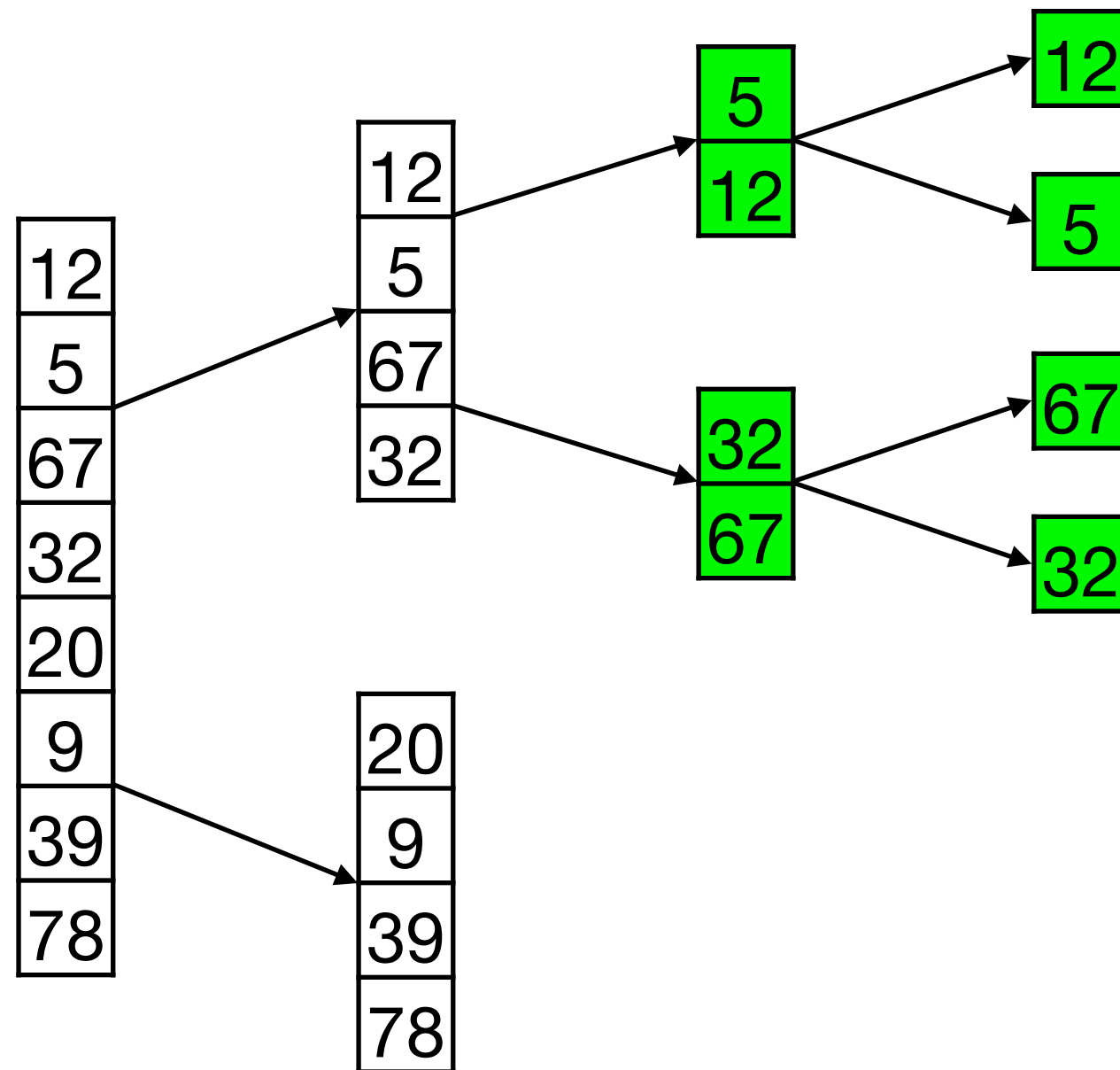


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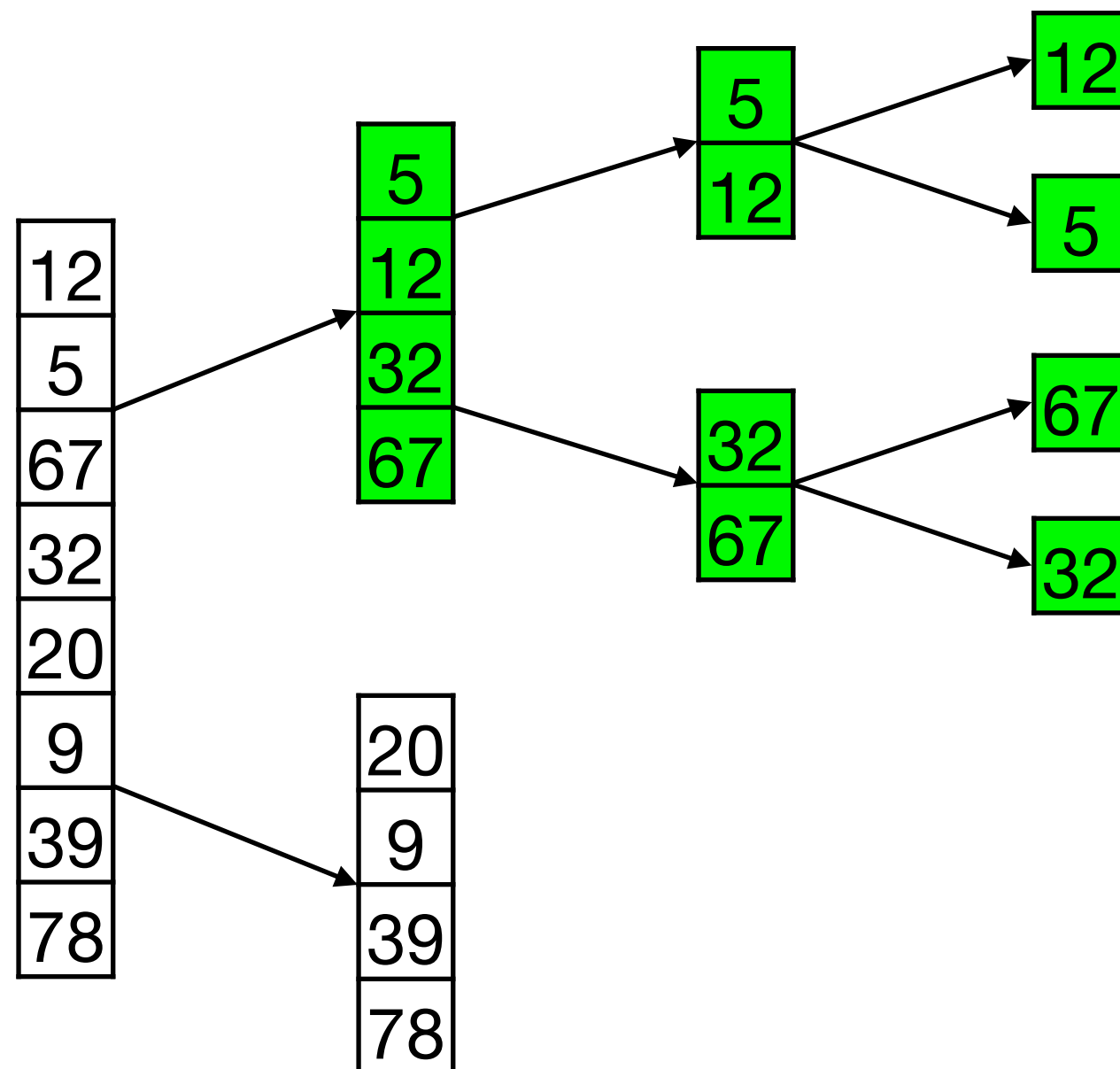




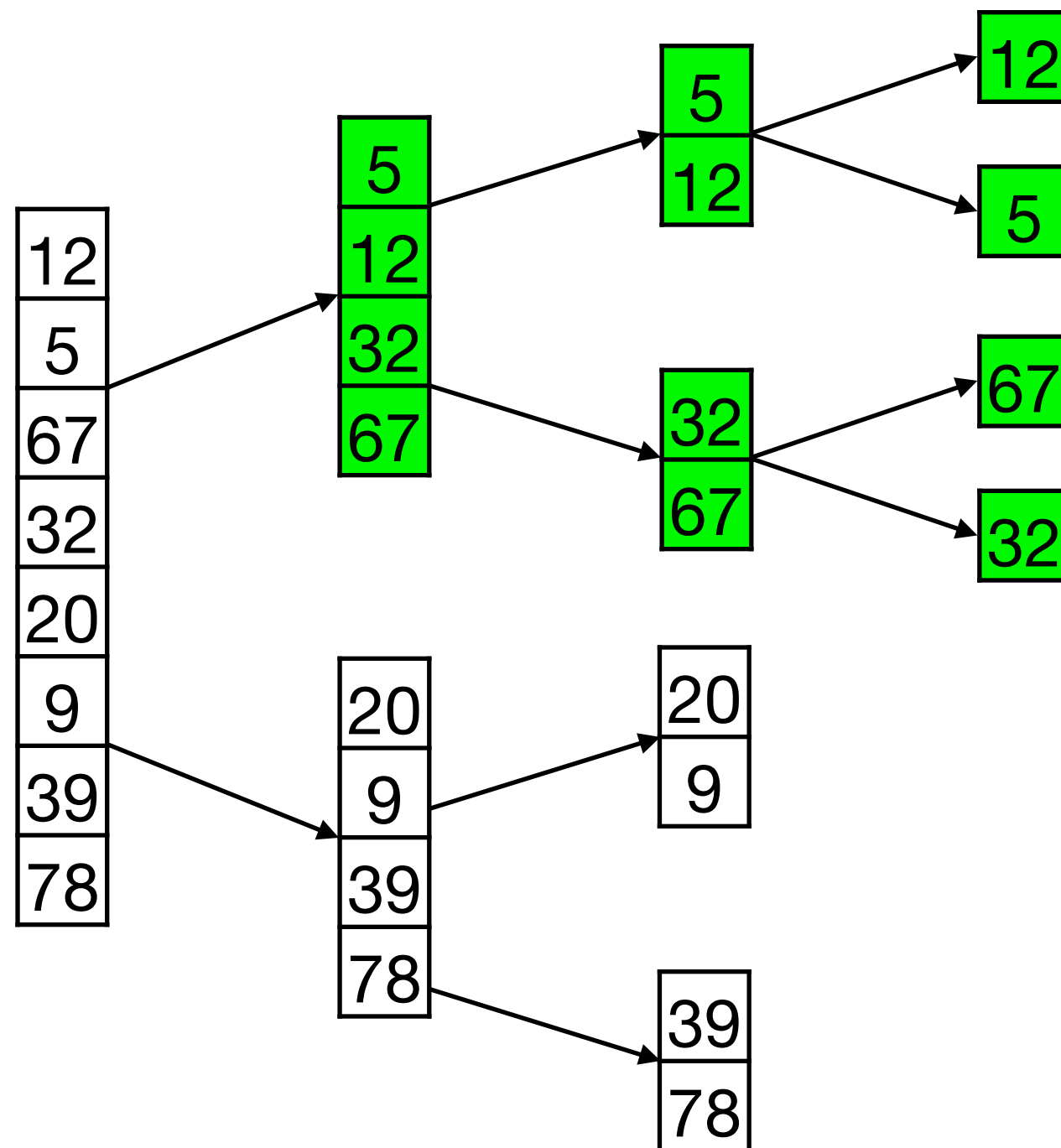
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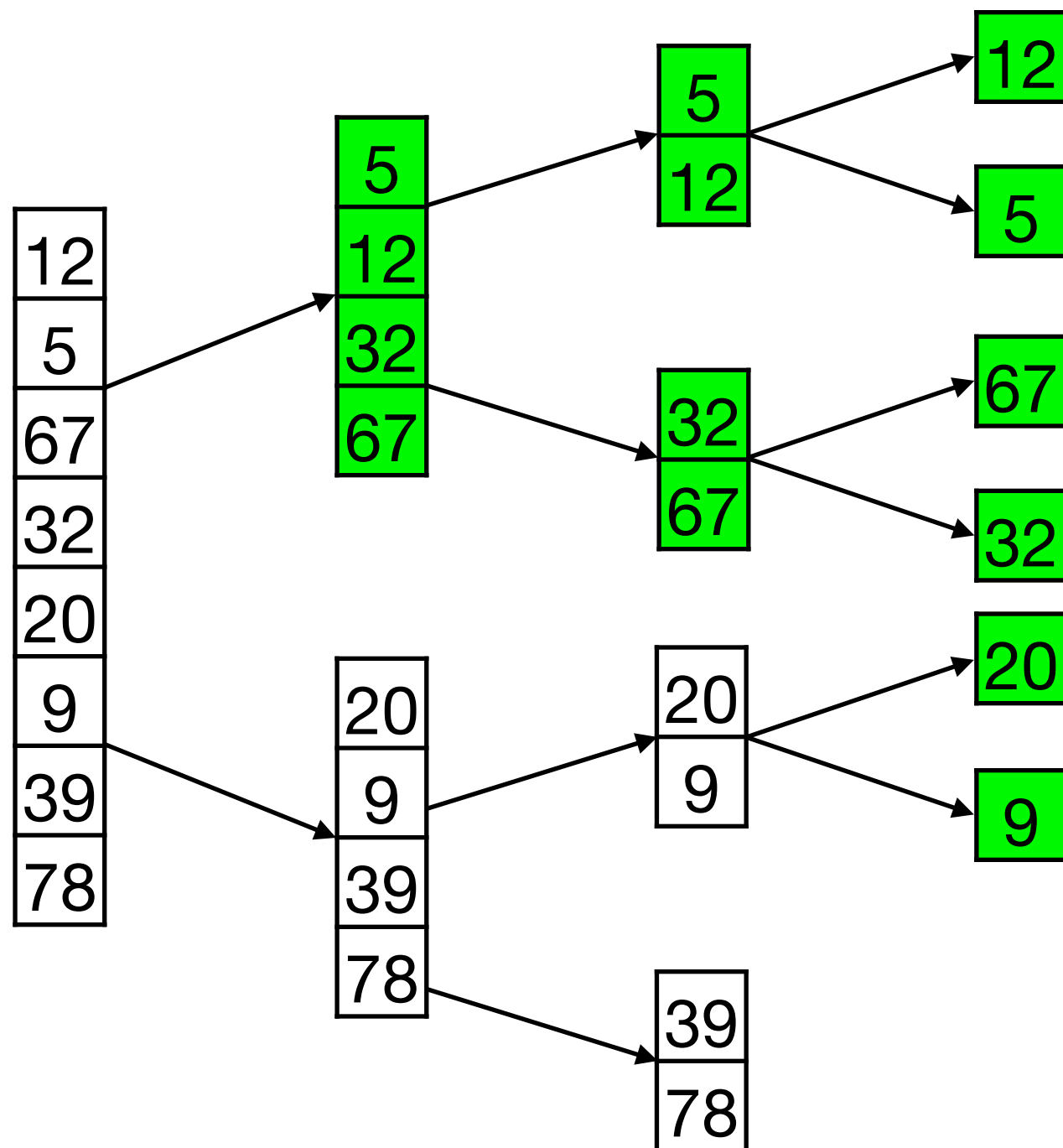
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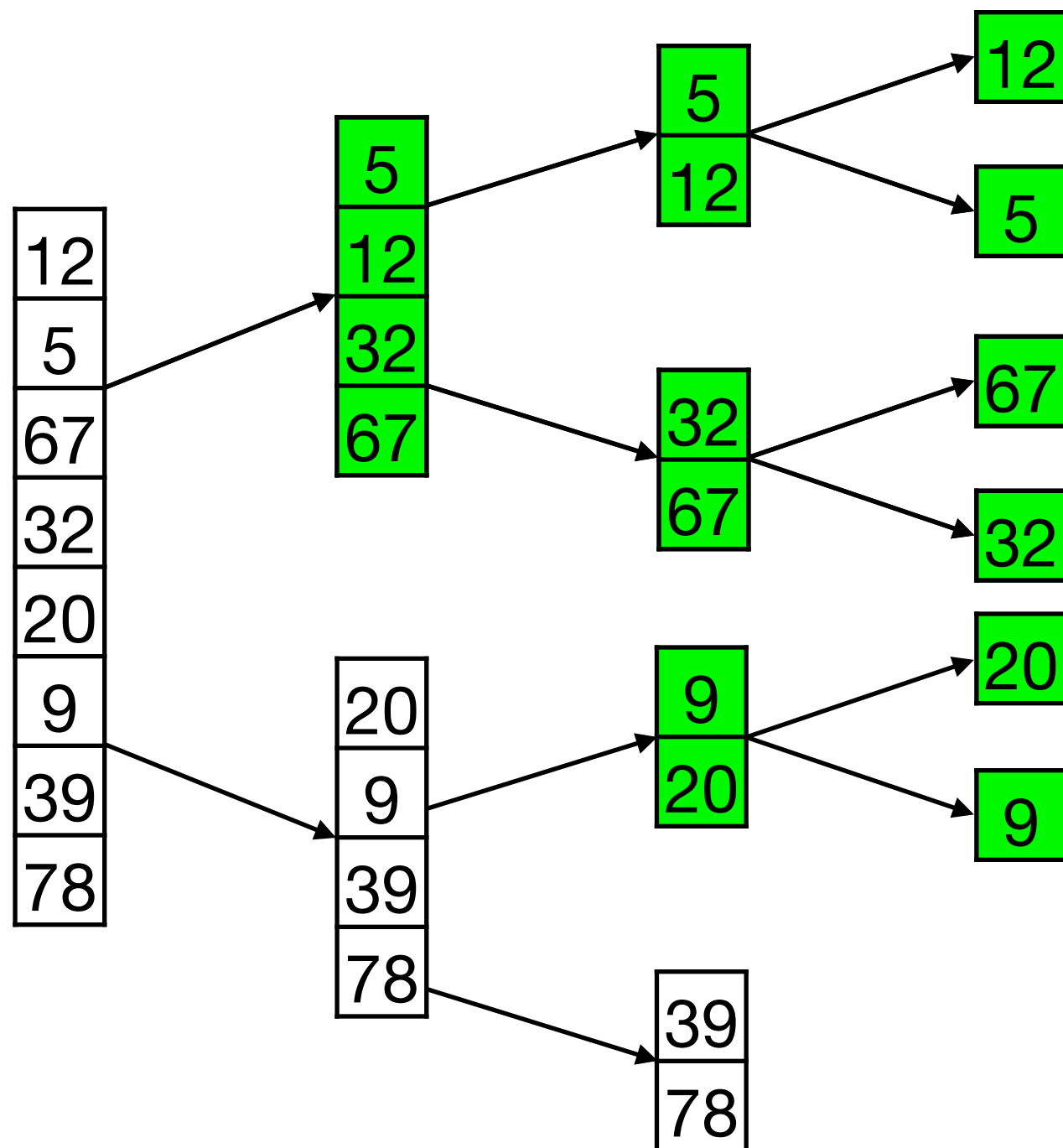
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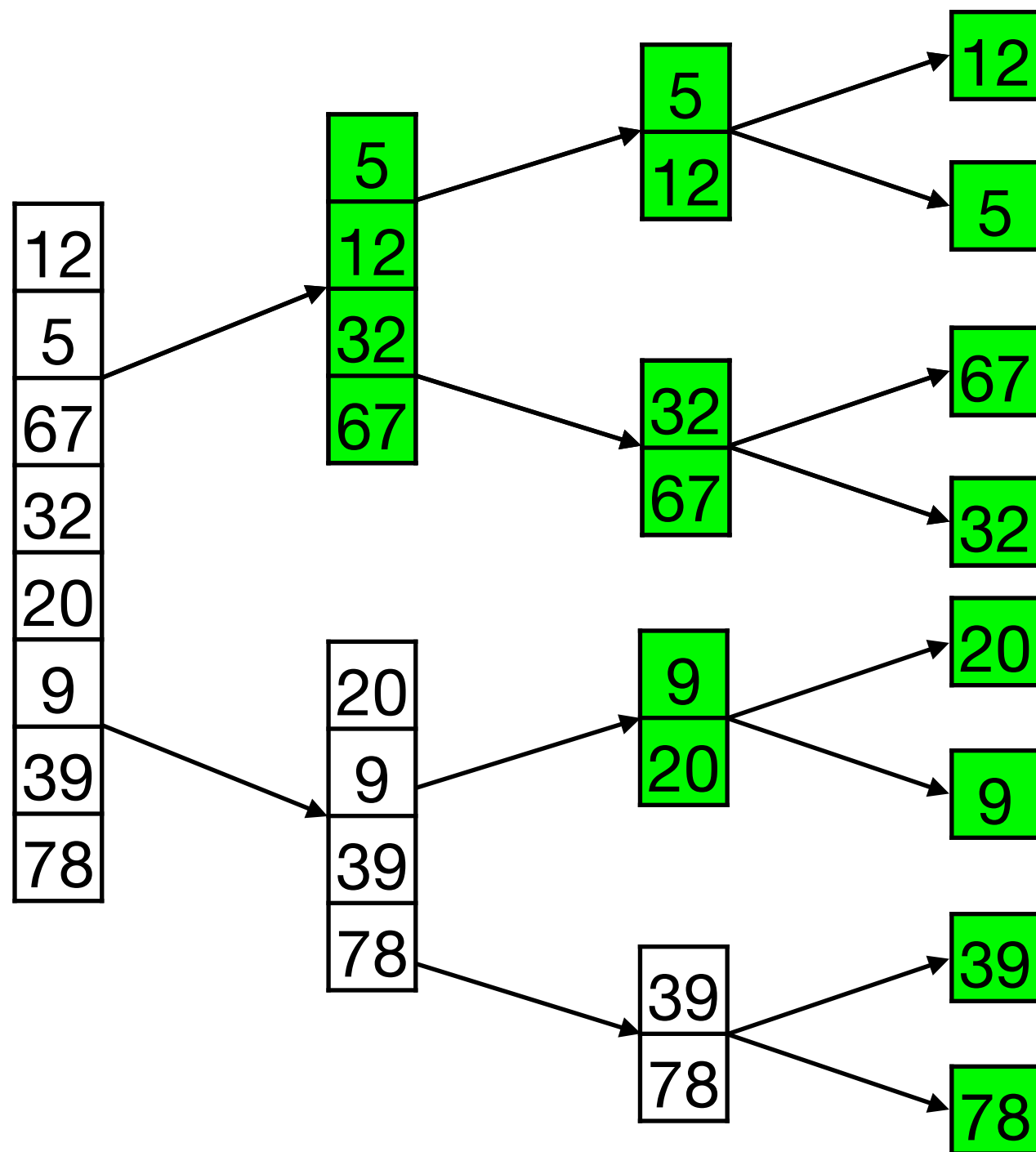
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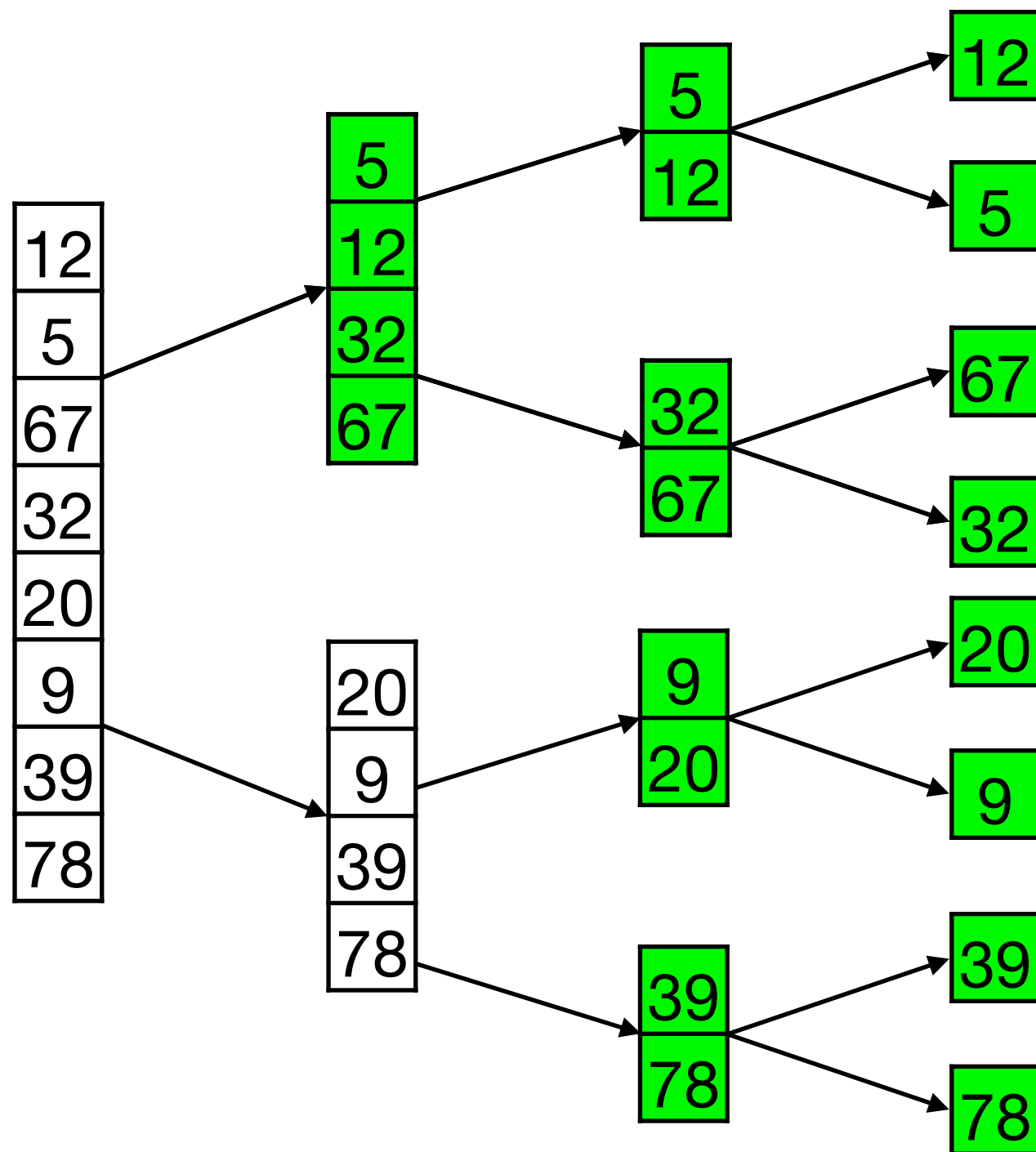
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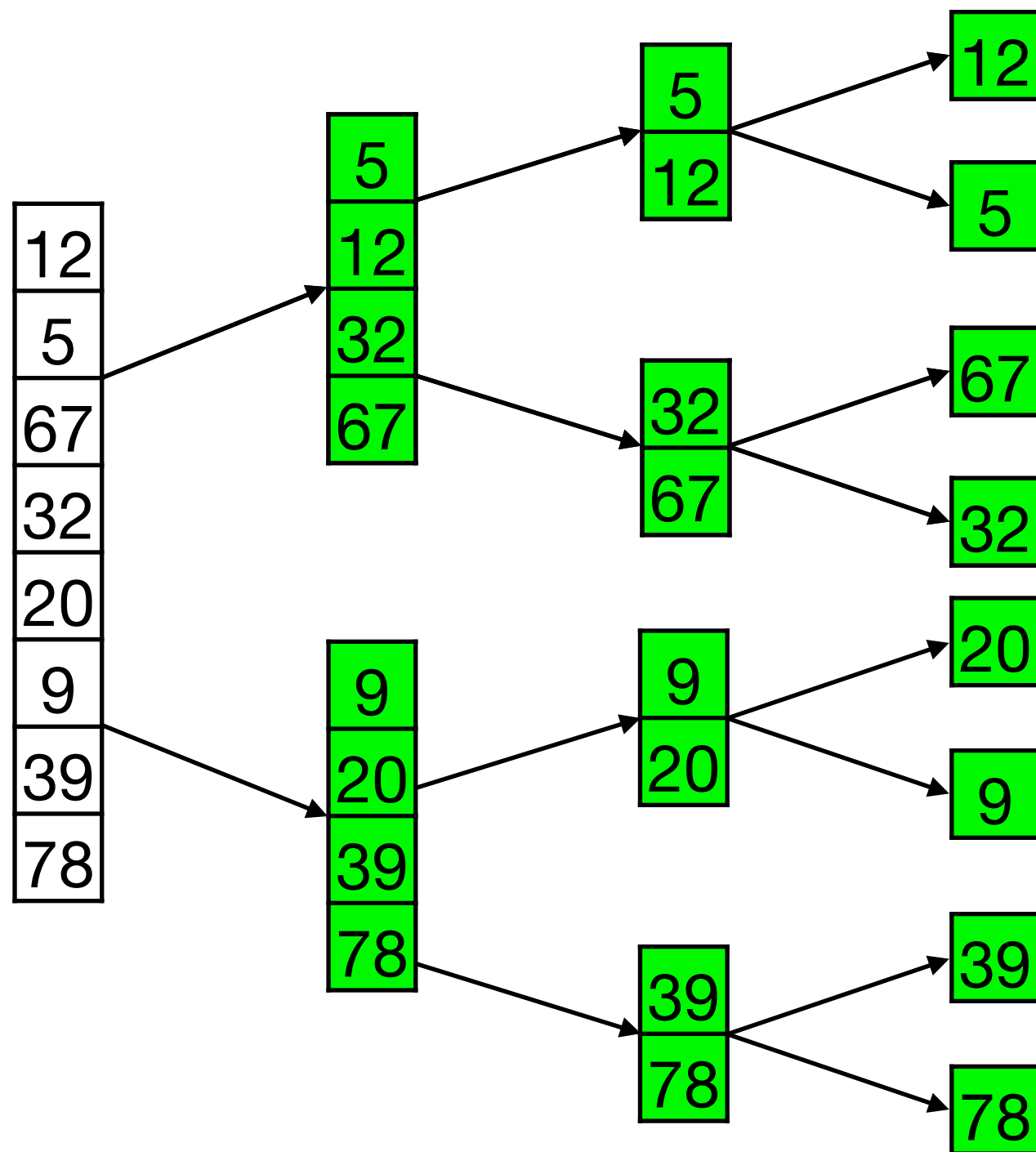
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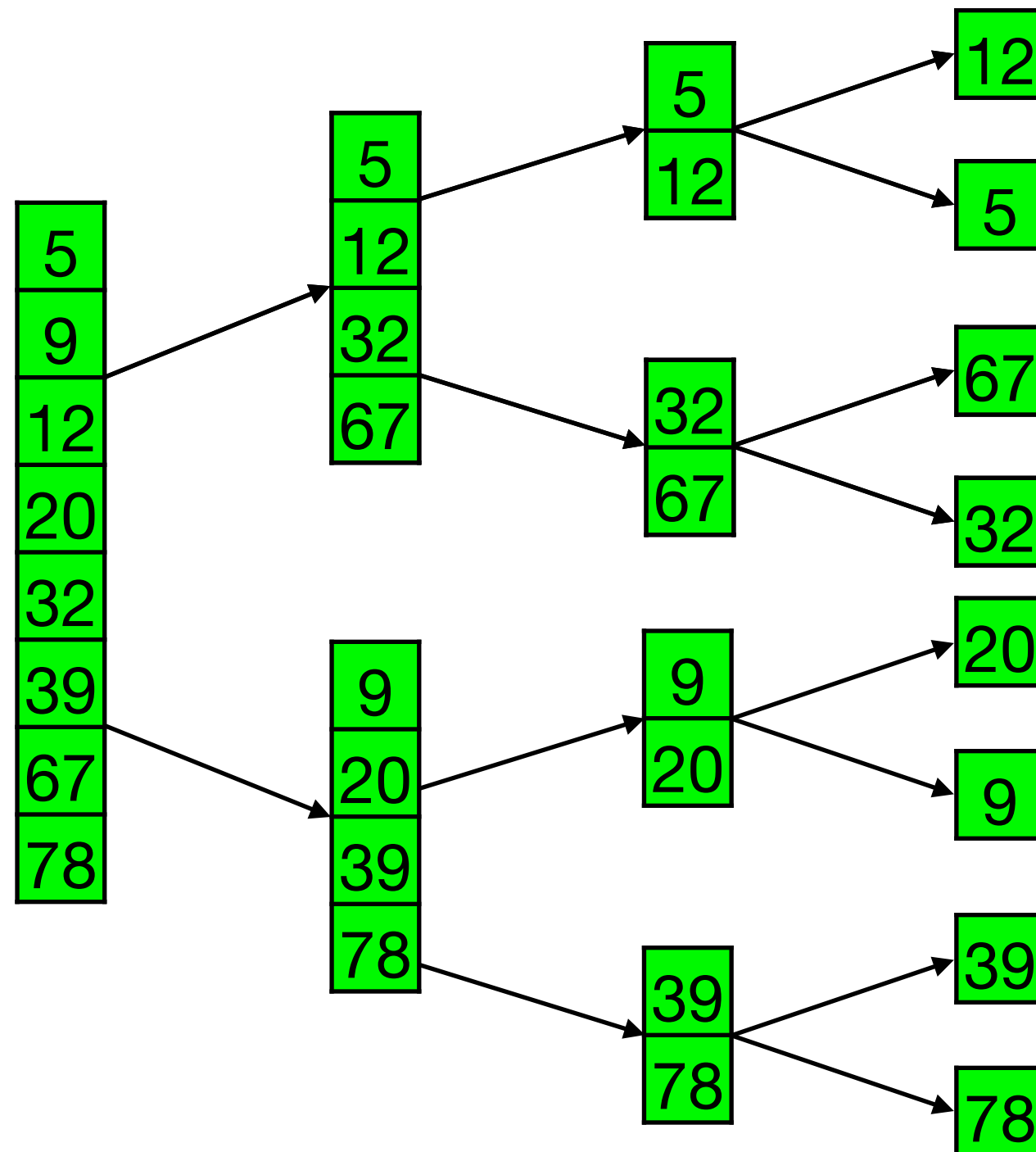


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# An Example of Merge Sort



### **Rest of this lecture:**

- 1) How to merge two sorted lists of integers?
- 2) Solving  $T(n)$ .

# Merging Two Sorted Lists of Numbers

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To merge sorted lists  $A = a_1, \dots, a_n$  and  $B = b_1, \dots, b_n$ :

Maintain a *Current* pointer into each list, initialized to point to the front elements

While both lists are nonempty:

Let  $a_i$  and  $b_j$  be the elements pointed to by the *Current* pointer

Append the smaller of these two to the output list

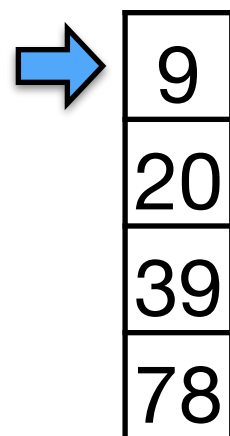
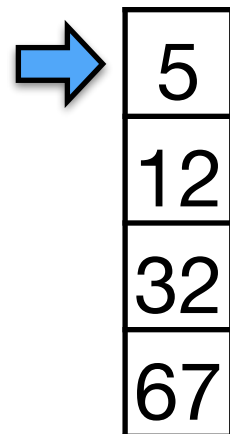
Advance the *Current* pointer in the list from which the smaller element was selected

EndWhile

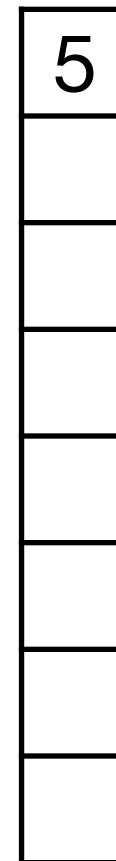
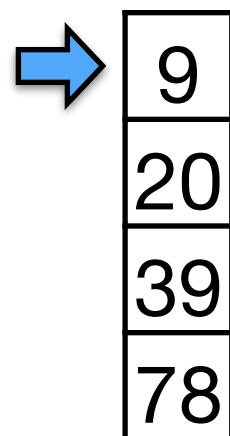
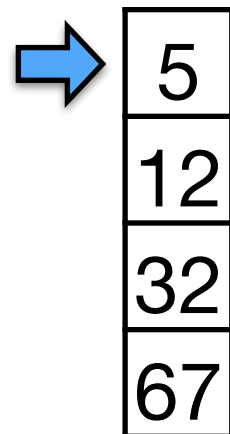
Once one list is empty, append the remainder of the other list to the output

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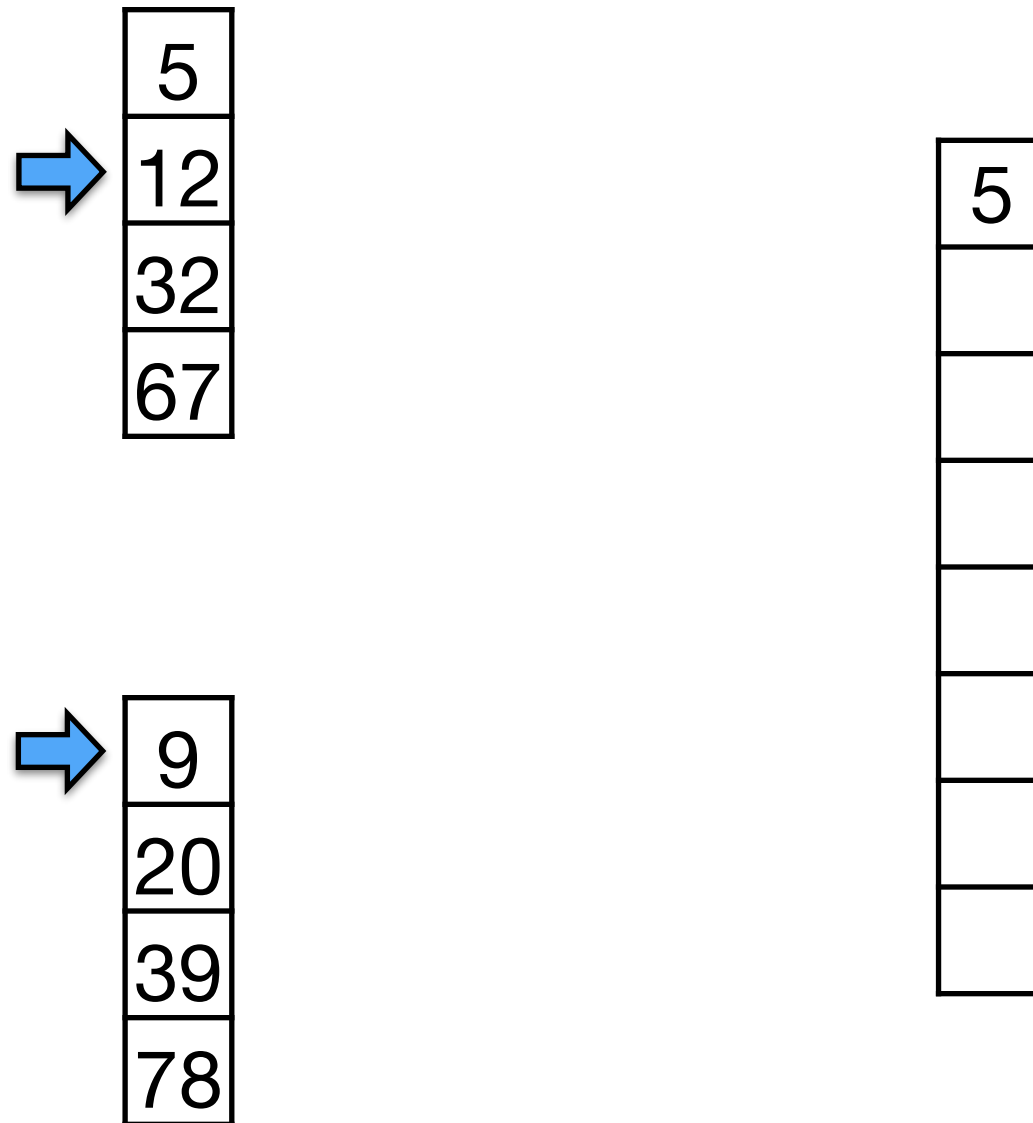
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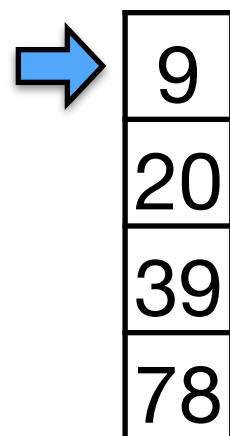
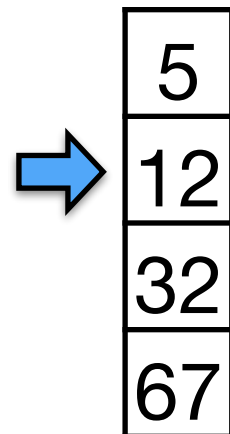
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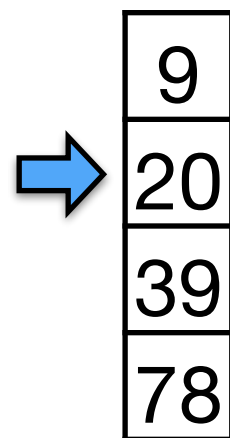
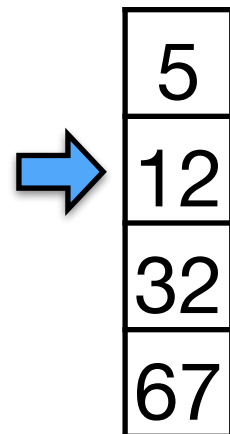
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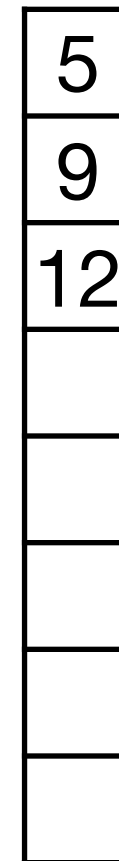
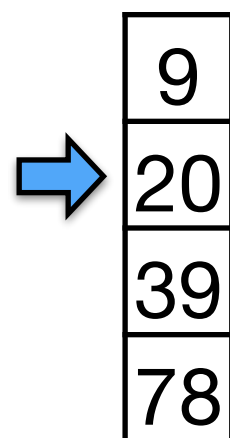
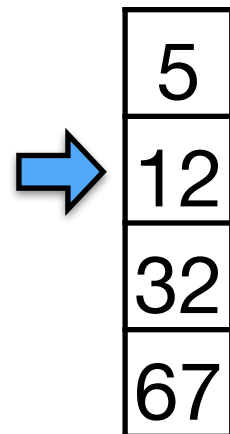


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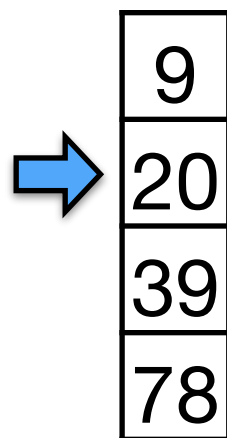
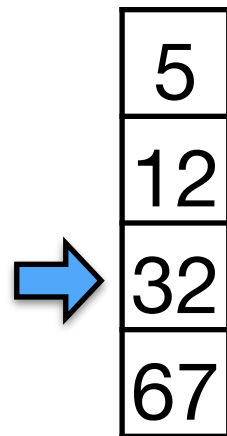




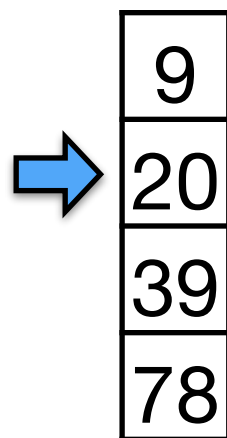
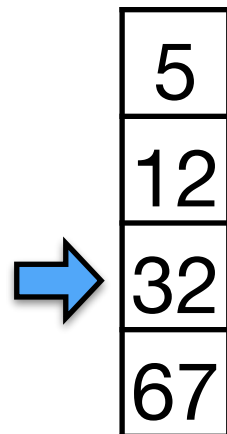
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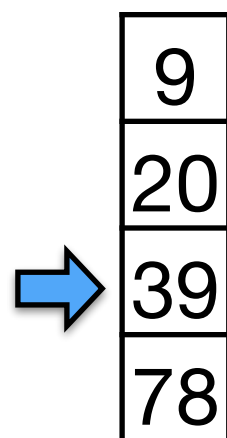
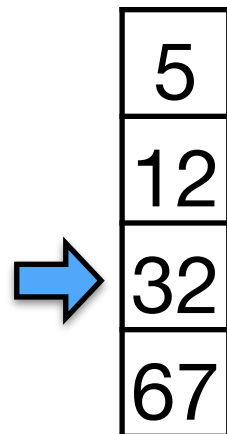
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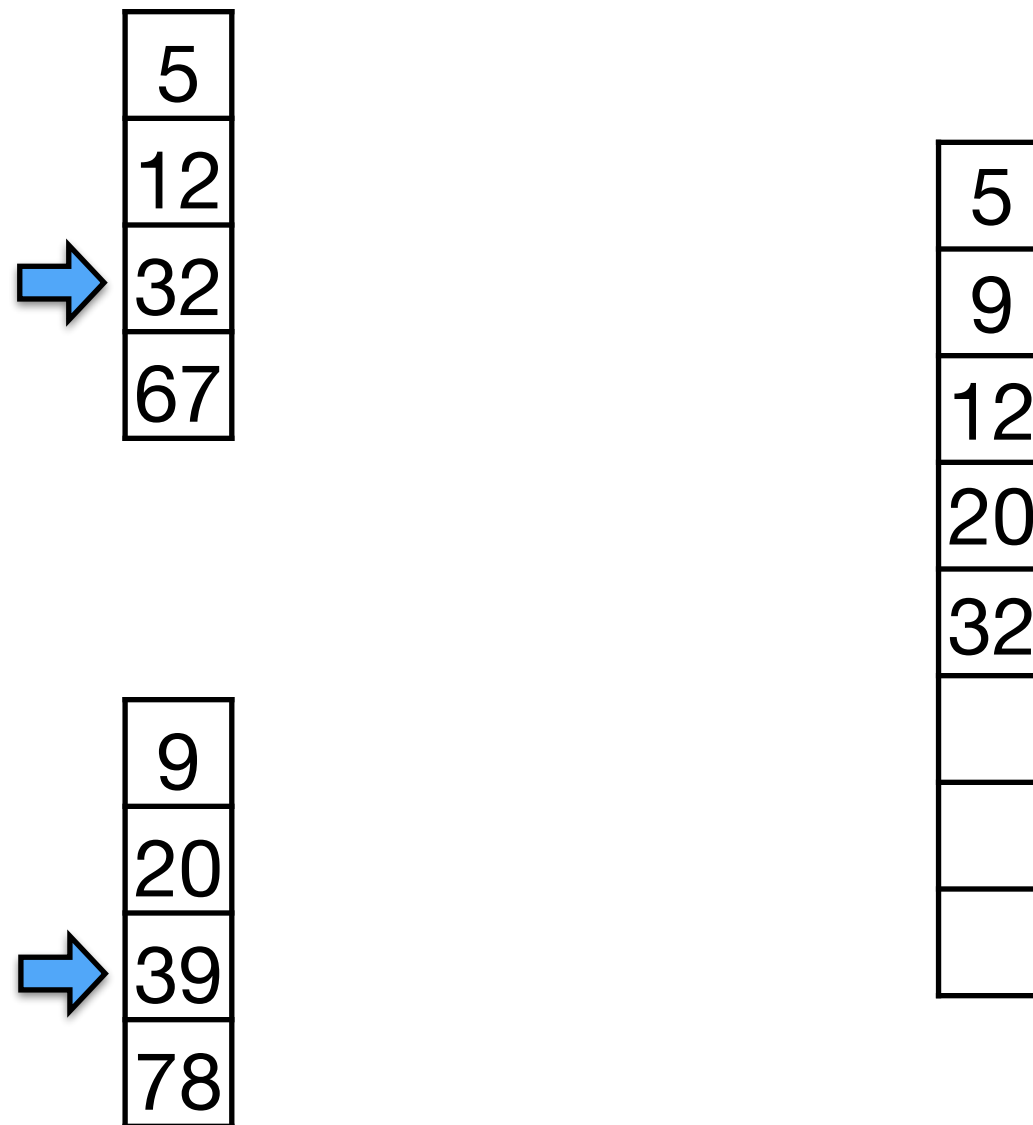
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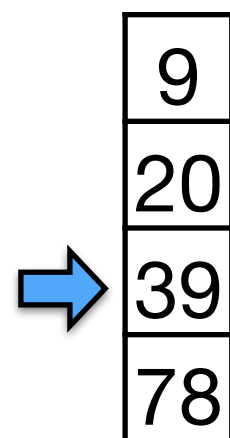
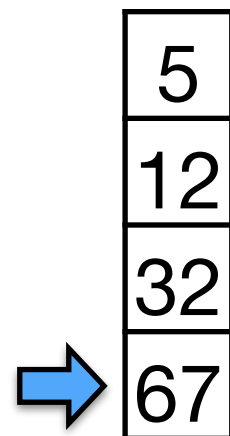
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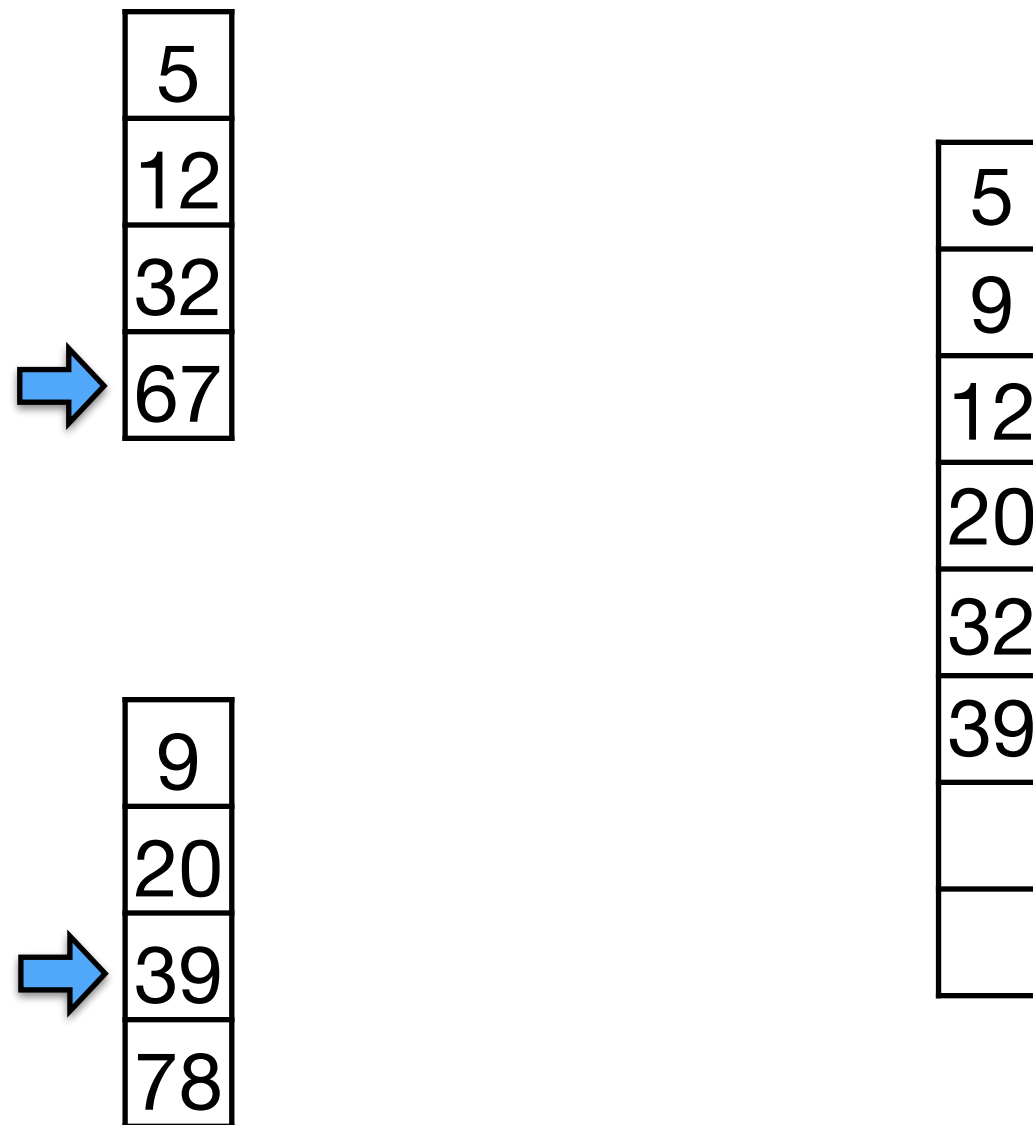
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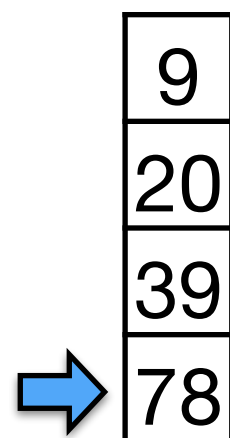
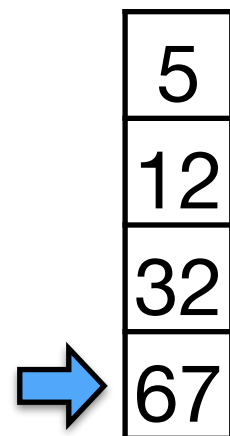
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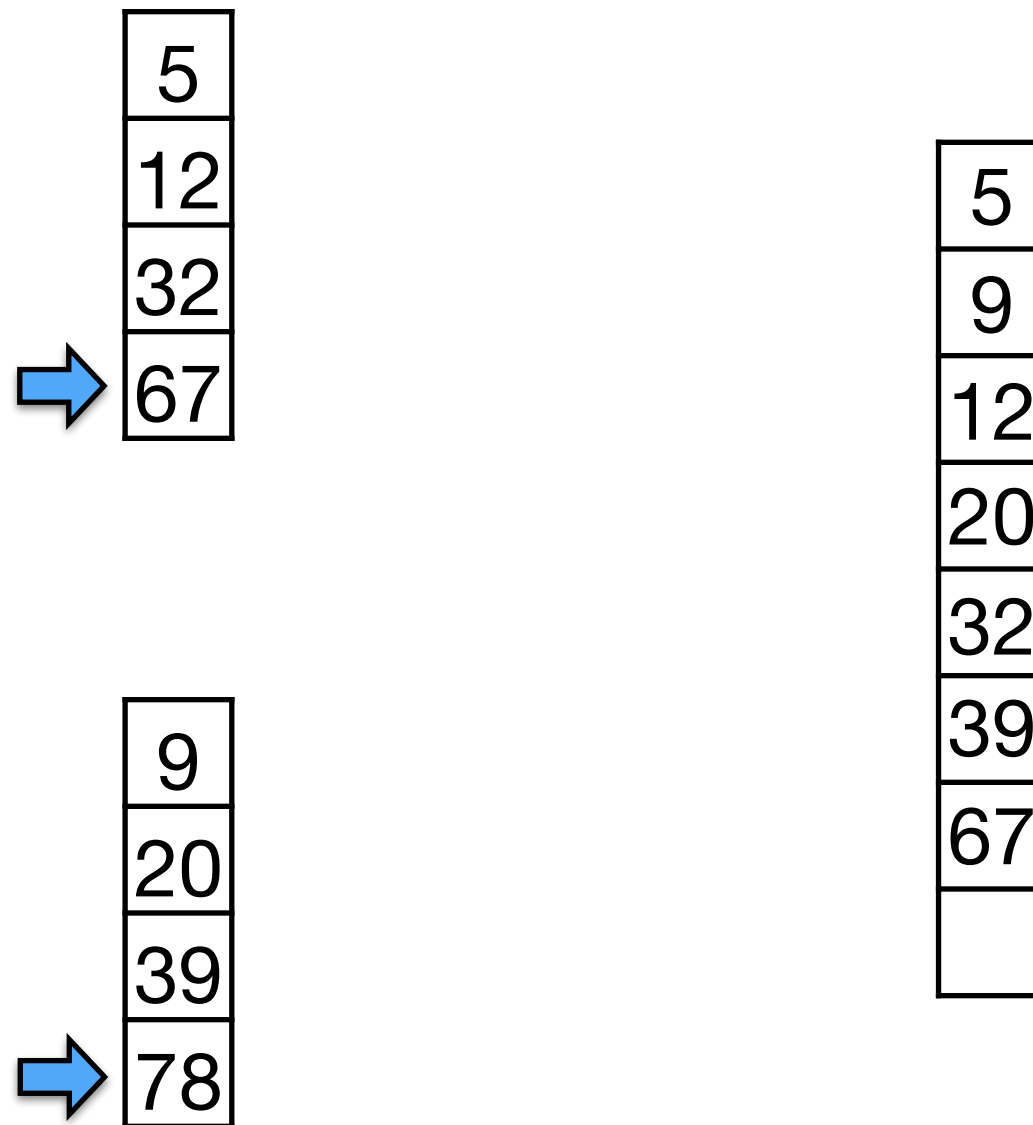


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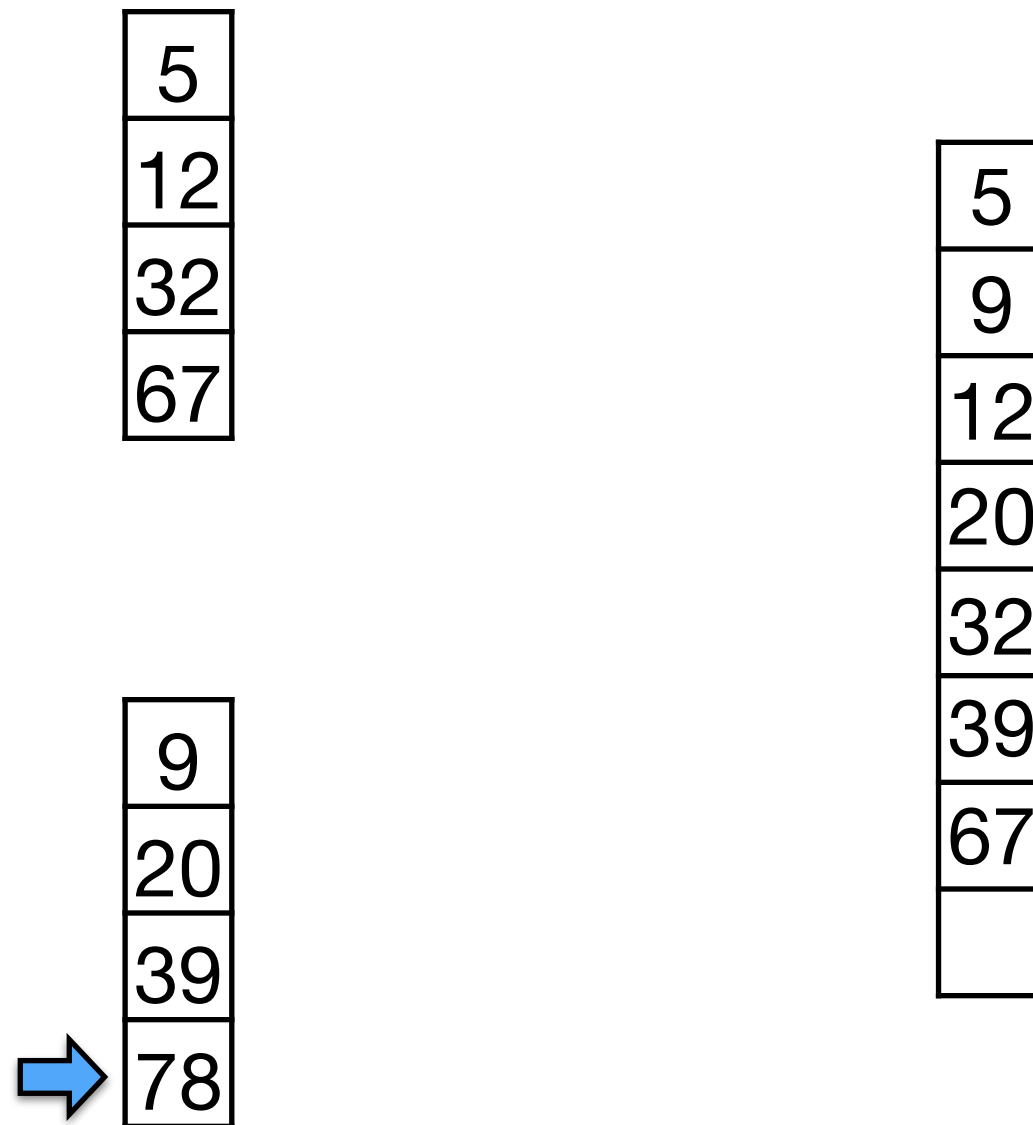




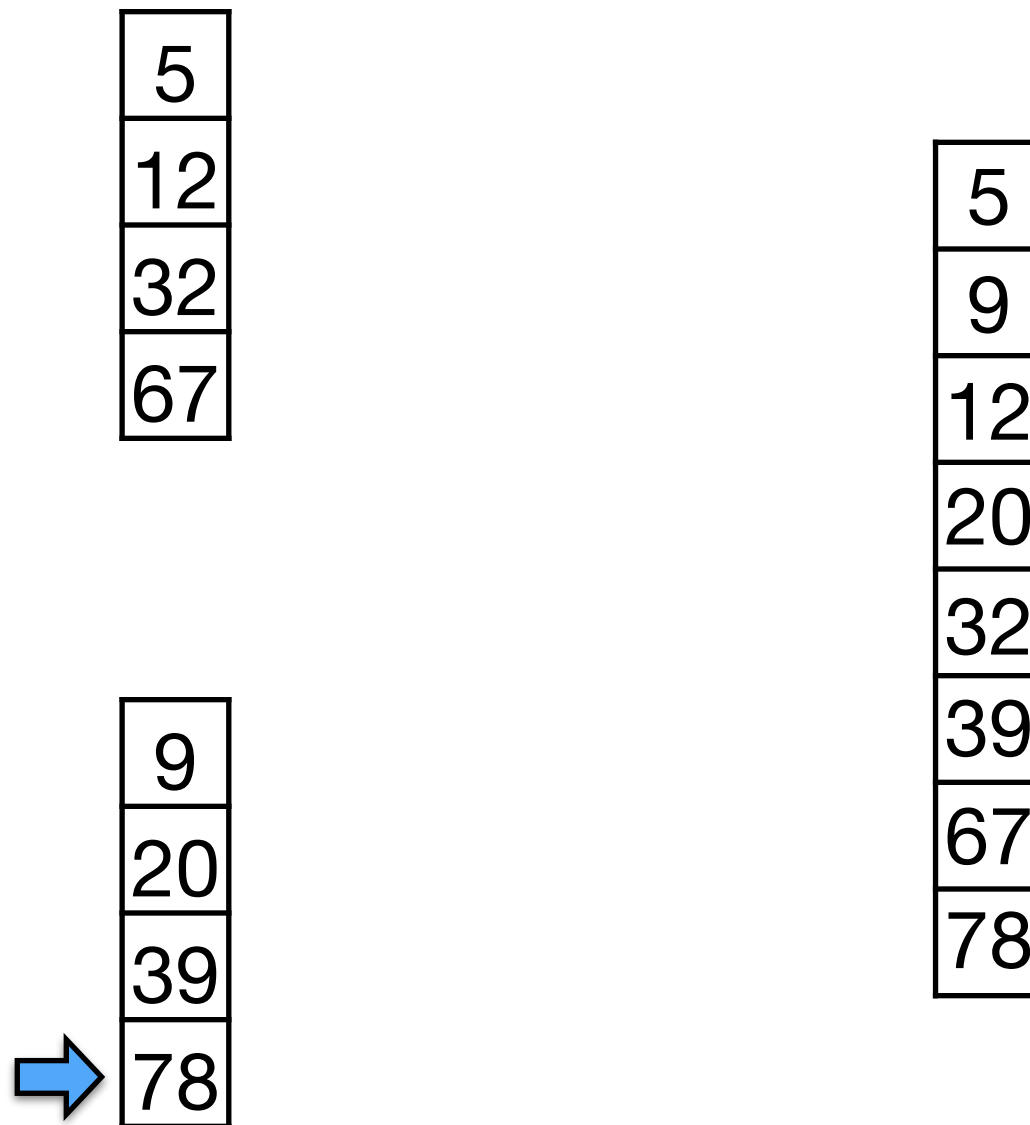
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- Thus, we have  $T(n) = 2T(n/2) + O(n)$

# Solving the Recursion

- Suppose  $T(n) = 2T(n/2) + O(n) \leq 2T(n/2) + cn$
- Keep expanding the RHS until it is a function of  $n$  and  $T(1)$ :

$$\begin{aligned} T(n) &\leq 2T(n/2) + cn \\ &\leq 2(2T(n/4) + c(n/2)) + cn \\ &= 4T(n/4) + 2cn \\ &\leq 4(2T(n/8) + c(n/4)) + 2cn \\ &= 8T(n/8) + 3cn \\ &\quad \dots \\ &\leq 2^i T(n/2^i) + icn \\ &\quad \dots \\ &= 2^{\log n} T(1) + cn \log n = O(n \log n) \end{aligned}$$