[illegible]

[illegible]

ESKO UKKONEN

[illegible]

The edit distance between strings [illegible] and [illegible] is the minimum cost of a [illegible]

[illegible] the order [illegible] well-known tabulating method computes [illegible] as well as the corresponding [illegible]

[illegible]

[illegible]

(ii) inserting a symbol [illegible] at position [illegible]

[illegible]

[illegible] (1968). The problem [illegible]

[illegible] sequence of editing steps has applications in various [illegible]

[illegible] assume that there is always an [illegible]

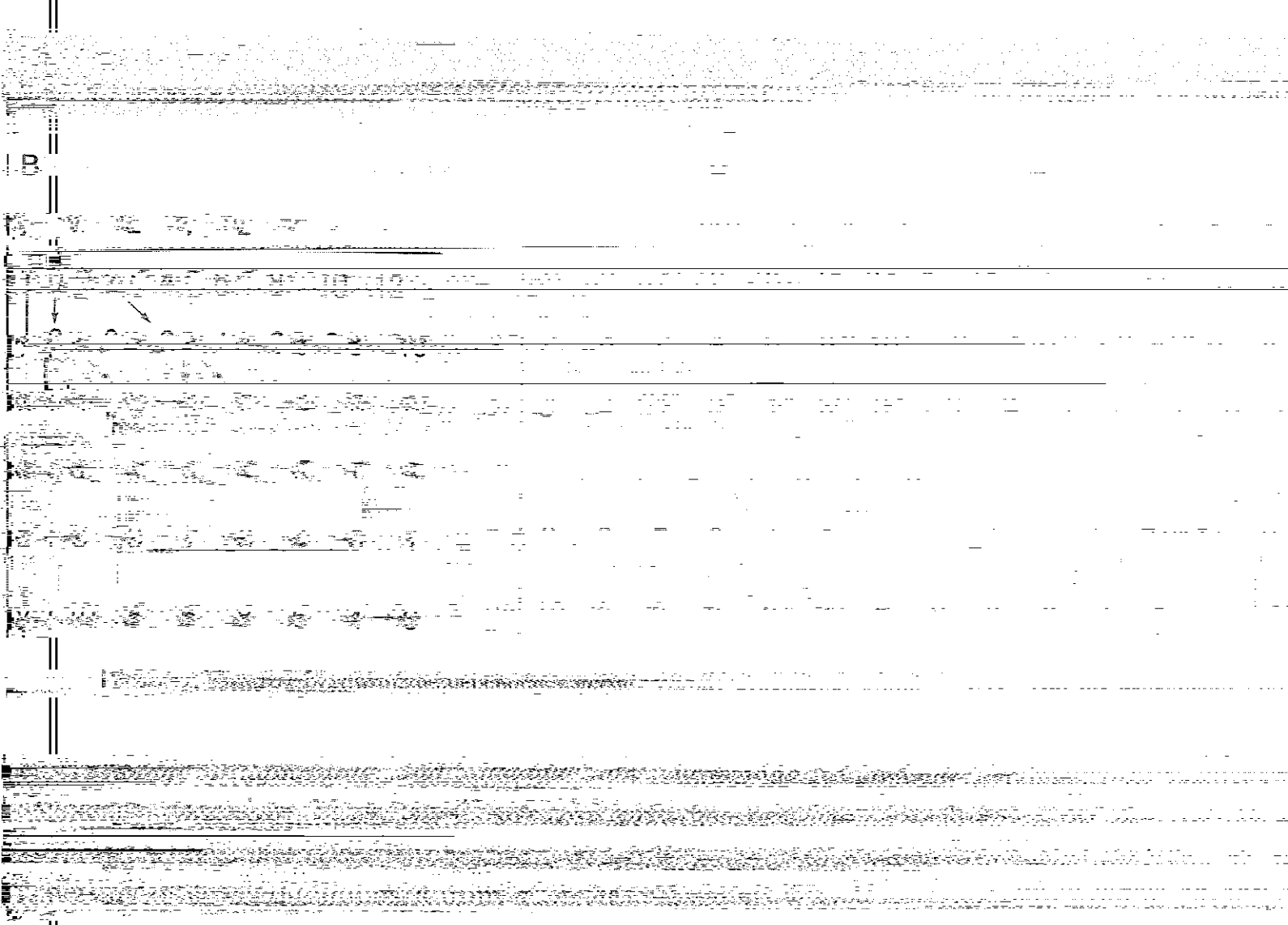
[illegible] not modified again. This means that all editing operations could be applied [illegible]

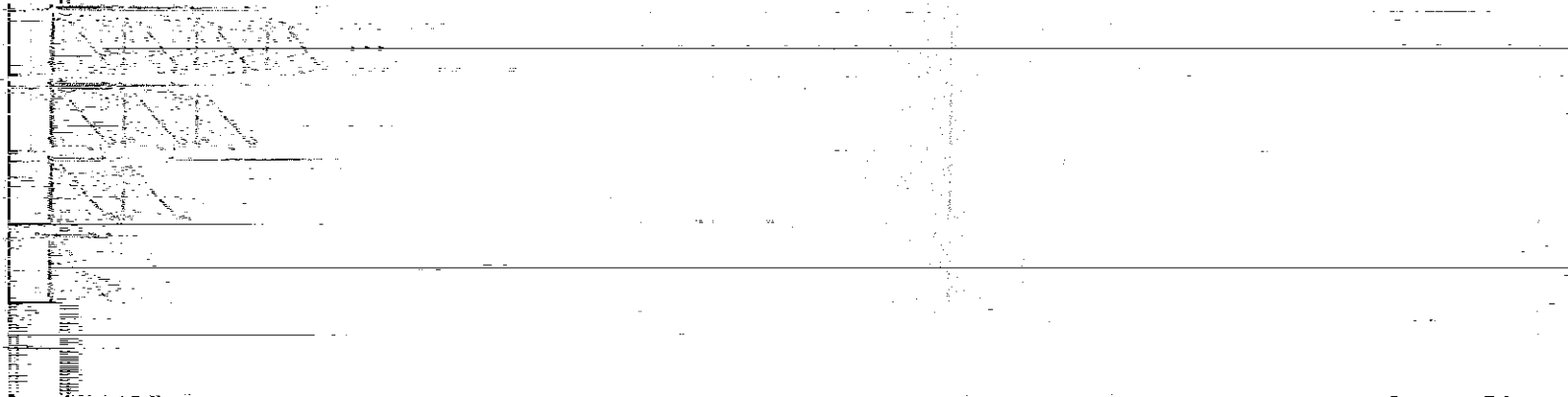
(1974).

As noted by Wagner and Fischer [illegible]

[illegible]

(3) [illegible]





[illegible]

[illegible] cannot belong to any path leading to [illegible]. Moreover [illegible]

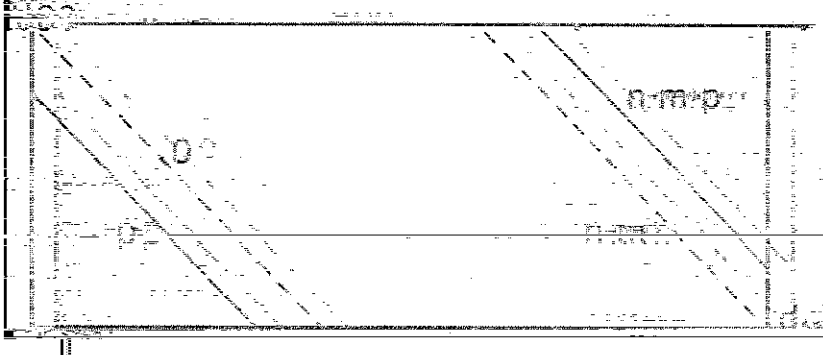
[illegible] the narrower the smaller [illegible]

[illegible]

(4)

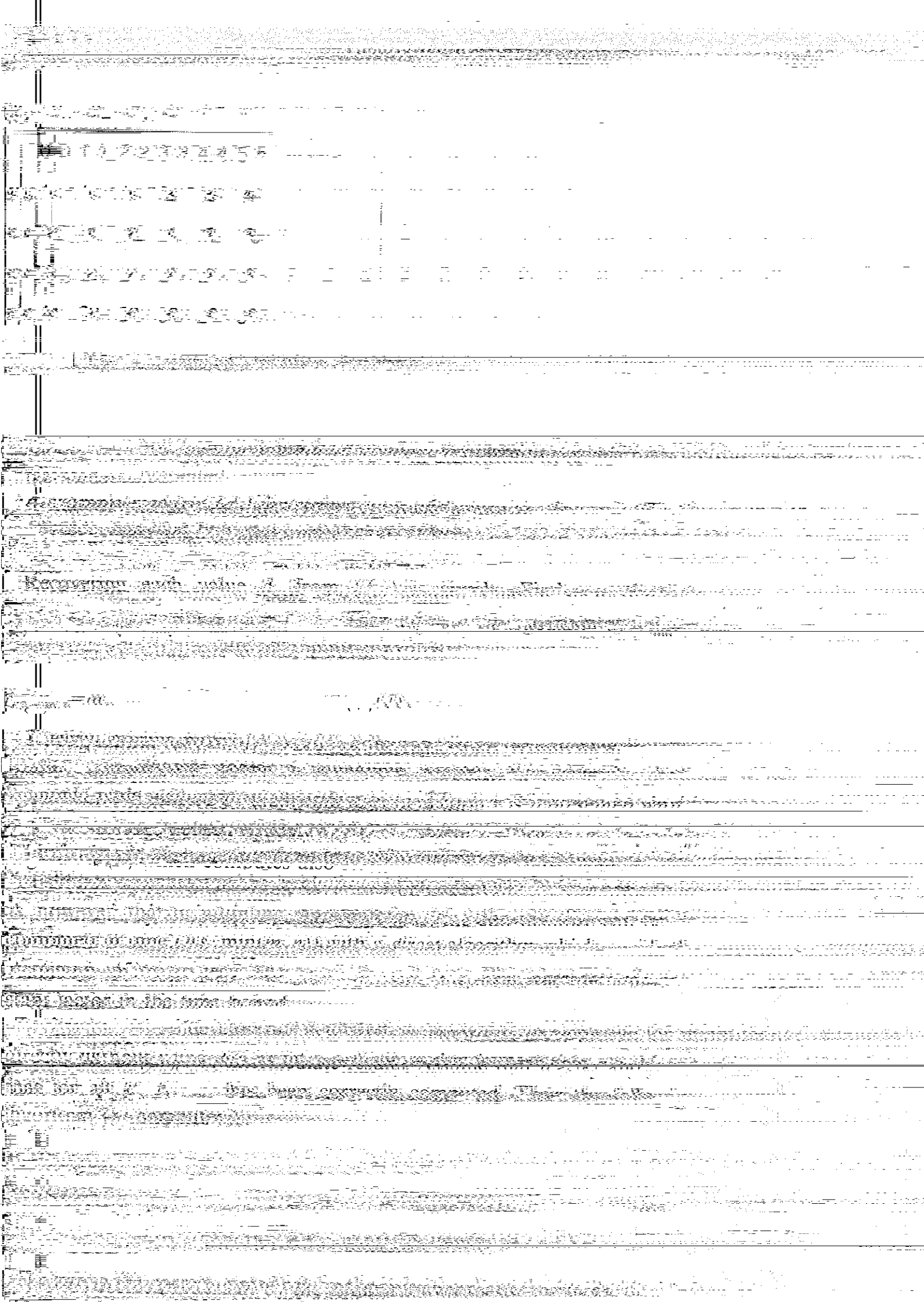
From Corollary 1 it follows that in testing whether $D(A, B) < t$, the

IF $t/\Delta < |n-m|$ THEN reject
ELSE
$k := \min(n-m, t/\Delta)$;
IF $n \geq m$ THEN $\lambda := p$ ELSE $\lambda := p-(n-m)$;
FOR $i := 0$ UPTO m DO
IF $i + k \geq 0$ THEN
ENDIF
IF $n-m+2p+k \leq t$ THEN accept ELSE reject ENDIF



(6)

The minimization step in (6) directly implies that γ_k cannot be larger



[illegible]

Denote by [illegible] the value of [illegible] after [illegible], that is, [illegible]

[illegible]

Assume, for example, that [illegible]

[illegible]

Consider then the problem of testing for a given message value c

most k. The method needs time $O(k \cdot \min(m,n))$ and space $O(\min(k,m)n)$.

for some constants c_1 and c_2 and for all $a \neq c$, $\delta(a \to c) = c_1$ and

[illegible]

[illegible] operation set [illegible] presented as

[illegible]

For an arbitrary [illegible]

[illegible]

There exists an equivalent

the results...

for similarities in the amino acid [illegible] general method applicable to the search