

Week 1-2: Logic and Truth Tables

< Section 2.1: Logical Form & Logical Equivalences >

[Warm Up] Consider the following.

"argument form"

Specific \rightarrow General

1, 2, 3, <u>4</u> , ..., <u>n</u>
, st , ^{2nd} , ^{3rd} , ^{4th} , ..., ^{nth}

"inductive logic"

General \rightarrow Specific

All men are mortal (premise 1)
Socrates is a man (premise 2)
 \therefore Socrates is mortal (conclusion)

"deductive reasoning"

Example 1

If P then Q

\hookrightarrow If all integers P are rational, then the number 1 is rational.
 \hookrightarrow All integers are rational $\xrightarrow{\text{?}} Q$

Therefore, the number 1 is rational

[Def^n] (statement)

A statement is a sentence that is true, or false, but not both

[Example 2] "Steve is a student." \rightarrow just a sentence
 \hookrightarrow could be both true and false
 \hookrightarrow not a statement

We can combine statements with "NOT", "AND", "OR"

p : = today is cold

q : = today is windy

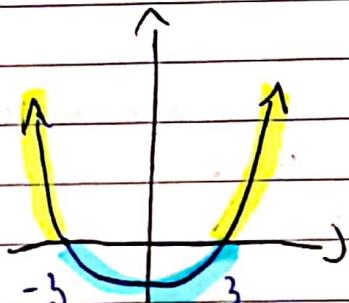
Meaning	Symbol	Name	Example
NOT	\sim	negation	$\sim p$: Today is <u>not</u> cold
AND	\wedge	conjunction	$p \wedge q$: Today is cold <u>and</u> today is windy
OR	\vee	disjunction	$p \vee q$: Today is cold <u>OR</u> today is windy.

Example 3 Write this sentence symbolically,

Today is neither cold nor windy

Today is not cold and not windy

$$[\sim p \quad \wedge \quad \sim q]$$



$x > 3$ means $[x > 3 \text{ or } x = 3]$

$-3 \leq x \leq 3$ means $[x > -3 \text{ and } x \leq 3]$

- Inequalities
- $x \leq a$ means $x < a$ or $x = a$
 - $a \leq x \leq b$ means $x > a$ and $x \leq b$

Example 4: let say we assign as follows:

$$\begin{aligned} p &:= 0 < x \\ q &:= x < 3 \\ r &:= x = 3 \end{aligned}$$

Write the following symbolically:

$$1. x \leq 3 \Rightarrow q \vee r$$

$$2. 0 < x < 3 \Rightarrow p \wedge q$$

$$3. 0 < x \leq 3 \Rightarrow p \wedge (q \vee r)$$

Truth Values of Truth Table

To find the truth values of statements, we use "truth table"

(Note: Each statement is either true or false.)

p	mp
T	F
F	T

Truth table for two statements: p, q

p	q	$p \wedge q$	$p \vee q$
T	T	T	T
T	F	F	T
F	T	F	T
F	F	F	F

$$(p \vee q) \wedge \neg(p \wedge q)$$

Example 5! Build the truth table for $(p \vee q) \wedge \sim(p \wedge q)$
 (Tip! start with $p, q, p \vee q, p \wedge q$, and even)

p	q	$p \vee q$	$p \wedge q$	$\sim(p \wedge q)$	$(p \vee q) \wedge \sim(p \wedge q)$
T	T	T	T	F	F
T	F	T	F	T	T
F	T	T	F	T	T
F	F	F	F	T	F

- This is the truth table for "exclusive or" $p \oplus q$
 - "one or the other, but not both!"

Table-Item Example | Build a truth table for the statement
 $(p \wedge q) \vee \sim r$

p	q	r	$p \wedge q$	$\sim r$	$(p \wedge q) \vee \sim r$
T	T	T	T	F	T
T	T	F	T	T	T
T	F	T	F	F	F
T	F	F	F	T	T
F	T	T	F	F	F
F	T	F	F	T	T
F	F	T	F	F	F
F	F	F	F	T	T

Week 1-3: Logical Equivalences of Conditional Statements

<sec 2.1: Logical Equivalences>

Basic idea: If two statement forms have identical truth values, then they are logically equivalent.

P	$\sim P$	$\sim(\sim P)$
T	F	T
F	T	P

↑ ↑
same truth values

$\therefore P \text{ is equivalent to } \sim(\sim P)$
 $P \equiv \sim(\sim P)$

"Useful Negation Laws: De Morgan's Laws"

De Morgan's Laws

$\sim(p_1 \&) \equiv \sim p_1 \vee \sim q$	warning! $\sim(p_1 \&) \equiv \sim p_1 \sim q$
$\sim(p \vee q) \equiv \sim p \& \sim q$	wrong!

P	Q	$\sim P$	$\sim Q$	$p \& q$	$\sim(p \& q)$	$\sim P \vee \sim Q$	$\sim P \& \sim Q$
F	T	T	F	F	T	F	F
T	F	F	T	F	T	T	F
F	T	T	F	F	T	T	F
F	F	T	T	F	T	F	T

This shows that
 $\sim(p \& q) \equiv \sim p \vee \sim q$ $\sim p \& \sim q$

Example 6: Use De Morgan's laws to write negation of $-1 < x \leq 4$.

$$\left. \begin{array}{l} P := -1 < x \\ Q := x \leq 4 \\ R := x = 4 \end{array} \right\} \quad \left. \begin{array}{l} -1 < x \leq 4 \text{ means} \\ P \wedge (Q \vee R) \end{array} \right.$$

$$\begin{aligned} \text{so } \neg [P \wedge (Q \vee R)] &\equiv \neg P \vee \neg (Q \vee R) && \text{By De Morgan's laws} \\ &\equiv \neg P \vee \neg Q \wedge \neg R && \neg\neg - \\ &&& \downarrow \quad \downarrow \quad \downarrow \\ && x \leq -1 \text{ or } (x \geq 4 \text{ and } x \neq 4) & \end{aligned}$$

Ans: $x \leq -1 \text{ or } x > 4$

Tautologies and Contradictions

- Tautology: is a statement form that is always true
- Contradiction: is a statement form that is always false

P	$\neg P$	$P \vee \neg P$	$P \wedge \neg P$
T	F	T	F
F	T	T	F

$\vee \quad \wedge$

tautology contradiction

$P \vee \neg P \equiv T$

$P \wedge \neg P \equiv F$

$X \cdot 1 = X$

- $P \wedge T \equiv P$
- $P \wedge C \equiv C$

Theorem 2.1.1.
Logical equivalences

statement variables: p, q, r tautology A , contradiction c

1. Commutative laws $p \wedge q \equiv q \wedge p$ $p \vee q \equiv q \vee p$

2. Associative laws $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$
 $(p \vee q) \vee r \equiv p \vee (q \vee r)$

3. Distributive laws $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$
 $p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$

4. Identity laws $p \wedge t \equiv p$
 $p \vee c \equiv p$

5. Negation laws $p \vee \sim p \equiv t$
 $p \wedge \sim p \equiv c$

6. Double negative law: $\sim(\sim p) \equiv p$

7. Idempotent laws: $p \wedge p \equiv p$ $p \vee p \equiv p$

8. Universal bound laws: $p \vee t \equiv t$ $p \wedge c \equiv c$

9. De Morgan's laws: $\sim(p \wedge q) \equiv \sim p \vee \sim q$
 $\sim(p \vee q) \equiv \sim p \wedge \sim q$

10. Absorption laws: $p \vee (p \wedge q) \equiv p$

$$p \wedge (p \vee q) \equiv p$$

11. Negations of t and c : $\sim t \equiv c$ $\sim c \equiv t$

Example 7: Use Theorem 2.1.1. to simplify $\sim(p \wedge q) \vee (p \vee q)$

$$\begin{aligned}
 \sim(\sim(p \wedge q)) \vee (p \vee q) &= [\sim(\sim p) \vee \sim q] \vee p \vee q && \text{By DeMorgan's Law} \\
 &\equiv (p \vee \sim q) \vee p \vee q && \text{By Double Negative Law} \\
 &\equiv p \vee (\sim q \vee q) && \text{By Distributive Laws} \\
 &\equiv p \vee c && \text{By Negation Law} \\
 &\equiv p && \text{by Identity Law}
 \end{aligned}$$

<Section 2.2 Conditional Statements>

statement form: if p then q
symbol: $p \rightarrow q$

This is conditional because the truth of statement q is conditioned on the truth of p .

Warm Up Example: If you study hard, then you will earn an A $\rightarrow q$

Question: When is this statement false?

Answer: You studied hard, but not getting an A

"us iteratur
cogit u cbe" Truth Table for $p \rightarrow q$

p	q	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

Warm Up Example: Determine true or false

if $1=2$ then $3=0$
 $\swarrow p$ $\swarrow q$

Note: Hypothesis is False. This means $p \rightarrow q$ is true by default! (i.e. vacuously true)

+ Order of Operation

1. \sim
2. $\wedge \vee$
3. $\rightarrow \leftarrow \leftrightarrow$

Example 1: Construct a truth table for $p \vee \sim q \rightarrow \sim p$

p	q	$\sim p$	$\sim q$	$p \vee \sim q$	$\sim p$	$p \vee \sim q \rightarrow \sim p$
T	T	F	F	T	F	F
T	F	F	T	T	F	F
F	T	T	F	F	T	T
F	F	T	T	T	T	T

hypothesis conclusion Answer

Your turn! Find the truth tables for

a) $p \rightarrow q$ b) $\sim p \vee q$

p	q	$p \rightarrow q$	$\sim p$	$\sim p \vee q$
T	T	T	F	T
T	F	F	F	F
F	T	T	T	T
F	F	T	T	T

$p \rightarrow q \equiv \sim p \vee q$ \because they are logically equivalent!

$$\begin{aligned}
 \sim(p \rightarrow q) &\equiv \sim(\sim p \vee q) \text{ By De Morgan's Law} \\
 &\equiv \sim(\sim p) \wedge \sim q \text{ by Double Negative Law} \\
 &\equiv p \wedge \sim q
 \end{aligned}$$

$\boxed{\sim(p \rightarrow q) \equiv p \wedge \sim q}$

Example 2: Negate the statement

If my car is in the repair shop, then I cannot get to class.

p

q

$$\neg(p \rightarrow q) \equiv p \wedge \neg q$$

Ans.

My car is in the repair shop AND I can't get to class.

Converse, Contrapositive, and Inverse of $p \rightarrow q$

Example 3: Consider the statement.

If today is Easter, then tomorrow is Monday.

p

q

Type	Symbol	Example
Contrapositive	$\neg q \rightarrow \neg p$	If tomorrow is <u>not</u> Monday, then today is <u>not</u> Easter.

Converse	$q \rightarrow p$	If tomorrow is Monday, then today is Easter.
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Inverse	$\neg p \rightarrow \neg q$	If today is <u>not</u> Easter, then tomorrow is <u>not</u> Monday.
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p is sufficient for q.

p is not necessary for q.

Biconditional Statements

Not tip! "only if" " \rightarrow " } "if and only if"
"if" " \leftarrow " } \leftrightarrow

$P \text{ only if } Q$

- P can happen only if we have Q
- if we don't have Q , then P can't happen
- [contrapositive] if $\neg P$, then $\neg Q$ $(P \rightarrow Q)$

Versus:

$P \text{ if } Q$

- P can happen if we have Q (sufficient)
- Q is sufficient to have P .
- if Q , then P .
- $Q \rightarrow P$

Combine them

$$P \leftrightarrow Q$$

This is called "biconditional of P and Q "

Summary

- Symbol: \leftrightarrow (it stands for iff)

write: $p \leftrightarrow q$

$$\underline{p \rightarrow q \text{ and } q \rightarrow p}$$

Example 4: Construct truth table for $p \leftrightarrow q$
(Tip! $p \leftrightarrow q$ is equivalent to $p \rightarrow q$ and $q \rightarrow p$)

p	q	$p \rightarrow q$	$q \rightarrow p$	$(p \rightarrow q) \wedge (q \rightarrow p)$
T	T	T	T	T
T	F	F	T	F
F	T	T	F	F
F	F	T	T	T

Sufficient and Necessary Conditions

[Def] if p and q are statements:

- p is a sufficient condition for q : if p then q
- p is a necessary condition for q : if $\neg p$ then $\neg q$ [If $\neg q$ then $\neg p$]

Example 5 Consider $p \rightarrow q$ where

$p \rightarrow q$ X

p : = Today is Easter

$p \rightarrow q$ ✓

q : = Tomorrow is Monday

- Today being Easter is sufficient for tomorrow to be Monday
- Check for necessary condition:
if today is not Easter, then tomorrow is not Monday $\neg p$ X
NOT Today is not Easter is not necessary for tomorrow to be Not Monday