

---

# Deep into ethereum light client protocol

---

Gary Rong  
[garyrong@ethereum.org](mailto:garyrong@ethereum.org)  
rjl493456442

---

# Agenda

---

- ❖ Introduction

- ❖ Merkle Trie & Merkle Proof

- ❖ Checkpoint Synchronization

- ❖ What can light client do

- ❖ Flow control & Capacity management model

# 1. Introduction



# What's light client protocol

- ❖ The protocol used by “light client”
  - ❖ Low resource requirement
  - ❖ Data verification capability
- ❖ Download and check the validity of the block headers
- ❖ Don't check the validity of state transition
- ❖ Fetch and proof the other part of blockchain on-demand
- ❖ Regard the whole DHT as database, local database as cache

# Category of Light Client Protocol

- ❖ Les(Light Ethereum SubProtocol)
- ❖ PIP(Parity Light Protocol)



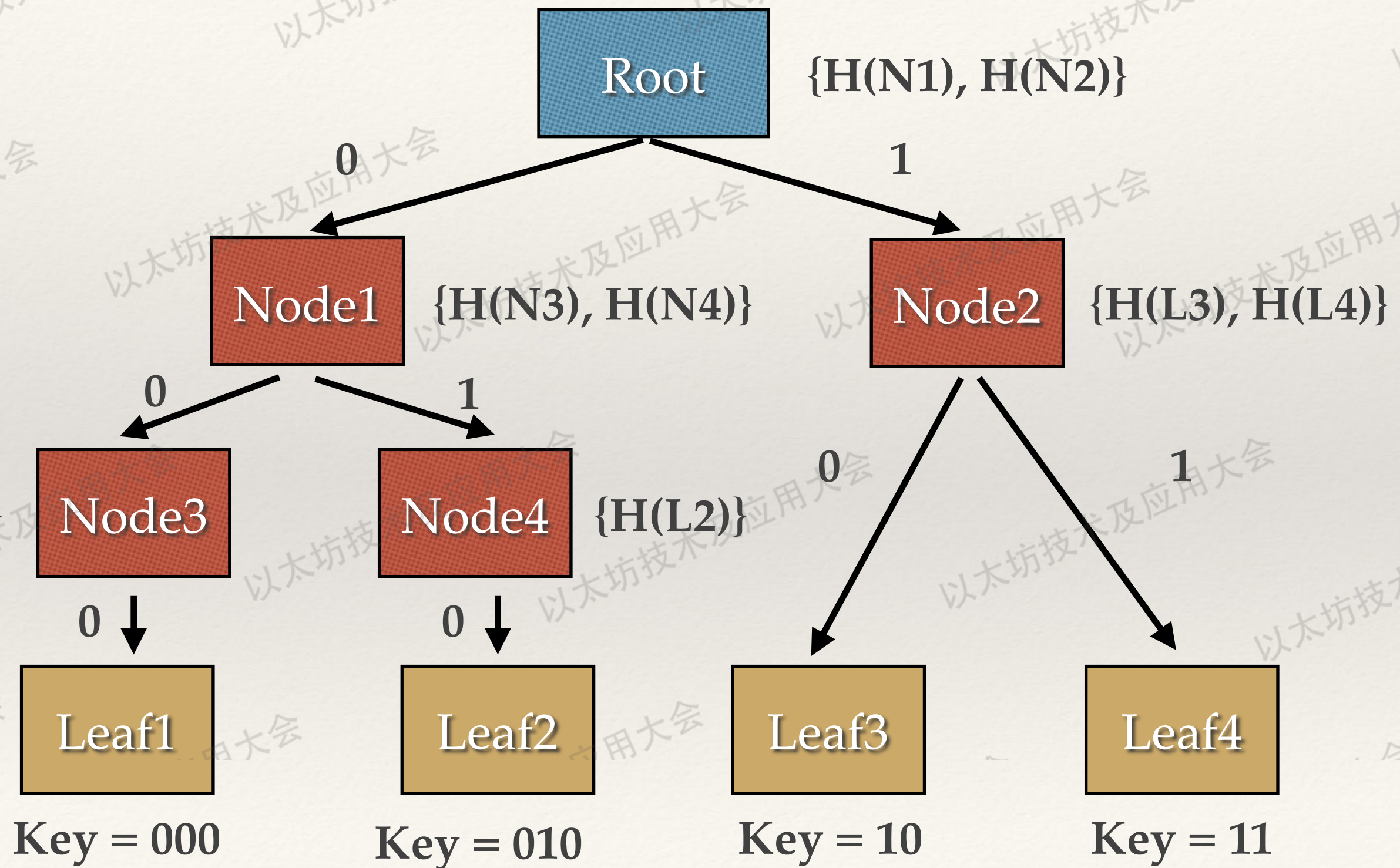
# Classification of nodes

Type	Database size	Sync Time	Security Guarantee
Archive node	~2.3TB	~13days	High
Full node	~131GB	~4hours	Medium
Light client	~50MB	~1minute	Low

## 2. Merkle Trie & Merkle Proof

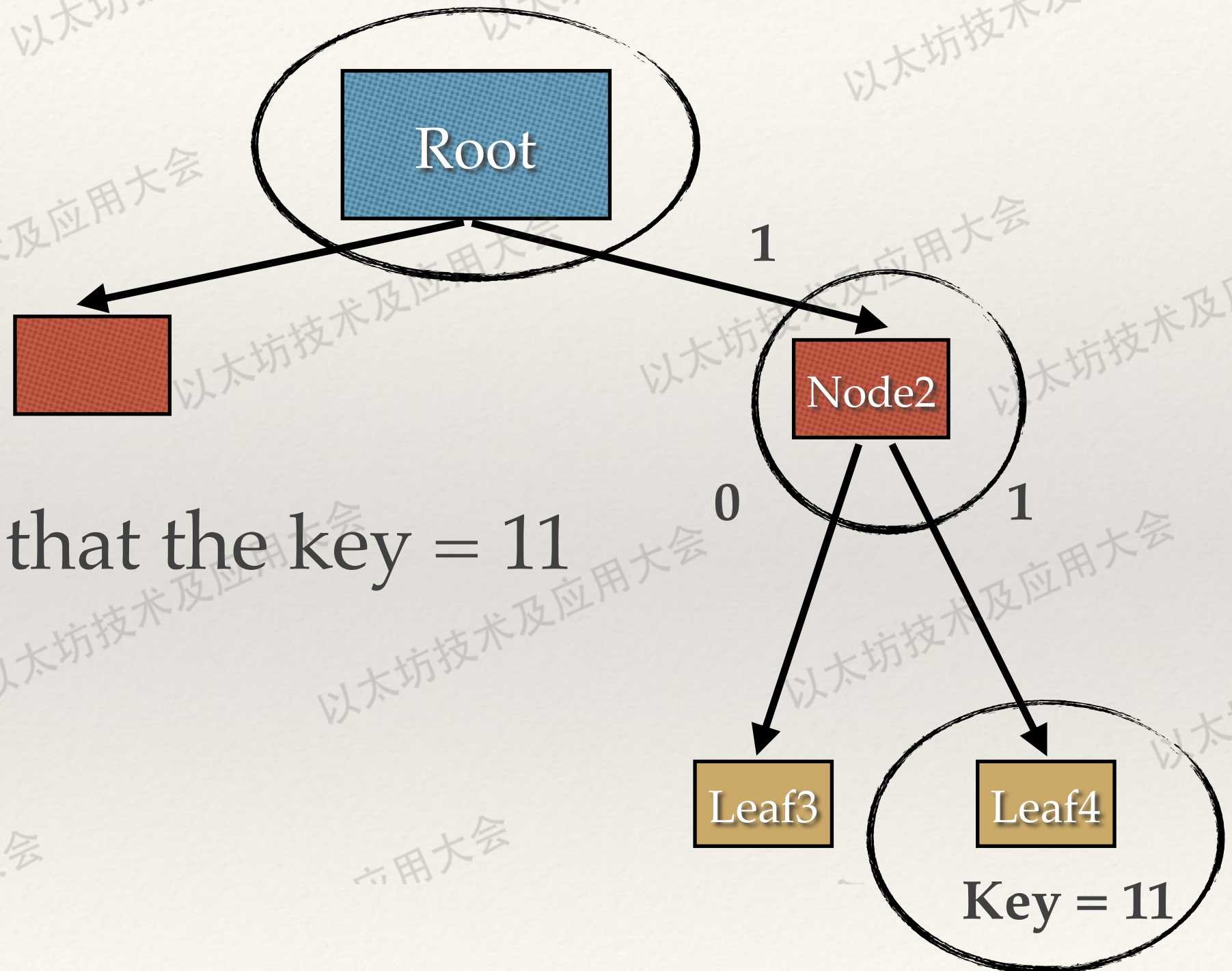


# Merkle Trie





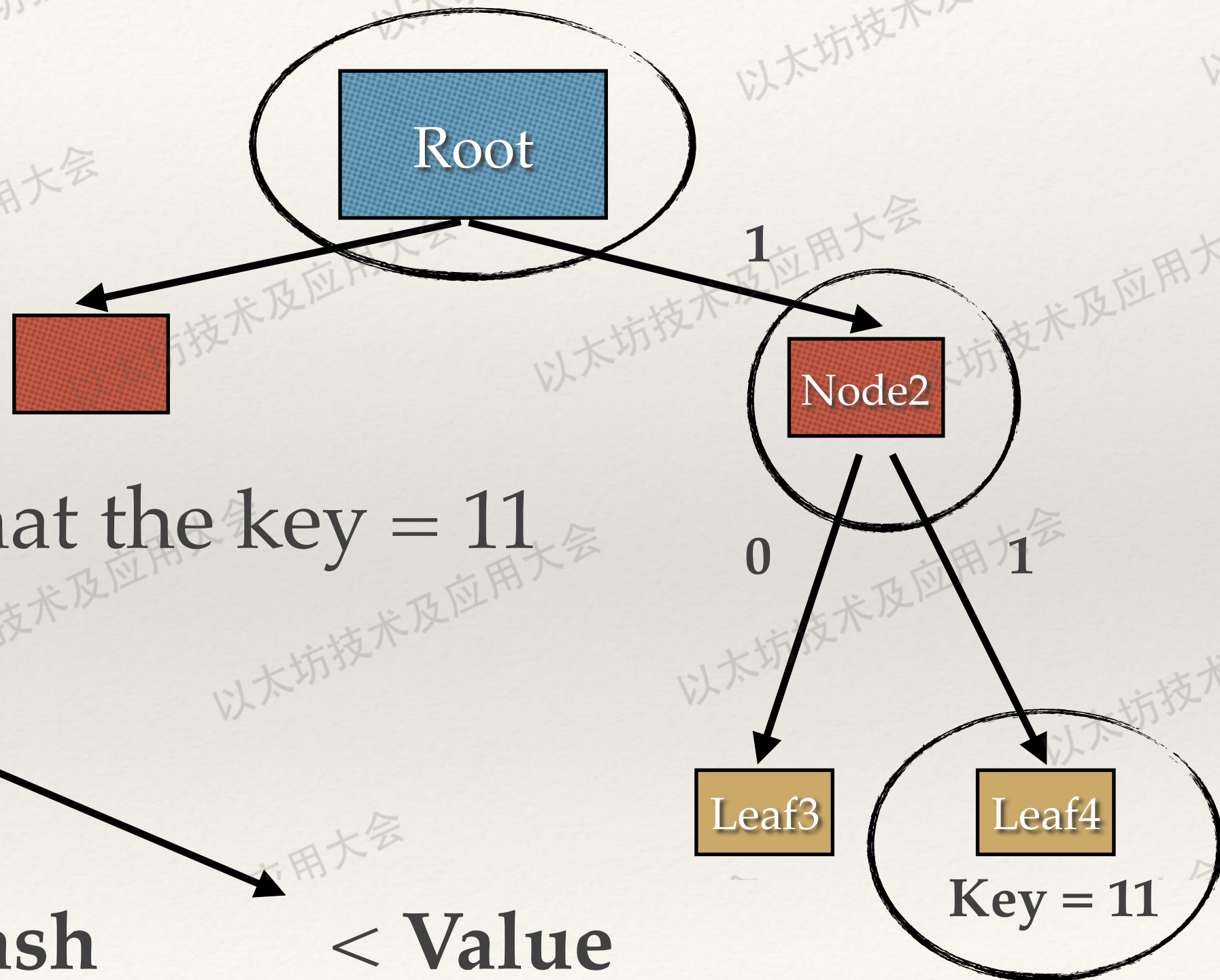
# Merkle Proof (generation)



> Proof node that the key = 11

< [R, N2, L4]

# Merkle Proof (verification)



> Verify node that the key = 11

> [R, N2, L4]

> Local Root Hash

< Value



# 3. Checkpoint Synchronization

# Checkpoint sync

- ❖ Start from a checkpoint
  - ❖ Trusted
  - ❖ Can verify the validity of historical chain
- ❖ Download block headers only
- ❖ One hundredth of the header is PoW verified(Network bandwidth  $\gg$  Ethash bandwidth)



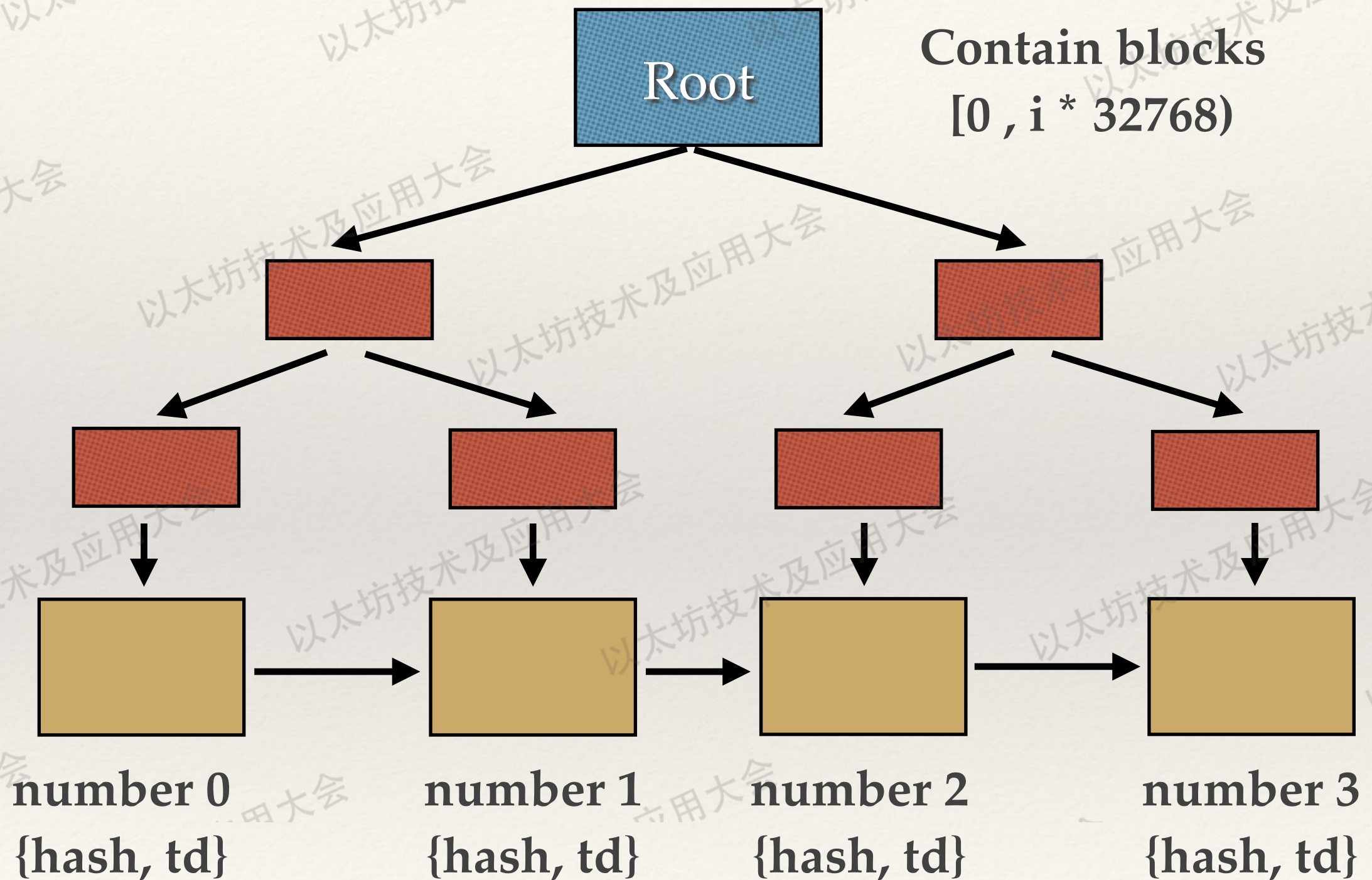
---

# Checkpoint

---

- ❖ Generate a checkpoint every 32768 blocks(server or light client)
- ❖ Section index
- ❖ Cumulative canonical hash trie root
- ❖ Cumulative bloom trie root

# Canonical Hash Trie





# Checkpoint registration

## Approach

## Pros

## Cons

Hardcode

Simple

1. Update according to release cycle
2. Old client can't use new checkpoint
3. Centralized

Checkpoint Oracle

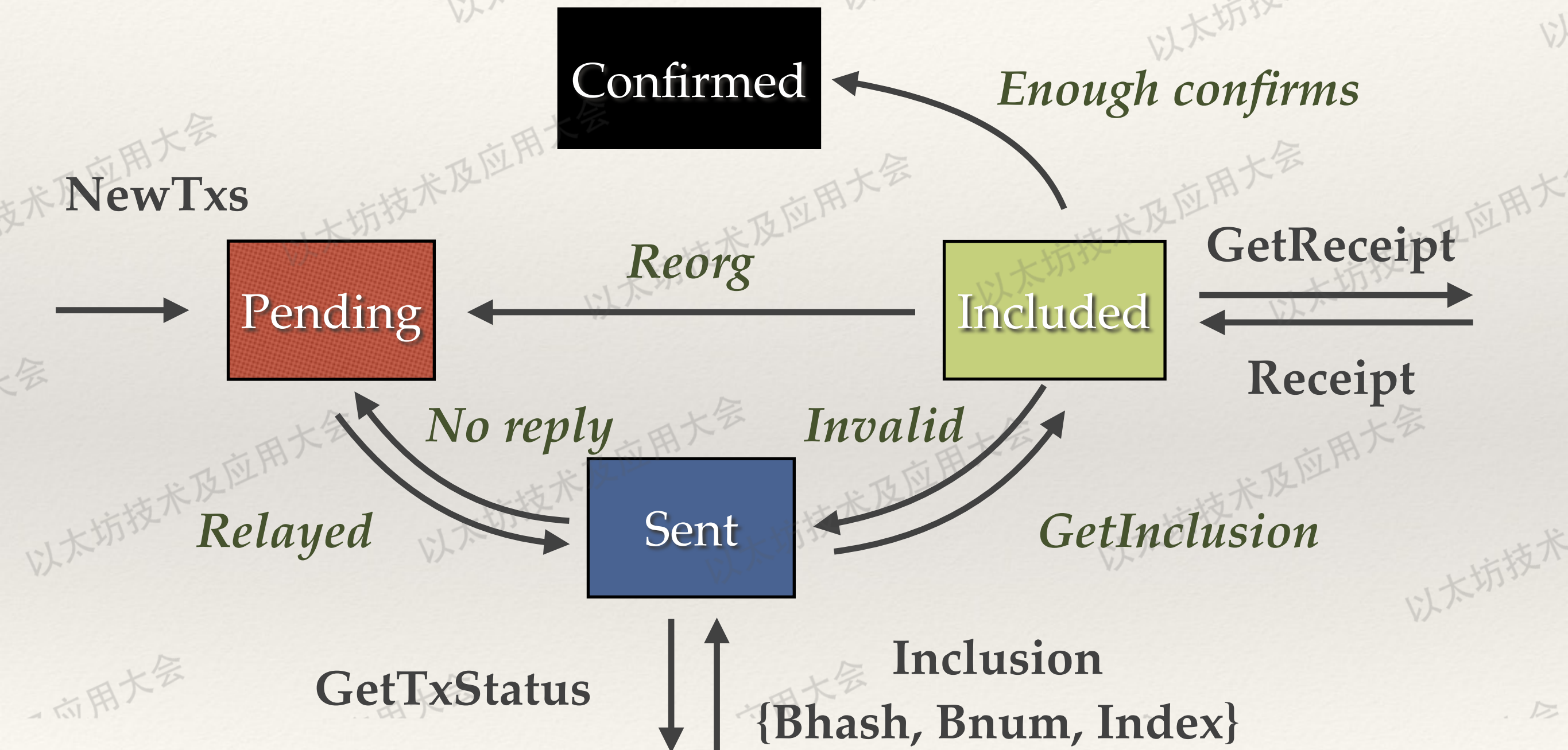
1. Update according to admin registration
2. Old client can always use latest checkpoint

1. Complicated
2. Still centralized

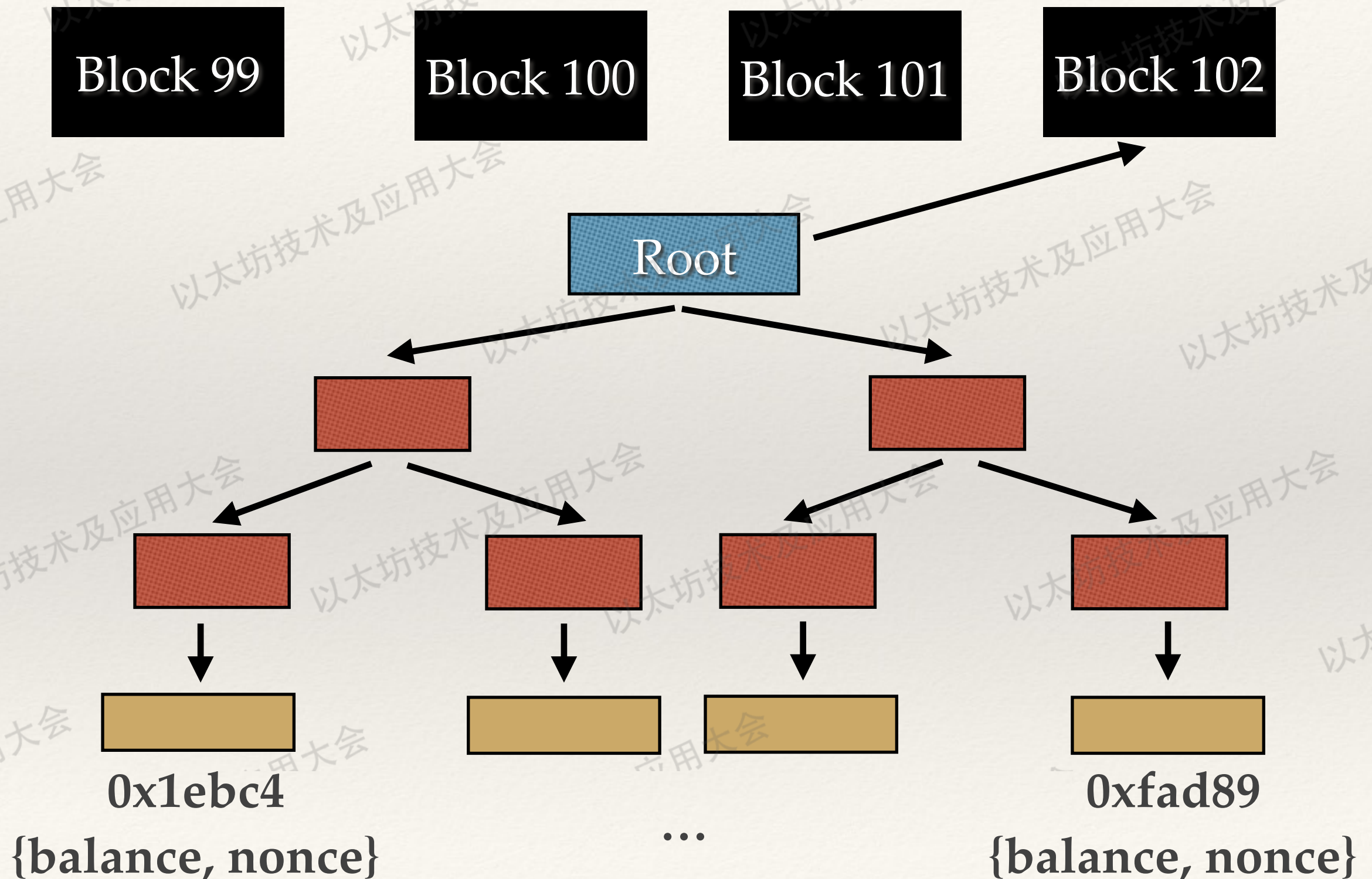
## 4. What can light client do



# Transaction relay

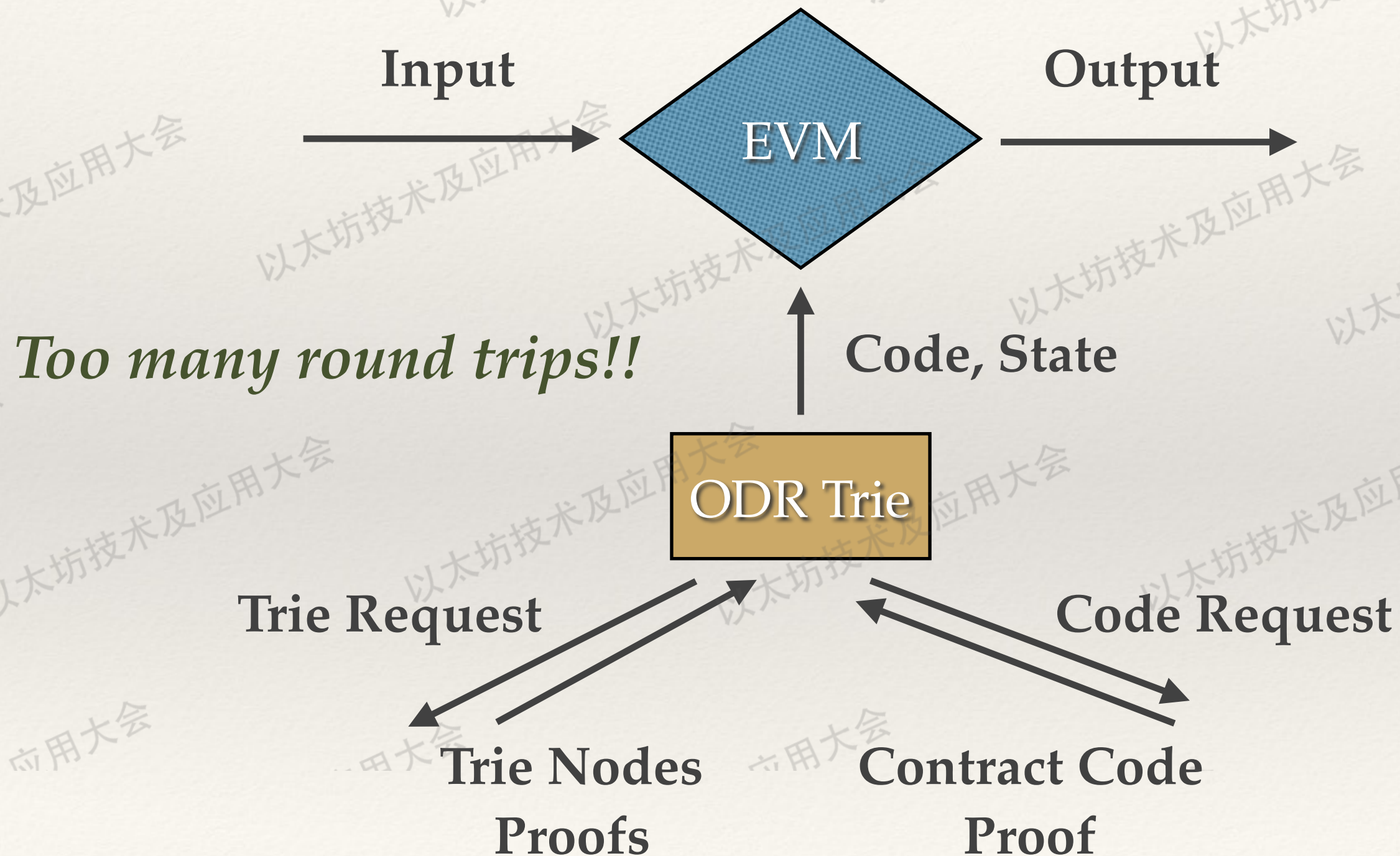


# State query





# Contract call



---

# Event subscription

---

- ❖ Event logs are embedded in the receipts
- ❖ The bloom for logs is included in headers
- ❖ Match the header bloom first
- ❖ Download and verify the target receipts
- ❖ Filter logs

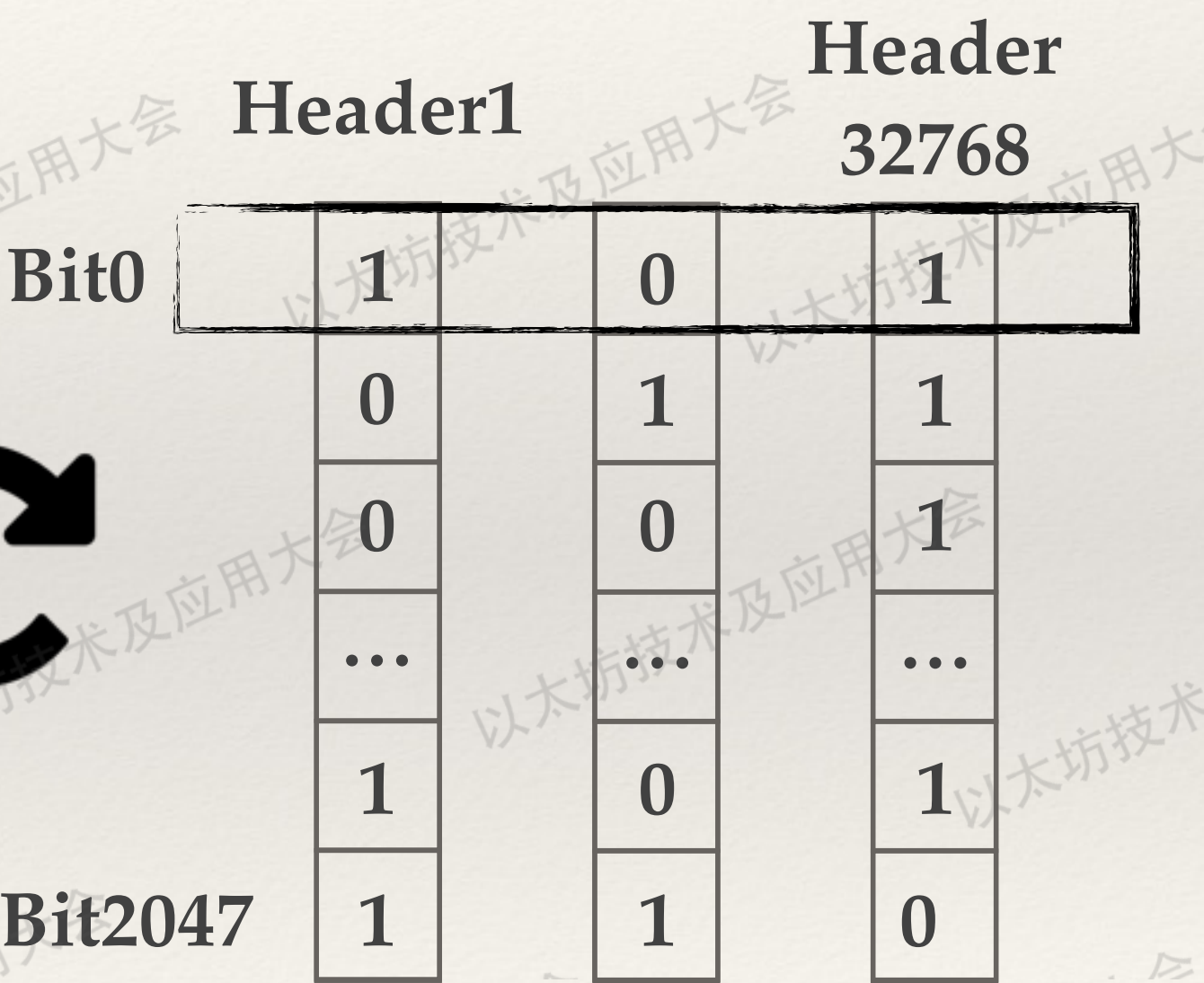


# Event searching(history)

*Bloom filter*



*Rotated Bloom filter*

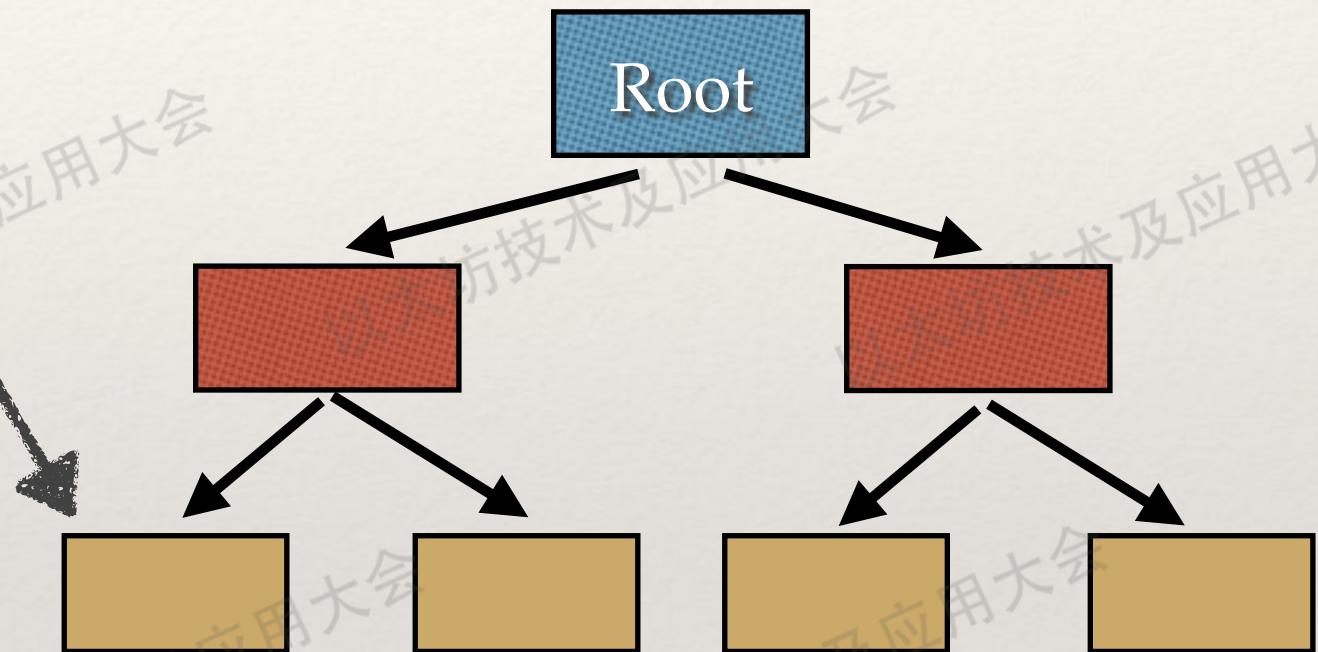


# Event searching contd.

## Bloom Section

Header1			Header		
			32768		
Bit0	1	0	1		
	0	1	1		
	0	0	1		
	...	...	...		
	1	0	1		
Bit2047	1	1	0		

## Bloom Trie



Key = BitIndex + SectionIndex  
Value = Rotated bloom

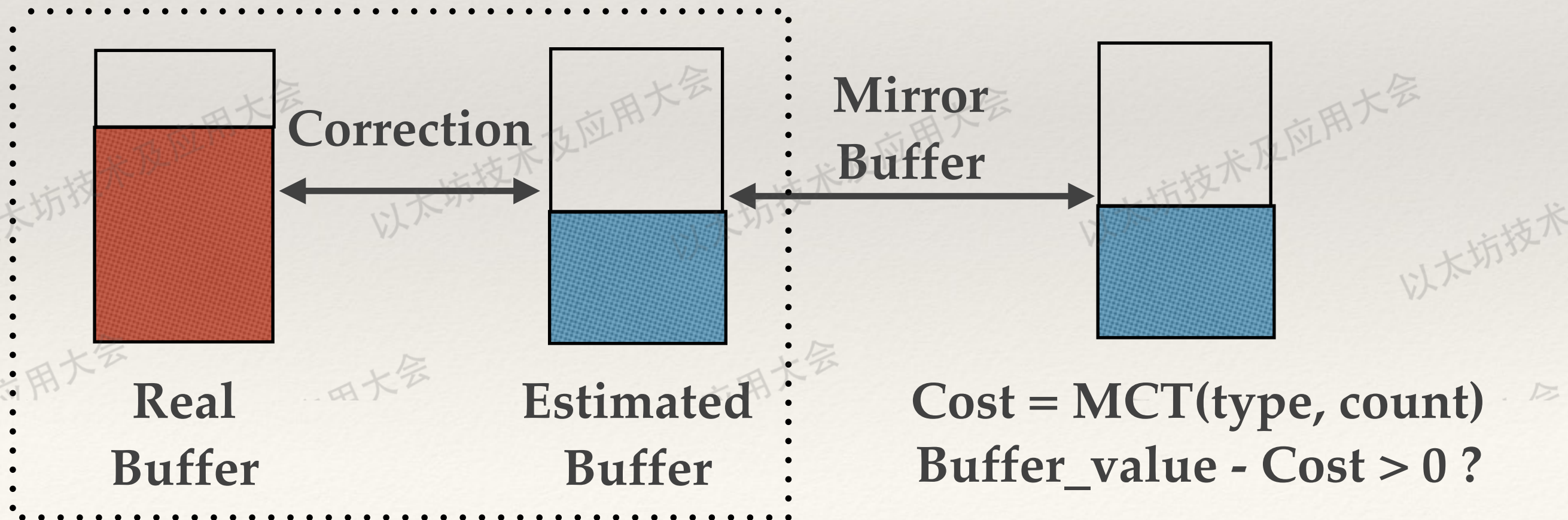


# 5. Flow Control & Capacity management

# Flow control model

*Server*  $\longleftrightarrow$  Handshake *Client*

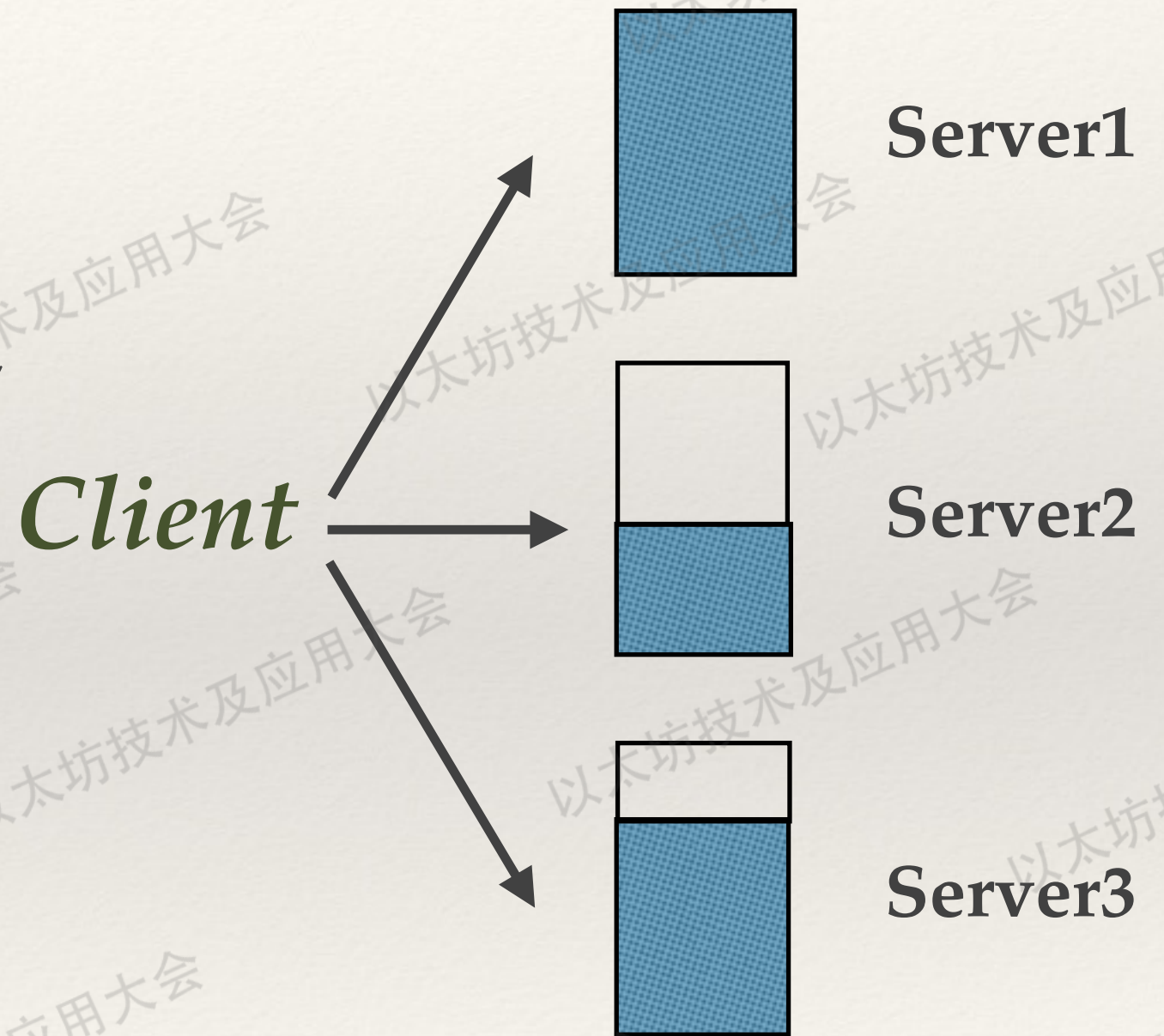
1. BufferLimit
2. MinRecharge
3. MaxCostTable





# Flow control model contd.

- ❖ Local Load Balancer
- ❖ Higher buffer remaining, higher priority
- ❖ Avoid low process priority and high response latency



# Capacity management

- ❖ —lightServ

- ❖ Percentage of serving light client requests per second

- ❖ Multi-threaded processing allows value over 100

- ❖ —lightbwout, —lightbwin

- ❖ Network bandwidth limitation

- ❖ —lightpeers

- ❖ Maximum number of light clients



# Capacity management contd.

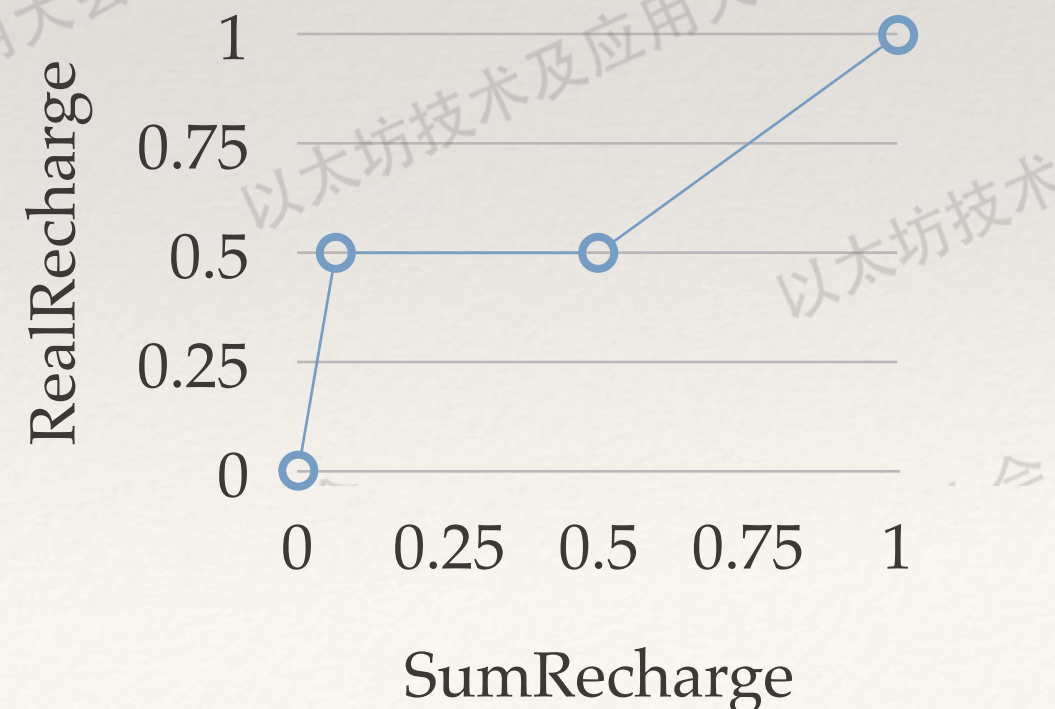
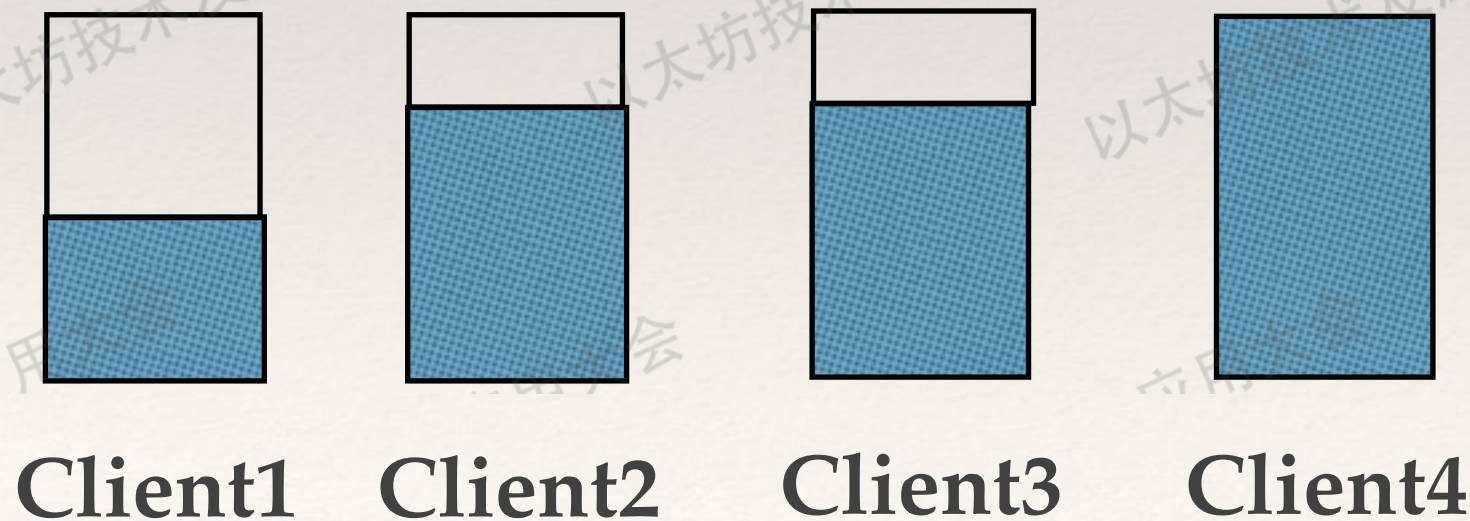
## *Cost calculation*

- ❖  $\text{TimeCost} = \text{ServingTime}$
- ❖  $\text{TrafficCost} = \text{PacketSize} * \text{TrafficFactor}$
- ❖  $\text{Cost} = \text{Max}(\text{TimeCost}, \text{InTrafficCost}, \text{OutTrafficCost})$

# Capacity management contd.

## *Buffer recharging*

- ❖  $\text{TotalRecharge} = \text{LightServ} * 1000,000,000 * \text{Factor} / 100$
- ❖  $\text{Bonus} = \text{Curve}(\text{SumRecharge}) / \text{SumRecharge}$
- ❖  $\text{RealRecharge} = \text{MinRecharge} * \text{Bonus}$
- ❖ Enjoy recharge bonus if the server is idle





# Capacity management contd.

## *Overbooking and Freezing*

- ❖ The maximum capacity of system is determined by the maximum number of peers and the minimal capacity.
- ❖ The maximum capacity of system can be much higher than total recharge.
- ❖ Freeze and kick out some clients if the **recent serving time** plus the **queued estimated time** exceeds the limitation.

# Capacity management contd.

## *Free client and Priority client*

- ❖ Higher buffer limit and recharge speed
- ❖ Kick out free client first
- ❖ Economic incentive for full node



用大会

以太坊技术及应用大会

以太坊技术及应用大会

以太坊技术及应用大会

以太坊技术及应用大会

技术及应用大会

以太坊技术及应用大会

以太坊技术及应用大会

以太坊技术及应用大会

以太坊技术及应用大会

大会

以太坊技术及应用大会

以太坊技术及应用大会

以太坊技术及应用大会

以太坊技术及应用大会

以太坊技术

应用大会

应用大会

应用大会

应用大会

Thanks