Paradigma P2P

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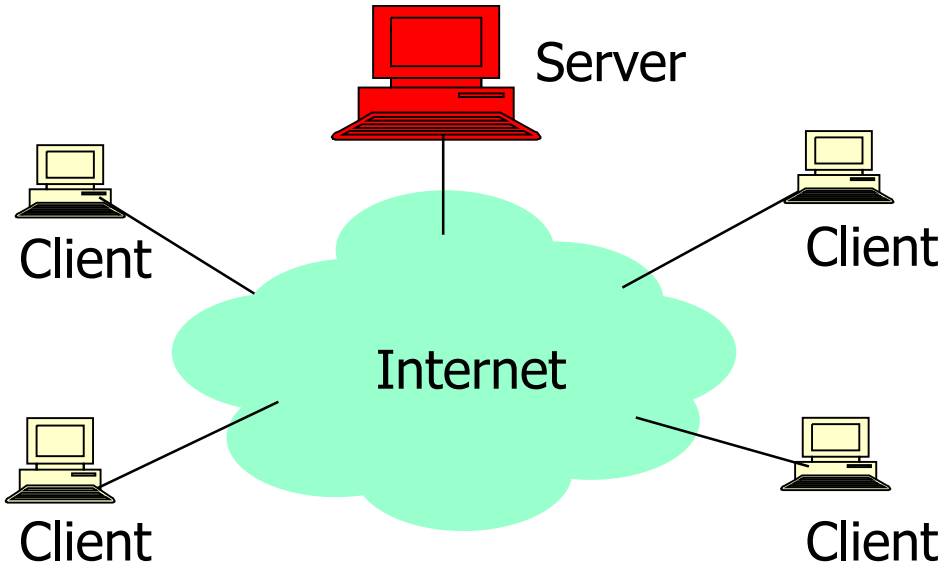
Cuprins

- **Paradigma peer-to-peer (P2P)**
  - Preliminarii
  - Definitii
  - Caracterizare
  - Tipuri de aplicatii
  - Infrastructuri
  - Instrumente

Preliminarii

...sa ne reamintim modelul client/server



Preliminarii

...sa ne reamintim modelul client/server

- Uzual, privim clientul ca fiind o componenta avand capacitati computationale reduse
- Serverul este mentinut si administrat in mod centralizat

Probleme ale arhitecturii client/server:

- Lipsa robustetii
- Lipsa sigurantei (eng. *reliability*)
- Lipsa scalabilitatii
- Vulnerabilitate la atac

Definitii

Peer *= one that is of equal standing with another* (conform Webster)

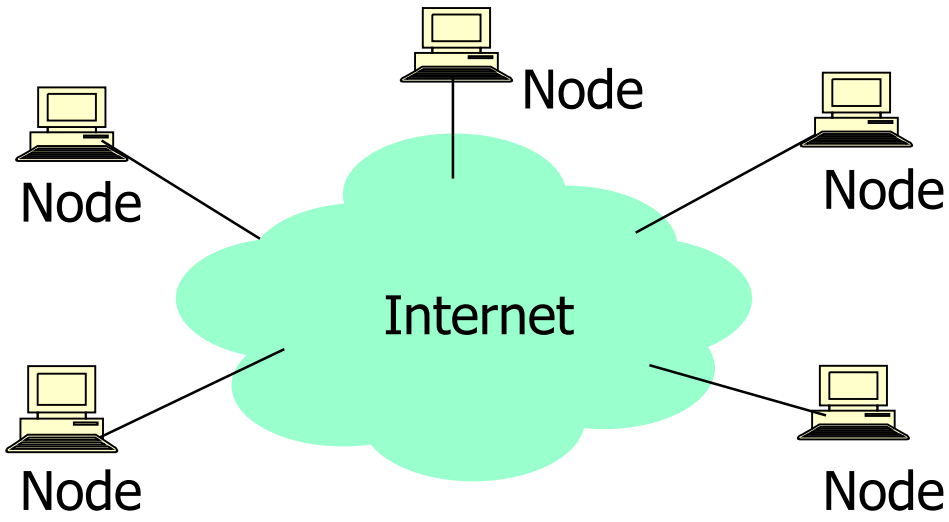
Peer-to-peer (P2P) = arhitectura de retea in care nodurile sunt relativ egale

- In sensul ca fiecare nod este, in principiu, capabil sa realizeze functii specifice retelei

Definitii

Sistemele P2P, in sens strict, sunt sisteme complet distribuite

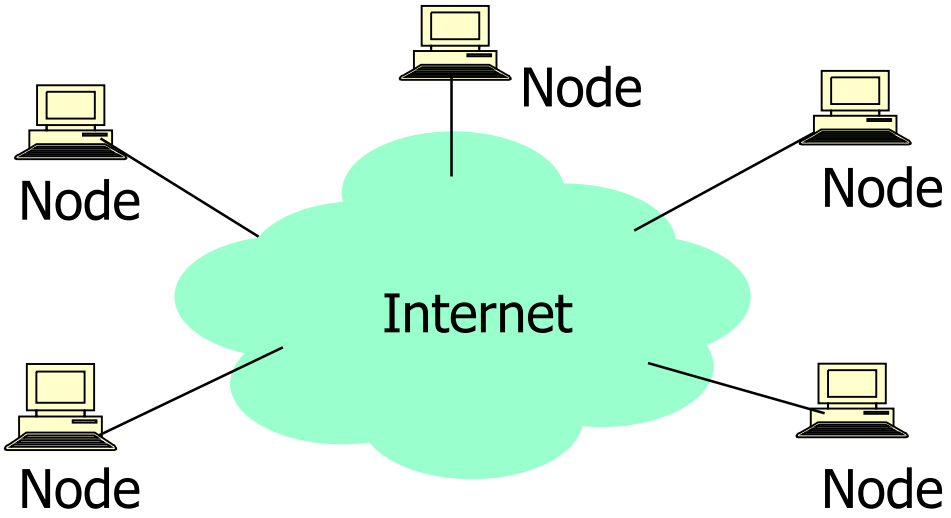
- Toate nodurile sunt total echivalente, in termeni de functionalitate si a activitatilor pe care le pot desfasura



Obs.: Sistemele P2P pure sunt rare (de ex. Gnutella); majoritatea sunt hibride, avand supernoduri sau servere cu diferite roluri (de ex. cautare de date, control, etc.)

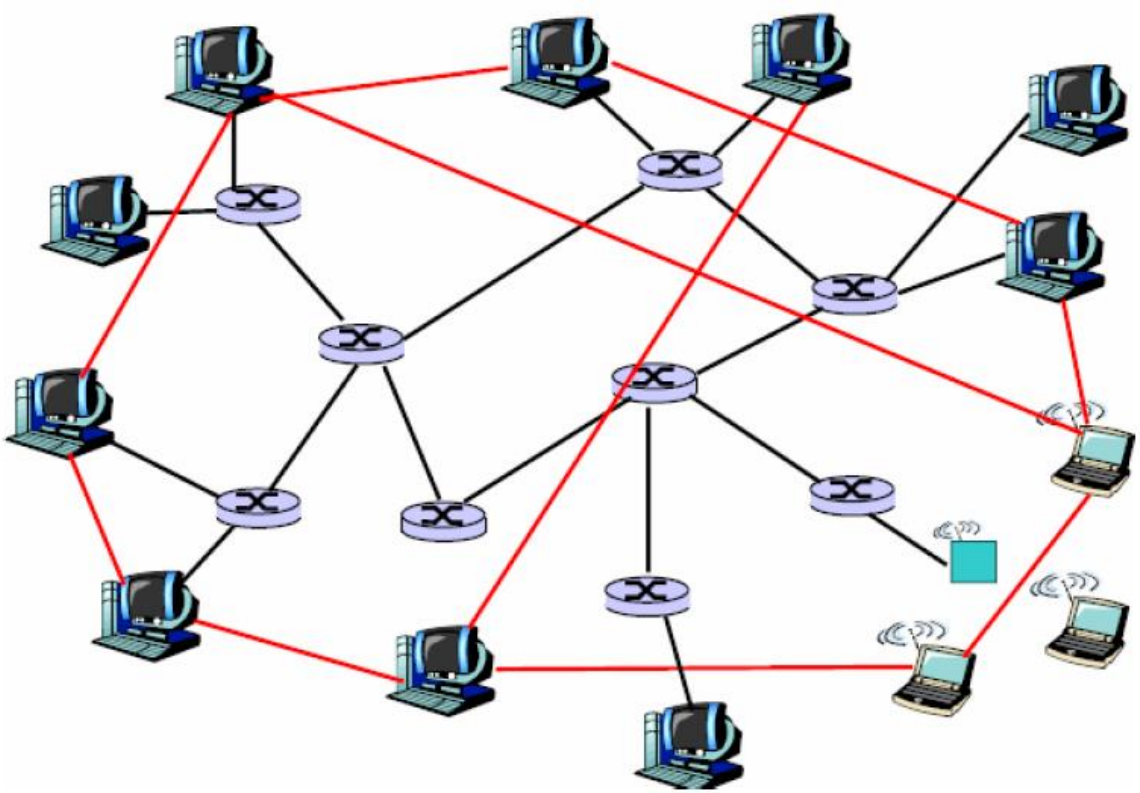
Definitii

- Nodurile
  - Pot consuma si oferi date
  - Orice nod poate initia o conexiune
- Nu exista sursa de date centralizata =>
  - Forma de democratie pe Internet
  - Protectie *copyright* amenintata



Definitii

- “**P2P** este clasa de aplicatii care se bazeaza pe resursele (de stocare, de procesare, continut, prezente umane) disponibile la marginile (edges) Internet-ului



Edges of the Internet
(*overlay networks*)

Definitii

"**P2P** is a class of applications that take advantage of resources – storage, cycles, content, human presence – available at the edges of the Internet. Because accessing these decentralized resources means operating in an environment of unstable and unpredictable IP addresses, P2P nodes must operate outside the DNS system and have significant, or total autonomy from central servers"

"A distributed network architecture may be called a **P2P** network if the participants share a part of their own resources. These shared resources are necessary to provide the service offered by the network. The participants of such a network are both resource providers and resource consumers"

Caracterizare

Caracteristici definitorii:

- Partajarea resurselor computationale prin interschimb direct si mai putin prin intermedieri oferite de o autoritate centralizata (server)
  - Serverele centralizate pot fi folosite insa pentru a realiza activitati specifice (initializarea retelei P2P, adaugarea de noi noduri in retea,…)
  - Ideal, nodurile participa activ si unilateral la realizarea de operatii ca localizarea & caching-ul nodurilor/continutului, dirijarea informatiilor, managementul resurselor transferate etc.

Caracterizare

Caracteristici definitorii:

- Abilitatea de a trata instabilitatea si variatiile conctivitatii retelei, adaptandu-se automat la erorile survenite sau la dinamicitatea nodurilor
  - Topologia retelei P2P e adaptiva si toleranta la defecte, nodurile auto-organizandu-se in vederea mentinerii conectivitatii si performantei retelei

Caracterizare

Reteaua P2P este una suprapusa (*overlay*) peste cea fizica

- Se situeaza la nivelul aplicatie => flexibilitate
- Muchiile virtuale sunt conexiuni TCP sau pointeri la adrese IP
- Mentinerea retelei P2P se face prin verificarea periodica a conectivitatii (*ping*) ori a existentei (mesaje "*still alive*?")
- Cand un nod pica, sistemul P2P ar putea stabili noi muchii
- Proximitatea (fizica) a nodurilor nu e importanta
- Reteaua P2P poate fi structurata sau nu

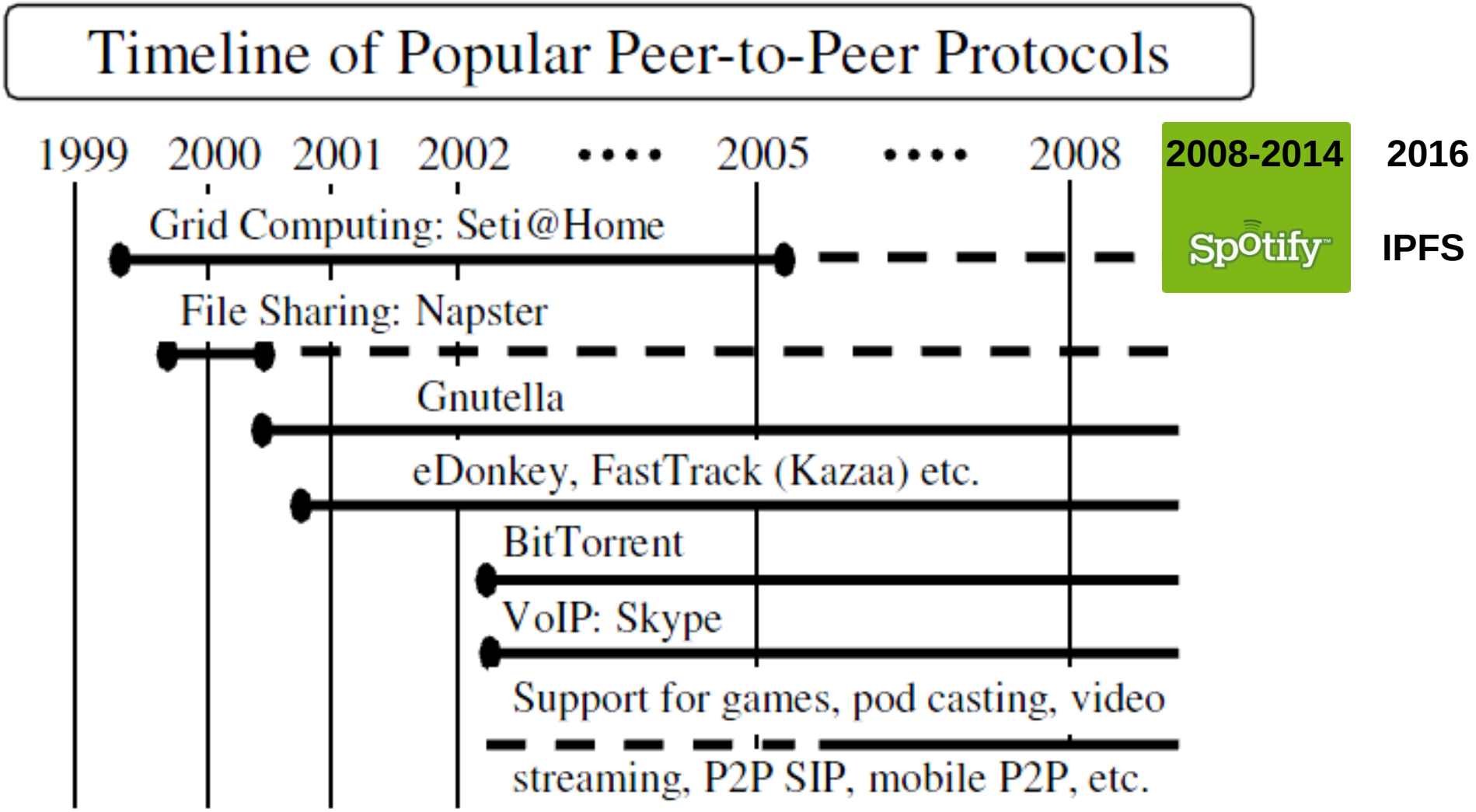
Scopuri si beneficii

- **Utilizarea eficienta a resurselor**
  - Latimea de banda neutilizata, resurse de stocare, putere de procesare disponibile la marginile (*edges*) retelei
- **Scalabilitate**
  - Fara informatii centralizate, fara *bottleneck*-uri (de comunicare si de calcul)
  - Agregarea resurselor se face in mod natural odata cu utilizarea sistemului
- **Siguranta** (eng. *reliability*)
  - Existenta de copii a datelor
  - Distribuire geografica
  - Nu mai exista "*single point of failure*"
- **Administrare usoara**
  - Nodurile se auto-organizeaza
  - Cresterea tolerantei la erori si a echilibrarii incarcarii
  - Cresterea autonomiei
- **Anonimitatea**
  - Greu de realizat intr-un sistem centralizat
- **Dinamism**
  - Mediu dinamic
  - Colaborare si comunicare ad-hoc

Dezavantaje/Probleme

- Arhitecturile P2P sunt probabilistice
  - Localizare impredictibila a resurselor
  - Resursele sunt volatile
- Inexistenta unui control centralizat
  - Probleme privind impunerea unei autoritati asupra aplicatiilor, continutului si utilizatorilor
  - Dificultati in detectarea si identificarea utilizatorilor (aspecte anti-sociale)
- Incurajarea folosirii sistemelor P2P in scop abuziv si ilegal (e.g. drepturile de autor asupra continutului digital)
- Lipsa increderii la nivel comercial, de afaceri
- Probleme de securitate (curs viitor)

Evolutie…


Timeline of Popular Peer-to-Peer Protocols
1999
2000
2001
2002
2005
2008
Grid Computing: Seti@Home
File Sharing: Napster
Gnutella
eDonkey, FastTrack (Kazaa) etc.
BitTorrent
VoIP: Skype
Support for games, pod casting, video
streaming, P2P SIP, mobile P2P, etc.
2008-2014
Spotify
2016
IPFS

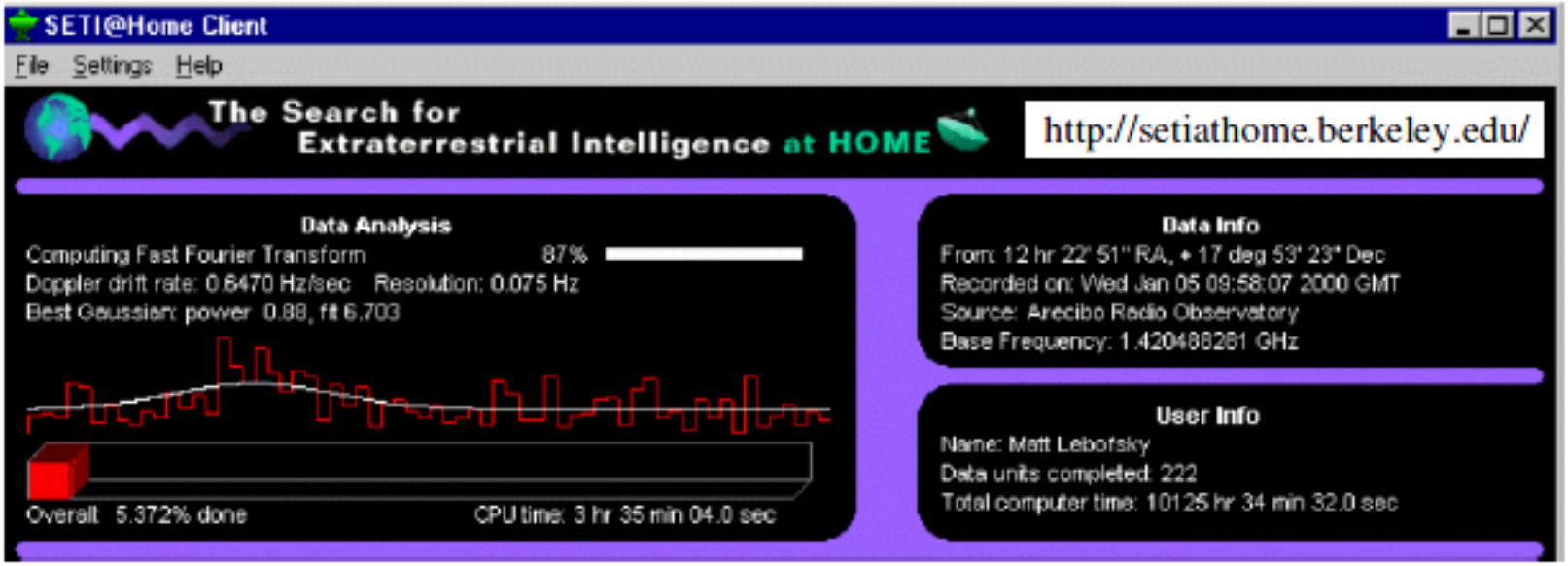
Tipuri de aplicatii

- **Comunicare & colaborare**
  - Sisteme ce ofera o infrastructura pentru facilitarea comunicarii & colaborarii directe, deseori in timp real, intre noduri
    - Sisteme conversationale (chat, mesagerie instantanee):

    IRC (Internet Relay Chat), ICQ (1996), YM!, MSN Messenger, Skype, Sisteme multicast P2P (e.g. Cirrus – Adobe Flash http://labs.adobe.com/technologies/cirrus/; WebRTC), …
- **Calcul distribuit**
  - Sisteme ce folosesc puterea computationala a nodurilor disponibile (cicli de procesor)
    - Rezolvarea unor probleme prin *divide-et-impera*: SETI@home (*Search for Extra-Terrestrial Intelligence*-Berkeley), genome@home
    - Reteaua P2P reprezinta un gen de Grid computational (…curs master)

Tipuri de aplicatii| Calcul distribuit - Exemplu

SETI@Home: A Public-Resource Computing Experiment



- Radio telescope signal analysis has insatiable appetite for computing power
- Usage of computers in homes and offices around the world has provided unprecedented computing power
- Grid computing application via peer-to-peer approach under central control

Tipuri de aplicatii

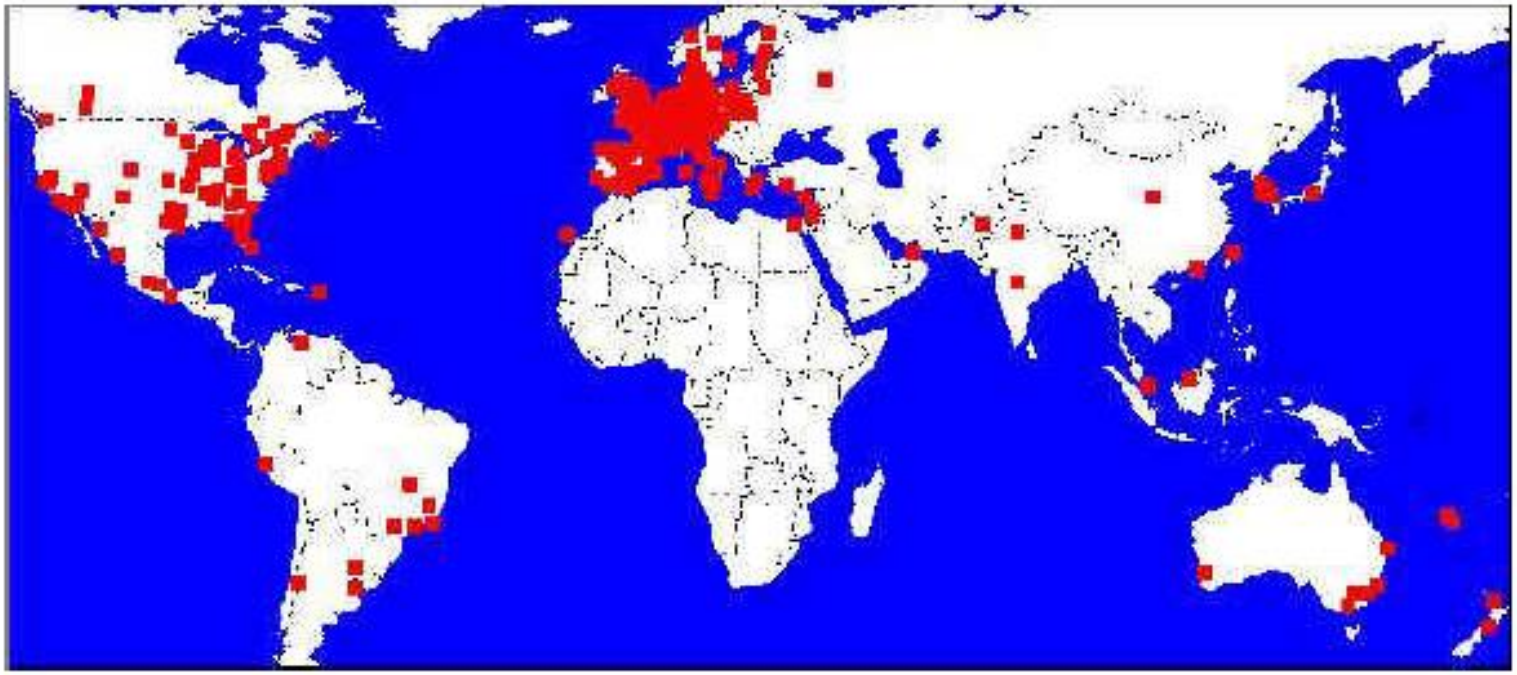
- **Sisteme de stocare (baze de date)**
  - Proiectare de sisteme de baze de date distribuite bazate pe infrastructuri P2P
    - PIER – motor scalabil de interogare distribuita (http://pier.cs.berkeley.edu/)
    - Edutella – proiect open-source pentru interogari si stocare de *meta-date* (P2P pentru Semantic Web)
- **Distribuire de continut digital**
  - Sisteme & infrastructuri pentru partajarea resurselor digitale (multimedia si alte date) intre utilizatori
    - Aplicatii pentru partajarea fisierelor (e.g. Napster, Gnutella, KaZaA, Freenet, BitTorrent, eDonkey etc.)
    - Medii de stocare distribuita pentru publicarea, organizarea, indexarea, cautarea si regasirea datelor in maniera securizata & eficienta

      (PAST, Chord, Groove, Mnemosyne, Avalanche,…)

Tipuri de aplicatii

- **Distribuire de continut digital | Exemplu**

P2P File-Sharing: Fast distribution of large files



Example: Harry Potter III early propagation after 2 hours on May 28th 2004 (Source: *www.itic.ca/DIC/News/archive.html*)

Tipuri de aplicatii

- **Distribuire de continut prin P2P**
  - Sisteme P2P de “interschimb de fisiere”
    - Nodurile transfera un fisier la un moment dat
    - Se ofera facilitati pentru realizarea unei retele P2P si pentru cautarea&transferul de fisiere intre noduri
    - Nu se ofera suport pentru securitate, disponibilitate si persistenta
    - Exemple: Napster, KaZaA, Gnutella

Tipuri de aplicatii

- **Distribuire de continut prin P2P**
  - Sisteme P2P pentru publicarea & stocarea continutului
    - Utilizatorii pot publica, stoca si distribui continut digital, pe baza unor drepturi de acces (privilegii)
    - Se focalizeaza asupra securitatii si persistentei
    - Unele ofera si facilitati privind colaborarea intre utilizatori
    - Exemple: Scan, Groove, Freenet, MojoNation, Tangler

Tipuri de aplicatii

- **Distribuire de continut prin P2P**
  - Infrastructuri pentru:
    - Dirijare & Localizare:

      Chord, Can, Pastry, Tapestry, Kademila
    - Anonimitate:

      Onion Routing, ZeroKnowledge, Freedom, Tarzan
    - Managementul reputatiei:

      Eigentrust, PeerTrust

Infrastruc.(localizare & dirijare)

- Mecanismele de localizare si dirijare ce pot fi adoptate depind de:
  - Topologia
  - Structura
  - Gradul de centralizare

  ale retelei suprapuse, acoperitoare (*overlay network*)

Infrastruc.(localizare & dirijare)

- **Aspecte privind centralizarea**
  - **Arhitecturi pur descentralizate**: toate nodurile realizeaza exact aceleasi activitati, jucand simultan roluri de servere si clienti, fara a beneficia de o coordonare centrala
    - Nodurile se numesc servents (SERVers + clieENTS)

Infrastruc.(localizare & dirijare)

- **Aspecte privind centralizarea**
  - **Arhitecturi partial centralizate**: unele noduri au un rol mai important (de ex. stocand indecsi locali pentru fisierele partajate)
    - Nodurile devin supernoduri conform politicilor fiecarui sistem P2P
    - Rolul de supernod este stabilit dinamic
  - **Arhitecturi descentralizate hibride**: exista un server central facilitand interactiunea intre noduri, mentinand cataloage de meta-date ale fisierelor
    - Serverele pot identifica si verifica nodurile de stocare
    - Sistemele se mai numesc *broker mediated*

Infrastruc.(localizare & dirijare)

- **Aspecte privind structura retelei:**
  - **Nestructurata:** plasarea continutului este complet independenta de topologia retelei suprapuse
    - Continutul trebuie localizat
    - Strategii de cautare prin "*forta bruta*": inundarea retelei – cereri propagate via BFS/DFS
    - Strategii mai sofisticate: drumuri aleatorii, probabilistice, etc.
  - **Slab structurata** (*loosely structured*): desi localizarea continutului nu e complet specificata, aceasta este afectata de dirijare
    - Categorie aflata intre retele structurate si cele nestructurate

Infrastruc.(localizare & dirijare)

- **Aspecte privind structura retelei:**
  - **Structurata:** topologia este controlata, iar fisierele (sau pointerii la ele) sunt plasate in locatii precise
    - Se realizeaza o asociere (*mapping)* intre continut (identificatorul de fisier) si locatie (adresa nodului)
      - In genul unei tabele de rutare distribuita
    - Cautarile exacte (*exact-match queries*) pot fi realizate in mod scalabil
    - Structura folosita la dirijarea eficienta a mesajelor este dificil de mentinut in cazul unor noduri tranziente, cu rata mare de atasare si deconectare de la retea

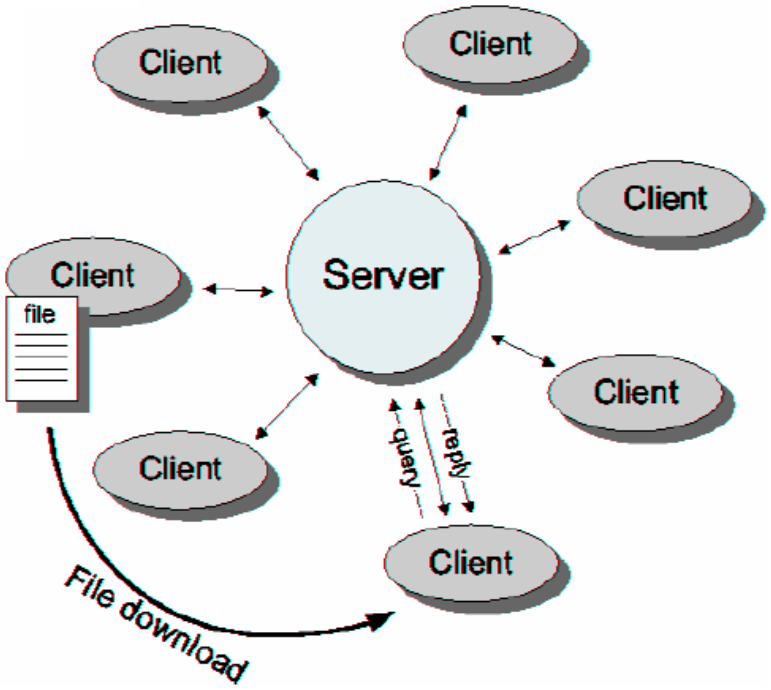
Infrastruc.(localizare & dirijare)

	Centralizare		
	Hibridă	Parțială	Absentă
Nestructurată	Napster Publius	KaZaA Morpheus Edutella	Gnutella FreeHaven
Infrastructură structurată			Chord, CAN, Tapestry, Pastry
Sisteme structurate			OceanStore Scan, PAST, Kademlia, Tarzan

Arhitecturi nestructurate

Descentralizate hibride

- Fiecare calculator client stocheaza continut (fisiere) partajat(e)
- Serverul central mentine o tabela cu conexiunile utilizatorilor inregistrati (IP, latime de banda,…) + o tabela cu lista fisierelor fiecarui utilizator &meta-date
- Exemple: **Napster, Publius**



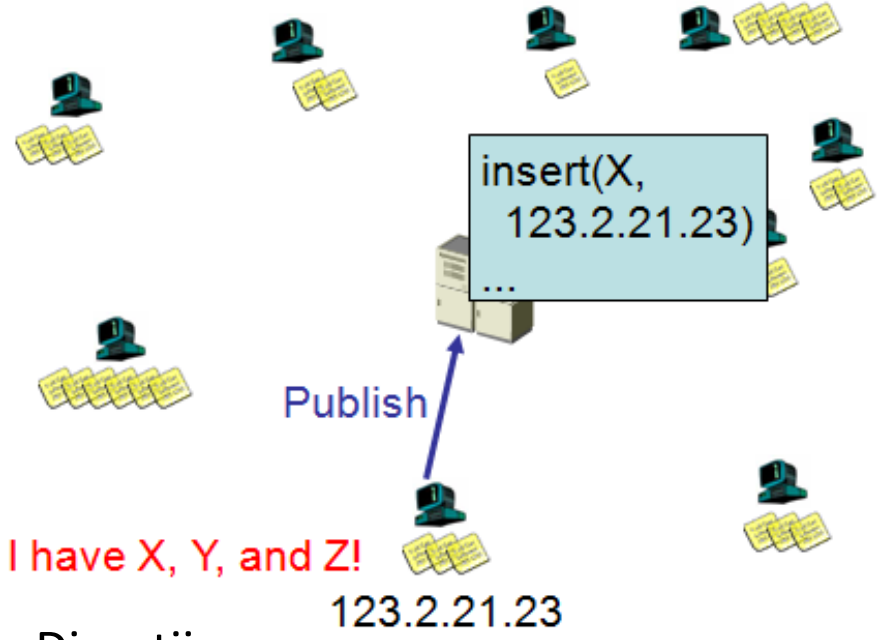
Arhitecturi nestructurate

Napster

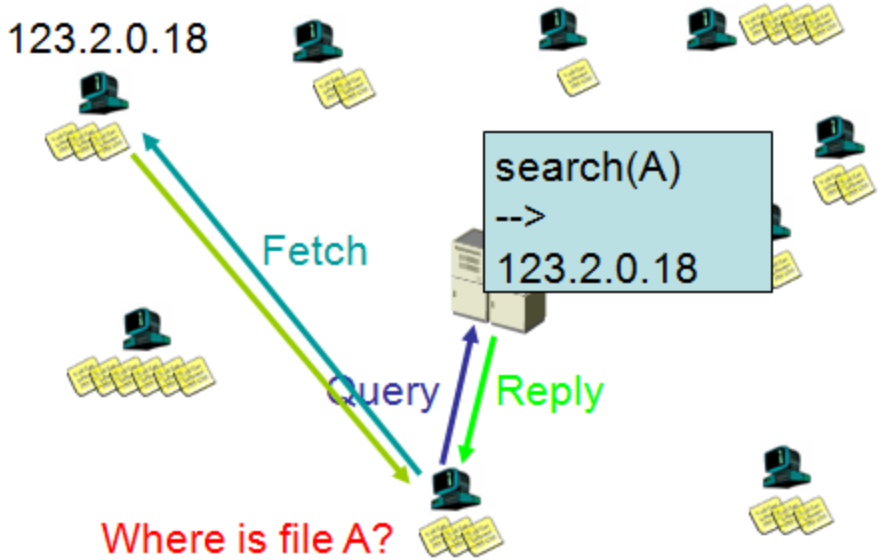
- 1999: Sean Fanning lanseaza Napster
- A atins cota de 1.5 milioane de utilizatori simultan
- Baza de date centralizata - operatii:
  - **Join**: clientul contacteaza serverul central (via TCP)
  - **Publish**: raportarea unei liste de fisiere serverului central
  - **Search**: interogarea serverului => se intoarce cineva care stocheaza fisierul cerut
  - **Fetch**: ia fisierul direct de la *peer* (cel cu cea mai buna rata de transfer)
- Iulie 2001: Napster a fost inchis

Arhitecturi nestructurate

Napster: Publish



Napster: Search



Discutii:

- Serverul face toate procesarile
- Avem "*single point of failure*"
- Probleme de scalabilitate, unele sisteme nu permit adaugarea altor servere (lista serverelor disponibile este statica)

Arhitecturi nestructurate

Descentralizate pure

- Se construieste o retea acoperitoare (*overlay*) cu propriile mecanisme de rutare prin IP
- Nu exista o coordonare centrala
- Utilizatorii se conecteaza via o aplicatie care are rol dublu – *servent*
- Comunicarea intre *serventi* se bazeaza pe un protocol la nivel de aplicatie, cu 4 tipuri de mesaje:
  - Ping – cere ca un nod sa se anunte
  - Pong – replica la mesajul *ping* (IP, port, numarul & marimea fisierelor)
  - Query – cerere de cautare (sir de cautare + viteza minima de transfer)
  - Query hints – raspuns (IP, port, viteza, dim. fis., index fis.)

Arhitecturi nestructurate

Descentralizate pure

- Cautarea se realizeaza prin inundare (flooding)
  - Daca nu ai fisierul dorit, intreaba pe *n* vecini
  - Daca nici ei nu au fisierul, vor intreba pe vecinii lor in maxim *m* hop-uri
  - Pe calea de intoarcere se vor intoarce raspunsurile (nu continutul fisierelor)
- Fiecare mesaj are un TTL atasat
- Exemplu: **Gnutella**

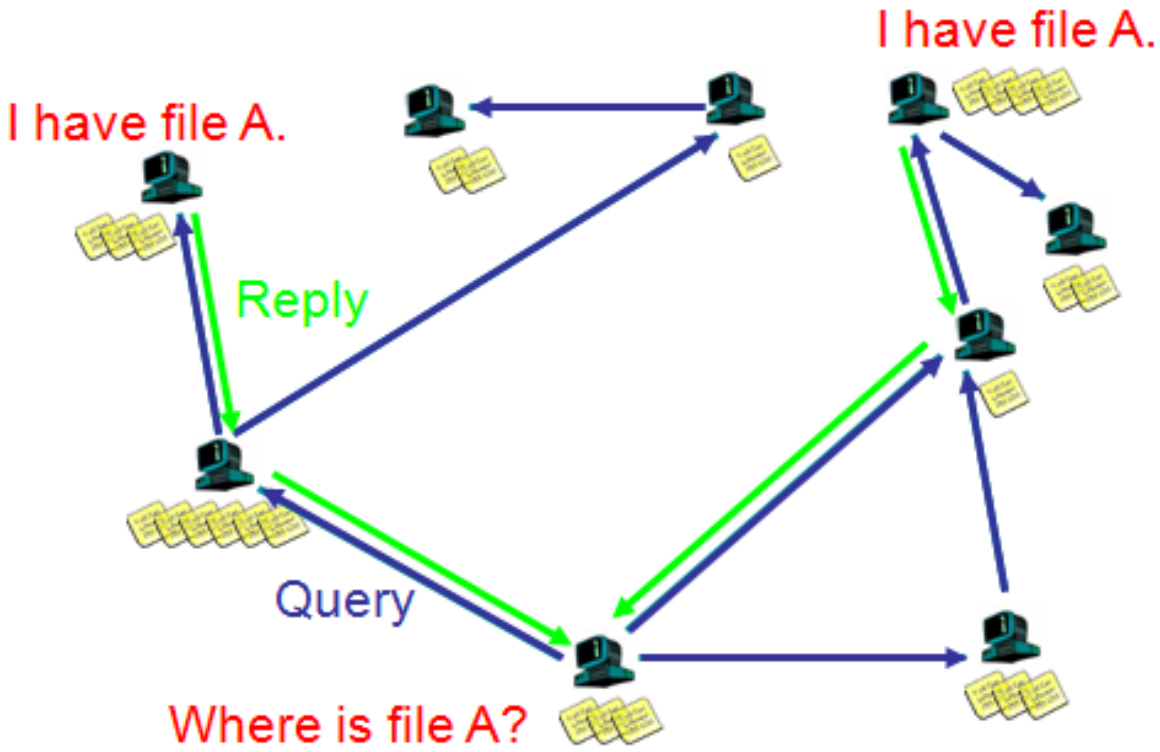
Arhitecturi nestructurate

Gnutella

- 2000: L. Frankel si T. Pepper(Nullsoft) lanseaza Gnutella
- Apar clienti: Bearshare, Morpheus, LimeWire
- *Query Flooding*:
  - **Join**: la intrare in sistem, clientul contacteaza cateva noduri care devin “vecinii” sai
  - **Publish**: nu este necesar
  - **Search**: se intreaba vecinii, care isi intreaba vecinii lor, …
    - Exista un TTL ce limiteaza propagarea
  - **Fetch**: preia fisierul direct de la *peer*

Arhitecturi nestructurate

Gnutella



Aspecte:

- Timpul de cautare este… O(?)
- Nodurile pleaca adesea => reteaua instabila

Arhitecturi nestructurate

Partial centralizate

- Folosesc conceptul de supernod: are activitati de servire a unei sub-retele P2P (indexare, *caching*)
- Nodurile sunt alese automat ca fiind supernoduri daca au suficienta latime de banda si putere computationala
- Toate cererile sunt trimise initial la supernoduri
- Avantaje: timpul descoperirii resurselor e mai redus + eterogenitatea este exploatata
- Exemplu: **KaZaA**

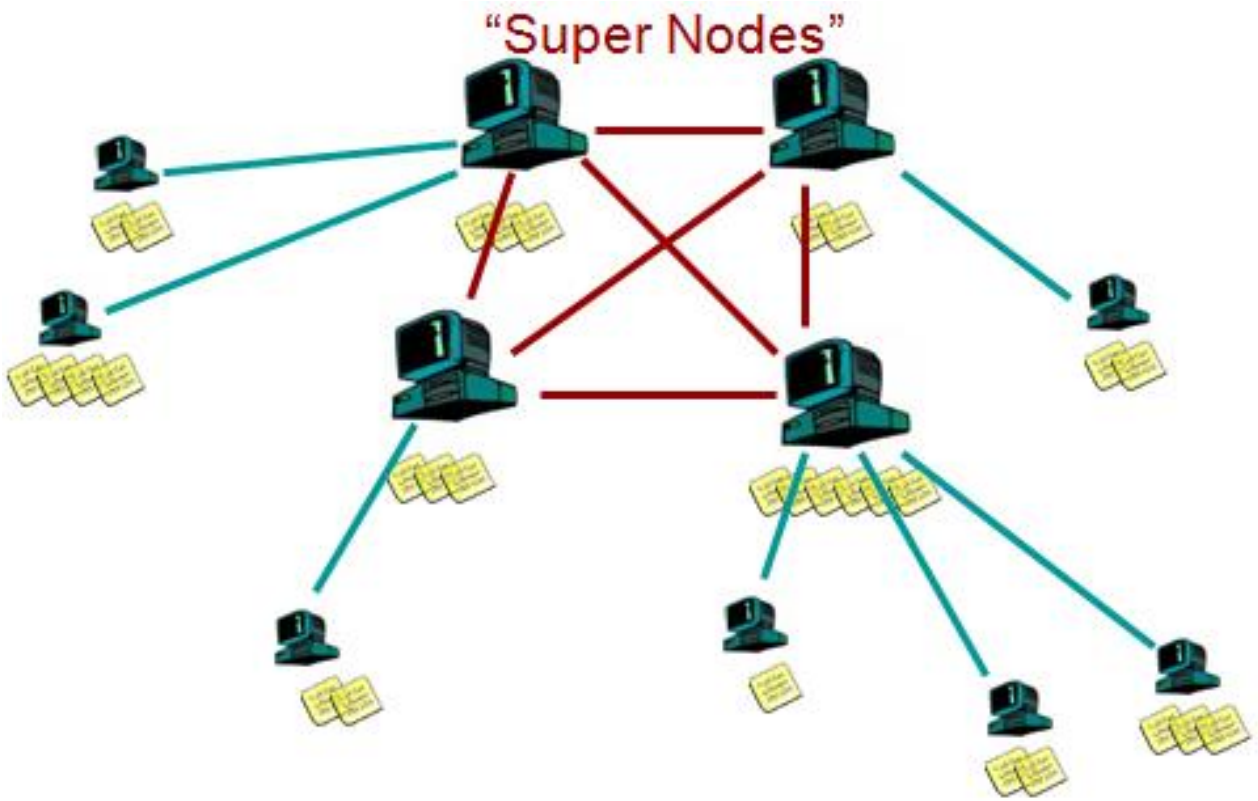
Arhitecturi nestructurate

KaZaA

- 2001: Se lanseaza KaZaA
- Apar clienti: Morpheus, giFT
- Se utilizeaza un mecanism de tip "*smart*" *query flooding*:
  - **Join**: la intrare in sistem, clientul contacteaza un "*supernode*" (poate deveni si el *supernod* la un moment dat)
  - **Publish**: trimite lista de fisiere *supernodului*
  - **Search**: trimite interogarea *supernodului*, si *supernodurile* se interogheaza intre ele
  - **Fetch**: ia fisierul direct de la *peer(s);* poate prelua fisierul simultan de la mai multe *peer*-uri

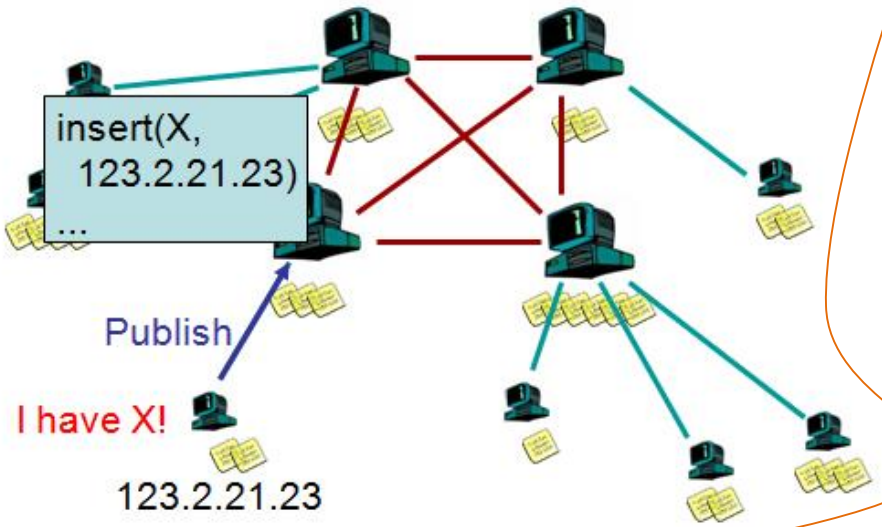
Arhitecturi nestructurate

KaZaA: Designul retelei

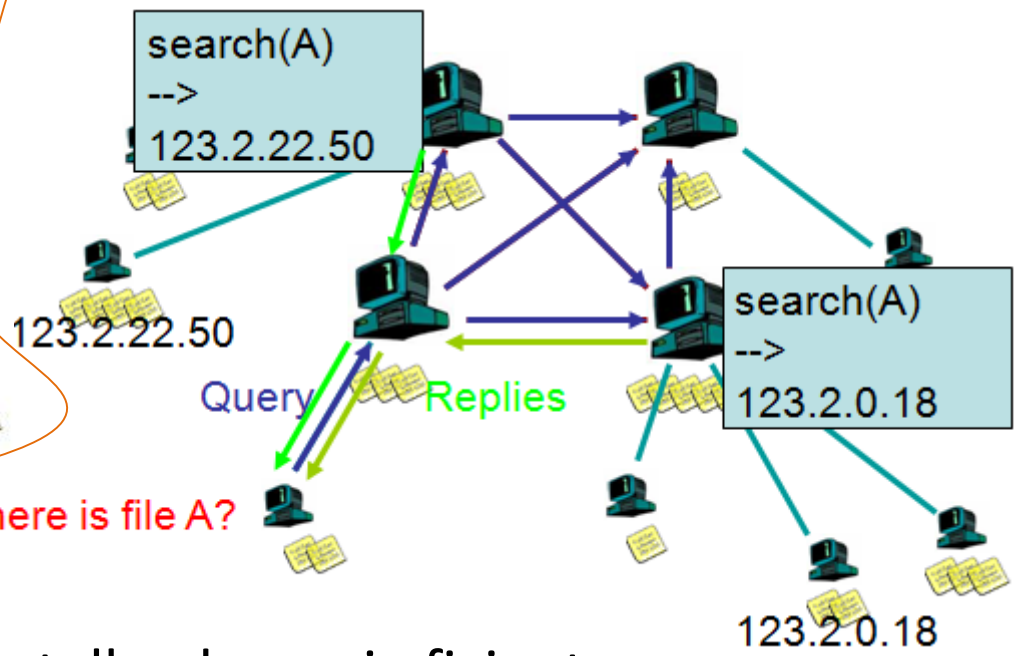


Arhitecturi nestructurate

KaZaA: Inserarea de Fisiere



KaZaA: Cautarea de Fisiere



Discutii:

- Comportament similar cu Gnutella, dar mai eficient
- Nu este nici o garantie asupra timpului de cautare sau a domeniului de cautare

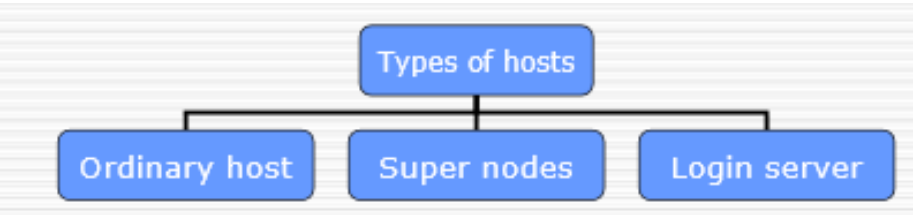
Arhitecturi nestructurate

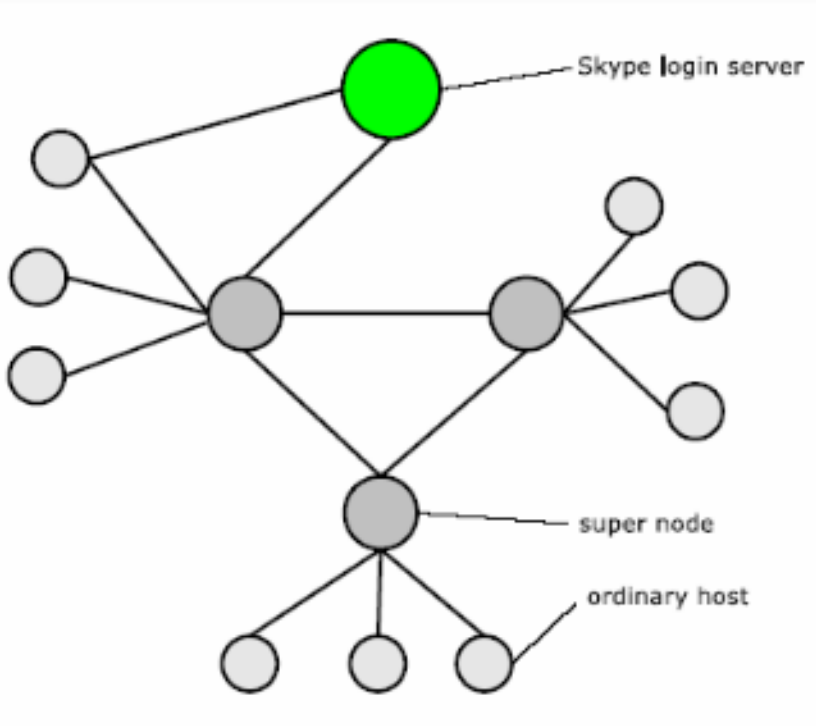
Partial centralizate

- Software-ul KaZaA este proprietar
- Datele de control P2P sunt criptate
- Mesajele folosesc HTTP
- Un nod e fie un supernod, fie asignat unui supernod
- Un supernod are 100-150 noduri-copil
- O retea poate avea ~30000 supernoduri
- Fiecare supernod are conexiuni TCP cu 30-50 supernoduri
- Pentru fiecare fisier se mentin meta-date (nume, dimensiune, *content hash*, descriptor de fisier)
- *Content hash*-ul este folosit pentru cautarea altei copii a unui fisier partial transferat
- Varianta fara *spyware* si *pop-up*-uri: **KaZaA-lite**

Arhitecturi nestructurate

Skype



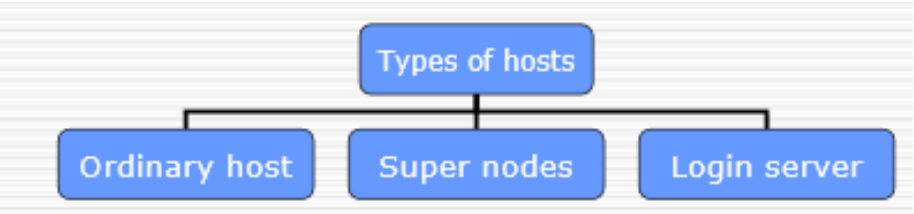


- Prima retea de telefonie p2p bazata pe IP
- din Iunie 2014, Microsoft a anuntat incompatibilitatea cu protocolul anterior Skype
- foloseste Microsoft Notification Protocol 24 (prima utilizare -> MSN Messenger in 1999)
- arhitectura era similara cu KaZaA

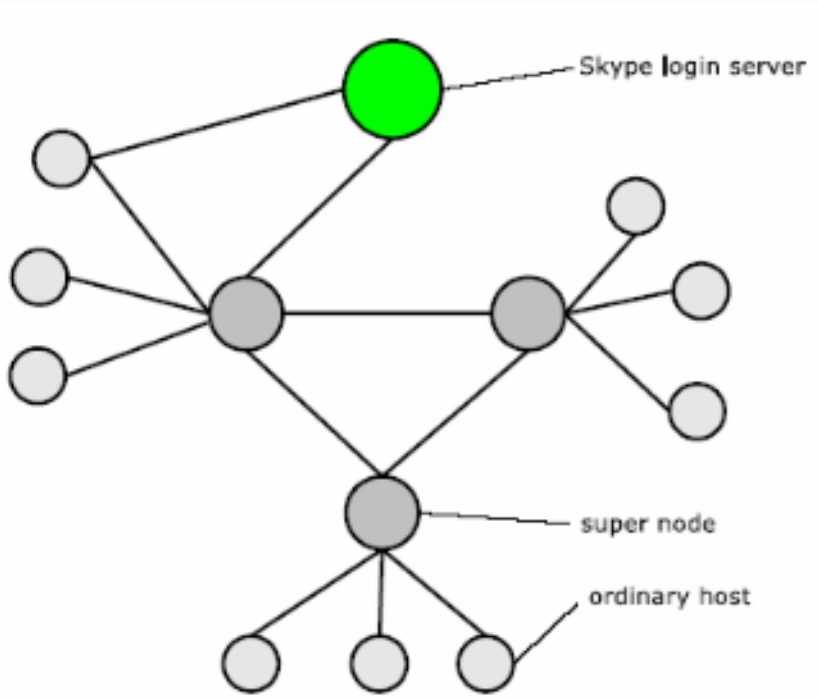
http://www1.cs.columbia.edu/~salman/publications/skype1_4.pdf

Arhitecturi nestructurate

Skype



- Fiecare client mentinea un *host cache* cu adresele IP si numerele de port ale supernodurilor accesibile
- Orice client cu latime de banda (si fara restrictii de firewall sau NAT) putea deveni supernode
- din 2012, Microsoft a inceput gazduirea supernodurilor in servere din centrele sale de date



http://en.wikipedia.org/wiki/PRISM_%28surveillance_program%29

Arhitecturi nestructurate

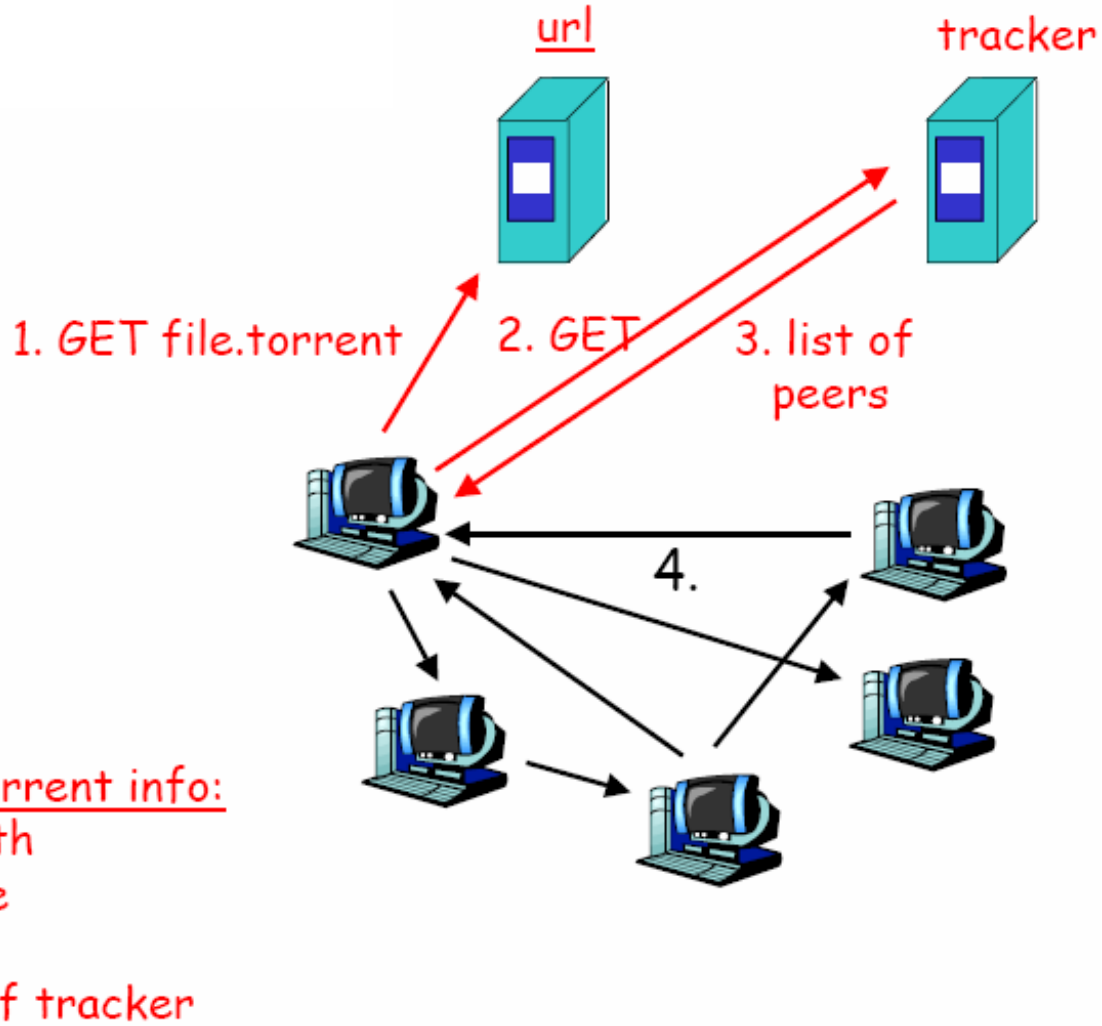
Partial centralizate

- Daca un fisier este gasit pe mai multe noduri, transferul poate fi realizat in paralel
  - Copiile identice se identifica via *content hash*
- Diferite portiuni din fisier sunt transferate de pe noduri diferite
- Pentru transferurile intrerupte, se face o recuperare automata *(automatic recovery)*
- Exemplu: **BitTorrent**
  - In 2002, B. Cohen a lansat BitTorrent
  - Si-a propus concentrarea pe problema legata de obtinerea eficienta a resurselor (*efficient fetching)* si nu pe cautare (*searching*)
  - Sustinatori inca de la aparitie
    - Blizzard Entertainment folosea BitTorrent pentru distributia versiunilor beta a noilor jocuri

Arhitecturi nestructurate

Partial centralizate

BitTorrent - arhitectura



Arhitecturi nestructurate

Partial centralizate

BitTorrent

- Se bazeaza pe mecanismul de *swarming*:
  - Join: contacteaza un server centralizat (*tracker*) si obtine o lista de *peer*-uri
  - Publish: ruleaza un server *tracker*
  - Search: de ex. foloseste Google pentru a gasi un *tracker* pentru fisierul dorit
  - Fetch: Preia bucati de fisiere de la *peer*-uri; incarca bucatile de fisier pe care le ai
- Obs. Diferenta fata de Napster
  - *Downlod-ul* de bucati (*chunk*) de fisiere
  - Utilizarea strategiei “tit-for-tat”: daca A face *download* de la alte noduri, atunci A trebuie sa permita si *download-ul* de la el (*free-rider problem*)

Arhitecturi nestructurate

Probleme

- **Noduri ale caror adrese IP sunt disponibile via NAT (cu restrictii)**
  - Nu pot fi servere TCP pentru reteaua P2P
  - Solutie partiala: *reverse call*
    - A vrea sa transfere de la B, iar B foloseste NAT
    - A si B stabilesc conexiuni TCP cu serverul C (IP rutabil)
    - A poate cere lui B, via C, sa realizeze o conexiune TCP de la B la A
    - A poate trimite o cerere lui B, iar B ii ofera raspunsul
  - Daca A si B utilizeaza NAT?
- **Flash crowd: o crestere neasteptata de cereri pentru o anumita resursa**
  - Pentru continutul dorit nu exista suficiente copii incarcate
  - Cat timp ia unui utilizator sa localizeze fisierul?
  - Cate mesaje va primi un nod datorita cautarilor realizate de alte noduri?
  - Se poate folosi un protocol de cautare generic, bazat pe TTL

Arhitecturi structurate

- Reprezinta solutia **academica** pentru P2P
- Scop:
  - Cautare cu succes
  - Timp de cautare in limite cunoscute
  - Scalabilitate demonstrata
- Abordare: **DHT (Distributed Hash Table)**
  - Se stocheaza perechi (key, value)
    - Key – nume de fisiere
    - Value – continut de fisier sau pointer la o locatie
  - Fiecare *peer* stocheaza o multime de (key, value)
  - *Operatii: gaseste nodul responsabil cu un Key*
    - Mapare *key – node*
    - Rutarea eficienta a cererilor de *insert/lookup/delete* asociate cu acest nod
  - Se permite o mare fluctuatie a nodurilor

Arhitecturi structurate

- Aspect de interes: **localizarea continutului**
- Idee: Responsabilitatea este distribuita mai multor noduri ale retelei de acoperire, intr-un mod adaptiv
- Fiecarei resurse i se asociaza o cheie unica via o functie *hash*: h(“Curs Retele”)->7929; Intervalul de valori ale functiei *hash* se distribuie in reteaua P2P
- Fiecare nod trebuie sa “cunoasca” locatia macar a unei singure copii a fiecarei resurse pentru care functia sa *hash* ia valori in intervalul lui
- Nodurile pot mentine in cache-ul propriu o copie a fiecarei resurse pe care trebuie sa o “cunoasca”

Arhitecturi structurate

- Aspect de interes: **dirijarea**
- Pentru fiecare resursa, un nod ce “cunoaste” resursa trebuie sa fie accesat pe calea cea mai “scurta”
- Abordarile de sisteme P2P structurate difera prin strategia de dirijare
- Nodurile din sistem formeaza o structura de date distribuita care poate fi: inel, arbore, hypercub, skip list, etc.
- Se ofera un API pentru tabelele distribuite de hash-uri (DHT – Distributed Hash Table)
  - Dand o cheie *k*, API-ul va returna adresa IP a nodului responsabil pentru valoarea cheii k

Arhitecturi structurate

- **Implementari**
  - Chord [MIT]
  - Pastry [Microsoft Research UK, Rice University]
  - Tapestry [UC Berkeley]
  - Content Addressable Network (CAN) [UC Berkeley]
  - SkipNet [Microsoft Research US, Univ. of Washington]
  - Kademlia [New York University]
  - Viceroy [Israel, UC Berkeley]
  - P-Grid [EPFL Switzerland]

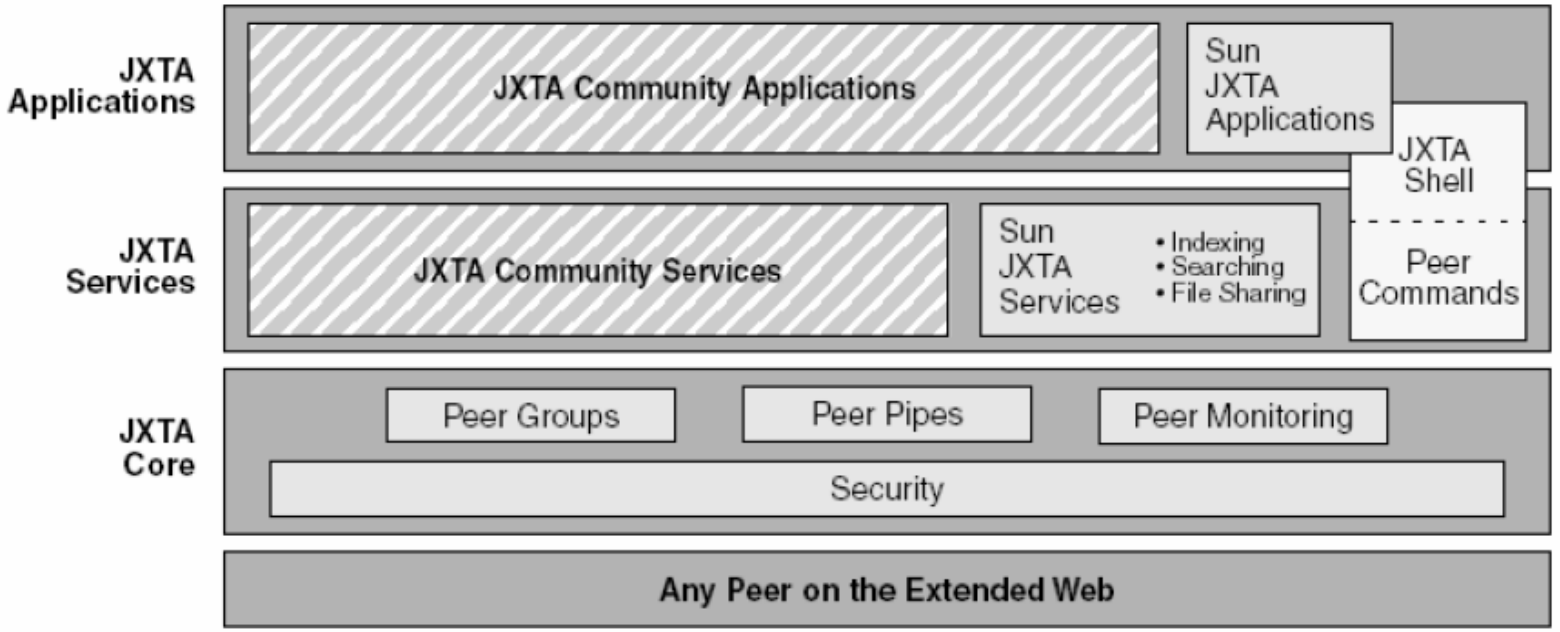
Arhitecturi structurate

- **Slab structurate**
  - Nodurile pot estima ce noduri stocheaza resursele cautate
    - Se evita broadcasturile oarbe
    - Se foloseste o propagare in lant (*chain mode propagation*): fiecare nod ia decizii locale privitoare la care va fi nodul urmator interogat
  - Cautarea unui fisier presupune utilizarea unei chei si a unui mecanism de *timeout*
  - Exemplu: **Freenet**

[https://en.wikipedia.org/wiki/Freenet]

Instrumente

- JXTA
  - Mediu de dezvoltare a sistemelor & aplicatiilor P2P
  - Bazat pe java, disponibil in regim open source
  - S-a renuntat la dezvoltarea acestuia



Instrumente

- P2P – framework pentru Android

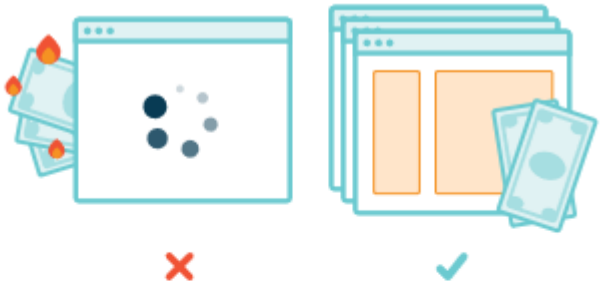
  https://code.google.com/p/p2p-communication-framework-for-android/
- p2psim – simulator pentru protocoalele p2p

  http://pdos.csail.mit.edu/p2psim/
- Instrumente si protocoale pentru P2P:

  http://en.wikibooks.org/wiki/The_World_of_Peer-to-Peer_%28P2P%29/Networks_and_Protocols/Other_Software_Implementations
- “ IPFS is the Distributed Web” - https://ipfs.io/

Instrumente

- " IPFS is the Distributed Web" - https://ipfs.io/
  - **–A peer-to-peer hypermedia protocol to make the web faster, safer, and more open**



HTTP is inefficient and expensive

HTTP downloads a file from a single computer at a time, instead of getting pieces from multiple computers simultaneously. With video delivery, a P2P approach could save 60% in bandwidth costs.

IPFS makes it possible to distribute high volumes of data with high efficiency. And zero duplication means savings in storage.

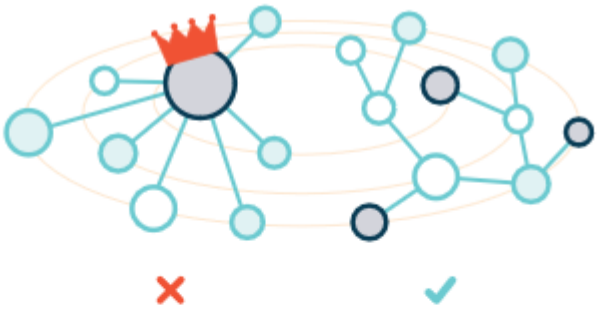
Humanity's history is deleted daily

The average lifespan of a web page is 100 days. Remember GeoCities? The web doesn't anymore. It's not good enough for the primary medium of our era to be so fragile.

IPFS provides historic versioning (like git) and makes it simple to set up resilient networks for mirroring of data.

Instrumente

- " IPFS is the Distributed Web" - https://ipfs.io/
  - **–A peer-to-peer hypermedia protocol to make the web faster, safer, and more open**



The web's centralization limits opportunity

The Internet has been one of the great equalizers in human history and a real accelerator of innovation. But the increasing consolidation of control is a threat to that.

IPFS remains true to the original vision of the open and flat web, but delivers the technology which makes that vision a reality.

Our apps are addicted to the backbone

Developing world. Offline. Natural disasters. Intermittent connections. All trivial compared to interplanetary networking. The networks we're using are so 20th Century. We can do better.

IPFS powers the creation of diversely resilient networks which enable persistent availability with or without Internet backbone connectivity.

Instrumente

- “ IPFS is the Distributed Web” - https://ipfs.io/

Let's take a look at what happens when you add files to IPFS:

Each file and all of the **blocks within it** are given a **unique fingerprint** called a **cryptographic hash**.

Each **network node** stores only content it is interested in, and some indexing information that helps figure out who is storing what.

IPFS **removes duplications** across the network and tracks **version history** for every file.

When **looking up files**, you're asking the network to find nodes storing the content behind a unique hash.

Every file can be found by **human-readable names** using a decentralized naming system called **IPNS**.

- https://github.com/ipfs/papers/raw/master/ipfs-cap2pfs/ipfs-p2p-file-system.pdf

Statistici

Global Consumer Internet Traffic 2005–2011

Consumer Internet Traffic 2005–2011							
	2005	2006	2007	2008	2009	2010	2011
By Sub-Segment (terabytes per month)							
Web, e-mail, file transfer	362,084	505,996	692,812	948,425	1,233,172	1,603,615	2,756,415
P2P	1,060,226	1,329,770	1,772,403	2,379,025	3,111,891	4,040,403	5,269,360
Gaming	66,844	91,943	133,367	188,680	250,574	318,212	386,832
Video Communications	11,629	15,575	24,932	36,638	47,173	66,101	92,453
VoIP	10,965	23,035	39,339	57,653	75,575	92,815	110,456
Internet Video to PC	53,074	174,427	484,027	838,154	1,232,461	1,726,114	2,331,908
Internet Video to TV	0	12,727	110,692	353,095	620,197	936,580	1,342,482
By Geography (TB per month)							
North America	534,236	618,765	917,365	1,287,026	1,698,700	2,242,841	2,861,772
Western Europe	334,600	505,329	814,015	1,281,041	1,856,310	2,515,070	3,458,721
Asia Pacific	565,782	819,072	1,201,277	1,742,834	2,315,755	3,049,294	4,663,774
Japan	60,080	98,747	147,733	223,120	319,788	436,057	556,631
Latin America	19,917	33,755	57,083	90,765	130,466	189,992	268,559
Central Eastern Europe	40,773	59,097	86,196	122,272	165,387	222,895	294,901
Middle East and Africa	9,435	18,708	33,904	54,613	84,637	127,689	185,549

Statistici

Table 8. Global Consumer Internet Traffic, 2011–2016

Consumer Internet Traffic, 2011–2016	2011	2012	2013	2014	2015	2016	CAGR 2011–2016
By Network (PB per Month)							
Fixed	20,121	29,095	35,943	45,372	57,991	74,247	30%
Mobile	402	879	1,717	3,116	5,213	8,313	83%
By Subsegment (PB per Month)							
File sharing	6,013	7,403	9,153	11,569	14,758	18,892	26%
Internet video	10,423	16,880	20,904	26,722	34,755	45,280	34%
Web, email, and data	3,863	5,422	7,274	9,783	13,119	17,583	35%
Online gaming	77	115	170	251	404	630	52%
Voice over IP (VoIP)	147	154	159	163	169	174	3%
By Geography (PB per Month)							
North America	5,394	8,041	9,022	10,294	12,499	15,646	24%
Western Europe	5,132	7,463	9,311	11,822	14,796	18,233	29%
Asia Pacific	8,220	11,795	15,266	20,204	26,515	34,553	33%
Latin America	780	1,196	1,833	2,835	4,352	6,506	53%
Central and Eastern Europe	817	1,211	1,686	2,398	3,392	4,699	42%
Middle East and Africa	180	328	542	935	1,649	2,783	73%
Total (PB per Month)							
Consumer Internet traffic	20,523	29,974	37,660	48,488	63,204	82,560	32%

Source: Cisco VNI, 2012

Statistici

Table 15. Global consumer internet traffic, 2017–2022

Consumer Internet Traffic,2017–2022	2017	2018	2019	2020	2021	2022	CAGR 2017–2022
By Network (EB per Month)							
Fixed	67	86	111	141	179	225	27%
Mobile	10	16	25	36	50	68	47%
By Subsegment (EB per Month)							
Internet video	56	77	105	140	184	240	34%
Web, email, and data	12	15	19	23	27	31	22%
Online gaming	1	3	4	7	11	15	59%
File sharing	8	7	7	7	7	7	-3%
By Geography (EB per Month)							
Asia Pacific	27	38	54	74	99	130	37%
North America	25	32	40	50	61	74	25%
Western Europe	12	15	19	24	30	37	25%
Central and Eastern Europe	5	7	9	11	15	19	29%
Middle East and Africa	3	4	6	9	13	18	46%
Latin America	5	6	8	9	11	14	21%
Total (EB per Month)							
Consumer Internet traffic	77	102	136	177	229	293	31%

Source: Cisco VNI, 2018

Statistici

Global Consumer Peer-to-Peer Traffic 2005–2011

Consumer Peer-to-Peer Traffic 2005–2011							
	2005	2006	2007	2008	2009	2010	2011
By Geography (TB per month)							
North America	381,746	378,538	462,356	560,817	673,083	852,483	1,080,979
Western Europe	223,519	304,988	411,057	540,032	757,818	991,817	1,330,885
Asia Pacific	391,235	550,664	762,276	1,074,759	1,401,028	1,811,094	2,327,648
Japan	28,621	42,883	58,463	87,446	117,967	154,868	206,803
Latin America	8,732	14,358	23,247	37,284	53,587	80,043	117,731
Central Eastern Europe	22,075	31,009	43,117	59,928	79,589	106,543	141,282
Middle East and Africa	4,297	7,329	11,886	18,759	28,819	43,553	64,033
Total (TB per month)							
Peer-to-Peer Traffic	1,060,226	1,329,770	1,772,403	2,379,025	3,111,891	4,040,403	5,269,360

Statistici

Table 10. Global Consumer File-Sharing Traffic, 2011–2016

Consumer File Sharing, 2011–2016	2011	2012	2013	2014	2015	2016	CAGR 2011–2016
By Network (PB per Month)							
Fixed	5,967	7,337	9,093	11,524	14,658	18,698	26%
Mobile	46	66	84	106	142	194	33%
By Subsegment (PB per Month)							
P2P file transfer	4,656	5,401	6,234	7,314	8,627	10,215	17%
Other file transfer	1,357	2,002	2,942	4,315	6,174	8,677	45%
By Geography (PB per Month)							
North America	785	935	1,119	1,349	1,636	2,006	21%
Western Europe	1,609	1,845	2,154	2,552	2,918	3,181	15%
Asia Pacific	2,997	3,789	4,794	6,201	8,100	10,660	29%
Latin America	210	266	347	485	786	1,276	43%
Central and Eastern Europe	363	504	690	937	1,266	1,699	36%
Middle East and Africa	49	64	73	105	96	71	8%
Total (PB per Month)							
Consumer file sharing	6,013	7,403	9,177	11,629	14,801	18,892	26%

Source: Cisco VNI, 2012

Table 10. Global Consumer File-Sharing Traffic, 2015–2020

Consumer File Sharing, 2015–2020							
	2015	2016	2017	2018	2019	2020	CAGR 2015–2020
By Network (PB per Month)							
Fixed	5,942	5,909	5,829	5,713	5,616	5,939	0%
Mobile	22	28	29	29	29	35	9%
By Subsegment (PB per Month)							
P2P file transfer	4,798	4,550	4,224	3,840	3,438	3,633	-5%
Other file transfer	1,166	1,388	1,634	1,902	2,207	2,340	15%
By Geography (PB per Month)							
Asia Pacific	2,335	2,269	2,186	2,098	2,004	2,098	-2%
North America	1,015	1,137	1,260	1,371	1,478	1,576	9%
Western Europe	1,124	1,105	1,096	1,075	1,053	1,131	0%
Central and Eastern Europe	829	763	691	646	621	666	-4%
Latin America	554	573	558	514	454	463	-4%
Middle East and Africa	107	91	68	39	34	39	-18%
Total (PB per Month)							

File Sharing
This category includes traffic from P2P applications such as BitTorrent and eDonkey, as well as web-based file sharing. Note that a large portion of P2P traffic is due to the exchange of video files, so a total view of the impact of video on the network should count P2P video traffic in addition to the traffic counted in the Internet video-to-PC and Internet video-to-TV categories. Table 10 shows the forecast for consumer P2P traffic from 2015 to 2020. Note that the P2P category is limited to traditional file exchange and does not include commercial video-streaming applications that are delivered through P2P, such as PPStream or PPLive.

- http://www.cisco.com/c/en/us/solutions/collateral/service-provider/visual-networking-index-vni/complete-white-paper-c11-481360.html

Rezumat

- **Paradigma peer-to-peer(P2P)**
  - Preliminarii
  - Definitii
  - Caracterizare
  - Tipuri de aplicatii
  - Infrastructuri
  - Instrumente

Bibliografie

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- https://www.kirsle.net/blog/entry/skype-switched-to-the-msn-messenger-protocol

Intrebari?