

COERCION IN FCOM

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1. PRELIMINARIES

The typing rules for coercion, homogeneous composition and heterogeneous composition are:

$$\frac{\Gamma, i : \mathbb{I} \vdash A \quad \Gamma \vdash r : \mathbb{I} \quad \Gamma \vdash s : \mathbb{I} \quad \Gamma \vdash u : A(r/i)}{\Gamma \vdash \text{coe}_i^{r \rightarrow s} A u : A(s/i)[r = s \mapsto u]}$$

$$\frac{\Gamma \vdash A \quad \Gamma \vdash r : \mathbb{I} \quad \Gamma \vdash s : \mathbb{I} \quad \Gamma \vdash \varphi : \mathbb{F} \quad \Gamma, \varphi, i : \mathbb{I} \vdash u : A \quad \Gamma \vdash u_0 : A[\varphi \mapsto u(r/i)]}{\Gamma \vdash \text{hcom}_i^{r \rightarrow s} A [\varphi \mapsto u] u_0 : A[\varphi \mapsto u(s/i), r = s \mapsto u_0]}$$

$$\frac{\Gamma, i : \mathbb{I} \vdash A \quad \Gamma \vdash r : \mathbb{I} \quad \Gamma \vdash s : \mathbb{I} \quad \Gamma \vdash \varphi : \mathbb{F} \quad \Gamma, \varphi, i : \mathbb{I} \vdash u : A \quad \Gamma \vdash u_0 : A(r/i)[\varphi \mapsto u(r/i)]}{\Gamma \vdash \text{com}_i^{r \rightarrow s} A [\varphi \mapsto u] u_0 : A(s/i)[\varphi \mapsto u(s/i), r = s \mapsto u_0]}$$

The rules for fcom types:

$$\frac{\Gamma \vdash r : \mathbb{I} \quad \Gamma \vdash s : \mathbb{I} \quad \Gamma, \varphi, i : \mathbb{I} \vdash B \quad \Gamma \vdash A[\varphi \mapsto B(r/i)]}{\Gamma \vdash \text{fcom}_i^{r \rightarrow s} [\varphi \mapsto B] A [r = s \mapsto A, \varphi \mapsto B(s/i)]}$$

$$\frac{\Gamma \vdash r : \mathbb{I} \quad \Gamma \vdash s : \mathbb{I} \quad \Gamma, \varphi, i : \mathbb{I} \vdash B \quad \Gamma \vdash A[\varphi \mapsto B(r/i)] \quad \Gamma, \varphi \vdash u : B(s/i) \quad \Gamma \vdash u_0 : A[\varphi \mapsto \text{coe}_i^{s \rightarrow r} B u]}{\Gamma \vdash \text{box}^{r \rightarrow s} [\varphi \mapsto u] u_0 : \text{fcom}_i^{r \rightarrow s} [\varphi \mapsto B] A [r = s \mapsto u_0, \varphi \mapsto u]}$$

$$\frac{\Gamma \vdash r : \mathbb{I} \quad \Gamma \vdash s : \mathbb{I} \quad \Gamma, \varphi, i : \mathbb{I} \vdash B \quad \Gamma \vdash A[\varphi \mapsto B(r/i)] \quad \Gamma \vdash u : \text{fcom}_i^{r \rightarrow s} [\varphi \mapsto B] A}{\Gamma \vdash \text{cap}_i^{r \leftarrow s} [\varphi \mapsto B] u : A[r = s \mapsto u, \varphi \mapsto \text{coe}_i^{s \rightarrow r} B u]}$$

$$\frac{\Gamma \vdash r : \mathbb{I} \quad \Gamma \vdash s : \mathbb{I} \quad \Gamma, \varphi, i : \mathbb{I} \vdash B \quad \Gamma \vdash A[\varphi \mapsto B(r/i)] \quad \Gamma, \varphi \vdash u : B(s/i) \quad \Gamma \vdash u_0 : A[\varphi \mapsto \text{coe}_i^{s \rightarrow r} B u]}{\Gamma \vdash \text{cap}_i^{r \leftarrow s} [\varphi \mapsto B] (\text{box}^{r \rightarrow s} [\varphi \mapsto u] u_0) = u_0 : A}$$

$$\frac{\Gamma \vdash r : \mathbb{I} \quad \Gamma \vdash s : \mathbb{I} \quad \Gamma, \varphi, i : \mathbb{I} \vdash B \quad \Gamma \vdash A[\varphi \mapsto B(r/i)] \quad \Gamma \vdash u : \text{fcom}_i^{r \rightarrow s} [\varphi \mapsto B] A}{\Gamma \vdash \text{box}^{r \rightarrow s} [\varphi \mapsto u] (\text{cap}_i^{r \leftarrow s} [\varphi \mapsto 1B] u) = u : \text{fcom}_i^{r \rightarrow s} [\varphi \mapsto B] A}$$

The fiber of a map $e : A \rightarrow B$ over $x : B$ is defined as:

$$\text{Fiber } e \, x := (y : A) \times \text{Path } B \, (e \, y) \, x$$

Lemma (“Dagstuhl lemma”). *Given*

$$i : \mathbb{I} \vdash A \quad \vdash r, s : \mathbb{I} \quad \vdash y : A(s/i) \quad \vdash \varphi : \mathbb{F}$$

we can extend

$$\varphi \vdash (x, \langle _ \rangle \text{coe}_i^{r \rightarrow s} A \, x) : \text{Fiber}(\text{coe}_i^{r \rightarrow s} A) \, y$$

to a total element

$$\vdash (a, p) : (\text{Fiber}(\text{coe}_i^{r \rightarrow s} A) \, y) \, [\varphi \mapsto (x, \langle _ \rangle \text{coe}_i^{r \rightarrow s} A \, x)]$$

Proof. Note that $\varphi \vdash x : A(r/i)$ and $\varphi \vdash \text{coe}_i^{r \rightarrow s} A \, x = y$ by assumption so that we can first let

$$\vdash a := \text{com}_i^{s \rightarrow r} A \, [\varphi \mapsto \text{coe}_i^{r \rightarrow i} A \, x] \, y : A(r/i)$$

We then let

$$i : \mathbb{I} \vdash \alpha_0 := \text{coe}_i^{r \rightarrow i} A \, a : A$$

$$i : \mathbb{I} \vdash \alpha_1 := \text{com}_i^{s \rightarrow i} A \, [\varphi \vdash \text{coe}_i^{r \rightarrow i} A \, x] \, y : A$$

so that we can take

$$\vdash p := \langle j \rangle \text{com}_i^{r \rightarrow s} A \, [(j = 0) \mapsto \alpha_0, (j = 1) \mapsto \alpha_1, \varphi \mapsto \text{coe}_i^{r \rightarrow i} A \, x] \, a$$

for a fresh dimension $j : \mathbb{I}$. Note that

$$\vdash p \, 0 = \alpha_0(s/i) = \text{coe}_i^{r \rightarrow s} A \, a$$

$$\vdash p \, 1 = \alpha_1(s/i) = y$$

as desired. □

2. COERCION IN FCOM

The goal of this section is to explain how to compute:

$$\text{coe}_i^{r \rightarrow r'} (\text{fcom}_j^{s \rightarrow s'} [\varphi \mapsto B] A) \, u_0 : (\text{fcom}_j^{s \rightarrow s'} [\varphi \mapsto B] A)(r'/i)$$

We omit the ambient context Γ and are given:

$$i : \mathbb{I} \vdash \varphi : \mathbb{F}$$

$$i : \mathbb{I} \vdash s, s' : \mathbb{I}$$

$$i : \mathbb{I}, \varphi, j : \mathbb{I} \vdash B$$

$$i : \mathbb{I} \vdash A[\varphi \mapsto B(s/j)]$$

so that

$$i : \mathbb{I} \vdash F := \text{fcom}_j^{s \rightarrow s'} [\varphi \mapsto B] A$$

is a well-formed fcom type such that:

$$i : \mathbb{I}, \varphi \vdash F = B(s'/j)$$

$$i : \mathbb{I}, (s = s') \vdash F = A$$

Given $\vdash r, r' : \mathbb{I}$ and $\vdash u_0 : F(r/i)$ the goal is to define

$$\vdash w := \text{coe}_i^{r \rightarrow r'} F \, u_0 : F(r'/i)$$

such that

$$\begin{aligned} \forall i. \varphi \vdash w &= \text{coe}_i^{r \rightarrow r'} B(s'/j) u_0 : B(s'/j)(r'/i) \\ (r = r') \vdash w &= u_0 : F(r/i) \\ \forall i. (s = s') \vdash w &= \text{coe}_i^{r \rightarrow r'} A u_0 : A(r'/i) \end{aligned}$$

Step 1. Compute the cap of u_0 :

$$\begin{aligned} \vdash a_0 &:= \text{cap}_j^{s(r/i) \leftarrow s'(r/i)} [\varphi(r/i) \mapsto B(r/i)] u_0 \\ &: A(r/i) [(s(r/i) = s'(r/i)) \mapsto u_0 \\ &\quad, \varphi(r/i) \mapsto \text{coe}_j^{s'(r/i) \rightarrow s(r/i)} B(r/i) u_0] \end{aligned}$$

Note that if i does not occur in s and s' then the first restriction reads like $(s = s') \mapsto u_0$.

Step 2. Note that $\varphi \vdash u_0 : B(s'/j)(r/i)$ and we can coerce in $B(s'/j)$ to get a line along i :

$$\forall i. \varphi, i : \mathbb{I} \vdash \tilde{b} := \text{coe}_i^{r \rightarrow i} B(s'/j) u_0 : B(s'/j)$$

so that

$$\begin{aligned} \forall i. \varphi \vdash \tilde{b}(r/i) &= u_0 : B(s'/j)(r/i) \\ \forall i. \varphi \vdash b_1 &:= \tilde{b}(r'/i) = \text{coe}_i^{r \rightarrow r'} B(s'/j) u_0 : B(s'/j)(r'/i) \end{aligned}$$

Note that $(r = r') \vdash b_1 = u_0$.

Step 3. We now compute a composition of a_0 with the tubes given by \tilde{b} and a suitable correction:

$$\vdash a_1 := \text{com}_i^{r \rightarrow r'} A [\forall i. \varphi \mapsto \text{coe}_j^{s' \rightarrow s} B \tilde{b}, \forall i. (s = s') \mapsto \text{coe}_i^{r \rightarrow i} A u_0] a_0 : A(r'/i)$$

which is well-formed as $\forall i. \varphi \leq \varphi(r/i)$ so that

$$\forall i. \varphi \vdash (\text{coe}_j^{s' \rightarrow s} B \tilde{b})(r/i) = \text{coe}_j^{s'(r/i) \rightarrow s(r/i)} B(r/i) u_0 = a_0$$

and

$$\forall i. (s = s') \vdash a_1 = \text{coe}_i^{r \rightarrow r} A u_0 = u_0 = a_0$$

by the remark at the end of Step 1.

Furthermore this satisfies:

$$\begin{aligned} \forall i. \varphi \vdash a_1 &= \text{coe}_j^{s'(r'/i) \rightarrow s(r'/i)} B(r'/i) (\text{coe}_i^{r \rightarrow r'} B(s'/j) u_0) : B(r'/i)(s(r'/i)/j) \\ (r = r') \vdash a_1 &= a_0 : A(r/i) \\ \forall i. (s = s') \vdash a_1 &= \text{coe}_i^{r \rightarrow r'} A u_0 : A(r'/i) \end{aligned}$$

where the first equation makes sense as $\varphi(r'/i) \vdash A(r'/i) = B(s/j)(r'/i) = B(r'/i)(s(r'/i)/j)$.

Step 4. Note that we have an element in the fiber of

$$e := \text{coe}_j^{s'(r'/i) \rightarrow s(r'/i)} B(r'/i) : B(r'/i)(s'(r'/i)/j) \rightarrow B(r'/i)(s(r'/i)/j)$$

over a_1 given by:

$$\varphi(r'/i), \forall i. \varphi \vdash (b_1, \langle _ \rangle a_1) : \text{Fiber } e a_1$$

where

$$\text{Fiber } e a_1 := (x : B(r'/i)(s'(r'/i)/j)) \times \text{Path } B(r'/i)(s(r'/i)/j) (e x) a_1$$

which makes sense by the remark in the end of Step 3 above.

The fact that the pair $(b_1, \langle _ \rangle a_1)$ has type $\text{Fiber } e a_1$ on this restriction is justified by:

$$\varphi(r'/i), \forall i. \varphi \vdash e b_1 = \text{coe}_j^{s'(r'/i) \rightarrow s(r'/i)} B(r'/i) (\text{coe}_i^{r \rightarrow r'} B(s'/j) u_0) = a_1$$

Furthermore, we have:

$$\varphi(r'/i), (r = r') \vdash e b_1 = \text{coe}_j^{s'(r/i) \rightarrow s(r/i)} B(r/i) u_0 = a_0 = a_1$$

The second equality holds as $\varphi(r/i) \vdash \text{coe}_j^{s'(r/i) \rightarrow s(r/i)} B(r/i) u_0 = a_0$, and the other equalities hold by the remarks about what happens on $(r = r')$ in Step 2 and 3. We hence get

$$\varphi(r'/i), \forall i. \varphi \vee (r = r') \vdash (b_1, \langle _ \rangle a_1) : \text{Fiber } e a_1$$

Step 5. As $\varphi(r'/i), \forall i. \varphi \vee (r = r') \vdash a_1 = e b_1 = \text{coe}_j^{s'(r'/i) \rightarrow s(r'/i)} B(r'/i) b_1$ we can use Lemma 1 to extend

$$\varphi(r'/i), \forall i. \varphi \vee (r = r') \vdash (b_1, \langle _ \rangle e b_1) : \text{Fiber } e a_1$$

to an element $\varphi(r'/i) \vdash (a, p) : (\text{Fiber } e a_1)[\forall i. \varphi \vee (r = r') \mapsto (b_1, \langle _ \rangle e b_1)]$.

Step 6. We can then use p to extend this to a total element in $A(r'/i)$:

$$\begin{aligned} \vdash a'_1 &:= \text{hcom}_k^{1 \rightarrow 0} A(r'/i) [\varphi(r'/i) \mapsto p k \\ &\quad, (r = r') \mapsto a_0 \\ &\quad, \forall i. (s = s') \mapsto a_1] a_1 : A(r'/i) \end{aligned}$$

Note that on $\varphi(r'/i)$ we have $p 1 = a_1$ (as $(a, p) : \text{Fiber } e a_1$) and on $(r = r')$ we have $a_1 = a_0$, making the homogeneous composition well-formed.

Step 7. Finally we combine all the parts we have computed:

$$\vdash w := \text{box}^{s(r'/i) \rightarrow s'(r'/i)} [\varphi(r'/i) \mapsto a] a'_1 : F(r'/i)$$

To see that w is well-formed note that:

$$\varphi(r'/i) \vdash a'_1 = p 0 = e a = \text{coe}_j^{s'(r'/i) \rightarrow s(r'/i)} B(r'/i) a$$

where the second equality holds as $(a, p) : \text{Fiber } e a_1$.

To see that w satisfies the three necessary properties note that $\forall i. \varphi \leq \varphi(r'/i)$ so that

$$\forall i. \varphi \vdash w = a = b_1 = \text{coe}_i^{r \rightarrow r'} B(s'/j) u_0$$

as desired. Furthermore,

$$\begin{aligned} \varphi(r'/i), (r = r') \vdash a = b_1 = u_0 \\ (r = r') \vdash a'_1 = a_0 = \mathbf{cap}_j^{s(r/i) \leftarrow s'(r'/i)} [\varphi(r/i) \mapsto B(r/i)] u_0 \end{aligned}$$

so that

$$\begin{aligned} (r = r') \vdash w &= \mathbf{box}^{s(r/i) \rightarrow s'(r'/i)} [\varphi(r/i) \mapsto u_0] \\ &(\mathbf{cap}_j^{s(r/i) \leftarrow s'(r'/i)} [\varphi(r/i) \mapsto B(r/i)] u_0) \\ &= u_0 \end{aligned}$$

by the η -rule for box and cap.

Finally, we have:

$$\begin{aligned} \forall i. (s = s') \vdash w &= \mathbf{box}^{s(r'/i) \rightarrow s'(r'/i)} [\varphi(r/i) \mapsto a] a'_1 \\ &= a'_1 \\ &= a_1 \\ &= \mathbf{coe}_i^{r \rightarrow r'} A u_0 \end{aligned}$$

3. SUMMARY OF THE ALGORITHM

Using the variable names above the goal is to define:

$$\vdash w := \text{coe}_i^{r \rightarrow r'} F u_0 : F(r'/i)$$

The algorithm to compute w is then:

$$\begin{aligned} a_0 &= \text{cap}_j^{s(r/i) \leftarrow s'(r'/i)} [\varphi(r/i) \mapsto B(r/i)] u_0 \\ a_1 &= \text{com}_i^{r \rightarrow r'} A [\forall i. \varphi \mapsto \text{coe}_j^{s' \rightarrow s} B(\text{coe}_i^{r \rightarrow i} B(s'/j) u_0), \forall i. (s = s') \mapsto \text{coe}_i^{r \rightarrow i} A u_0] a_0 \\ b_1 &= \text{coe}_i^{r \rightarrow r'} B(s'/j) u_0 \\ \bar{a} &= \text{com}_j^{s(r'/i) \rightarrow j} B(r'/i) [\forall i. \varphi \vee (r = r') \mapsto \text{coe}_j^{s'(r'/i) \rightarrow j} B(r'/i) b_1] a_1 \\ a &= \bar{a}(s'(r'/i)/j) \\ p &= \text{com}_j^{s'(r'/i) \rightarrow s(r'/i)} B(r'/i) [(k = 0) \mapsto \text{coe}_j^{s'(r'/i) \rightarrow j} B(r'/i) a \\ &\quad, (k = 1) \mapsto \bar{a} \\ &\quad, \forall i. \varphi \vee (r = r') \mapsto \text{coe}_j^{s'(r'/i) \rightarrow j} B(r'/i) b_1] a \\ a'_1 &= \text{hcom}_k^{1 \rightarrow 0} A(r'/i) [\varphi(r'/i) \mapsto p, (r = r') \mapsto a_0, \forall i. (s = s') \mapsto a_1] a_1 \\ w &= \text{box}^{s(r'/i) \rightarrow s'(r'/i)} [\varphi(r'/i) \mapsto a] a'_1 \end{aligned}$$

Here is the algorithm from Part III:

$$\begin{aligned} a_0 &= \text{cap}_j^{s(r/i) \leftarrow s'(r'/i)} [\varphi(r/i) \mapsto B(r/i)] u_0 \\ a'_0 &= \text{hcom}_j^{s'(r'/i) \rightarrow j} A(r/i) [\varphi(r/i) \mapsto \text{coe}_j^{j \rightarrow s(r/i)} B(r/i) (\text{coe}_j^{s'(r'/i) \rightarrow j} B(r/i) u_0)] a_0 \\ a_1 &= \text{com}_i^{r \rightarrow r'} A [\forall i. \varphi \mapsto \text{coe}_j^{s' \rightarrow s} B(\text{coe}_i^{r \rightarrow i} B(s'/j) u_0), \forall i. (s = s') \mapsto \text{coe}_i^{r \rightarrow i} A u_0] a'_0(s(r/i)/j) \\ b_1 &= \text{coe}_i^{r \rightarrow r'} B(s'/j) u_0 \\ \bar{a} &= \text{com}_j^{s(r'/i) \rightarrow j} B(r'/i) [\forall i. \varphi \vee (r = r') \mapsto \text{coe}_j^{s'(r'/i) \rightarrow j} B(r'/i) b_1] a_1 \\ p &= \text{coe}_j^{j \rightarrow s(r'/i)} B(r'/i) \bar{a} \\ a'_1 &= \text{hcom}_j^{s(r'/i) \rightarrow s'(r'/i)} A(r'/i) [\varphi(r'/i) \mapsto p, (r = r') \mapsto a'_0] a_1 \\ w &= \text{box}^{s(r'/i) \rightarrow s'(r'/i)} [\varphi(r'/i) \mapsto \bar{a}(s'(r'/i)/j)] a'_1 \end{aligned}$$

The major difference is the simplification in p , which in turn leads to an extra hcom in a'_0 .