A Generalized Logical Framework

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Overview

1. Two-level type theories (2LTT):
   - metaprogramming over a **single model** of a **single type theory**.
   - the chosen model is defined **outside the system**.
   - **only a second-order (“internal”)** view on the model.

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   - *No substructural modalities*.

In this talk:

1. A syntax of GLF + examples + increasing amount of syntactic sugar.
2. A short overview of semantics.

GLF: basic rules

$\mathbf{U}$	$: \mathbf{U}$	A universe of that supports ETT.
$\mathbf{Base}$	$: \mathbf{U}$	Type of "base categories".
$\mathbf{1}$	$: \mathbf{Base}$	The terminal category as a base category.
$\mathbf{PSh}$	$: \mathbf{Base} \to \mathbf{U}$	Universes of presheaves. Cumulativity: $\mathsf{PSh}_i \subseteq \mathsf{U}$. Supports ETT.
		We can only eliminate from PSh_i to PSh_i.
$\mathbf{Cat}_i$	$: \mathbf{PSh}_i$	$:=$ *type of categories in* PSh_i
$\mathbf{In}$	$: \mathbf{Cat}_i \to \mathbf{U}$	"Permission token" for working in presheaves over some $\mathbb{C} : \mathsf{Cat}_i$.
$\mathbf{base}$	$: \mathbf{In}\,\mathbb{C} \to \mathbf{Base}$	"Using the permission".

We use type-in-type everywhere for simplicity, i.e. $\mathsf{U} : \mathsf{U}$ and $\mathsf{PSh}_i : \mathsf{PSh}_i$.

Basic things we can do

$$\mathsf{U} : \mathsf{U} \qquad \mathsf{Base} : \mathsf{U} \qquad 1 : \mathsf{Base} \qquad \mathsf{PSh} : \mathsf{Base} \to \mathsf{U}$$

$$\mathsf{Cat}_i : \mathsf{PSh}_i := \textit{type of cats in } \mathsf{PSh}_i \qquad \mathsf{In} : \mathsf{Cat}_i \to \mathsf{U} \qquad \mathsf{base} : \mathsf{In}\,\mathbb{C} \to \mathsf{Base}$$

PSh_1 is a universe supporting ETT. Semantically, PSh_1 is a universe of sets.

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At this point, we have no interesting interaction between PSh_1 and PSh_i.

Syntactic sugar: we'll omit "base" in the following.

Example: embedding pure lambda calculus

A **second-order model of pure LC** in PSh_i consists of:

$$
\begin{aligned}
&\mathsf{Tm} : \mathsf{PSh}_i \\
&\mathsf{lam} : (\mathsf{Tm} \to \mathsf{Tm}) \to \mathsf{Tm} \\
&\_\$\_ : \mathsf{Tm} \to \mathsf{Tm} \to \mathsf{Tm} \\
&\beta \quad : \mathsf{lam}\, f \,\$\, t = f\ t \\
&\eta \quad : \mathsf{lam}\, (\lambda x.\, t \,\$\, x) = t
\end{aligned}
$$

We define $\mathsf{SMod}_i : \mathsf{PSh}_i$ as the above Σ-type.

Example: embedding pure lambda calculus

A **first-order model of pure LC** consists of:

- A category of contexts and substitutions with $\mathsf{Con} : \mathsf{PSh}_i$, $\mathsf{Sub} : \mathsf{Con} \to \mathsf{Con} \to \mathsf{PSh}_i$ and terminal object $\bullet$.
- $\mathsf{Tm} : \mathsf{Con} \to \mathsf{PSh}_i$, plus a term substitution operation.
- A context extension operation $- \triangleright : \mathsf{Con} \to \mathsf{Con}$ such that $\mathsf{Sub}\,\Gamma\,(\Delta \triangleright) \simeq \mathsf{Sub}\,\Gamma\,\Delta \times \mathsf{Tm}\,\Gamma$.
- A natural isomorphism $\mathsf{Tm}\,(\Gamma \triangleright) \simeq \mathsf{Tm}\,\Gamma$ whose components are λ and application.

We define $\mathsf{FMod}_i : \mathsf{PSh}_i$ as the above Σ-type.

FMod is mechanically derivable from SMod.[1]

[1] Ambrus Kaposi & Szumi Xie: *Second-Order Generalised Algebraic Theories*.

Example: embedding pure lambda calculus

GLF rule

Assuming $M : \mathsf{FMod}_i$ and $j : \mathsf{In}\, M$, we have $\mathsf{S}_j : \mathsf{SMod}_j$.
(In "In M" we implicitly convert M to its underlying category.)

Now we have a 2LTT inside PSh_j:

- ETT type formers in PSh_j comprise the outer level.
- S_j comprises the inner level.

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Y-combinator as example:

$$\mathsf{YC} : \mathsf{Tm}_{\mathsf{S}_j}$$

$$\mathsf{YC} := \mathsf{lam}_{\mathsf{S}_j}(\lambda\, f.\, (\mathsf{lam}_{\mathsf{S}_j}(\lambda x.\, x \,\$_{\mathsf{S}_j}\, x))\,\$_{\mathsf{S}_j}\, (\mathsf{lam}_{\mathsf{S}_j}(\lambda f.\, \mathsf{lam}_{\mathsf{S}_j}(\lambda x.\, f \,\$_{\mathsf{S}_j}\, (x \,\$_{\mathsf{S}_j}\, x)))))$$

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With a reasonable amount of sugar:

$$\mathsf{YC} : \mathsf{Tm}_{\mathsf{S}_j}$$
$$\mathsf{YC} := \mathsf{lam}\, f.\,(\mathsf{lam}\, x.\, x\, x)\,(\mathsf{lam}\, f.\,\mathsf{lam}\, x.\, f\,(x\,x))$$

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- Hence: all 2LTTs are syntactic fragments of GLF.
- (For each 2LTT, the semantics of GLF restricts to the standard presheaf semantics of the 2LTT.)

Yoneda: conversion between internal & external views

GLF rule: Yoneda embedding for pure LC

Assuming $M : \mathsf{FMod}_i$ and writing $\simeq$ for definitional isomorphism, we have

$$
\begin{aligned}
\mathsf{Y} : \mathsf{Con}_M &\to ((j : \mathsf{In}_M) \to \mathsf{PSh}_j)\\
\mathsf{Y} : \mathsf{Sub}_M\,\Gamma\,\Delta &\simeq ((j : \mathsf{In}_M) \to \mathsf{Y}\,\Gamma\,j \to \mathsf{Y}\,\Delta\,j)\\
\mathsf{Y} : \mathsf{Tm}_M\,\Gamma &\simeq ((j : \mathsf{In}_M) \to \mathsf{Y}\,\Gamma\,j \to \mathsf{Tms}_j)
\end{aligned}
$$

such that Y preserves empty context and context extension:

$$
\begin{aligned}
\mathsf{Y}\,\bullet\,j &\simeq \top\\
\mathsf{Y}\,(\Gamma\,\triangleright)\,j &\simeq \mathsf{Y}\,\Gamma\,j \times \mathsf{Tms}_j
\end{aligned}
$$

and Y preserves all other structure strictly.

Notation: we write Λ for inverses of Y.

LC examples, sugar

Υ and Λ allow ad-hoc switching between first-order and second-order notation. Let's redefine some operations using second-order notation:

$$\begin{array}{ll} \mathsf{id} : \mathsf{Sub}_M\,\Gamma\,\Gamma & \mathsf{comp} : \mathsf{Sub}_M\,\Delta\,\Theta \to \mathsf{Sub}_M\,\Gamma\,\Delta \to \mathsf{Sub}_M\,\Gamma\,\Theta \\ \mathsf{id} := \Lambda\,(\lambda\, j\,\gamma.\,\gamma) & \mathsf{comp}\,\sigma\,\delta := \Lambda\,(\lambda\, j\,\gamma.\,\Upsilon\,\sigma\,(\Upsilon\,\delta\,\gamma\,j)\,j) \end{array}$$

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Or even:

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With reasonable amount of sugar:

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Or even:

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Example for "pattern matching" notation:

$$\begin{aligned} &\mathsf{p} : \mathsf{Sub}_M\,(\Gamma\,\triangleright)\,\Gamma \\ &\mathsf{p} := \Lambda\,(\gamma,\,\alpha).\,\gamma \qquad \textit{Note:}\ \mathsf{Y}\,(\Gamma\,\triangleright) \simeq \mathsf{Y}\,\Gamma \times \mathsf{Tm}_{\mathsf{S}_j} \end{aligned}$$

Second-order notation

- When working with CwF-s, De Bruijn indices and substitutions can be hard to read.
- Handwaved "named" binders in CwFs have been used in literature (e.g. by me).
- GLF provides a rigorous implementation of such notation.
- For many use cases, we can use second-order notation and just forget about the first-order combinators.

Embedding dependent type theories

In a first order model, we have:

$$
\begin{aligned}
\mathsf{Con} &: \mathsf{PSh}_i \\
\mathsf{Sub} &: \mathsf{Con} \to \mathsf{Con} \to \mathsf{PSh}_i \\
\mathsf{Ty} &: \mathsf{Con} \to \mathsf{PSh}_i \\
\mathsf{Tm} &: (\Gamma : \mathsf{Con}) \to \mathsf{Ty}\,\Gamma \to \mathsf{PSh}_i \\
&\ldots
\end{aligned}
$$

In a second order model, we have

$$
\begin{aligned}
\mathsf{Ty} &: \mathsf{PSh}_i \\
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Yoneda embedding:

$$\begin{aligned} \mathsf{Y} : \mathsf{Con}_M &\to ((j : \mathsf{In}\,M) \to \mathsf{PSh}_j) \\ \mathsf{Y} : \mathsf{Sub}_M\,\Gamma\,\Delta &\simeq ((j : \mathsf{In}\,M) \to \mathsf{Y}\,\Gamma\,j \to \mathsf{Y}\,\Delta\,j) \\ \mathsf{Y} : \mathsf{Ty}_M\,\Gamma &\simeq ((j : \mathsf{In}\,M) \to \mathsf{Y}\,\Gamma\,j \to \mathsf{Ty}_{\mathsf{S}_j}) \\ \mathsf{Y} : \mathsf{Tm}_M\,\Gamma\,A &\simeq ((j : \mathsf{In}\,M) \to (\gamma : \mathsf{Y}\,\Gamma\,j) \to \mathsf{Tm}_{\mathsf{S}_j}\,(\mathsf{Y}\,A\,j\,\gamma)) \end{aligned}$$

Embedding dependent type theories

Sugar for contexts:

$$(\Gamma \triangleright A \triangleright B) : \mathsf{Con}_M \quad \text{is equal to} \quad \Gamma \triangleright (\Lambda\, \gamma . \mathsf{Y} A\, \gamma) \triangleright (\Lambda\, (\gamma, \alpha) . \mathsf{Y} B\, (\gamma, \alpha))$$

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This suggests the notation:

$$(\gamma : \Gamma,\ \alpha : \mathsf{Y} A\,\gamma,\ \beta : \mathsf{Y} B\,(\gamma,\,\alpha)) : \mathsf{Con}_M$$

With implicit Y:

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Sugar for $\mathbf{Tm}_M$. We have

$$\mathsf{Tm}_M\,(\Gamma \triangleright A \triangleright B)\,C \;=\; \mathsf{Tm}_M\,(\Gamma \triangleright A \triangleright B)\,(\Lambda\,(\gamma,\,\alpha,\,\beta).\,B\,(\gamma,\,\alpha,\,\beta))$$

which suggests the notation

$$\mathsf{Tm}_M\,(\gamma : \Gamma,\, \alpha : A\,\gamma,\, \beta : B\,(\gamma,\,\alpha))\,(B\,(\gamma,\,\alpha,\,\beta))$$

Embedding dependent type theories

Example: a construction which looks awful in explicit CwF notation[2]

$$
\begin{aligned}
\mathsf{Con}^\circ\,\Gamma \quad &:= \mathsf{Ty}\,(F\,\Gamma) \\
\mathsf{Ty}^\circ\,\Gamma^\circ\,A \quad &:= \mathsf{Ty}\,(F\,\Gamma \triangleright \Gamma^\circ \triangleright F\,A[\mathsf{p}]) \\
\mathsf{Tm}^\circ\,\Gamma^\circ\,A^\circ\,t &:= \mathsf{Tm}\,(F\,\Gamma \triangleright \Gamma^\circ)\,(A^\circ[\mathsf{id},\,F\,t[\mathsf{p}]]) \\
\Gamma^\circ \triangleright^\circ A^\circ \quad &:= \Sigma(\Gamma^\circ[\mathsf{p} \circ F_{\triangleright.1}])(A^\circ[\mathsf{p} \circ F_{\triangleright.1} \circ \mathsf{p},\,\mathsf{q},\,\mathsf{q}[F_{\triangleright.1} \circ \mathsf{p}]]) \\
\ldots
\end{aligned}
$$

but is reasonable in sugary GLF notation:

$$
\begin{aligned}
\mathsf{Con}^\circ\,\Gamma \quad &:= \mathsf{Ty}\,(\gamma : F\,\Gamma) \\
\mathsf{Ty}^\circ\,\Gamma^\circ\,A \quad &:= \mathsf{Ty}\,(\gamma : F\,\Gamma,\,\gamma^\circ : \Gamma^\circ\,\gamma,\,\alpha : F\,A\,\gamma) \\
\mathsf{Tm}^\circ\,\Gamma^\circ\,A^\circ\,t &:= \mathsf{Tm}\,(\gamma : F\,\Gamma,\,\gamma^\circ : \Gamma^\circ\,\gamma)\,(A^\circ\,(\gamma,\,\gamma^\circ,\,F\,t\,\gamma)) \\
\Gamma^\circ \triangleright^\circ A^\circ \quad &:= \Lambda\,(F_{\triangleright.2}(\gamma,\,\alpha)).\,\Sigma(\gamma^\circ : \Gamma^\circ\,\gamma) \times A^\circ\,(\gamma,\,\gamma^\circ,\,\alpha)
\end{aligned}
$$

It's a fair amount of sugar, but we can always rigorously desugar when it doubt!

[2] Kaposi, Huber, Sattler: *Gluing for Type Theory*, Section 5

General GLF rules

For every second-order generalized algebraic signature $\mathbb{T}$:

- We compute (externally to GLF) $\mathsf{FMod}_{(\mathbb{T},\,i)}$ and $\mathsf{SMod}_{(\mathbb{T},\,i)}$.
- We specify that GLF has $\mathsf{S}_{(\mathbb{T},\,i)}$.
- We specify that GLF has Yoneda embedding.

It's not simple compute the specification of Yoneda embedding from $\mathbb{T}$! Doing this is part of future work.

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It's not simple compute the specification of Yoneda embedding from $\mathbb{T}$! Doing this is part of future work.

Also, these are not all rules that we might want to have!

- *For example:* conversion between internal and external natural numbers, i.e. $\mathbb{N}_i \simeq ((j : \mathsf{In}_M) \to \mathbb{N}_j)$ where $M : \mathsf{Cat}_i$.
- This can be broadly generalized to an isomorphism of "external" and "internal" 2LTT models.
- But we're not sure yet which rules are the best to enshrine in GLF syntax.

Semantics

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We should work with **Cat** somehow, but there are issues with that:

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GLF contexts are modeled as certain *trees of categories*:

- Each node represents a presheaf universe, each edge represents an internal/external switch.
- Tree morphisms only have non-trivial action on discrete data in trees.

Semantics

Notation:

- For a category C and a split fibration A over it, we write $C \triangleright A$ for the total category.
- For a presheaf A, we write $\mathrm{Disc}\, A$ for the derived discrete fibration.

Definition. A *category telescope* is either the terminal category, or it is (inductively) of the form $C \triangleright \mathrm{Disc}\, A \triangleright B$ where C is a category telescope. We write $C : \mathrm{CatTel}$ for a category telescope.

Definition. A tree of categories is inductively defined as:

$$
\begin{aligned}
&\textbf{data}\ \mathrm{Tree}\,(B : \mathrm{CatTel}) : \mathrm{Set}\ \textbf{where} \\
&\quad \mathrm{node} : (\Gamma : \mathrm{PSh}\, B) \\
&\qquad \to (n : \mathbb{N}) \\
&\qquad \to (C : \mathrm{Fin}\, n \to \mathrm{Fib}\,(B \triangleright \mathrm{Disc}\,\Gamma)) \\
&\qquad \to ((i : \mathrm{Fin}\, n) \to \mathrm{Tree}\,(B \triangleright \mathrm{Disc}\,\Gamma \triangleright C\, i)) \\
&\qquad \to \mathrm{Tree}\, B
\end{aligned}
$$

Semantics

$$\mathsf{node} : (\Gamma : \mathsf{PSh}\, B)(n : \mathbb{N})(C : \mathsf{Fin}\, n \to \mathsf{Fib}\,(B \triangleright \mathsf{Disc}\,\Gamma)) \to ((i : \mathsf{Fin}\, n) \to \mathsf{Tree}\,(B \triangleright \mathsf{Disc}\,\Gamma \triangleright C\, i)) \\ \to \mathsf{Tree}\, B$$

A GLF context is an element of Tree 1. Some examples for semantic contexts. We have $\mathbb{N}_i : \mathsf{PSh}_i$. We use $- \triangleright -$ for "context extension" in presheaves as well.

$$\begin{aligned}
\bullet \quad & := \mathsf{node}\, 1\, 0\, [\,]\, [\,] \\
(\bullet \triangleright \mathbb{N}_1) \quad & := \mathsf{node}\,(1 \triangleright \mathbb{N})\, 0\, [\,]\, [\,] \\
(\bullet \triangleright \mathbb{N}_1 \triangleright \mathsf{In}\ C) \quad & := \mathsf{node}\,(1 \triangleright \mathbb{N})\, 1\, [C]\, [\mathsf{node}\, 1\, 0\, [\,]\, [\,]] \\
(\bullet \triangleright \mathbb{N}_1 \triangleright i : \mathsf{In}\ C \triangleright \mathbb{N}_{(\mathsf{base}\, i)}) & := \mathsf{node}\,(1 \triangleright \mathbb{N})\, 1\, [C]\, [\mathsf{node}\,(1 \triangleright \mathbb{N})\, 0\, [\,]\, [\,]]
\end{aligned}$$

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\bullet \quad & := \mathrm{node}\, 1\, 0\, [\,]\, [\,] \\
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(\bullet \triangleright \mathbb{N}_1 \triangleright \mathrm{In}\, C) \quad & := \mathrm{node}\,(1 \triangleright \mathbb{N})\, 1\, [C]\, [\mathrm{node}\, 1\, 0\, [\,]\, [\,]] \\
(\bullet \triangleright \mathbb{N}_1 \triangleright i : \mathrm{In}\, C \triangleright \mathbb{N}_{(\mathrm{base}\, i)}) & := \mathrm{node}\,(1 \triangleright \mathbb{N})\, 1\, [C]\, [\mathrm{node}\,(1 \triangleright \mathbb{N})\, 0\, [\,]\, [\,]]
\end{aligned}$$

- A Base in context Γ points to a node in Γ.
- An In C in context Γ points to a subtree of a node.
- Extending a context with $A : \mathrm{PSh}_i$ extends the presheaf in node i.
- Extending a context with $j : \mathrm{In}\, C$ for $C : \mathrm{Cat}_j$ adds a new subtree at node j.

Semantics

$$\mathsf{node} : (\Gamma : \mathsf{PSh}\,B)(n : \mathbb{N})(C : \mathsf{Fin}\,n \to \mathsf{Fib}\,(B \triangleright \mathsf{Disc}\,\Gamma)) \to ((i : \mathsf{Fin}\,n) \to \mathsf{Tree}\,(B \triangleright \mathsf{Disc}\,\Gamma \triangleright C\,i)) \\ \to \mathsf{Tree}\,B$$

Tree morphisms are defined inductively & levelwise, containing

- natural transformations between $\Gamma : \mathsf{PSh}\,B$ components
- functions for reindexing subtrees of type $\mathsf{Fin}\,n \to \mathsf{Fin}\,m$

such that the non-discrete fibrations are preserved.

A semantic $\mathbf{PSh}_{\boldsymbol{i}}$ in context Γ is a presheaf over the category given by the path from the root of Γ to the node i.

Further work

- Decide on the exact rules of GLF.
- Compute the specification of Yoneda embedding from SOGAT signatures, define semantics in this generality.
- Investigate syntactic metatheory.
  - For computer implementation, we need to wean ourselves off extensional TT!
    - (but informal extensional GLF is already useful)
  - Definitional isos for Y are unusual in syntax.
  - Simpler syntactic fragments of GLF could be useful & easier to implement.

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Thank you!

Shameless bonus advertisement: 40th Agda implementors' meeting, Budapest, May 26-31, free participation, `https://wiki.portal.chalmers.se/agda/Main/AIMXXXX`