

In this TP we consider a compression function  $h$  that takes an  $n$ -bit message block  $m$  and an index  $i$ , and returns an  $n$ -bit string:

$$h : \mathbb{N} \times \{0, 1\}^n \rightarrow \{0, 1\}^n .$$

The function  $h$  is defined in the file `tp4_code.py`, where  $n := 48$ . For compactness, both input and output values are encoded as python integers.

For a fixed integer  $k$ , we use  $h$  to define a  $k$ -block hash function as follows:

$$\begin{cases} H : & (\{0, 1\}^n)^k & \rightarrow & \{0, 1\}^n \\ & m_0, \dots, m_{k-1} & \mapsto & \left( \bigoplus_{i=0}^{k-1} h(i, m_i) \right) \end{cases} \quad (1)$$

## k-Collisions

Let  $c, m$  be two integers. Let  $L_i, L_j$  be two lists of pairs  $(h(i, x), x)$  and  $(h(j, y), y)$  respectively, which are sorted using lexicographic ordering of the bit-strings  $h(i, x)$  and  $h(j, y)$  (LSBs first). In the file `tp4_code.py` we provide several functions to facilitate the manipulation of integers instead of bit-strings:

- Of course, the  $\wedge$  operator on integers takes the XOR of the bit-strings;
- The function `lower(x, y, c)` returns True iff the value of the  $c$  LSBs of  $x$  is strictly lower (in lexicographic ordering) than the one of  $y$ ;
- The function `eq(x, y, c)` returns True iff they are equal;
- The function `sort(l)` sorts a list of integer tuples according to the lexicographic ordering of the bit-string of the first value.

If the  $c$  LSBs of a bit-string / integer  $x$  are zero, we write  $x|_c = 0$ .

Let  $L_i \bowtie_c L_j$  be the set of tuples:

$$L_i \bowtie_c L_j = \{(h(i, x) \oplus h(j, y), x, y), x \in L_i, y \in L_j, (h(i, x) \oplus h(j, y))|_c = 0\} .$$

also represented as a *sorted list*. The list  $L_i \bowtie_c L_j$  will be called “merged list”. In other words, we are computing a list of *partial collision* pairs, where the  $c$  LSBs of the pairs collide.

**Question 1.** Show that there exists an algorithm of complexity:  $\max(|L_i|, |L_j|, |L_i \bowtie_c L_j|)$  that, on input  $L_i, L_j$ , returns  $L_i \bowtie_c L_j$ .

**Question 2.** Implement this algorithm in a general setting, when the two lists are lists of tuples of integers, lexicographically sorted according to the first value (see the description in `tp4_code.py`).

**Question 3.** Fix  $i, j = 1, 2$  and  $c = n/3$ . Asymptotically, what is the expected size of  $L_1 \bowtie_c L_2$ ?

**Question 4.** Let  $k = 4$ . Find an algorithm of complexity  $\mathcal{O}(2^{n/3})$  that finds a preimage of 0 by  $H$ . Implement this attack.

**Question 5.** Can we improve this complexity? How does it depend on the value of  $k$ ?