

avocado: A Variant Caller, Distributed

Frank Austin Nothaft, Peter Jin, Brielin Brown
Department of Electrical Engineering and Computer Science
University of California, Berkeley
{fnothaft,phj,brielin}@berkeley.edu

ABSTRACT

In this paper, we present *avocado*, a distributed variant caller built on top of ADAM and Spark. *avocado*'s goal is to provide both high performance and high accuracy in an open source variant calling framework. To achieve this, we implement both local assembly and pileup-based single nucleotide polymorphism (SNP) calling. A key innovation presented in our work involves the development of heuristics for when to choose more expensive assembly-based methods instead of pileup-based methods. Additionally, we introduce the concept of "significant statistics," a tool for performing incremental joint variant calling.

Categories and Subject Descriptors

Applied Computing [Life and Medical Sciences]: Computational Biology—*Sequencing and Genotyping Technologies*; Applied Computing [Genomics]: Computational Genomics; Computing Methodologies [Distributed Computing Methodologies]: Distributed Algorithms—*MapReduce Algorithms*

General Terms

Algorithms, Performance

Keywords

Variant Calling, Genotyping, Genomics Pipeline, Local Assembly, Distributed Computing, MapReduce

1. INTRODUCTION

Modern genomics processing pipelines can be divided into four primary ordered stages:

1. **Sequencing:** Gathering of read data from DNA
2. **Alignment:** Alignment of read data against reference genome

3. **Variant Calling:** Statistical determination of differences against reference genome
4. **Variant Annotation:** Annotation of impact of variation

Currently, to run a genomics pipeline end-to-end for a single high coverage genome¹ consumes approximately 100 hours [15]. Of this 100 hour figure, both alignment and variant calling consume approximately 50 hours each.

Although some applications that use genomic data are latency insensitive (for example, population genomics), many medical applications like genomic medicine, or genomic classification of viral outbreaks [14] are latency sensitive. However, it is unacceptable to sacrifice accuracy in the pursuit of speed. Recent work has focused on the problem of accelerating alignment [17]; in this paper, we address accelerating variant calling.

As noted above, it is unacceptable to sacrifice accuracy for performance. To achieve improved performance, we implement several enhancements:

- Current pipelines are penalized by I/O performance, we address this by using an in-memory MapReduce framework [18] to reduce I/O pressure
- Additionally, we leverage the new ADAM data format [10], a high performance file format for distributed genomics
- Finally, we achieve high accuracy at a low performance cost by using high fidelity assembly-based methods only on complex segments of the genome

In this paper, we discuss this system, related work, and perform a performance analysis. We start with a discussion of the related work in §2. We then describe our architecture and algorithms in §3. Finally, we analyze the performance of our system in §4, and propose future research directions in §5.

¹High coverage refers to having on average $>30\times$ bases aligned to each location in the reference genome.

2. RELATED WORK

There has been significant work related to variant calling, and towards accelerating the genomic processing pipeline. In this section, we discuss other variant callers, and tools that we use in our evaluation.

2.1 ADAM

ADAM [10] is a new data format for genomics meant to replace the Sequence/Binary Alignment/Map (SAM/BAM) formats for read data [9], and the Variant Call Format (VCF) for variant/genotype data [2]. The original SAM/BAM/VCF formats were designed for single-node processing, and do not easily distribute across several machines. Although a library was designed for processing BAM/VCF data in Hadoop [12], this API does not scale well past 8 nodes. ADAM achieves scalability beyond 100 machines by eliminating the central file header, and by using the Parquet data store which is optimized for parallel data access [16].

In the process of developing *avocado*, we contributed 3,500 lines of code (LOC) to the ADAM project. This contribution comprised the variant and genotype format, code for calculating normalized variant data from genotypes, and converters to/from the VCF format. Additionally, this contribution included code for translating between read and reference oriented views of data.

2.2 The Genome Analysis Toolkit

[11, 3]

2.3 Samtools Mpileup

[7]

2.4 FreeBayes

[4]

2.5 SNAP

SNAP is a high performance short-read aligner that is optimized for longer read lengths, and for distributed computing [17]. At the current moment, we have not integrated with SNAP, but long term, we plan to incorporate SNAP as the aligner in our read-alignment and variant calling pipeline. This is significant, as variant callers typically optimize to correct for the error models of the aligners that they coexist with.

SNAP leverages the increasing length of reads to build a large alignment index, which is similar to the method used by BLAST [1], and which is dissimilar to the Burrows-Wheeler transform based methods used by BWA [8]. Aligners which use the Burrows-Wheeler transform perform very well in the absence of mismatching data—however, they cannot handle mismatches or inserts well. BWA has bad failure modes for reads with mismatches within the first 20 bases of the read, and Bowtie [6] does not handle insertions when aligning. As SNAP will have better performance when aligning indels, it is likely that we will be able to omit the local realignment stage from ADAM [10]—this is significant as local realignment is the most expensive phase of read processing before variant calling.

2.6 SMaSH

SMASH is a benchmarking suite for alignment and variant calling pipelines [15], and was a key tool used for the evaluation of *avocado*.

Traditionally, variant calling pipelines have been evaluated on concordance², and through using venn diagrams [5]. This is because genomics rarely has access to *ground truth*: typical sequencing methods have insufficient fidelity to detect all variants clearly, and extremely high fidelity sequencing methods are too expensive/slow to be used in clinical practice. However, this method is fraught with risk: concordance is not a good metric to use if variant callers or aligners are making similar systemic errors.

To address this problem, SMASH leverages synthetic data which by definition has known ground truth, and rigorously verified mouse and human genomes. The human genomes and validation data come from the 1000 Genomes project, a project which surveyed the genomes of 1000 individuals using multiple sequencing technologies [13]. On top of this curated data, SMASH provides a VCF based interface for determining the precision of a variant calling pipeline. Novel to this benchmarking suite, the authors introduced a “rescue” phase, which is used to resolve ambiguities that are caused by the VCF specification.

It is worth noting that SMASH does not include any datasets that are designed for joint variant calling. Because of this, we fall back on concordance as a metric for evaluating our joint variant calling algorithms which are described in §3.3.

3. ARCHITECTURE

When architecting *avocado*, we made a conscientious decision to prioritize modularity and extensibility. There are several reasons behind this design choice:

- Current variant calling pipelines are not meant to be extended. Because of this, anyone looking to prototype a new variant calling algorithm must implement their own variant calling infrastructure, which is a significant impediment to variant calling research.
- Although we have limited specialization in our current pipeline (we specialize towards local assembly and pileup based SNP calling), long term we plan to add increasingly specialized variant calling algorithms.
- Similarly, it is known that modern variant callers perform poorly when calling structural variants (SVs). This is a critical area that we hope to attack through specialized variant calling algorithms.

To improve modularity, we pushed as much functionality into the ADAM stack as possible [10]. ADAM implements several important transformations that are used on the read processing frontend, including sort-by-reference position, duplicate marking, base quality score recalibration (BQSR), and local realignment (see §2.5). Additionally, we have

²Identifying overlap between the call set of multiple variant calling pipelines.

moved portions of the variant calling stack into ADAM—after genotyping samples, we use the ADAM pipeline to transform genotypes into variant calls. We include a brief discussion of the merits of this approach in Appendix B.

A diagram of the `avocado` architecture is included in Figure 1.

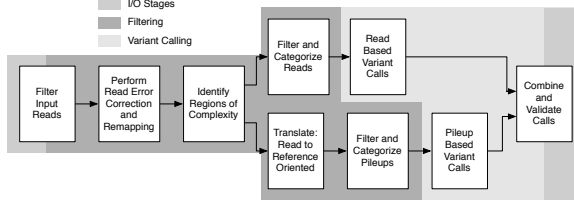


Figure 1: System Architecture

`avocado` is roughly divided into four stages:

1. **Read pre-processing:** Applies transformations from the ADAM pipeline, including sort, duplicate marking, BQSR, and local realignment. Operates strictly on read data.
2. **Filtering:** Selects the best algorithm for calling variants on a segment of data; transforms data to be reference-oriented if necessary. This is discussed in §3.4.
3. **Variant calling:** Calls genotypes for samples in input set on either read-oriented or reference-oriented data. These algorithms are discussed in §3.1 and §3.2.
4. **Variant filtering:** Here, variant data is calculated from genotypes using the ADAM API [10]. Additionally, variants can be filtered for quality.

All stages are designed to be configurable, and easy to replace. Specifically, the filtering and variant calling stages are designed with a specific eye towards modularity. For variant calls, we provide a base class, and two specialized subclasses for variant calls that operate on either read or reference oriented data. A variant calling algorithm can be added by implementing one of these class interfaces, and then registering the call with a filter. For filters, we provide two filtering stages: the first operates on read data, and is used to implement the filter described in §3.4. The second filtering stage operates on pileup (reference oriented) data—currently, we use this to filter out pileups that have no mismatch evidence. For both of these filters, we provide a simple interface for developers to implement.

To improve pipeline performance, we made a significant optimization to the reads-to-pileup transformation step. At the start of transforming reads into reference oriented data, the reads are sorted by reference position. However, due to the partitioning used by Spark [18], after a naïve transformation of reads to single reference oriented bases, the reference oriented bases are no longer sorted. This makes the ensuing grouping of bases by position very expensive due to significant disk shuffle I/O. Instead, we perform a quasi-streaming

transform on the read data. Here, the sorted read blocks are chunked together. All of the reads in this window are then converted into pileup bases, and then grouped together into rods. This leads to significantly better performance, but has tricky edge cases: reads that border the end of a group must be duplicated into both read groups.

3.1 Local Assembly

Given a partition of reads, we can group them by their starting locus in intervals of W , creating regions of length $W + L - 1$ where L is the read length. Within each region, we can evaluate the likelihood of observing the reads given the reference haplotype:

$$\begin{aligned} \mathcal{L}(H^{\text{ref}}) &\equiv \mathbf{P}(\{r_i\} | H^{\text{ref}}) \\ &= \prod_i \mathbf{P}(r_i | H^{\text{ref}}) \end{aligned} \quad (1)$$

where $\mathbf{P}(r|H)$ is obtained from aligning the read and the candidate haplotype by a pairwise HMM alignment. Note that, in practice, all probabilities are computed in units of logarithm base 10, so products become sums, etc.

If a reference haplotype likelihood is below a fixed threshold, the region corresponding to the haplotype is marked *active*. Each *active region* is assembled independently and in parallel.

The assembly of an active region starts by splitting all reads assigned to the region into k -mers, where k is a fixed parameter for all assemblies. A read generates $L - k + 1$ total k -mers. Each k -mer is uniquely identified by the substring of its originating read. Because of coverage overlap and sequence repeats, some k -mers will be duplicates; these are consolidated, and the duplication factor is recorded as the k -mer multiplicity.

The k -mers define edges in the completed k -mer assembly graph. Within a read, each adjacent pair of k -mers have an overlapping substring of length $k - 1$; these are seeded as the initial vertices in the k -mer graph. Because there are duplicated k -mers, some vertices will be “merged,” connecting the graph. Unlike an exact de Bruijn graph, which connects all overlaps between k -mers, we only connect the overlaps found in the reads, performing a simple form of read threading.

Once the k -mer graph is complete, we perform a depth-first traversal with an upper bound on the total path multiplicity, defined as the sum of the edge multiplicities, to enumerate a set of possible paths. The traversal begins at a graph source, and a completed path must also end at a sink. Each path is an assembled haplotype to be evaluated.

Since assembly depends on the exact sequences of the input data, the quality of the reads is critical for performing an accurate assembly. One error correction method is spectral filtering, which depends on the observation that the distribution of k -mers from all reads with respect to multiplicity is bimodal, one due to the Poisson sampling of reads (with high multiplicity), the other due to errors (with low multiplicity), so that splitting the reads between the two modes and keeping the reads near the mode with higher multiplicity serves to prune the poor quality reads. Errors in the

read data lead to *spurs*, which are spurious short sections of the graph connected to the source or the sink. We do not remove spurs, although it may slightly improve performance. Empirically, we have found during other work that utilizing mate pair information greatly improves the quality of an assembly. We also do not employ mate pair threading, which requires collecting the insert size distribution of the data, but we expect that implementing it has the potential to vastly improve the accuracy of variant calls.

From the assembled haplotypes, we order them according to the haplotype likelihood:

$$\mathcal{L}(H_j) = \prod_i \mathbf{P}(r_i | H_j). \quad (2)$$

Among the ordered haplotypes, we pick the top scoring haplotypes and ignore the low scoring ones. The likelihood of observing the reads $\{r_i\}$, given a pair of haplotypes H_j and $H_{j'}$, is defined to be:

$$\begin{aligned} \mathcal{L}(H_j, H_{j'}) &\equiv \mathbf{P}(\{r_i\} | H_j, H_{j'}) \\ &= \prod_i \left[\frac{\mathbf{P}(r_i | H_j)}{2} + \frac{\mathbf{P}(r_i | H_{j'})}{2} \right]. \end{aligned} \quad (3)$$

We compute the posterior probability of observing the pair of haplotypes H_j and $H_{j'}$ from the haplotype pair likelihood and a haplotype pair prior probability:

$$\mathbf{P}(H_j, H_{j'} | \{r_i\}) = \frac{1}{Z} \mathcal{L}(H_j, H_{j'}) \mathbf{P}(H_j, H_{j'}) \quad (4)$$

where Z is a normalization:

$$Z = \sum_j \sum_{j'} \mathcal{L}(H_j, H_{j'}) \mathbf{P}(H_j, H_{j'})$$

and where we obtain the prior $\mathbf{P}(H_j, H_{j'})$ by aligning the haplotype pair with the same pairwise HMM alignment as above, and taking the product of the prior probabilities for each SNP and indel event.

We choose the maximum a priori estimate among haplotypes with any variants as the called non-reference maternal and paternal haplotype pair:

$$(H_{\text{mat}}^{\text{nonref}}, H_{\text{pat}}^{\text{nonref}}) = \arg \max_{H_j, H_{j'} : n_{\text{var}}(H_j, H_{j'}) > 0} \mathbf{P}(H_j, H_{j'} | \{r_i\}). \quad (5)$$

Similarly, we may define the reference haplotype pair as $(H^{\text{ref}}, H^{\text{ref}})$. The error probability of calling the non-reference haplotype pair is:

$$\begin{aligned} \mathbf{P}_{\text{error}}(H_{\text{mat}}^{\text{nonref}}, H_{\text{pat}}^{\text{nonref}}) \\ = \frac{\mathbf{P}(H^{\text{ref}}, H^{\text{ref}})}{\mathbf{P}(H_{\text{mat}}^{\text{nonref}}, H_{\text{pat}}^{\text{nonref}}) + \mathbf{P}(H^{\text{ref}}, H^{\text{ref}})}. \end{aligned} \quad (6)$$

The quality score of all variants present in the nonreference haplotype pair is defined as the Phred scaling of $\mathbf{P}_{\text{error}}$.

3.2 SNP Calling

3.3 Joint Variant Calling

3.4 Algorithm Selection

As discussed in §3, we seek to improve the performance of our variant caller without reducing accuracy by directing the variant caller to use higher accuracy methods in areas that

show increased complexity. Loosely, we define a complex region as an area that is highly similar to other areas of the genome, or where it is likely that a complex variant (such as an indel) is segregating. To identify those regions, we use the following heuristics:

- Areas that are highly similar to other areas of the genome can be distinguished by low mapping quality. Reduced mapping quality indicates that alignment found several areas where the read could map with similar quality.
- Complex variants lead to a change in coverage over the effected area. Deletions will see reduced coverage, and insertions lead to increased coverage.

We implemented our filter by stepping a window across the reference genome. The window was stepped by 1000 base pairs. In each 1000 base pair window, we would compute the average coverage, and mapping quality. If a window violated either of the mapping quality or the coverage threshold, we would flag it as high complexity and pass it to the assembler. If a window did not violate either of those two thresholds, we built pileups out of the data on this interval, and then used the SNP calling methods described in §3.2.

We provided several tuning parameters to the end user. Specifically, the end user could set the target mapping quality and coverage deviation percentage (percent change from mean coverage). Table 1 summarizes the sensitivity of these parameters.

Table 1: % Reads in High Complexity Region

MapQ, Cov	0.2	0.4	0.6
40	88%	54%	20%
50	90%	56%	22%
60	98%	91%	88%

As can be noted, mapping complexity begins to saturate as the mapping quality threshold increases to 60³, and as the coverage variation coefficient decreases to 0.2.

4. EVALUATION

4.1 Accuracy

4.2 Performance

5. FUTURE WORK

6. CONCLUSION

APPENDIX

A. AVAILABILITY

avocado is open source and is licensed under the Apache 2 license. The source code is available at:

<http://www.github.com/bigdatagenomics/avocado>

³Phred-scaled 60 is equivalent to a probability of $P = 0.999999$ that the read is correctly mapped.

B. GENOTYPE/VARIANT REFACTORING

Traditionally, variant calling has involved genotyping n samples, and then determining which of the genotypes of these samples contain true variants. This involves the calculation of a *variant quality score*⁴, and filtering variants.

Semantically, all of the information necessary for the creation of variant data can be extracted from the genotypes of the samples. Therefore, the step of packing/unpacking genotype data to create variant calls is unnecessary. To reduce the likelihood of making errors, we have migrated code to do this into the ADAM [10] framework. Then, to get variant calls, we must just genotype the samples we see, and call the ADAM library.

This does not preclude joint variant calling—although joint variant calling implies that multi-sample data is being used to influence the filtering of variants, this is a bit of a misnomer. Practically, as discussed in §3.3, joint variant calling involves the use of data from multiple samples to refine genotype probabilities across a population to support genotypes that are frequently seen in this population.

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⁴Loosely defined as the likelihood that there is *at least* one genotype with this variant out of all haplotypes seen in the call set.