



**POLITECNICO**  
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# BLOCKCHAIN NOTARIZATION: EXTENSIONS TO THE OPENTIMESTAMPS PROTOCOL

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# INTRODUCTION

- **Notarization** of digital documents is traditionally achieved by a trusted certification authority.
- A notary public could be malevolent and represents a **single point of failure**.
- An application built upon a **distributed system** reaching consensus in a **decentralized way** can greatly strengthen the notarization process.
- The Bitcoin blockchain, the most reliable decentralized system, can be used as a notary in a procedure called **timestamping**.
- The open-source protocol **OpenTimestamps** aims to be a standard for timestamping and solves a scalability issue.

# OUTLINE

- 1 Distributed Consensus
  - Distributed Systems
  - The Consensus Problem
- 2 Nakamoto Consensus in Bitcoin
- 3 Blockchain Timestamping
- 4 OpenTimestamps
- 5 Conclusions

## Definition (distributed system)

*A distributed system is a system whose components are located on different networked computers, which communicate and coordinate their actions by passing messages to one another. The components interact with one another in order to achieve a common goal.*

## Examples:

- ATM points.
- Computers in cloud.
- Modern smartphones.

# DISTRIBUTED SYSTEMS: PROS VS. CONS

## PROS:

- Scalability.
- Performance.
- Fault tolerance.
- Reliability.
- Availability

## CONS:

- Security.
- Consensus.

Thus, the main challenge of distributed systems is to achieve **distributed consensus** in the presence of a number of **faulty processes**.

# THE CONSENSUS PROBLEM

- Networks are composed by many agents, called **nodes**.
- Nodes are either **honest** or **byzantine**.
- They could also be **correct** and **faulty**, but not as alternatives to honest and byzantine.

## Definition (consensus)

*There are  $n$  nodes, of which at most  $f$  might crash, i.e., at least  $n - f$  nodes are correct. Node  $i$  starts with a proposed value  $v_i$ . The nodes must decide for one of those values, satisfying the following properties:*

- *Agreement: All correct nodes decide for the same proposal.*
- *Termination: All correct nodes terminate in finite time.*
- *Validity: The decision value must be the proposal of some proposer.*

# BYZANTINE AGREEMENT IN AN ASYNCHRONOUS NETWORK

- Is it possible to achieve **distributed consensus** in an **asynchronous** network in presence of **byzantine** nodes ?
- The **Byzantine General's Problem** is the abstraction of such situation.

More precisely:

## **Definition** (byzantine general's problem)

*A commanding general must send an order to his  $n - 1$  lieutenant generals such that the following conditions are satisfied:*

- *All loyal lieutenants obey the same order.*
- *If the commanding general is loyal, then every loyal lieutenant obeys the order he sends.*

# IMPOSSIBILITY RESULT

- Achieving **distributed consensus** in an asynchronous network in presence of byzantine nodes is a very **hard problem**.
- Fisher, Lynch and Paterson [12] formally proved that it is **impossible** to reach distributed consensus in an asynchronous network with **one faulty process**, even for a **binary** variable.
- So how Satoshi Nakamoto reaches consensus in **Bitcoin**, a decentralized, distributed, peer-to-peer network ?



# OUTLINE

- 1 Distributed Consensus
- 2 Nakamoto Consensus in Bitcoin
  - Eventual Consistency
  - Bitcoin Design
- 3 Blockchain Timestamping
- 4 OpenTimestamps
- 5 Conclusions

# BITCOIN AS AN EXAMPLE OF EVENTUAL CONSISTENCY

- **Bitcoin** is a distributed, decentralized, peer-to-peer electronic payment system based on **cryptographic proof**.
- The Bitcoin network is subjected to **network partition**.
- It is inherently characterized by a trade-off between:
  - **Consistency**: All nodes agree on the current state.
  - **Availability**: The system is operational and instantly processing incoming requests.
  - **Partition tolerance**: The ability to continue operating correctly even in the presence of a network partition.
- CAP theorem proves that is **impossible** to achieve **simultaneously** the three properties.
- Nakamoto's intuition has been to relax the **agreement** property in the consensus definition to hold **probabilistically**, to reach **eventual consistency**.

# DOUBLE SPENDING PROBLEM

- In a digital cash scheme, a **single digital token**, being just a file that can be duplicated, can be **spent twice**.

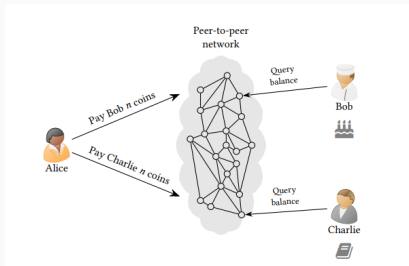


Figure 1: Illustration of the double spending problem.

- Nakamoto solves the **double spending problem** combining **cryptography** and **social incentives**. But first, some protocol's mechanics.

# TRANSACTIONS IN BITCOIN

- A **bitcoin** is defined as a chain of digital signatures.
- Coins cannot be combined, subdivided or transferred, but only **entirely consumed** as transaction inputs (TxIn) to create new output coins (TxOut).
- A TxOut can be in two states: **spent** or **unspent**.
- The **UTXO** is the set of current unspent transactions.
- A **TxOut** consists of an amount of bitcoins and a **locking script**, a cryptographic puzzle which determines the spending conditions.
- A **TxIn** consists of a pointer referencing the UTXO that it consumes and an **unlocking script**, the solution to solve the locking one.

# THE BLOCKCHAIN

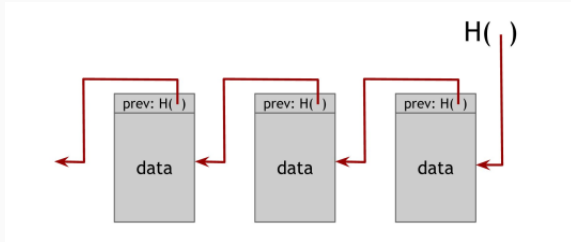


Figure 2: Blockchain as a hash pointer linked list.

- Transactions are recorded in a **distributed ledger**, the **blockchain**, formally a hash pointer linked list.
- The **cryptographic link** between blocks requires computing power to be created.
- A block is **valid** only if it includes valid transactions.

# MINING & PROOF OF WORK

- All network nodes **validate** and **clear** all transactions.
- Special nodes, called **miners**, compete to finalize a new block of transactions, providing **proof-of-work**, which consists of finding a special number **x** called **nonce** s.t.

$$\text{SHA256}(\underbrace{\text{SHA256}(\dots\|\text{prev\_block\_header\_hash}\|\dots\|x))}_{\text{Candidate Block Header}} < \frac{2^{224}}{d}.$$

- The miner who **first** find it is rewarded with the issuance of new bitcoins in a special **coinbase** transaction.
- Miners solve the **double spending problem**:
  - A double spending transaction invalidates a block.
  - The bitcoin reward would have removed.
  - The winning miner would have wasted his work.
  - Proof-of-work ensures **blockchain immutability**.

# OUTLINE

- 1 Distributed Consensus
- 2 Nakamoto Consensus in Bitcoin
- 3 Blockchain Timestamping
  - Notarization
  - Commitment Operations
  - Timestamping Procedure
- 4 OpenTimestamps
- 5 Conclusions

# NOTARIZATION

- **Notarization** is the official fraud-deterrend process that assures the parties of a transaction that a document is **authentic** and can be **trusted**.
- Traditionally, it is achieved by a trusted certification authority which represents a **single point of failure**.
- Need to **decentralize** the source of trust to be reliable even if the security of such central authority is violated, using the Bitcoin blockchain as a **notary** for **timestamping**.

## **Definition** (timestamp)

*A timestamp is a proof that some data  $d$  existed prior to a certain time  $t$  and in a certain state.*



# COMMITMENT OPERATIONS

- To create a timestamp, data  $d$  has to **cause** an event that could not have been generated without the existence of  $d$ , must be **attested** to time  $t$  and can be **publicly observed**.
- A **good** timestamp must become **invalid** even if a **single bit** of the input data is modified.
- What is published on the blockchain is a **commitment** to  $d$ .

## Definition (commitment operation)

A function  $C : X \rightarrow Y$  is a commitment operation if given  $x_1 \in X$  it is not feasible to compute  $x_2 \in X$  s.t.  $x_1 \neq x_2, C(x_1) = C(x_2)$ .

## Examples:

- append: "hello"  $\xrightarrow{\text{append}(\text{"world"})}$  "helloworld".
- prepend: "world"  $\xrightarrow{\text{prepend}(\text{"hello"})}$  "helloworld".

# HASH FUNCTIONS

- However, both “append” and “prepend” commitment operations are **not good** for timestamping:
  - They do not **hide** input data.
  - Outputs are always **bigger** in size.
- **Hash functions** map input data of **arbitrary length** into hash values, i.e., outputs of **fixed length**:
  - **Preimage resistant** (one-wayness): given  $h(x)$  it is not feasible to compute  $x$ .
  - **Second-preimage resistant**: given  $x$  it is not feasible to compute  $y$  s.t.  $x \neq y$ ,  $h(x) = h(y)$ .
  - **Collision resistant**: it is not feasible to find  $x, y$  s.t.  $x \neq y$ ,  $h(x) = h(y)$ .
- An hash value represents a **digital fingerprint**.
- Bitcoin uses **SHA256**: outputs have a fixed size of 256-bit.

# TIMESTAMPING PROCEDURE

- A generic data file can be **hashed** to produce a short unique identifier.
- Such a **digital fingerprint** can be associated to a particular Bitcoin transaction, called **null data transaction**.
- These transactions make use of **OP\_RETURN**, an opcode script which allows to add up to 80 bytes of arbitrary data in a **provably unspendable** transaction.
- Such transaction **commits** to a specific field of the block header, called **merkleroot**, which is the root of a binary tree that summarize all the transactions in a block.
- Blockchain **immutability** provides **timestamping**, proving the data file existence at that moment in time in that specific status.

# OUTLINE

- 1 Distributed Consensus
- 2 Nakamoto Consensus in Bitcoin
- 3 Blockchain Timestamping
- 4 OpenTimestamps**
  - OpenTimestamps Protocol
  - A Practical Use Case
- 5 Conclusions

- Blockchain timestamping have some **downsides**:
  - Not efficient: it requires one transaction per document.
  - Lack of standardization.
- **OpenTimestamps** is an open-source protocol that aims to be a **standard** for blockchain timestamping.
- It defines a set of rules for **conveniently** creating **provable timestamps** and later **independently verifying** them.
- A **proof** made with OpenTimestamps consists in a list of **commitment operations** applied in sequence to the document, ending with one or more **time attestations**.
- It proposes a **solution** to the **scalability problem**.

# OPENTIMESTAMPS AS A SCALABILITY SOLUTION

- OpenTimestamps allows to aggregate up to an infinite number of document in a single transaction.
- The trick is to compose documents in a merkle tree and to timestamp only the root of that tree.

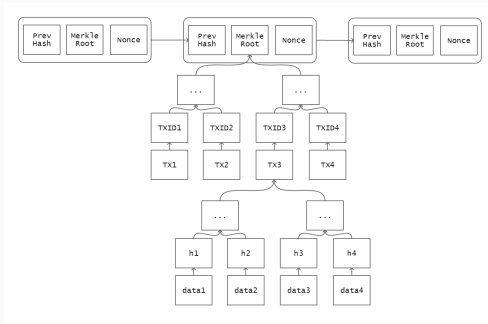


Figure 3: OTS scalability solution.

# OPENTIMESTAMPS AS A SCALABILITY SOLUTION

- Moreover, **aggregation servers** collect hash values and periodically **aggregate** them in a single merkle tree.
- Aggregation servers are **efficient** but **not convenient**: to obtain a proof you must wait for the transaction to be included in a block.
- To provide proofs almost **instantly**, OpenTimestamps makes use of **public calendar servers**.
- Aggregation servers submit the merkleroot to a calendar server, which **promise** that every submitted digest will be timestamped with Bitcoin.
- Proofs made this way are **incomplete**, but they can be **upgraded** once the blockchain has completed the timestamp, **adding** the path up the the **block header**.

## A PRACTICAL USE CASE: AIM OF THE PROJECT

- The author, in partnership with **DGI** (Digital Gold Institute) and **ANIA** (Associazione Nazionale fra le Imprese Assicuratrici) worked on a **practical use case** of blockchain timestamping.
- The aim of the project is to provide a fully operating **timestamping service**.
- It is build upon OpenTimestamps, while marginally **improving** and **extending** such protocol with additional features.
- For example, any insurance company could make use of it to grant **authenticity** of its associates policies.
- Simply, it could be useful whenever a **digital signature** is involved.



# ARCHITECTURE OF THE SOLUTION

- The **architecture** of the solution consists of **three cloud-servers**, accessible from outside via traditional network elements (DNS server, reverse proxy, firewall, etc.).
- Such elements transparently remap host names under the domain “aniasafe.it”.
- The three servers are divided in a **front-end** server and two **back-end** servers:
  - **Front-end**: hosts the public **web interface** (<https://timestamp.aniasafe.it>).
  - **Back-end**: they both run a **calendar server** connected to a local Bitcoin node: (<https://calendar.aniasafe.it>) for mainnet, (<https://test-calendar.aniasafe.it>) for testnet.

# ARCHITECTURE OF THE SOLUTION

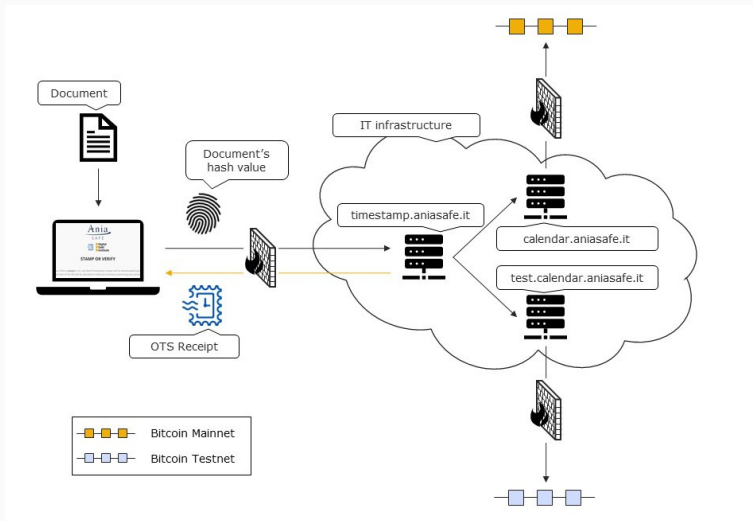


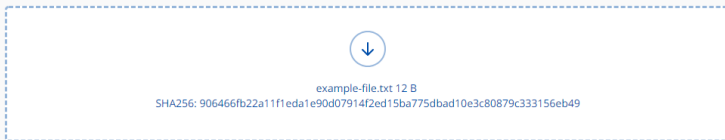
Figure 4: Architecture of the solution

- The front-end server hosts the **web interface**.
- The underlying **JavaScript library** defines the **main functions** invoked when a user interacts with the web interface: “**stamp**”, “**verifyTimestamp**” and “**upgradeTimestamp**”.
- The **JS library** has been extended for **additional features**:
  - **Testnet** as an alternative chain.
  - **Multi-validation**: the original protocol supports timestamping on multiple chains, but the verification of one attestation at a time.
  - **Esplora** (block-explorer): a new class has been created to query Esplora’s API correctly.

## STAMP OR VERIFY

Drop below a file to **stamp** it: its `.ots` OpenTimestamps receipt will be downloaded automatically.  
The hash of the file will be calculated inside your browser preserving your privacy.

Alternatively, drop below an `.ots` receipt/proof file to **verify** it.



**SUCCESS!**

OpenTimestamps receipt created and download started


Figure 5: Stamping

# WEB INTERFACE: VERIFYING



## STAMP OR VERIFY

Drop below a file to **stamp** it: its .ots OpenTimestamps receipt will be downloaded automatically.  
The hash of the file will be calculated inside your browser preserving your privacy.


Alternatively, drop below an .ots receipt/proof file to **verify** it.




example-file.txt.ots 4.4 kB  
Stamped SHA256 hash: 906466fb22a11f1eda1e90d07914f2ed15ba775dbad10e3c80879c333156eb49

 **VERIFIED!** 

BitcoinTestnet block 1486048 attests existence as of 2019-03-26 CET  
Bitcoin block 568877 attests existence as of 2019-03-26 CET



example-file.txt 12 B  
SHA256: 906466fb22a11f1eda1e90d07914f2ed15ba775dbad10e3c80879c333156eb49

 **SUCCESS!**

The provided file is the stamped one: its hash value matches the stamped one

Figure 6: Multi-validation

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## CONCLUSIONS AND FUTURE WORK

- The **notarization** process can be strengthened by means of a **decentralized solution** built upon a **blockchain technology**.
- The **Bitcoin blockchain** can be used as a **notary** for timestamping.
- **Timestamping** provides only **proof-of-existence** at a given date; it does not convey authorship, non-repudiation, veracity.
- OpenTimestamps provides a **standardization** for timestamping and it is **scalable**.
- **Elliptic curve commitments** could improve the timestamping process, allowing to push data into the payee public key or in a digital signature, **avoiding fees**.

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