

# An Introduction to Quantum Computing

## Lecture 07 *The Deutsch-Jozsa Algorithm*

Paolo Zuliani



SAPIENZA  
UNIVERSITÀ DI ROMA

[zuliani@di.uniroma1.it](mailto:zuliani@di.uniroma1.it)

# Outline

- The Deutsch-Jozsa Problem
- The quantum algorithm

# The Deutsch-Jozsa Problem

- For  $n \in \mathbb{N}$ , consider  $f: \mathcal{B}^n \rightarrow \mathcal{B}$

$$\mathcal{B} = \{0,1\}$$

- $f$  is **constant**:  $\forall x \in \mathcal{B}^n \quad f(x) = 0$  (or 1)

- $f$  is **balanced**:  $\sum_{x \in \mathcal{B}^n} f(x) = \frac{2^n}{2} = 2^{n-1}$

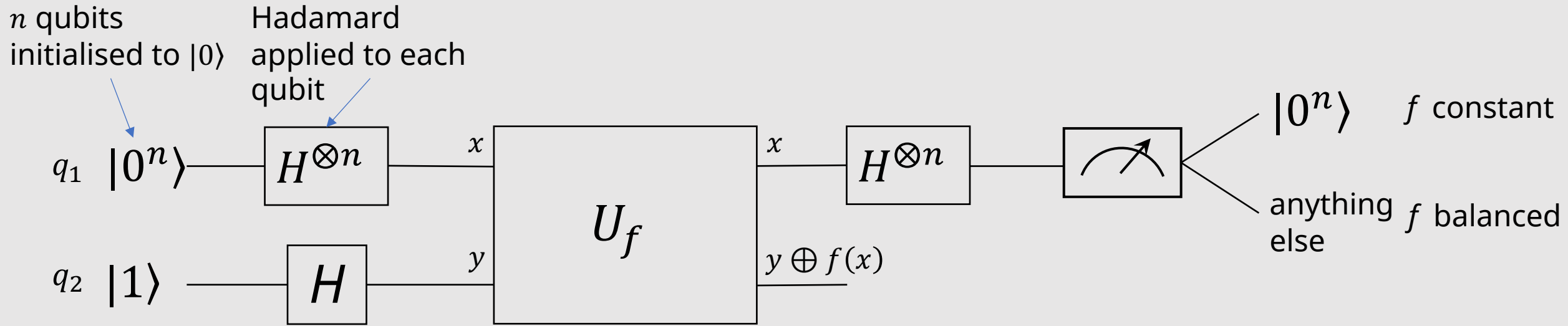
- A constant function cannot be balanced, and a balanced function cannot be constant
- Most Boolean functions are neither constant nor balanced



# The Deutsch-Jozsa Problem

- $f: \mathcal{B}^n \rightarrow \mathcal{B}$  is either **constant** or **balanced** (we don't know which one)
- **D-J Problem:** decide whether  $f$  is constant or  $f$  is balanced
- Complexity (*classical*): in the worst case, we need  $(2^{n-1}+1)$  evaluations of  $f$  to make a correct decision
- Complexity (*quantum*): **one** evaluation of  $f$  suffices!!

# The Deutsch-Jozsa Algorithm



Using a “programming” notation:

```

 $q_1, q_2 = |0^n 1\rangle;$ 
 $q_1, q_2 = H^{\otimes n} \otimes H(q_1, q_2);$ 
 $q_1, q_2 = U_f(q_1, q_2);$ 
 $q_1 = H^{\otimes n}(q_1);$ 
 $b = \text{Measure}(q_1)$ 
    
```

//  $q_1$  is a quantum register of  $n$  qubits!

# Preliminary: The Hadamard Transform

$$\begin{aligned} H|0\rangle &= \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \\ H|1\rangle &= \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) \end{aligned} \quad \longrightarrow \quad H|a\rangle = \frac{1}{\sqrt{2}} \sum_{b \in \mathcal{B}} (-1)^{a \cdot b} |b\rangle$$

Now two qubits:

$$\begin{aligned} H \otimes H |x\rangle \otimes |y\rangle &= H \otimes H |xy\rangle \\ &= \left( \frac{1}{\sqrt{2}} \sum_{a \in \mathcal{B}} (-1)^{x \cdot a} |a\rangle \right) \otimes \left( \frac{1}{\sqrt{2}} \sum_{b \in \mathcal{B}} (-1)^{y \cdot b} |b\rangle \right) \\ &= \frac{1}{2} \sum_{a, b \in \mathcal{B}} (-1)^{x \cdot a + y \cdot b} |a\rangle |b\rangle = \frac{1}{2} \sum_{a \in \mathcal{B}^2} (-1)^{xy \cdot a} |a\rangle \end{aligned}$$

# Preliminary: The Hadamard Transform

So, for two qubits:

$$H \otimes H |xy\rangle = \frac{1}{2} \sum_{a \in \mathcal{B}^2} (-1)^{xy \cdot a} |a\rangle$$

In general, for  $x \in \mathcal{B}^n$  (i.e., a string of  $n$  bits)

$$H^{\otimes n} |x\rangle = \frac{1}{\sqrt{2^n}} \sum_{a \in \mathcal{B}^n} (-1)^{x \cdot a} |a\rangle \quad \text{where } x \cdot a = \left( \sum_{i=0}^{n-1} x_i \cdot a_i \right) \bmod 2$$

# The Deutsch-Jozsa Algorithm

$$\begin{aligned} q_1, q_2 &= |0^n 1\rangle; \\ q_1, q_2 &= H^{\otimes n} \otimes H(q_1, q_2); \\ q_1, q_2 &= U_f(q_1, q_2); \\ q_1 &= H^{\otimes n}(q_1); \\ b &= \text{Measure}(q_1) \end{aligned}$$

$$|0^n 1\rangle$$

$$\text{Apply } q_1, q_2 = H^{\otimes n} \otimes H(q_1, q_2)$$

$$= \sum_{x \in \mathcal{B}^n} \frac{|x\rangle}{\sqrt{2^n}} \otimes \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$$

$$\text{Apply } q_1, q_2 = U_f(q_1, q_2)$$

$$= \sum_{x \in \mathcal{B}^n} \frac{(-1)^{f(x)} |x\rangle}{\sqrt{2^n}} \otimes \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$$

Recall that, for  $a \in \mathcal{B}^n$

$$U_f(|a\rangle \frac{(|0\rangle - |1\rangle)}{\sqrt{2}}) = (-1)^{f(a)} |a\rangle \frac{(|0\rangle - |1\rangle)}{\sqrt{2}}$$





# The Deutsch-Jozsa Algorithm

~~$q_1, q_2 = |0^n 1\rangle;$~~   
 ~~$q_1, q_2 = H^{\otimes n} \otimes H(q_1, q_2);$~~   
 ~~$q_1, q_2 = U_f(q_1, q_2);$~~   
 $q_1 = H^{\otimes n}(q_1);$   
 $b = \text{Measure}(q_1)$

$$= \sum_{x \in \mathcal{B}^n} \frac{(-1)^{f(x)} |x\rangle}{\sqrt{2^n}} \otimes \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$$

Apply  $q_1 = H^{\otimes n}(q_1)$

$$= \sum_{x \in \mathcal{B}^n} \sum_{a \in \mathcal{B}^n} \frac{(-1)^{x \cdot a + f(x)} |a\rangle}{2^n} \otimes \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$$

$$H^{\otimes n} |x\rangle = \frac{1}{\sqrt{2^n}} \sum_{a \in \mathcal{B}^n} (-1)^{x \cdot a} |a\rangle$$

# The Deutsch-Jozsa Algorithm

~~$q_1, q_2 = |0^n 1\rangle;$~~   
 ~~$q_1, q_2 = H^{\otimes n} \otimes H(q_1, q_2);$~~   
 ~~$q_1, q_2 = U_f(q_1, q_2);$~~   
 ~~$q_1 = H^{\otimes n}(q_1);$~~   
 ~~$b = \text{Measure}(q_1)$~~

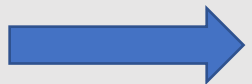
$$= \sum_{x \in \mathcal{B}^n} \sum_{a \in \mathcal{B}^n} \frac{(-1)^{x \cdot a + f(x)}}{2^n} |a\rangle \otimes \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$$

$$= \sum_{a \in \mathcal{B}^n} \sum_{x \in \mathcal{B}^n} \frac{(-1)^{x \cdot a + f(x)}}{2^n} |a\rangle \otimes \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$$

$$= \sum_{a \in \mathcal{B}^n} |a\rangle \sum_{x \in \mathcal{B}^n} \frac{(-1)^{x \cdot a + f(x)}}{2^n} \otimes \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$$

Apply  $b = \text{Measure}(q_1)$

The amplitude for  $|0^n\rangle$  is computed from the sum above by setting  $a = 0^n$



$$\frac{1}{2^n} \sum_{x \in \mathcal{B}^n} (-1)^{x \cdot a + f(x)} = \frac{1}{2^n} \sum_{x \in \mathcal{B}^n} (-1)^{f(x)}$$



# The Deutsch-Jozsa Algorithm

$$\begin{aligned} q_{\pm}, q_z &= |0^n 1\rangle; \\ q_{\pm}, q_z &= H^{\otimes n} \otimes H(q_{\pm}, q_z); \\ q_{\pm}, q_z &= U_f(q_{\pm}, q_z); \\ q_{\pm} &= H^{\otimes n}(q_{\pm}); \\ b &= \text{Measure}(q_{\pm}) \end{aligned}$$

The amplitude for  $|0^n\rangle$  is thus

$$\frac{1}{2^n} \sum_{x \in \mathcal{B}^n} (-1)^{f(x)} = \begin{cases} \pm 1 & \text{if } f \text{ is constant} \\ 0 & \text{if } f \text{ is balanced} \end{cases}$$

If we measure  $0^n$  we know precisely that  $f$  is constant.  
If we measure any other state, then  $f$  must be balanced.

Only one evaluation of  $f$  is needed!

