

OPTIMAL AREA LAYOUT OF THE BUTTERFLY NETWORK



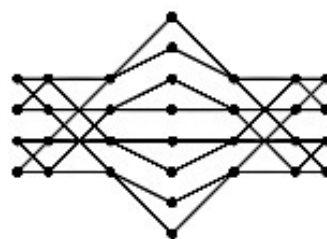
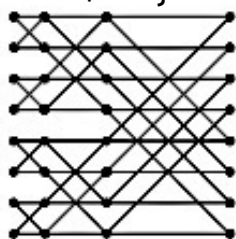
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OPTIMAL AREA LAYOUT - IDEA (1)

The two papers that provide an optimal area layout base their results on the following lemma:

Lemma: For any non-negative integers j, k , $0 \leq j \leq j+k \leq n$, the subgraph of the n -dim. Butterfly induced by the nodes of levels $j, j+1, \dots, j+k$ is the disjoint union of 2^{n-k} copies of k -dimensional butterflies.

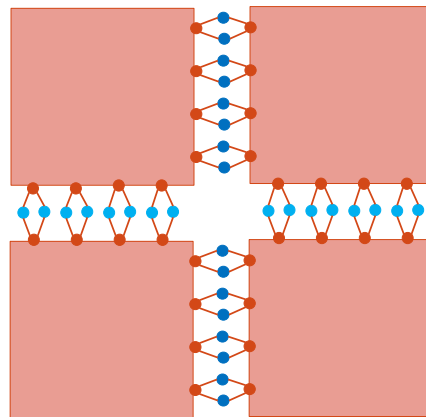
In particular, if $j=0$ and $k=n-1$:



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OPTIMAL AREA LAYOUT - IDEA (2)

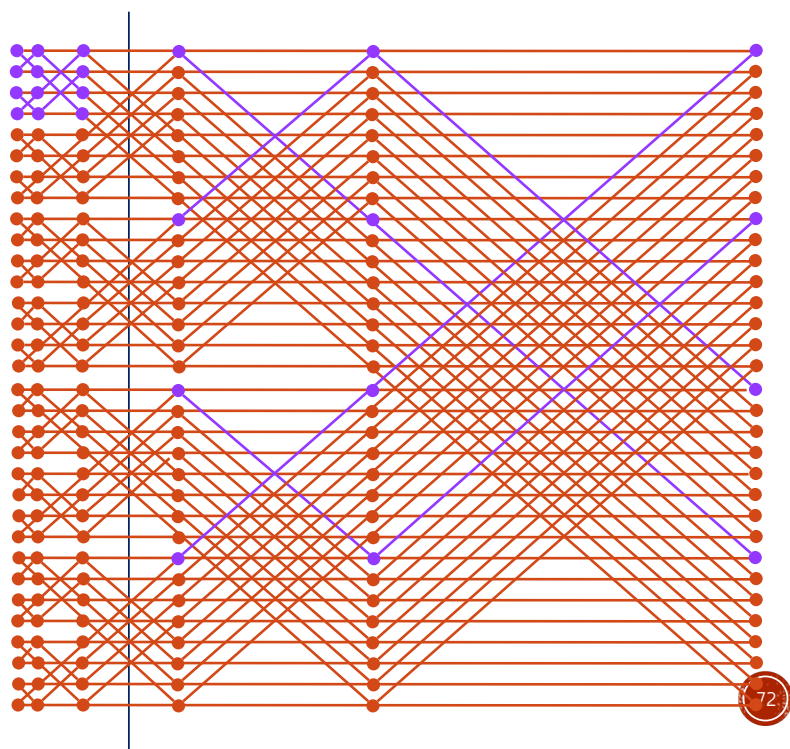
Hence, an $(n-1)$ -dimensional butterfly can be built as a pair of $(n-2)$ -dim. butterflies connected by one node layer and one edge layer. If we cut out the input and output nodes from an n -dim. Butterfly, we get:



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OPTIMAL AREA LAYOUT - IDEA (3)

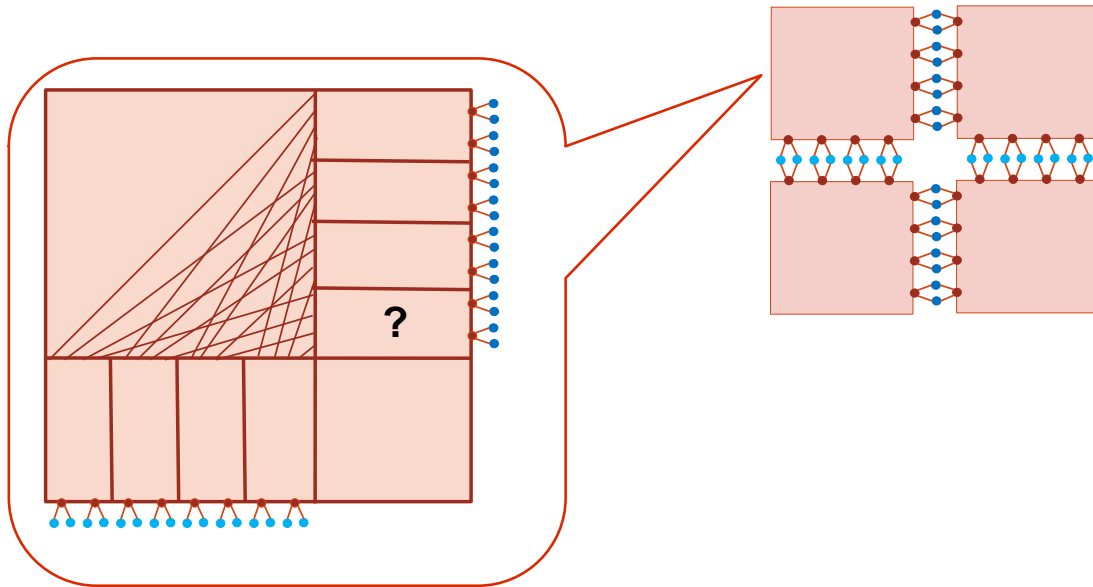
Each one of these $(n-2)$ -dim. Butterflies can be, in turn, cut into many smaller butterflies:



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OPTIMAL AREA LAYOUT - IDEA (4)

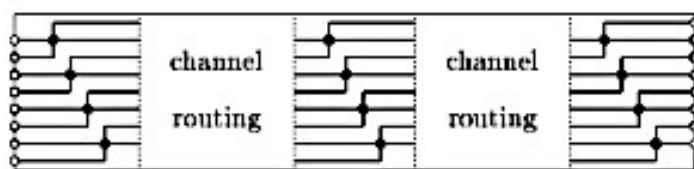
The previous layout can be better specified as follows:



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OPTIMAL AREA LAYOUT - IDEA (5)

Each rectangle contains a Butterfly that can be represented, either horizontally or vertically, layer by layer as follows:

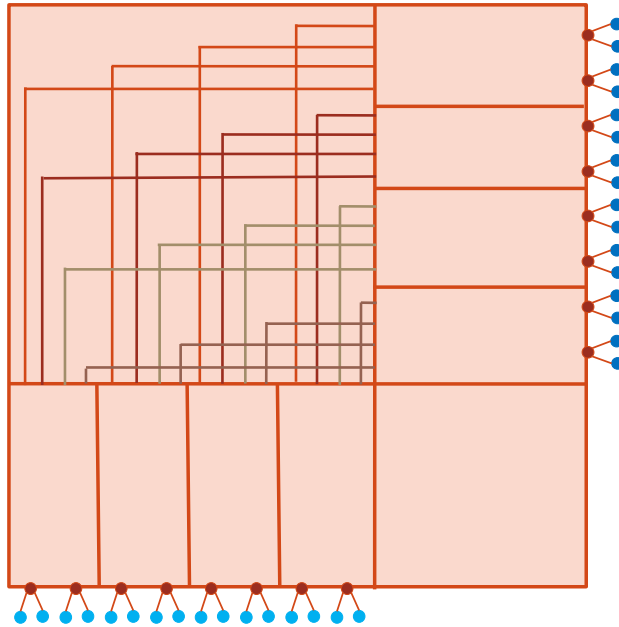


Obs.: this layout is far from being optimal; nevertheless it allows to produce a final optimal layout.

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OPTIMAL AREA LAYOUT - IDEA (6)

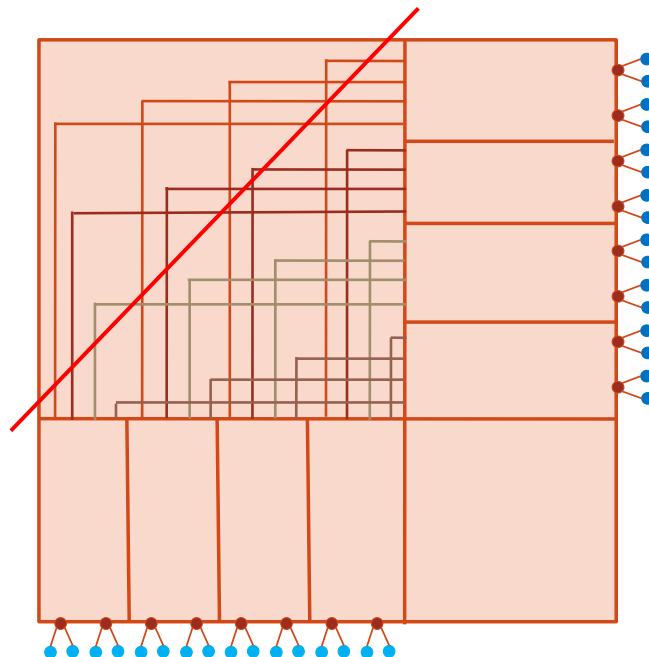
It remains to connect the small rectangular butterflies:



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OPTIMAL AREA LAYOUT - IDEA (7)

In the case of slanted layout, it can be bent along the line:



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OPTIMAL AREA LAYOUT - IDEA (8)

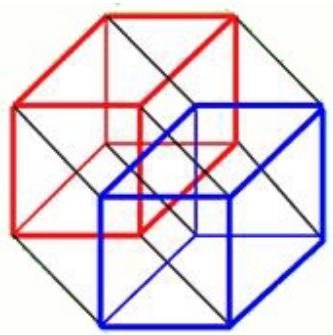
It is possible to prove tight lower and upper bounds on the layout area for both the models (usual and slanted).

The interested students can look at:

- A. Avior, T.C., S. Even, A. Litman, A.L. Rosenberg: A Tight Layout of the Butterfly Network. Theory of Computing Systems 31, 1998.
- Y. Dinitz, S. Even, M. Zapolotsky: A Compact Layout of the Butterfly. J. of Interconnection Networks 4, 2003.

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LAYOUT OF THE HYPERCUBE NETWORK



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