

An Introduction to Quantum Computing

Lecture 09:

The Quantum Fourier Transform and Phase Estimation - Towards Shor's Algorithm (I)

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Agenda

- Discrete Fourier Transform
- Quantum Fourier Transform
- Quantum Algorithm for Phase Estimation

The Discrete Fourier Transform

It maps N complex numbers x_0, \dots, x_{N-1} to N complex numbers y_0, \dots, y_{N-1} :

$$y_k = \frac{1}{\sqrt{N}} \sum_{j=0}^{N-1} x_j e^{2\pi i j k / N}$$

where i is the imaginary unit.

- The Fourier Transform is much used for signal processing (e.g., speech recognition, audio compression)
- It is sometimes easier to study a signal in a different domain (e.g., frequency instead of time)
- The Discrete Fourier Transform¹ “filters” the input sequence through a sinusoidal wave of frequency k/N .

¹For a deeper and clear treatment of the DFT see Chapter 7 of “The Design and Analysis of Computer Algorithms” by Aho, Hopcroft, and Ullman.

The Quantum Fourier Transform

Definition

The QFT maps each basis state $|0\rangle, |1\rangle, \dots, |N-1\rangle$ as follows

$$|j\rangle \xrightarrow{QFT} \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} e^{2\pi i j k / N} |k\rangle$$

Equivalently, for a generic vector:

$$\sum_{j=0}^{N-1} x_j |j\rangle \rightarrow \sum_{k=0}^{N-1} y_k |k\rangle$$

where $y_k = \frac{1}{\sqrt{N}} \sum_{j=0}^{N-1} x_j e^{2\pi i j k / N}$.

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so it can be written as

$$QFT = \sum_{j=0}^{N-1} \left(\frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} e^{2\pi i j k / N} |k\rangle \right) \langle j|$$

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Proposition

The QFT is a unitary operator, i.e., $QFT QFT^\dagger = QFT^\dagger QFT = I$.

$$QFT^\dagger = \sum_{j=0}^{N-1} |j\rangle \left(\frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} e^{-2\pi i j k / N} \langle k| \right)$$

The Quantum Fourier Transform: Unitarity

Let us show that the QFT is unitary:

$$QFT^\dagger QFT = \sum_{j=0}^{N-1} |j\rangle \left(\frac{1}{\sqrt{N}} \sum_{r=0}^{N-1} e^{-2\pi i jr/N} \langle r| \right) \sum_{k=0}^{N-1} \left(\frac{1}{\sqrt{N}} \sum_{s=0}^{N-1} e^{2\pi i sk/N} |s\rangle \right) \langle k|$$

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$$= \frac{1}{N} \sum_{j,k=0}^{N-1} |j\rangle \left(\sum_{r=0}^{N-1} e^{2\pi i r(k-j)/N} \right) \langle k|$$

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Exercise: prove $QFT QFT^\dagger = I$.

The Quantum Fourier Transform: Quantum Circuit

An **equivalent** QFT definition (assuming $N = 2^n$, hence n qubits):

$$|j_1 \dots j_n\rangle \xrightarrow{QFT} \frac{(|0\rangle + e^{2\pi i 0.j_n} |1\rangle) \otimes (|0\rangle + e^{2\pi i 0.j_{n-1}j_n} |1\rangle) \otimes \dots \otimes (|0\rangle + e^{2\pi i 0.j_1 \dots j_n} |1\rangle)}{2^{n/2}}$$

where j_1, \dots, j_n are bits, and the **binary fraction**

$$0.j_l j_{l+1} j_m = \frac{j_l}{2} + \frac{j_{l+1}}{4} + \dots + \frac{j_m}{2^{m-l+1}}$$

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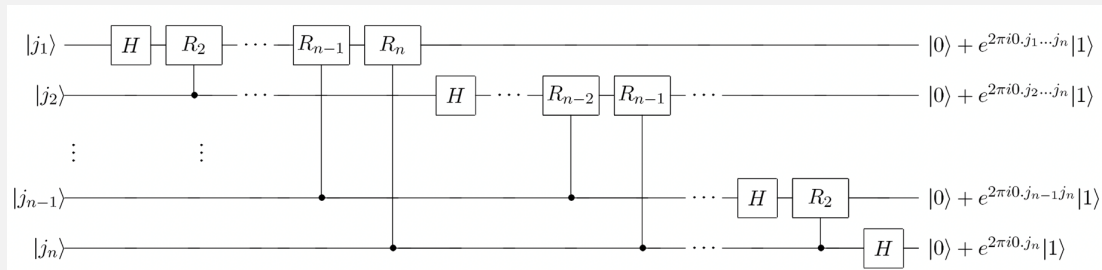
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Example:

$$0.1101 = \frac{1}{2} + \frac{1}{4} + \frac{1}{2^4}$$

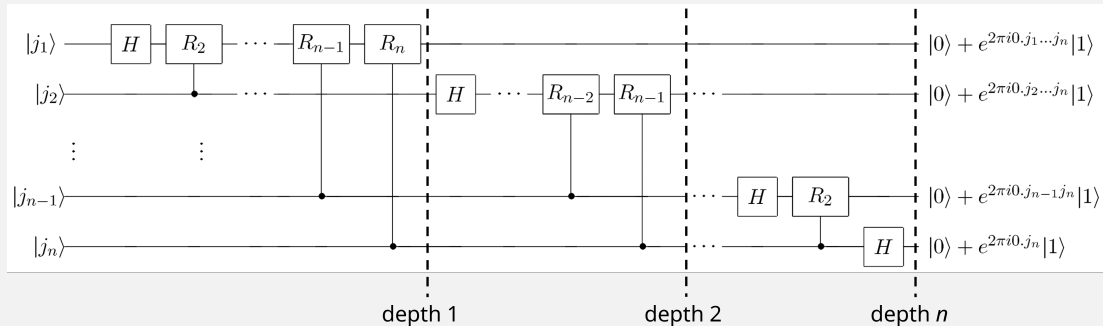
The Quantum Fourier Transform: Quantum Circuit



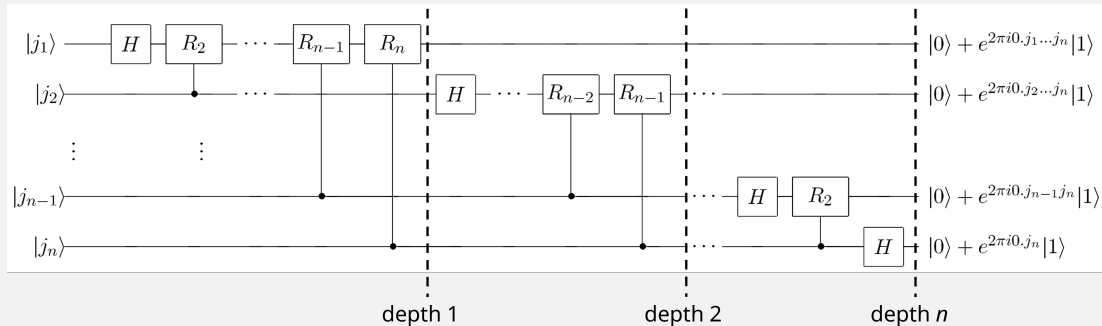
where H is the usual Hadamard and the controlled- R_k gates are defined on

$$R_k = \begin{pmatrix} 1 & 0 \\ 0 & e^{2\pi i / 2^k} \end{pmatrix}$$

The Quantum Fourier Transform: Quantum Circuit

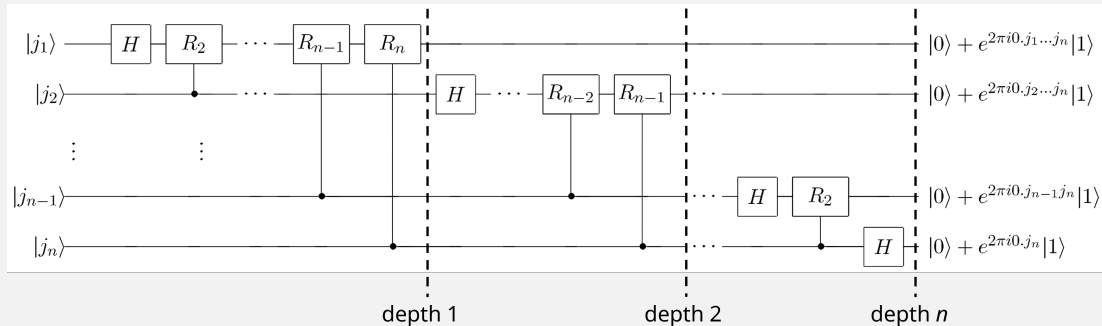


The Quantum Fourier Transform: Quantum Circuit



State at depth 1: $\frac{1}{2^{1/2}}(|0\rangle + e^{2\pi i 0 \cdot j_1 \dots j_n} |1\rangle) |j_2 \dots j_n\rangle$

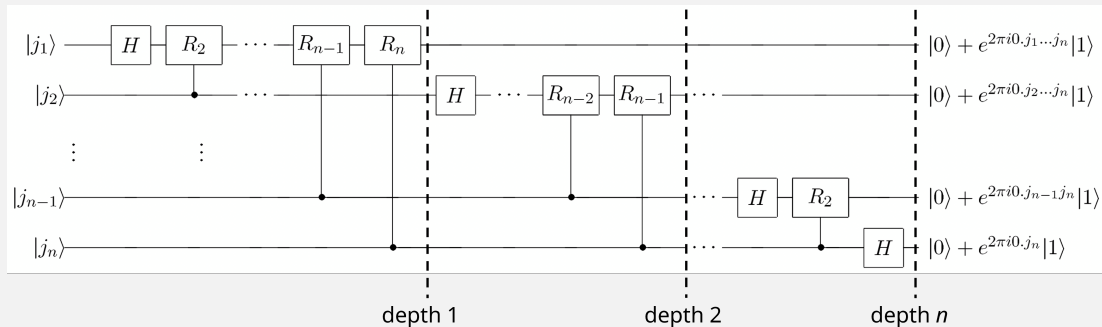
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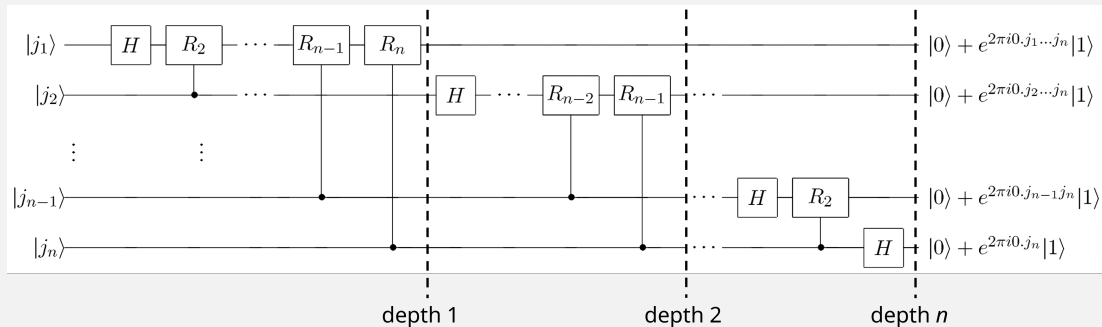
State at depth 2: $\frac{1}{2^{2/2}}(|0\rangle + e^{2\pi i 0 \cdot j_1 \dots j_n} |1\rangle) \otimes (|0\rangle + e^{2\pi i 0 \cdot j_2 \dots j_n} |1\rangle) |j_3 \dots j_n\rangle$

The Quantum Fourier Transform: Quantum Circuit



The state at depth n has the correct terms, but in the wrong order! (Remember the tensor product is NOT commutative.)

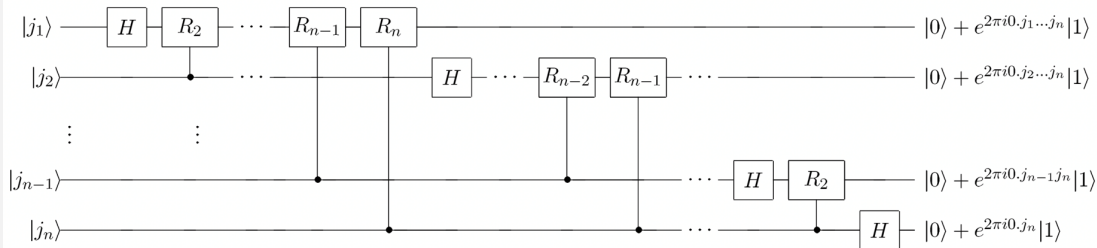
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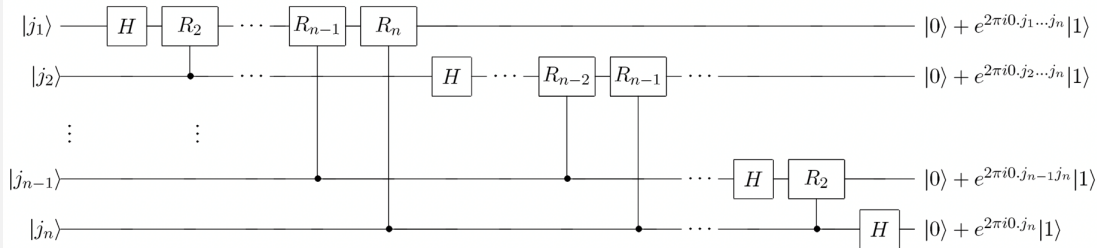
We need to **swap** the qubits, which can be done unitarily, of course.

The Quantum Fourier Transform: Complexity



- Quantum circuit has $O(n^2)$ gates
- Best classical circuit needs $O(n2^n)$ gates
- Looks great! Is it?

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$$\sum_{j=0}^{N-1} x_j |j\rangle \xrightarrow{QFT} \sum_{k=0}^{N-1} y_k |k\rangle \quad \left(\text{where } y_k = \frac{1}{\sqrt{N}} \sum_{j=0}^{N-1} x_j e^{2\pi i j k / N} \right)$$

- We want the DFT coefficients y_k 's, but they are encoded in the amplitudes!

Phase Estimation

Let's see an application of the QFT.

A previous exercise: the eigenvalues of a unitary operator are complex numbers of **modulus 1**.

This means that any eigenvalue of a unitary operator can be written as $e^{2\pi i\varphi}$ for some real $\varphi \in [0, 1]$.

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Definition (Phase Estimation Problem)

Let $\lambda = e^{2\pi i\varphi}$ be an eigenvalue of a unitary operator U . Find φ .

This problem can be solved quite easily with the QFT.

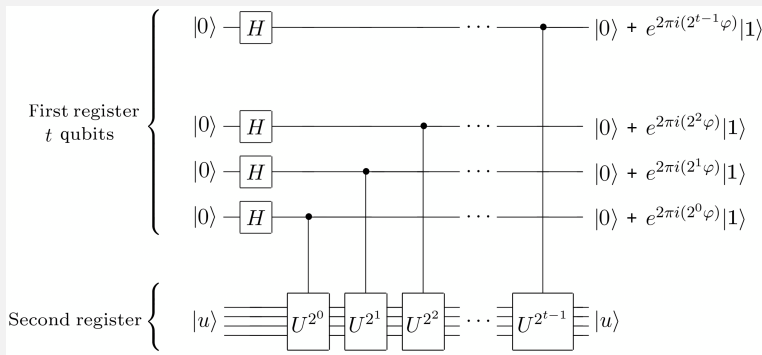
Quantum Phase Estimation Algorithm

Let u be an eigenvector associated to the unknown eigenvalue $e^{2\pi i\varphi}$ of a unitary operator U , i.e., $U|u\rangle = e^{2\pi i\varphi}|u\rangle$.

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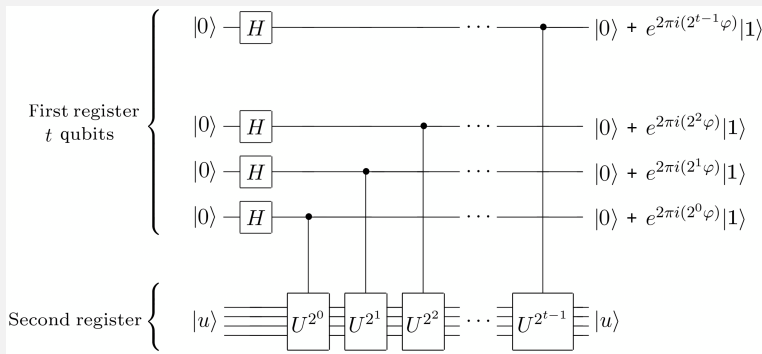
Consider the circuit below for some natural $t > 0$:



Quantum Phase Estimation Algorithm

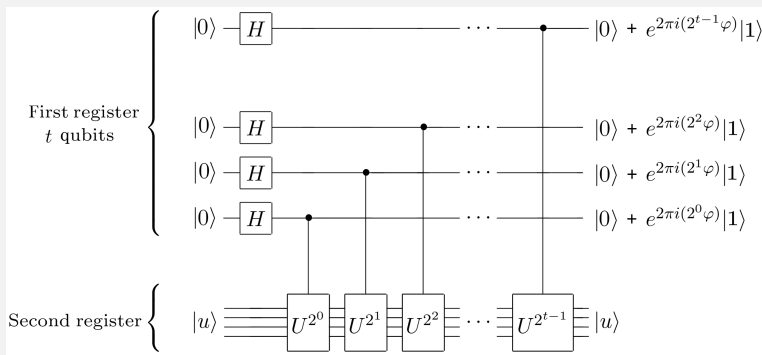
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A control- U^{2^k} gate conditionally applies $U^{2^k} = \underbrace{U \cdots U}_{2^k \text{ times}}$ to the second qubit register.

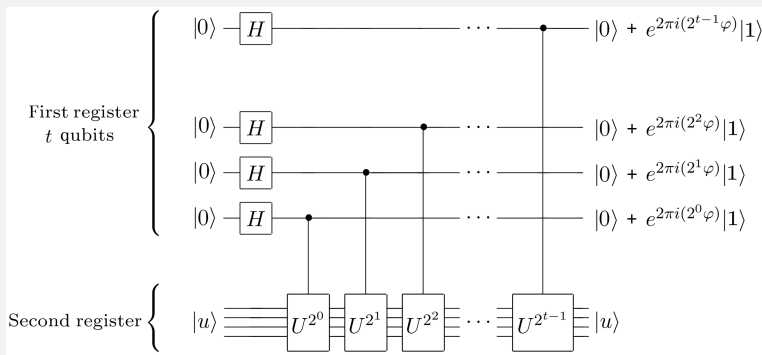
Quantum Phase Estimation Algorithm



The state of the t qubits at the end of the QPE circuit is:

$$\frac{1}{2^{t/2}} (|0\rangle + e^{2\pi i 2^{t-1}\varphi} |1\rangle) \otimes (|0\rangle + e^{2\pi i 2^{t-2}\varphi} |1\rangle) \otimes \dots \otimes (|0\rangle + e^{2\pi i 2^0\varphi} |1\rangle)$$

Quantum Phase Estimation Algorithm



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Quantum Phase Estimation Algorithm

Suppose now that φ can be written exactly with t bits:

$$\varphi = 0.\varphi_1 \dots \varphi_t$$

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which is *precisely* the final state of the QFT circuit (after the swap)!

Quantum Phase Estimation Algorithm

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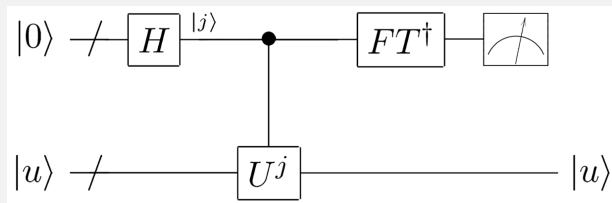
$$\frac{1}{2^{t/2}}(|0\rangle + e^{2\pi i 0.\varphi_t} |1\rangle) \otimes (|0\rangle + e^{2\pi i 0.\varphi_{t-1}\varphi_t} |1\rangle) \otimes \dots \otimes (|0\rangle + e^{2\pi i 0.\varphi_1\varphi_2\dots\varphi_t} |1\rangle)$$

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Therefore, we apply the inverse QFT circuit at the end of the QPE circuit and then measure to obtain the sought phase $|\varphi_1 \dots \varphi_t\rangle$ with probability 1!

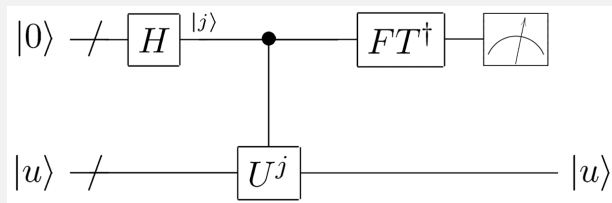
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The final quantum circuit for solving phase estimation is thus:



Quantum Phase Estimation Algorithm

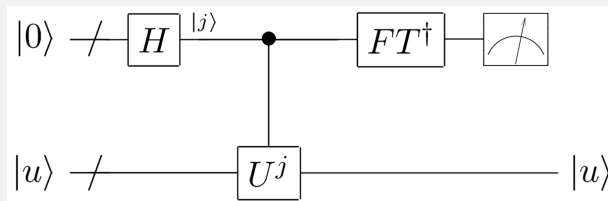
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Quantum Phase Estimation Algorithm

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Proposition

To estimate φ with n bits of precision and success probability at least $1 - \epsilon$, it is sufficient to use the QPE circuit with r qubits

$$r = n + \left\lceil \log \left(2 + \frac{1}{2\epsilon} \right) \right\rceil$$