## Performance Modeling of Computer Systems and Networks

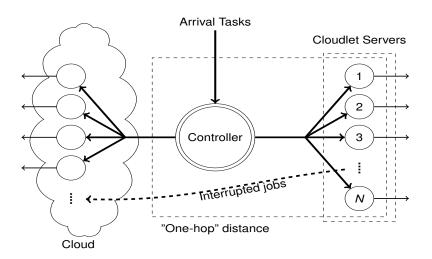
**Project 2018-2019** 

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### System description





## Specification model

#### • What are the state variables?

$$n_{x}^{(c)}( au)=$$
 Number of class  $c$  jobs currently running at  $x$  system's node at time  $au$ 

$$g_x^{(c)}(\tau) = \text{Number of class } c \text{ departed jobs from node } x \text{ at time } \tau$$

$$s_{x,i}^{(c)}$$
 = Service time of class c job i served on x node

$$i_{cloudlet}^{(2)}(\tau) =$$
 Number of class 2 interrupted jobs at time  $\tau$  which were running on cloudlet (Access control algorithm 2 only!)

Where:

$$c \in \{1,2\} = C$$
  
  $x \in \{cloudlet, cloud, global\} = X$   
  $\tau \in (t_0, t)$ 

# Specification model

 Are there any constrains regarding the values assumed by these variables?

$$\sum_{c \in C} n_{cloudlet}^{(c)}(\tau) \le N \quad \forall \tau \in (t_0, t)$$
(1)

If you are using Access Control Algorithm 2:

$$n_{cloudlet}^{(2)}( au) > 0 \quad \Rightarrow \quad \sum_{c \in C} n_{cloudlet}^{(c)}( au) \leq S \quad \forall au \in (t_0, t)$$
 $S + 1 \leq n_{cloudlet}^{(1)}( au) \leq N \quad \Leftrightarrow \quad n_{cloudlet}^{(2)}( au) = 0 \qquad \forall au \in (t_0, t)$ 

$$(2)$$

# Specification Model System State

- What's meant by system state?
- So, how we can define system state formally?

$$\omega(\tau) = (\omega_{cloudlet}(\tau), \omega_{cloud}(\tau)) \tag{3}$$

Where:

$$\omega_{cloudlet}(\tau) = (n_{cloudlet}^{(1)}(\tau), n_{cloudlet}^{(2)}(\tau)) 
\omega_{cloud}(\tau) = (n_{cloud}^{(1)}(\tau), n_{cloud}^{(2)}(\tau))$$
(4)

Thus:

$$\omega(\tau) = ((n_{cloudlet}^{(1)}(\tau), n_{cloudlet}^{(2)}(\tau)), (n_{cloud}^{(1)}(\tau), n_{cloud}^{(2)}(\tau))$$

$$\tag{5}$$

• How do they evolve in time?



## Specification Model Events

- What is an event?
- What is the first event that can occur? And the last one?

$$n_{x}^{(c)}(t_{0}) = 0 \qquad \forall c \in C, \forall x \in X$$

$$d_{x}^{(c)}(t_{0}) = 0 \qquad \forall c \in C, \forall x \in X$$

$$n_{x}^{(c)}(t) = 0 \qquad \forall c \in C, \forall x \in X$$

$$t_{0} = 0, t = \tau^{*}$$

$$(6)$$

So, what are the events capable to change system status?

# Specification Model Events

Event name that occurred at time $\tau$	Event's place	Event's effects
Class 1 job arrival	Cloudlet	$n_{\text{cloudlet}}^{(1)}(\tau)$ ++
	Cloud	$n_{\text{cloud}}^{(1)}(\tau)$ ++ $n_{\text{global}}^{(1)}(\tau)$ ++
	Controller	$n_{\text{global}}^{(1)}(\tau)$ ++
Class 2 job arrival	Cloudlet	$n_{\text{cloudlet}}^{(2)}(\tau)$ ++ $n_{\text{cloud}}^{(2)}(\tau)$ ++
	Cloud	$n_{\rm cloud}^{(2)}(\tau)$ ++
	Controller	l 'I'global (' ) TT
Class 1 job departure	Q1	$n_{\text{cloudlet}}^{(1)}(\tau)$ $n_{\text{global}}^{(1)}(\tau)$
	Cloudlet	$d_{\text{cloudlet}}^{(1)}(\tau)$ ++
		$d_{\text{global}}^{(1)}(\tau)$ ++
	Cloud	$n_{\text{cloud}}^{(1)}(\tau)$ $n_{\text{global}}^{(1)}(\tau)$
		$d_{\text{cloud}}^{(1)}(\tau)$ ++ $d_{\text{global}}^{(1)}(\tau)$ ++
Class 2 job departure	Cloudlet	$n_{\text{cloudlet}}^{(2)}(\tau)$ $n_{\text{global}}^{(2)}(\tau)$
		$d_{ ext{cloudlet}}^{(2)}( au)$ ++ $d_{ ext{global}}^{(2)}( au)$ ++
	Cloud	$n_{\text{cloud}}^{(2)}(\tau)$ $n_{\text{global}}^{(2)}(\tau)$
		$d_{\text{cloud}}^{(2)}(\tau)$ ++ $d_{\text{global}}^{(2)}(\tau)$ ++

#### Specification Model Events

- When you are using Access Control Algorithm 2:
  - If class 1 job arrival event on controller occurs:

$$n_{\text{global}}^{(1)}(\tau)$$
++ (7)

If following conditions are all true:

$$n_{cloudlet}^{(1)}( au) \neq N$$

$$n_{cloudlet}^{(1)}(\tau) + n_{cloudlet}^{(2)}(\tau) = S$$
 (8)

$$n_{cloudlet}^{(2)} > 0$$

#### Then:

- Class 1 job arrival on cloudlet node event is scheduled.
- Class 2 job arrival on cloud node event is scheduled.
- A Class 2 job departure on cloudlet node event is removed from the event list.

  •  $i_{cloudlet}^{(2)}(\tau) + +$



System description

• How to compute requested statistics?

$$E[N_x]^{(c)} = \frac{1}{t-t_0} \int_{t_0}^t n_x^{(c)}(\tau) d\tau \qquad \forall c \in C, \forall x \in X$$
 (9)

$$E[N_x] = \frac{1}{t - t_0} \cdot \int_{t_0}^t \left( n_x^{(1)}(\tau) + n_x^{(2)}(\tau) \right) d\tau$$

$$= \frac{1}{t - t_0} \cdot \int_{t_0}^t n_x^{(1)}(\tau) d\tau + \frac{1}{t - t_0} \cdot \int_{t_0}^t n_x^{(2)}(\tau) d\tau$$

$$= E[N_x]^{(1)} + E[N_x]^{(2)}$$

$$= \sum_{c \in C} E[N_x]^{(c)}$$

 $\forall x \in X$ 

## **Specification Model**

Statistics computation

$$E[T_x]^{(c)} = E[S_x]^{(c)} = \frac{1}{d_x^{(c)}(t)} \cdot \sum_{i=0}^{d_x^{(c)}(t)} s_{x,i}^{(c)} \quad \forall c \in C, \forall x \in X$$
(11)

$$E[T_x] = E[S_x] = \frac{1}{d_x^{(1)}(t) + d_x^{(2)}(t)} \cdot \left(\sum_{i=0}^{d_x^{(1)}(t)} s_{x,i}^{(1)} + \sum_{i=0}^{d_x^{(2)}(t)} s_{x,i}^{(2)}\right) \quad \forall x \in X$$
(12)

## Computational Model Next-event simulation logic

 Which are main features of a discrete-event simulation using a next-event approach?

- How we have implemented simulation clock and the event list?
- How we have implemented next-event simulation algorithm in our computational model?

Next-event simulation logic

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#### Listing 1: Computational Model.perform method

```
public void perform() {
   initializeSimulation();
   while (!this.simulationEventList.isEmpty()) {
      SimulationEvent actualEvent = this.simulationEventList.poll();
      if (actualEvent != null) {
          SimulationClock.getInstance().setCurrentEventTime(actualEvent.getStartTime());
          actualEvent.perform();
          actualEvent.scheduleFollowingEvent();
          SimulationEvent nextEvent = this.simulationEventList.peek();
          if (nextEvent != null)
              SimulationClock.getInstance().setNextEventTime(nextEvent.getStartTime());
      updateStatistics():
```

#### Listing 2: Computational Model.initialize Simulation method



#### Next-event simulation logic

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#### Listing 3: System class methods

```
Offverride
     protected void initializeSystemStateVariables() {
         this.cloud.initializeSystemStateVariables();
         this.cloudlet.initializeSystemStateVariables();
         this.globalNetwork.initializeSystemStateVariables();
     @Override
     protected void initializeSimulationClock() {
        SimulationClock.getInstance().setCurrentEventTime(0.0);
         SimulationClock.getInstance().setNextEventTime(0.0);
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     @Override
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     protected void scheduleInitialEvent() {
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         this.globalNetwork.scheduleInitialEvent();
```

#### Listing 4: GlobalNetwork.scheduleInitialEvent method

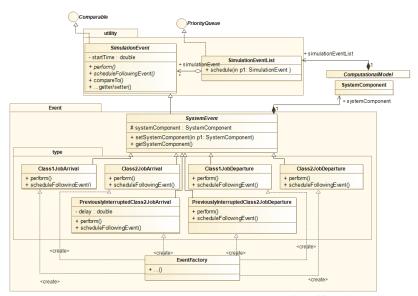
```
@Override
public void scheduleInitialEvent() {
    SystemEvent firstEventClass1 = SystemEventFactory.buildClass1JobArrival();
    SystemEvent firstEventClass2 = SystemEventFactory.buildClass2JobArrival();
    this.system.scheduleEventOnGlobalNetwork(firstEventClass1, this.
          getNextClass1JobInterArrivalTime());
    this.system.scheduleEventOnGlobalNetwork(firstEventClass2, this.
          getNextClass2JobInterArrivalTime());
```



# Computational Model Events

• How we have implemented events in our computational model?

## Computational Model Events





#### Computational Model **Events**

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#### Listing 5: SystemEvent class

```
public abstract class SystemEvent extends SimulationEvent {
         protected SystemComponent systemComponent;
         public void setSystemComponent(SystemComponent systemComponent) {
             this.systemComponent = systemComponent;
         public SystemComponent getSystemComponent() {
             return systemComponent:
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```

#### Listing 6: Class1JobArrival class

```
public class Class1JobArrival extends SystemEvent {
         Offverride
         public void perform() {
             this.systemComponent.updateStatusAfterClass1JobArrival():
         Offverride
         public void scheduleFollowingEvent() {
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             this.systemComponent.scheduleFollowingEventAfterClass1JobArrival();
```

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17 18 • PreviouslyInterruptedClass2JobArrival? What is it for?

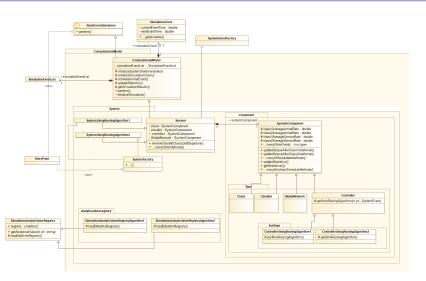
Listing 7: PreviouslyInterruptedClass2JobArrival class

```
public class PreviouslyInterruptedClass2JobArrival extends SystemEvent {
   private double delay;
   public PreviouslyInterruptedClass2JobArrival(double delay) {
        this.delay = delay;
   }
   @Override
   public void perform() {
        this.systemComponent.updateStatusAfterPreviouslyInterruptedClass2JobArrival(this.delay);
   }
   @Override
   public void scheduleFollowingEvent() {
        this.systemComponent.scheduleFollowingEventAfterPreviouslyInterruptedClass2JobArrival();
   }
}
```

System implementation

- How we have implemented our system?
- Which classes hold system state variables? Which compute output statistics?
- How can we start our simulator?

System implementation



Class 2 job interruption

• How we manage class 2 job interruption?

System description

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```
protected void performRoutingAlgorithm(SystemEvent event) {
    int n1 = this.system.getNumberOfClass1JobOnCloudlet();
    int n2 = this.system.getNumberOfClass2JobOnCloudlet();
        if (event instanceof Class1JobArrival) {
            if (n1 == this.system.getThreshold())
                this.system.scheduleEventOnCloud(event, 0);
            else if (n1 + n2 < this.system.getThreshold())
                this.system.scheduleEventOnCloudlet(event, 0);
            else if (n2 > 0) {
                double runningCloudletTimeOfInterruptedJob = this.system.
                      removeCloudletClass2JobDeparture();
                this.system.scheduleEventOnCloudlet(event, 0);
                double setupTime = RandomNumberGenerator.getInstance().getExponential(5, 0.8);
                this.system.scheduleEventOnCloud(SystemEventFactory.
                      buildPreviouslyInterruptedClass2JobArrival(setupTime +
                      runningCloudletTimeOfInterruptedJob), setupTime);
                this.numberOfInterruptedClass2Jobs++;
            } else
                this.system.scheduleEventOnCloudlet(event. 0):
       } else {
```

## Statistics Results A Finite-Horizon simulation

- What kind of simulation have we performed?
- What is generated at the end of each simulation? How we can provide an estimation for required metrics?
   Let i i-th replication, where i = 0, ..., n

$$\bar{x}_i(t) = \frac{1}{t} \int_0^t x_i(t) dt \qquad \forall i = 1, ..., n$$
 (13)

$$E[\bar{x}(t)] = \frac{1}{n} \sum_{i=1}^{n} \bar{x}_i(t)$$
 (14)

Obliviously, interval estimate for  $E[\bar{x}(t)]$  can be calculated from the ensemble mean and standard deviation.

## Statistics Results Il processo client

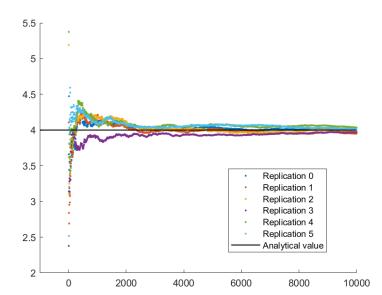
 $Avg = 6.627533 \pm 0.052390$  $I_{95\%} = [6.575143, 6.679922]$ = 6.622806V Analytical value Transient statistic value. Analytical value Computed mean value Confidence interval 95%. Effective sample width: [x-2s,x+2s] 6.45 6.5 6.55 6.6 6.65 6.7 6.75 6.8



## Statistics Results A Finite-Horizon simulation

- What are steady-state statistics? When they exist?
- So, do the steady-state statistics exist in our system?
- Is possible to obtain an interval estimation for steady state statistics using our simulator?

## Statistics Results Steady State statistics





#### Statistics Results Steady State statistics

- Finite-horizon interval estimates are accurate steady-state estimates, becoming increasingly more accurate as the number of jobs increases.
- The steady-state average wait indicator is interior to the finite-horizon interval estimate, suggesting that current system simulation parameters ( $\tau^*$ , n) allow us to achieve statistics that are close to their steady-state values.
- But there is another way...

Which routing policy is better?

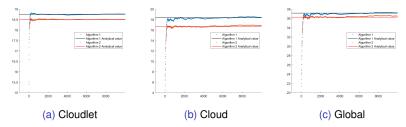


Figura: Time-Average Job Population

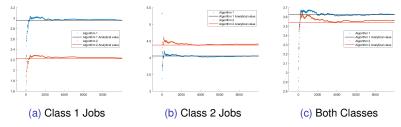


Figura: Global Time-Average service/response time

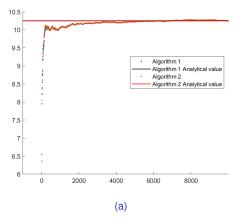


Figura: Global Throughput

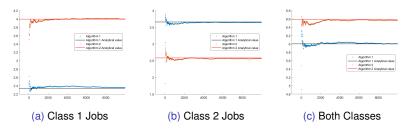


Figura: Cloudlet Throughput

Grazie per l'attenzione!