



TOR VERGATA
UNIVERSITÀ DEGLI STUDI DI ROMA

Macroarea di Ingegneria

A QoS-Aware Broker for Multi-Provider Serverless Applications

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My thesis is focused on “**serverless computing**”.

It is a development paradigm according to which:

- The provider takes care of all aspects of server management.
- Cloud application are abstracted as a group of so-called **serverless functions**, which are computation units implementing a business functionality.

A serverless function is executed inside a containerized environment: the so-called **function instance**.

The FaaS platform **automatically** scales the number of function instances.

FaaS platforms impose a **limit** on the number of function instances runnable at the **same time** called **concurrency limit**.

A delay is observed when a new function instance is started by the provider: this event is called **cold start**.

To invoke a **serverless function**, users have to specify a so-called **serverless function configuration**.

- Generally, the **amount of memory allocated** to a serverless function.

Configuration parameters **significantly** affect the **cost** and **response time** of serverless functions.

- Lack of support for applications whose functions are hosted on **multiple providers**.
- Lack of support for serverless function implementations abstraction, that is, for the so-called **concrete functions**.
- Fulfillment of **non-functional requirements** concerning the **quality of service** (QoS) levels that should be guaranteed for **multi-provider serverless applications**.
- Fulfillment of **functional requirements** concerning the **orchestration** of multi-provider applications.

Goal # 1

To guarantee the **satisfaction of QoS levels** for multi-provider serverless applications.

It was necessary to develop:

- 1 An **analytical model** to evaluate aforementioned class of applications.
- 2 A **methodological way** to find the “best” configuration to satisfy QoS constraints.

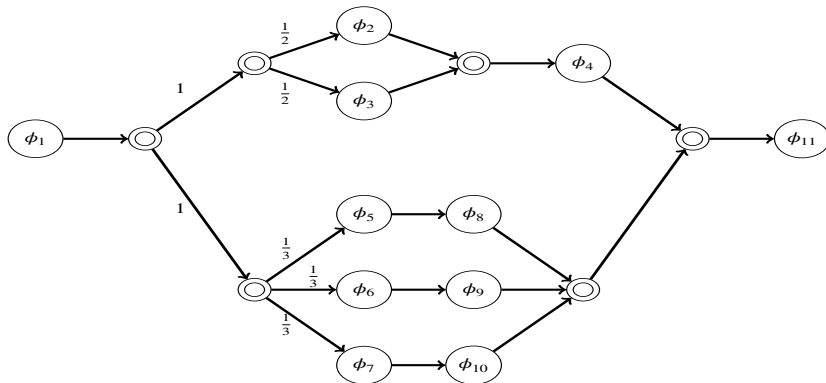
Goal # 2

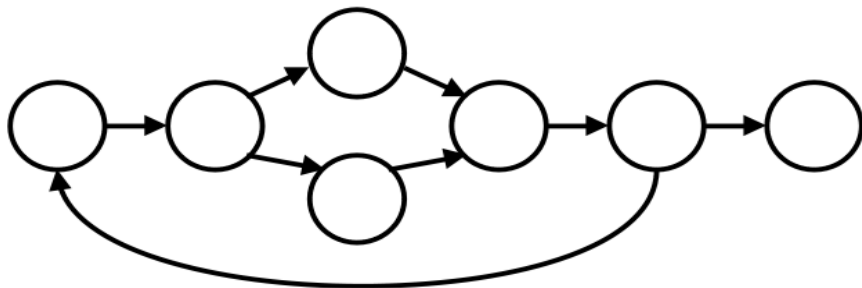
The **orchestration** for multi-provider serverless applications.

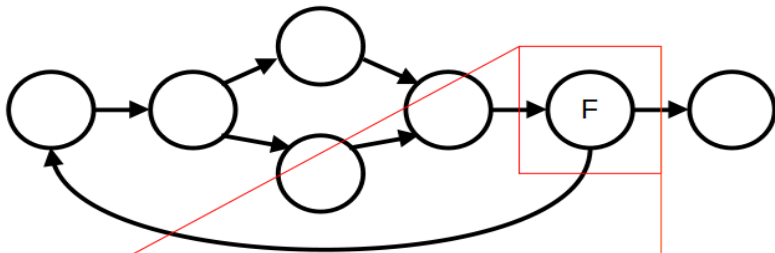
To achieve it, I had to build:

- 1 A **software framework**.
- 2 An extension to an already existent **representation scheme** to define a serverless application workflow.
 - Based on an existing language called **abstract function choreography language** (AFCL).

Serverless applications are abstracted to a **weakly connected weighted directed graph**.







AWS
Lambda

f_x Concrete Function 1

f_x Concrete Function 2

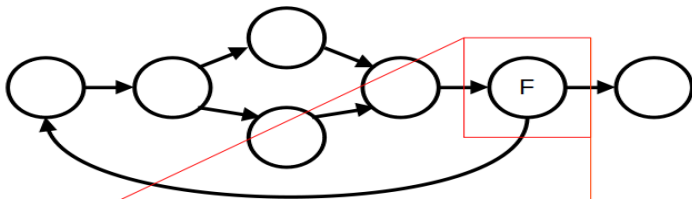

f_x Concrete Function 3



APACHE
OpenWhisk


f_x Concrete Function 1

f_x Concrete Function 2


AWS Lambda

- f_x Concrete Function 1
- f_x Concrete Function 2
- f_x Concrete Function 3



APACHE OpenWhisk

- f_x Concrete Function 1
- f_x Concrete Function 2

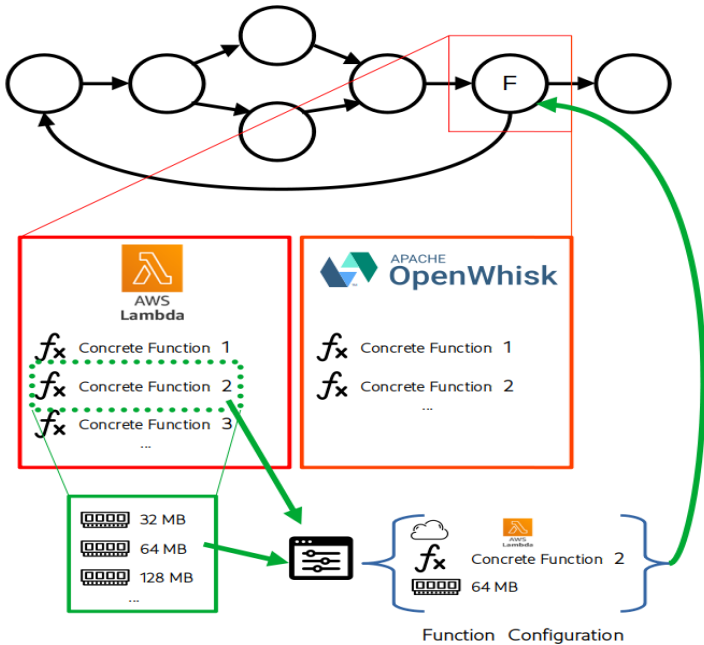


32 MB

64 MB

128 MB

...



- Estimations of **average response time** and **cost** under **any** possible configuration of all concrete functions is done using **exponential moving average** approach.
- An estimation of the **probability** according to which a request follows a cold start is required.
 - This is done using **Erlang-B** formula by modeling FaaS platform providers by **sets** of $M/G/K(t)/K(t)$ queueing systems.
 - $K(t)$: the number of function instances at time t .
- Performance estimations of the application depend on its workflow properties.

- To achieve our goal consisting in finding the best configuration to guarantee QoS constraints, we have to solve an **optimization problem**.
 - It is based on **multi-dimensional multi-choice knapsack problem formulation** (MMKP).

$$\max \sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(C)|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot p_{\phi_{i_j}}(t) \quad (5.10)$$

$$\text{subject to } \sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(C)|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot c_{\phi_{i_j}}(t) \leq C \quad (5.11)$$

$$\begin{aligned} & \sum_{\phi_i \in \mathcal{F}_{\mathcal{E}}(C) \setminus \mathcal{F}_{\mathcal{E}}(\tilde{\mathcal{P}})} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot rt_{\phi_{i_j}}(t) + \\ & + \sum_{\phi_h \in \delta_C} \sum_{j=1}^{|\mathbf{F}_{\phi_h} \times \mathbb{N}|} y_{\phi_{h_j}} \cdot rt_{\phi_{h_j}}(t) \leq RT \end{aligned} \quad \forall \delta_C \in \Delta_C \quad (5.12)$$

$$\sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(C)|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot a_{(\phi_{i_j}, \omega_P^{(l)})} \leq l - R(\mathbf{Q}_{\omega_P^{(l)}}, t) \quad \forall \omega_P^{(l)} \in \tilde{\mathbf{S}}_C \quad (5.13)$$

$$\sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} = 1 \quad \forall i \in \mathbb{N} \cap [1, |\mathcal{F}_{\mathcal{E}}|] \quad (5.14)$$

$$y_{\phi_{i_j}} \in \{0, 1\} \quad \forall i \in \mathbb{N} \cap [1, |\mathcal{F}_{\mathcal{E}}|] \\ \forall j \in \mathbb{N} \cap [1, |\mathbf{F}_{\phi_i} \times \mathbb{N}|]$$

$$\max \sum_{i=1}^{|\mathcal{F}_\mathcal{E}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot p_{\phi_{i_j}}(t) \quad (5.10)$$

$$\text{subject to } \sum_{i=1}^{|\mathcal{F}_\mathcal{E}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot c_{\phi_{i_j}}(t) \leq C \quad (5.11)$$

$$\begin{aligned} & \sum_{\phi_i \in \mathcal{F}_\mathcal{E}(\mathcal{C}) \setminus \mathcal{F}_\mathcal{E}(\tilde{\mathcal{P}})} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot rt_{\phi_{i_j}}(t) + \text{Cost Constraint} \\ & + \sum_{\phi_h \in \delta_C} \sum_{j=1}^{|\mathbf{F}_{\phi_h} \times \mathbb{N}|} y_{\phi_{h_j}} \cdot rt_{\phi_{h_j}}(t) \leq RT \end{aligned} \quad \forall \delta_C \in \Delta_C \quad (5.12)$$

$$\sum_{i=1}^{|\mathcal{F}_\mathcal{E}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot a_{(\phi_{i_j}, \omega_P^{(l)})} \leq l - R(\mathbf{Q}_{\omega_P^{(l)}}, t) \quad \forall \omega_P^{(l)} \in \tilde{\mathbf{S}}_C \quad (5.13)$$

$$\sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} = 1 \quad \forall i \in \mathbb{N} \cap [1, |\mathcal{F}_\mathcal{E}|] \quad (5.14)$$

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$$\max \sum_{i=1}^{|\mathcal{F}_\mathcal{E}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{ij}} \cdot p_{\phi_{ij}}(t) \quad (5.10)$$

$$\text{subject to } \sum_{i=1}^{|\mathcal{F}_\mathcal{E}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{ij}} \cdot c_{\phi_{ij}}(t) \leq C \quad (5.11)$$

$$\begin{aligned} & \sum_{\phi_i \in \mathcal{F}_\mathcal{E}(\mathcal{C}) \setminus \mathcal{F}_\mathcal{E}(\tilde{\mathcal{P}})} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{ij}} \cdot rt_{\phi_{ij}}(t) + \\ & + \sum_{\phi_h \in \delta_\mathcal{C}} \sum_{j=1}^{|\mathbf{F}_{\phi_h} \times \mathbb{N}|} y_{\phi_{hj}} \cdot rt_{\phi_{hj}}(t) \leq RT \end{aligned} \quad \begin{array}{l} \text{Response Time Constraint} \\ \forall \delta_\mathcal{C} \in \Delta_\mathcal{C} \end{array} \quad (5.12)$$

$$\sum_{i=1}^{|\mathcal{F}_\mathcal{E}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{ij}} \cdot a_{(\phi_{ij}, \omega_P^{(l)})} \leq l - R(\mathbf{Q}_{\omega_P^{(l)}}, t) \quad \forall \omega_P^{(l)} \in \tilde{\mathbf{S}}_\mathcal{C} \quad (5.13)$$

$$\sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{ij}} = 1 \quad \forall i \in \mathbb{N} \cap [1, |\mathcal{F}_\mathcal{E}|] \quad (5.14)$$

$$y_{\phi_{ij}} \in \{0, 1\} \quad \begin{array}{l} \forall i \in \mathbb{N} \cap [1, |\mathcal{F}_\mathcal{E}|] \\ \forall j \in \mathbb{N} \cap [1, |\mathbf{F}_{\phi_i} \times \mathbb{N}|] \end{array}$$

$$\max \sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(C)|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{ij}} \cdot p_{\phi_{ij}}(t) \quad (5.10)$$

$$\text{subject to} \quad \sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(C)|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{ij}} \cdot c_{\phi_{ij}}(t) \leq C \quad (5.11)$$

$$\begin{aligned} & \sum_{\phi_i \in \mathcal{F}_{\mathcal{E}}(C) \setminus \mathcal{F}_{\mathcal{E}}(\tilde{\mathcal{P}})} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{ij}} \cdot rt_{\phi_{ij}}(t) + \\ & + \sum_{\phi_h \in \delta_C} \sum_{j=1}^{|\mathbf{F}_{\phi_h} \times \mathbb{N}|} y_{\phi_{hj}} \cdot rt_{\phi_{hj}}(t) \leq RT \quad \forall \delta_C \in \Delta_C \end{aligned} \quad (5.12)$$

$$\sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(C)|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{ij}} \cdot a_{(\phi_{ij}, \omega_P^{(l)})} \leq l - R(\mathbf{Q}_{\omega_P^{(l)}}, t) \quad \forall \omega_P^{(l)} \in \tilde{\mathbf{S}}_C \quad (5.13)$$

$$\sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{ij}} = 1 \quad \text{Capacity Constraint} \quad \forall i \in \mathbb{N} \cap [1, |\mathcal{F}_{\mathcal{E}}|] \quad (5.14)$$

$$y_{\phi_{ij}} \in \{0, 1\} \quad \forall i \in \mathbb{N} \cap [1, |\mathcal{F}_{\mathcal{E}}|] \\ \forall j \in \mathbb{N} \cap [1, |\mathbf{F}_{\phi_i} \times \mathbb{N}|]$$

I develop a custom heuristic algorithm based on **ant colony optimization** (ACO).

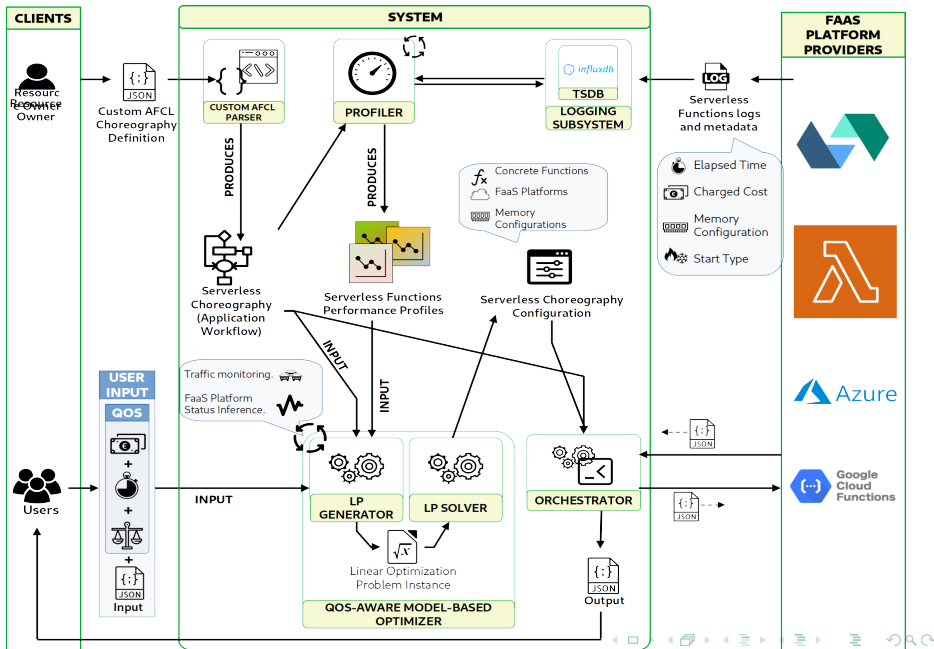
- It is based on a set of computational agents, called **artificial ants**, which **iteratively** construct a solution.
- At each iteration, each agent moves from a solution to another, applying a series of stochastic **local** decisions whose policy is based on following parameters:
 - **Attractiveness.**
 - **Pheromone trail.**

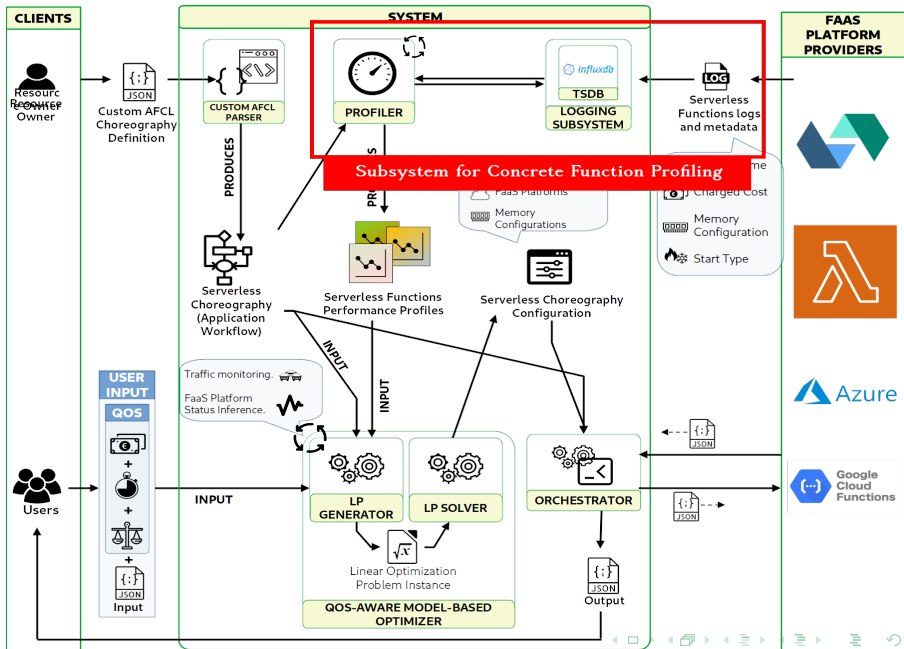
Pheromone trails are updated during **every iteration**.

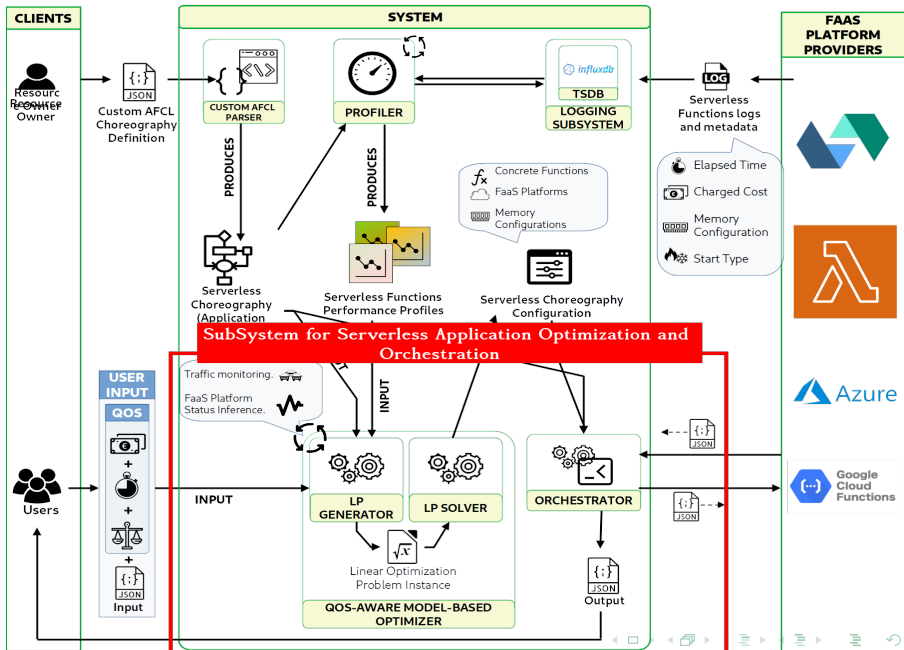
Pheromone trails are used to decide which solutions should be preferred during **next iterations**.

Main features of our software framework:

- **Client-server architecture.**
- **Cloud-native** application.
- Includes a set of **adapters** to interact with following FaaS providers:
 - AWS Lambda.
 - Apache OpenWhisk.
- **REST** architectural style.



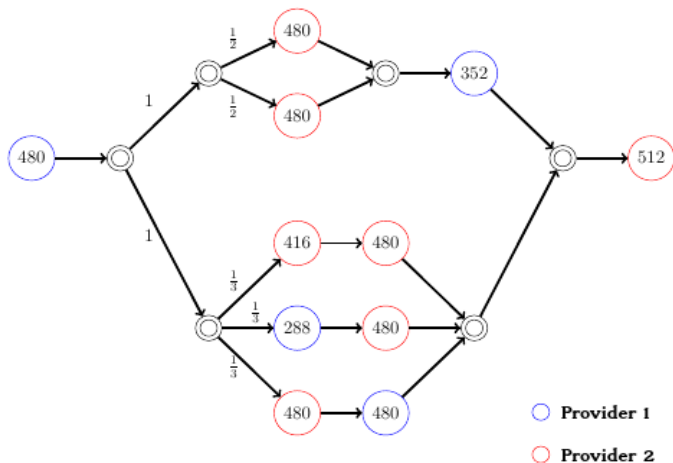




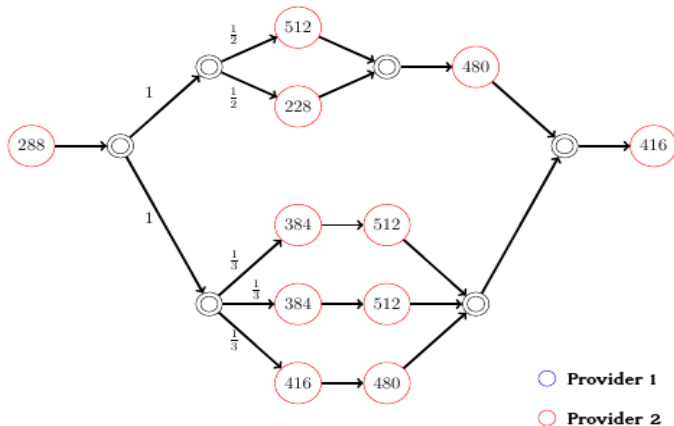
We validate our model through several experiments using an image-processing serverless application.

- Firstly, we check the respect of user specified QoS constraints in a static way thorough several **sequential invocations**.
- Then, we test the model in a dynamically way thorough **several concurrent and parallel invocations**.
 - That experiment aimed to **run out of capacity** on one FaaS provider and so force our prototype to schedule the concrete functions on the other.

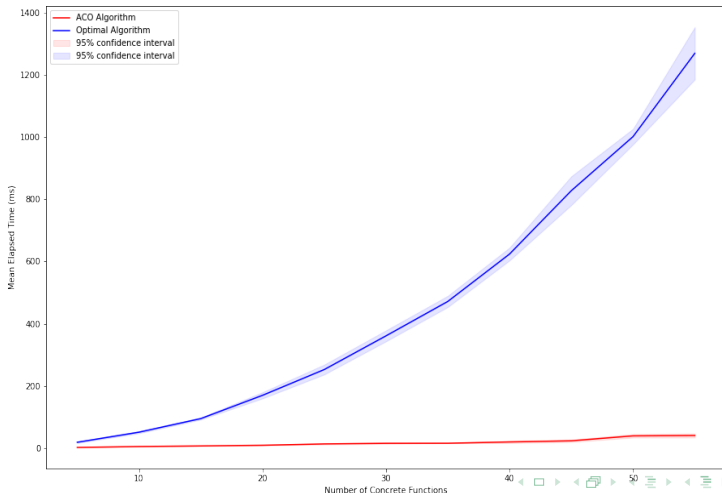
- Application configuration produced by our system **before** Provider 1 runs out of capacity...



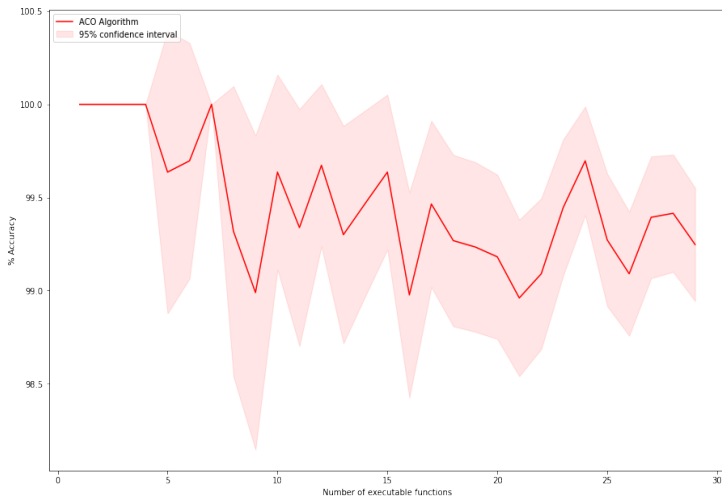
- ...and **after**, when Provider 1 runs out of capacity.



- We compared **execution time** of our heuristic algorithm with that of optimal algorithm through several experiments.



- According to my results, **accuracy** is above **98%**.

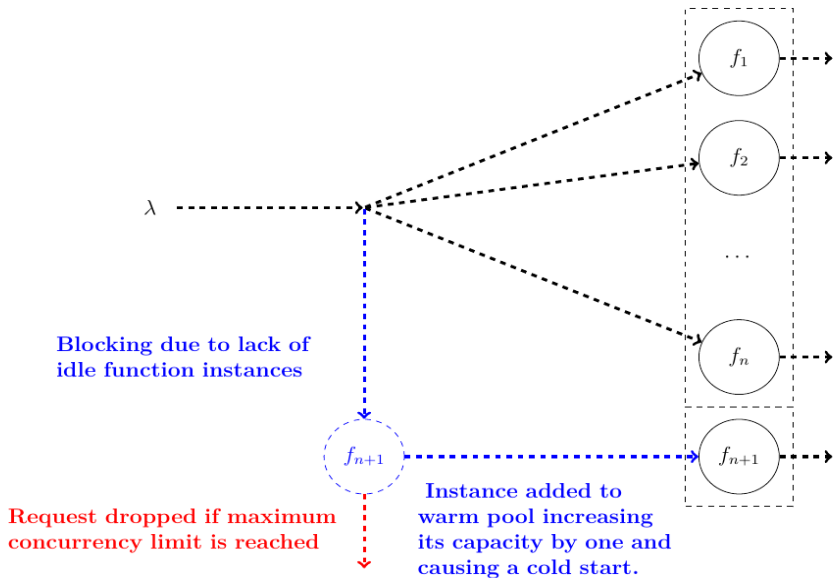


- 1 I presented an analytical model to evaluate the performance of multi-provider serverless application.
- 2 I defined an optimization problem formulation to address the problem of finding a suitable configuration in order to meet user defined QoS constraints.
- 3 I developed a heuristic algorithm to efficiently solve it.
- 4 I validated proposed solution through test and experimental evaluations using my prototype.



Thanks for your attention!
Questions?

Warm Pool having capacity n



Pheromone Trail Update

$$\tau(o_{i_a}, o_{j_b})_k = \tau(o_{i_a}, o_{j_b})_{k-1} \cdot \rho + \Delta\tau(o_{i_a}, o_{j_b})_k$$

where:

$$\Delta\tau(o_{i_a}, o_{j_b})_k \stackrel{\text{def}}{=} \sum_{l=1}^m \tau(o_{i_a}, o_{j_b})_k^{(l)} \quad (1)$$

and:

$$\tau(o_{i_j}, o_{a_b})_k \stackrel{\text{def}}{=} \begin{cases} 0 & \text{if } S_{k_l} = \emptyset \\ \frac{1}{1 + \text{profit}(S_{\text{best}}) - \text{profit}(S_{k_l})} & \text{if } \begin{matrix} S_{k_l} \neq \emptyset \\ (o_{i_j}, o_{a_b}) \in \pi_{k_l} \end{matrix} \end{cases} \quad (2)$$

Transition Probability

$$P(o_{i_j}, \mathbf{s}_{k_l}^{(z)}, \pi_{k_l}^{(z)}) \stackrel{\text{def}}{=} \frac{\left[\tau(o_{i_j}, \pi_{k_l}^{(z)}) \right]^\alpha \cdot \left[\eta(o_{i_j}, \mathbf{s}_{k_l}^{(z)}) \right]^\beta}{\sum_{o_{i_j} \in \mathcal{C}(\mathbf{G}_i, \mathbf{s}_{k_l}^{(z)})} \left[\tau(o_{i_j}, \pi_{k_l}^{(z)}) \right]^\alpha \cdot \left[\eta(o_{i_j}, \mathbf{s}_{k_l}^{(z)}) \right]^\beta} \quad (3)$$

Function Configuration

$$x_\phi = (f_\phi, m) \in f_\phi \times \mathbf{M}_{f_\phi} \subseteq \mathbf{F}_\phi \times \mathbb{N} \quad (4)$$

Choreography Configuration

$$\begin{aligned} \mathbf{x}_{\mathcal{C}} &\stackrel{\text{def}}{=} \left\{ x_{\phi_1}, \dots, x_{\phi_k} \right\} \\ &\in \left\{ \left\{ \bigcup_{j=1}^{|\mathbf{F}_{\phi_1}|} f_{\phi_{1j}} \times \mathbf{M}_{f_{\phi_{1j}}} \right\} \times \dots \times \left\{ \bigcup_{j=1}^{|\mathbf{F}_{\phi_k}|} f_{\phi_{kj}} \times \mathbf{M}_{f_{\phi_{kj}}} \right\} \right\} \\ &= \bigtimes_{i=1}^k \left\{ \bigcup_{j=1}^{|\mathbf{F}_{\phi_i}|} f_{\phi_{ij}} \times \mathbf{M}_{f_{\phi_{ij}}} \right\} \\ &\subseteq \bigtimes_{i=1}^k \left\{ \mathbf{F}_{\phi_i} \times \mathbb{N} \right\} = \mathbf{x}_{\mathcal{C}} \end{aligned} \quad (5)$$

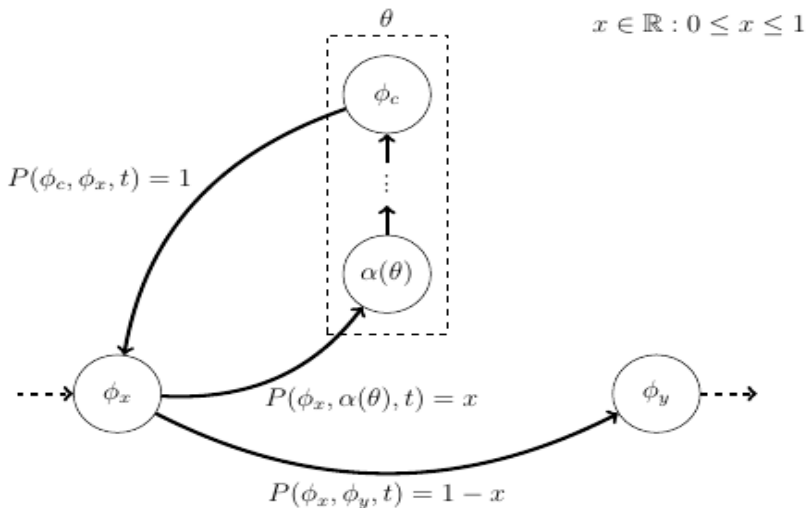


Figure 4.2: Conditional loop structure in a serverless workflow.

Performance Evaluation Loop Structure

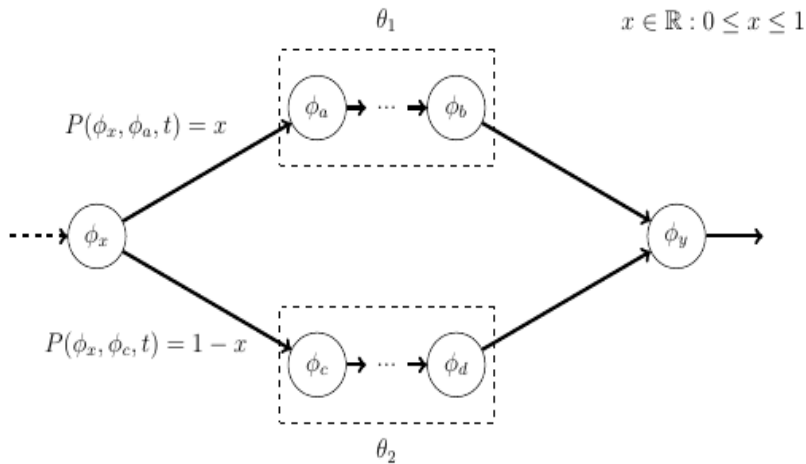
$$RT(\mathcal{L}, \mathbf{x}_{\mathcal{L}}, t) \stackrel{\text{def}}{=} E[I_{\theta}(t)] \cdot RT(\theta_i, \mathbf{x}_{\mathcal{L}}, t) \quad (6)$$

$$C(\mathcal{L}, \mathbf{x}_{\mathcal{L}}, t) \stackrel{\text{def}}{=} E[I_{\theta}(t)] \cdot C(\theta_i, \mathbf{x}_{\mathcal{L}}, t) \quad (7)$$

$$p = P(\phi_x, \phi_y, t) = 1 - P(\phi_x, \alpha(\theta), t) \quad (8)$$

$$\begin{aligned}
 P(I_\theta = k) &= p \cdot (1 - p)^{k-1} \\
 &= pq^{k-1} \\
 &= \left[1 - P(\phi_x, \alpha(\theta), t)\right] \cdot \left[P(\phi_x, \alpha(\theta), t)\right]^{k-1}
 \end{aligned} \tag{9}$$

$$\begin{aligned}
 E[I_\theta(t)] &= \sum_{k=1}^{\infty} (k-1)pq^{k-1} \\
 &= \frac{P(\phi_x, \alpha(\theta), t)}{1 - P(\phi_x, \alpha(\theta), t)}
 \end{aligned} \tag{10}$$



SLA definition

$$SLA \stackrel{\text{def}}{=} \langle (RT, w_{RT}), (C, w_C) \rangle$$

Function configuration score

$$p_{\phi_{ij}}(t) \stackrel{\text{def}}{=} w_{RT} \cdot p_{\phi_{ij}}(t)^{(rt)} + w_C \cdot p_{\phi_{ij}}(t)^{(c)}$$

$$p_{\phi_{ij}}(t)^{(rt)} \stackrel{\text{def}}{=} \begin{cases} 1 & \text{if } rt_{\phi_{i\text{MAX}}}(t) = rt_{\phi_{i\text{MIN}}}(t) \\ \frac{rt_{\phi_{i\text{MAX}}}(t) - rt_{\phi_{ij}}(t)}{rt_{\phi_{i\text{MAX}}}(t) - rt_{\phi_{i\text{MIN}}}(t)} & \text{otherwise} \end{cases} \quad (11)$$

$$p_{\phi_{ij}}(t)^{(c)} \stackrel{\text{def}}{=} \begin{cases} 1 & \text{if } c_{\phi_{i\text{MAX}}}(t) = c_{\phi_{i\text{MIN}}}(t) \\ \frac{c_{\phi_{i\text{MAX}}}(t) - c_{\phi_{ij}}(t)}{c_{\phi_{i\text{MAX}}}(t) - c_{\phi_{i\text{MIN}}}(t)} & \text{otherwise} \end{cases} \quad (12)$$

$$\max \sum_{i=1}^{|\mathbf{x}_C|} x_i \cdot S(\mathbf{x}_{C_i}, t) \quad (5.1)$$

$$\text{subject to } \sum_{\omega=1}^{|\mathbf{x}_C|} x_i \cdot C(\mathcal{C}, \mathbf{x}_{C_i}, t) \leq C \quad (5.2)$$

$$\sum_{i=1}^{|\mathbf{x}_C|} x_i \cdot RT(\mathcal{C}, \mathbf{x}_C, t) \leq RT \quad (5.3)$$

$$\sum_{i=1}^{|\mathbf{x}_C|} x_i \cdot N(\mathcal{C}, \mathbf{x}_C, \omega_P^{(l)}) \leq l - R(\mathbf{Q}_{\omega_P^{(l)}}, t) \quad \forall \omega_P^{(l)} \in \tilde{\mathbf{S}}_C \quad (5.4)$$

$$\sum_{i=1}^{|\mathbf{x}_C|} x_i = 1 \quad (5.5)$$

$$x_i \in \{0, 1\} \quad \forall i \in \mathbb{N} \cap [1, |\Omega|] \quad (5.6)$$

Concrete Function Average Response Time

$$\begin{aligned} RT(x_\phi, t) &\stackrel{\text{def}}{=} RT(f_\phi, m, t) \\ &= RT_{avg}^{(c)}(f_\phi, m, t) \cdot \mathbf{P}_{\omega_{PR}^{(l)}}(t) + \\ &\quad RT_{avg}^{(w)}(f_\phi, m, t) \cdot \left(1 - \mathbf{P}_{\omega_{PR}^{(l)}}(t)\right) \end{aligned} \quad (13)$$

where:

$$RT_{avg}^{(c)}(f_\phi, m, t) \stackrel{\text{def}}{=} \begin{cases} Y_o & \text{if } t = 0 \\ \alpha \cdot Y_t + (1 - \alpha) \cdot RT_{avg}^{(c)}(f_\phi, m, t - 1) & \text{if } t > 0 \end{cases} \quad (14)$$

Listing 3.1: Definition of a base function in AFCL

```
1 function: {
2     name: "name",
3     type: "type",
4     dataIns: [
5         {
6             name: "name", type: "type",
7             source: "source?", value: "value?",
8             properties: [{name: "name", value: "value"}+]?
9             constraints: [{name: "name", value: "value"}+]?
10        }+
11    ],
12    properties: [{name: "name", value: "value"}+]?
13    constraints: [{name: "name", value: "value"}+]?
14    dataOuts: [
15        {
16            name: "name", type: "type",
17            source: "source?", value: "value?",
18            properties: [{name: "name", value: "value"}+]?
19            constraints: [{name: "name", value: "value"}+]?
20        }+
21    ],
22 }
```

Listing 3.2: Definition of a serverless function according to our custom AFCL.

```
1      function: {
2          name: "name",
3          "implementations": [
4              {
5                  providerName: "providerName",
6                  concreteFunctionName: "concreteFunctionName"
7                  host: "host"?
8              }+
9          ],
10         "profilingPayloads": [
11             {...}]
12         ]
13         dataIns: [
14             {
15                 name: "name", type: "type",
16                 source: "source"
17             }+
18         ],
19         dataOuts: [
20             {
21                 name: "name", type: "type",
22                 source: "source"
23             }+
24         ]?,
25     }
```

Algorithm 2: Generic algorithmic skeleton for ACO algorithms

```
1 Initialization pheromone trails;  
2 while Termination conditions not met do  
3   | ConstructSolutions;  
4   | ApplyLocalSearch;  
5   | UpdatePheromoneTrails;  
6 end
```
