

A QoS-Aware Broker for Multi-Provider Serverless Applications

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Serverless Computing: Overview



My thesis is focused on "serverless computing".

It is a development paradigm according to which:

- The provider takes care of all aspect of server management.
- Small-granularity billing pricing model: pay-as-you-go.
- Cloud application are abstracted as a group of so-called serverless functions, which are computation units implementing a business functionality.

Serverless Computing: Overview



A serverless function is executed inside a containerized environment: the so-called **function instance**.

The FaaS platform automatically scales the number of function instances.

FaaS platforms impose a limit on the number of function instance runnable at the same time called **concurrency limit**.

A delay is observed when a new function instance is started by the provider: this event is called **cold start**.

Serverless Computing: Overview



To invoke a serverless function, users have to specify a so-called serverless function configuration.

• Generally, the amount of memory allocated to a serverless function.

Configuration parameters **significantly** affect the **cost** and **response time** of serverless functions.

Serverless Computing: Problems



- The lack of support for application whose functions are hosted on multiple providers.
- The lack of support for serverless function implementations abstraction, that is, for the so-called concrete functions.
- The fulfillment of non-functional requirements concerning the quality of service (QoS) levels that should be guaranteed for multi-provider serverless applications.
- The fulfillment of functional requirements concerning the orchestration of multi-provider applications.

State of Art



Solutions concerning QoS fulfillment already exist.

- Many solutions rely on QoS-aware scheduling algorithms while others rely on the formulation and solving of optimization problems.
 - Generally, proposed solution do not support Multi-Provider serverless applications.
- Despite there are some exceptions, current solution are unaware of both the current status of FaaS platforms and user traffic.

Thesis Goals



Goal #1

To guarantee the **satisfaction of QoS levels** for multi-provider serverless applications.

It was necessary to develop:

- An analytical model to evaluate aforementioned class of applications.
- A methodological way to find the "best" configuration to satisfy QoS constraints.
 - By solving an optimization problem (LP).
- A custom heuristic algorithm to rapidly resolve the aforementioned optimization problem.

Thesis Goals



Goal #2

The orchestration for multi-provider serverless applications.

To achieve it, I had to build:

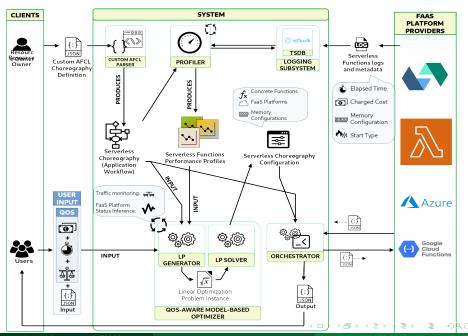
- A software framework
 - Users/Serverless application management (CRUD operations, profiling tasks, orchestration task, etc.)
- An extension to an already existent representation scheme to define a serverless application workflow.
 - Based on an existing language called abstract function choreography language (AFCL).

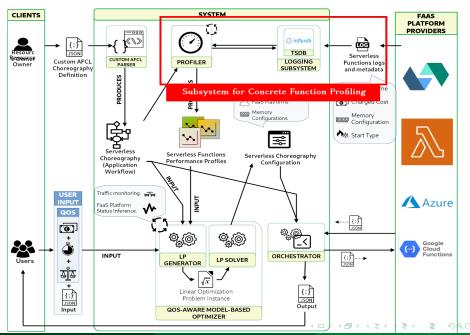
Prototype

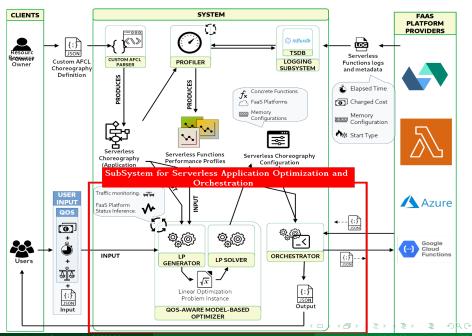


Main features of our software framework:

- Client-server architecture.
- Cloud-native application.
- Includes a set of adapters to interact with following FaaS providers:
 - AWS Lambda.
 - Apache OpenWhisk.
- REST architectural style.



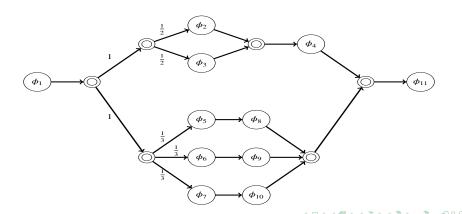




System Model



Serverless applications are abstracted to a weakly connected weighted directed graph.



System Model



- Estimations of average response time and cost under any possible configuration of all concrete functions is done using exponential moving average approach.
- An estimation of the probability according to which a request follows a cold start is required.
 - This is done using **Erlang-B** formula by modeling FaaS platform providers by **sets** of M/G/K(t)/K(t) queueing systems.
 - K(t): the number of function instances at time t.

Optimization Problem



- To achieve our goal consisting in finding the best configuration to guarantee QoS constraints, we have to solve an optimization problem.
 - It is based on multi-dimensional multi-choice knapsack problem formulation (MMKP).

$$max \sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_{i}} \times \mathbb{N}|} y_{\phi_{i_{j}}} \cdot p_{\phi_{i_{j}}}(t)$$

$$\text{subject to} \sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_{i}} \times \mathbb{N}|} y_{\phi_{i_{j}}} \cdot c_{\phi_{i_{j}}}(t) \leq C$$

$$\sum_{\phi_{i} \in \mathcal{F}_{\mathcal{E}}(\mathcal{C}) \setminus \mathcal{F}_{\mathcal{E}}(\widetilde{\mathcal{P}})} \sum_{j=1}^{|\mathbf{F}_{\phi_{i}} \times \mathbb{N}|} y_{\phi_{i_{j}}} \cdot rt_{\phi_{i_{j}}}(t) +$$

$$+ \sum_{\phi_{h} \in \mathcal{S}_{\mathcal{C}}} \sum_{j=1}^{|\mathbf{F}_{\phi_{h}} \times \mathbb{N}|} y_{\phi_{h_{j}}} \cdot rt_{\phi_{h_{j}}}(t) \leq RT$$

$$\forall \delta_{\mathcal{C}} \in \Delta_{\mathcal{C}}$$

$$(5.12)$$

$$\sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_{i}} \times \mathbb{N}|} y_{\phi_{i_{j}}} \cdot a_{(\phi_{i_{j}}, \omega_{P}^{(l)})} \leq l - R(\mathbf{Q}_{\omega_{P}^{(l)}}, t)$$

$$\forall \omega_{P}^{(l)} \in \widetilde{\mathbf{S}}_{\mathcal{C}}$$

$$(5.13)$$

$$i=1 \quad j=1$$

$$(5.13)$$

$$|\mathbf{F}_{\phi_i} \times \mathbb{N}|$$

$$\sum_{j=1}^{\mathbf{F}_{\phi_i}\times\mathbb{N}|}y_{\phi_{i_j}}=1 \qquad \qquad \forall i\in\mathbb{N}\cap[1,|\mathcal{F}_{\mathcal{E}}|]$$

(5.14) $\forall i \in \mathbb{N} \cap [1, |\mathcal{F}_{\mathcal{E}}|]$

$$\forall i \in \mathbb{N} \cap [1, |\mathcal{F}_{\ell_i}|]$$

$$\forall j \in \mathbb{N} \cap [1, |\mathbf{F}_{\phi_i} \times \mathbb{N}|]$$

 $y_{\phi_{i_j}} \in \{0,1\}$

$$\max \sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot p_{\phi_{i_j}}(t)$$

(5.10)

(5.11)

subject to
$$\sum_{i=1}^{|S_{\epsilon}(b)|} \sum_{i=1}^{|S_{\epsilon}(b)|}$$

$$\sum_{i=1}^{|i-1|}\sum_{j=1}^{|i-1|}y_{\phi_{i_j}}\cdot c_{\phi_{i_j}}(t)\leq C \qquad \text{Objective Function}$$

$$\sum_{\phi_i \in \mathcal{T}_{\mathcal{E}}(\mathcal{C}) \backslash \mathcal{T}_{\mathcal{E}}(\widetilde{\mathcal{P}})} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot rt_{\phi_{i_j}}(t) +$$

$$+ \sum_{\phi_h \in \delta_C} \sum_{j=1}^{|\mathbf{F}_{\phi_h} \times \mathbb{N}|} y_{\phi_{h_j}} \cdot rt_{\phi_{h_j}}(t) \le RT$$

 $\forall \delta_{\mathcal{C}} \in \Delta_{\mathcal{C}}$

$$\sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot a_{(\phi_{i_j}, \omega_P^{(l)})} \leq l - R(\mathbf{Q}_{\omega_P^{(l)}}, t)$$

 $\forall \omega_P^{(l)} \in \widetilde{\mathbf{S}}_C$

$$\sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} = 1$$

$$\forall i \in \mathbb{N} \cap [1, |\mathcal{F}_{\mathcal{E}}|]$$

$$y_{\phi_{i_i}} \in \{0, 1\}$$

$$\forall i \in \mathbb{N} \cap [1, |\mathcal{F}_{\mathcal{E}}|]$$

$$\forall j \in \mathbb{N} \cap [1, |\mathbf{F}_{\phi_i} \times \mathbb{N}|]$$

$$max \sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_{i}} \times \mathbb{N}|} y_{\phi_{i_{j}}} \cdot p_{\phi_{i_{j}}}(t)$$

$$\text{subject to} \sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_{i}} \times \mathbb{N}|} y_{\phi_{i_{j}}} \cdot c_{\phi_{i_{j}}}(t) \leq C$$

$$\sum_{\phi_{i} \in \mathcal{F}_{\mathcal{E}}(\mathcal{C}) \setminus \mathcal{F}_{\mathcal{E}}(\tilde{\mathcal{P}})} \sum_{j=1}^{|\mathbf{F}_{\phi_{i}} \times \mathbb{N}|} y_{\phi_{i_{j}}} \cdot rt_{\phi_{i_{j}}}(t) + \text{Cost Constraint}$$

$$+ \sum_{\phi_{h} \in \mathcal{S}_{\mathcal{E}}} \sum_{j=1}^{|\mathbf{F}_{\phi_{h}} \times \mathbb{N}|} y_{\phi_{h_{j}}} \cdot rt_{\phi_{h_{j}}}(t) \leq RT$$

$$\forall \delta_{\mathcal{C}} \in \Delta_{\mathcal{C}}$$

$$+ \sum_{\phi_h \in \delta_{\mathcal{C}}} \sum_{j=1} y_{\phi_{h_j}} \cdot rt_{\phi_{h_j}}(t) \le RT \qquad \forall \delta_{\mathcal{C}} \in \Delta_{\mathcal{C}}$$

$$|\mathcal{F}_{\mathcal{E}}(\mathcal{C})| |\mathbf{F}_{\phi_i} \times \mathbb{N}| \qquad (5.12)$$

$$\sum_{i=1}^{\mathcal{F}_{\mathcal{E}}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot a_{(\phi_{i_j}, \omega_P^{(l)})} \le l - R(\mathbf{Q}_{\omega_P^{(l)}}, t)$$

$$\forall \omega_P^{(l)} \in \widetilde{\mathbf{S}}_{\mathcal{C}}$$

$$(5.13)$$

$$\sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} = 1 \qquad \qquad \forall i \in \mathbb{N} \cap [1, |\mathcal{F}_{\mathcal{E}}|]$$

 $y_{\phi_{i_j}} \in \{0, 1\}$ $\forall i \in \mathbb{N} \cap [1, |\mathcal{F}_{\mathcal{E}}|]$ $\forall j \in \mathbb{N} \cap [1, |\mathbf{F}_{\phi_i} \times \mathbb{N}|]$

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(5.14)

$$\max \sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot p_{\phi_{i_j}}(t)$$
 (5.10)

subject to
$$\sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(\mathcal{C})|} \sum_{i=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot c_{\phi_{i_j}}(t) \leq C$$
 (5.11)

$$\sum_{\phi_{i} \in \mathcal{F}_{\mathcal{E}}(\mathcal{C}) \backslash \mathcal{F}_{\mathcal{E}}(\tilde{\mathcal{P}})} \sum_{j=1}^{|\mathbf{F}_{\phi_{i}} \times \mathbb{N}|} y_{\phi_{i_{j}}} \cdot rt_{\phi_{i_{j}}}(t) + \\ + \sum_{\phi_{b} \in \delta_{\mathcal{E}}} \sum_{j=1}^{|\mathbf{F}_{\phi_{b}} \times \mathbb{N}|} y_{\phi_{h_{j}}} \cdot rt_{\phi_{h_{j}}}(t) \leq RT \qquad \forall \delta_{\mathcal{C}} \in \Delta_{\mathcal{C}}$$

$$\sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot a_{(\phi_{i_j}, \omega_P^{(l)})} \le l - R(\mathbf{Q}_{\omega_P^{(l)}}, t)$$

$$\forall \omega_P^{(l)} \in \widetilde{\mathbf{S}}_{\mathcal{C}}$$

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$$\sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} = 1 \qquad \forall i \in \mathbb{N} \cap [1, |\mathcal{F}_{\mathcal{E}}|]$$

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$$\forall i \in \mathbb{N} \cap [1, |\mathcal{F}_{\mathcal{E}}|]$$

$$\forall j \in \mathbb{N} \cap [1, |\mathbf{F}_{\phi_i} \times \mathbb{N}|]$$

 $y_{\phi_{i_j}} \in \{0,1\}$

$$\max \sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot p_{\phi_{i_j}}(t)$$
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subject to
$$\sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot c_{\phi_{i_j}}(t) \le C$$

$$(5.11)$$

$$\sum_{\phi_i \in \mathcal{F}_{\mathcal{E}}(\mathcal{C}) \backslash \mathcal{F}_{\mathcal{E}}(\widetilde{\mathcal{P}})} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot rt_{\phi_{i_j}}(t) +$$

$$+\sum_{\phi_h \in \delta_C} \sum_{j=1}^{|\mathbf{F}_{\phi_h} \times \mathcal{N}|} y_{\phi_{h_j}} \cdot rt_{\phi_{h_j}}(t) \le RT$$

(5.12)

 $\forall \delta_{\mathcal{C}} \in \Delta_{\mathcal{C}}$

$$\sum_{i=1}^{|\mathcal{F}_{\mathcal{E}}(\mathcal{C})|} \sum_{j=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} \cdot a_{(\phi_{i_j}, \omega_P^{(l)})} \leq l - R(\mathbf{Q}_{\omega_P^{(l)}}, t)$$

$$\sum_{i=1}^{|\mathbf{F}_{\phi_i} \times \mathbb{N}|} y_{\phi_{i_j}} = 1$$

Capacity Constraint

 $\forall i \in \mathbb{N} \cap [1, |\mathcal{F}_{\varepsilon}|]$

(5.14)

 $\forall i \in \mathbb{N} \cap [1, |\mathcal{F}_{\mathcal{E}}|]$ $\forall j \in \mathbb{N} \cap [1, |\mathbf{F}_{\phi_i} \times \mathbb{N}|]$

$$y_{\phi_{i_j}} \in \{0,1\}$$

Heuristic Algorithm



I develop a custom heuristic algorithm based on ant colony optimization (ACO); it is called pre-provisioned colony optimization algorithm with lazy pheromone update.

- It is based on a set of computational agents, called artificial ants, which iteratively construct a solution.
- At each iteration, each agent moves from a solution to another, applying a series of stochastic local decisions whose policy is based on following parameters:
 - Attractiveness.
 - Pheromone trail.



Heuristic Algorithm



Pheromone trails are updated during every iteration.

Pheromone trails are used to decide which solutions should be preferred during **next iterations**.

- A lazy approach for pheromone trails update is used.
- A pre-provisioning tactic is used to anticipates data needs assuring a lower latency.

Validation Test



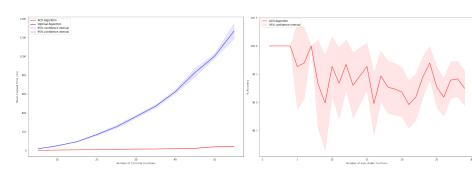
We validate our model through several experiments using an image-processing serverless application.

- Firstly, we check the respect of user specified QoS constraints in a static way thorough several sequential invocations.
- Then, we test the model in a dynamically way thorough several concurrent and parallel invocations.
 - That experiment aimed to run out of capacity on one FaaS provider and so force our prototype to schedule the concrete function on the other.

Heuristic Algorithm Evaluation



• We compared performance of our heuristic algorithm with that of optimal algorithm through several experiments.



Conclusions



- I presented an analytical model to evaluate the performance of multi-provider serverless application
- I defined an optimization problems formulation to address the problem of finding a suitable configuration in order to meet user defined QoS constraints
- I developed a heuristic algorithm to rapidly solve it.
- I validate proposed solution through test and experimental evaluations using my prototype.



Thanks for your attention! Questions?