# Vertically Scalable solutions illustrated via EPL and Esper

Emanuele Della Valle

Politecnico di Milano

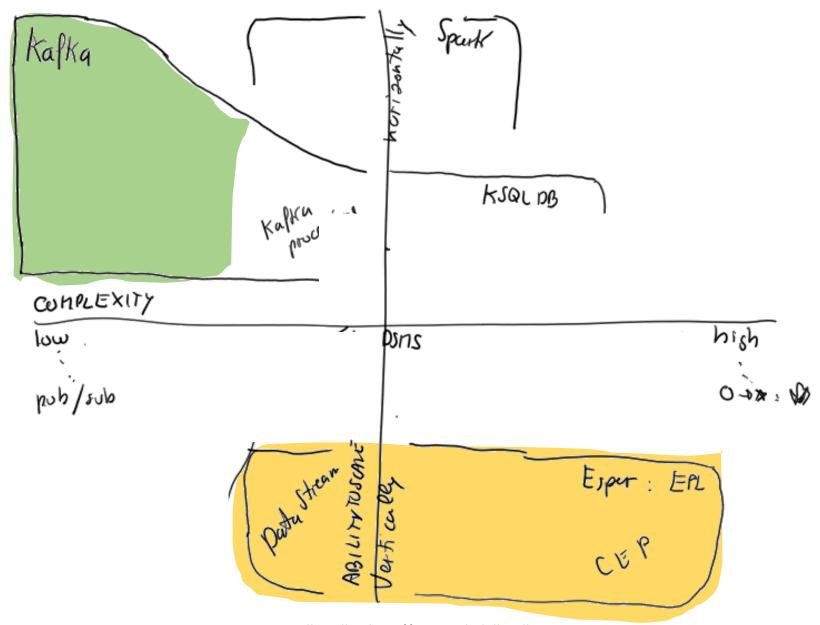
# Positioning this lecture

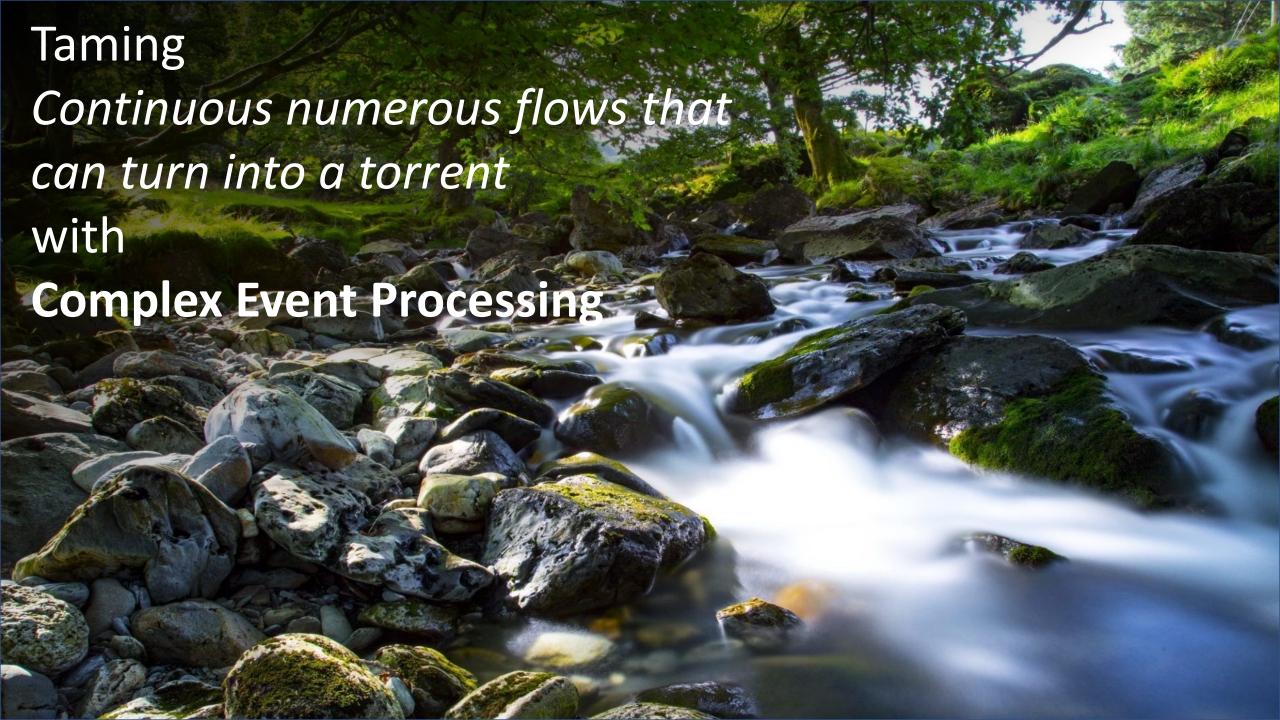
#### 1<sup>st</sup> premises: continuity matters

Traditional approach Velocity approach

Stop data to analyse

Analyse data in motion





# The Event Processing Language & its reference implementation Esper

#### Esper in a nutshell

- Implemented as a Java library
  - Can be embedded in any JVM application
- Designed for performance
  - High throughput
  - Low latency
- Tree-based recognition algorithm
- Inverted indexes to dispatch incoming events to EPL statements
- Builds indexes to quickly retrieve events with given properties among those stored

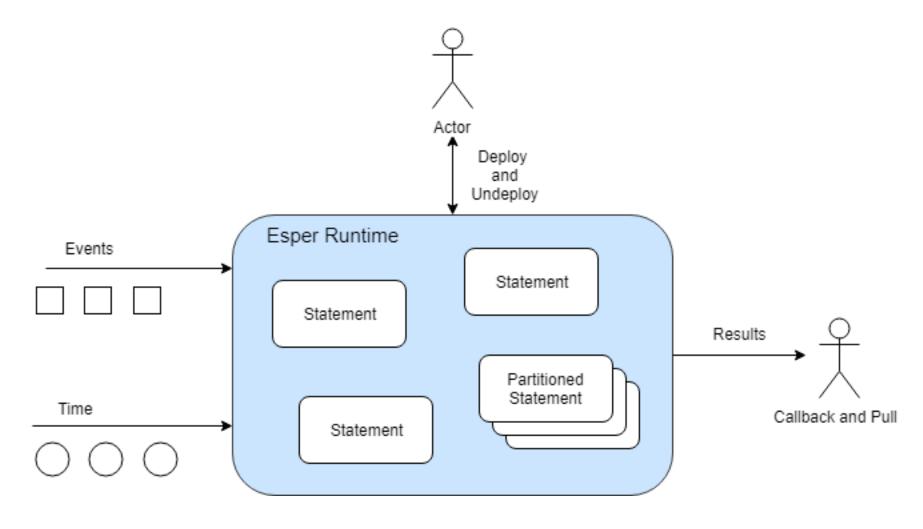
#### Esper in a nutshell (cont.)

- Interaction with static / historical data
- Configurable push or pull communication
- Several adapters for input/output
  - CSV, JMS in/out, API, DB, Socket, HTTP
- Esper HA
  - High Availability
  - Ensures that the state is recoverable in the case of failure

#### EPL in a nutshell

- EPL: rich language to express rules (a.k.a., statementa)
- Grounded on the DSMS approach
  - Windowing
  - Relational select, join, aggregate, ...
  - Relation-to-stream operators to produce output
- Queries can be combined to form a graph
- Includes complex event recognition abstractions
  - Pattern detection

### Esper & EPL



# The Basics of the Event Processing Language

#### Running example

 Count the number of fires detected using a set of smoke and temperature sensors in the last 10 minutes

#### Events

- Smoke event: String sensor, boolean state
- Temperature event: String sensor, double temperature
- Fire event: String sensor, boolean smoke, double temperature

#### Condition:

• Fire: at the same sensor smoke followed by temperature>50 within 2 minutes

#### Declare event types

- Two ways
  - EPL create schema clause
  - Runtime configuration API addEventType

#### Syntax

```
create schema
schema_name [as]
(property_name property_type
[,property_name property_type [,...])
[inherits inherited_event_type
[, inherited_event_type] [,...]]
```

#### Running example

```
create schema SmokeSensorEvent(
      sensor string,
      smoke boolean
);
create schema TemperatureSensorEvent(
      sensor string,
      temperature double
);
create schema FireComplexEvent(
      sensor string,
      smoke boolean,
      temperature double
);
```

#### Event Processing Language (EPL)

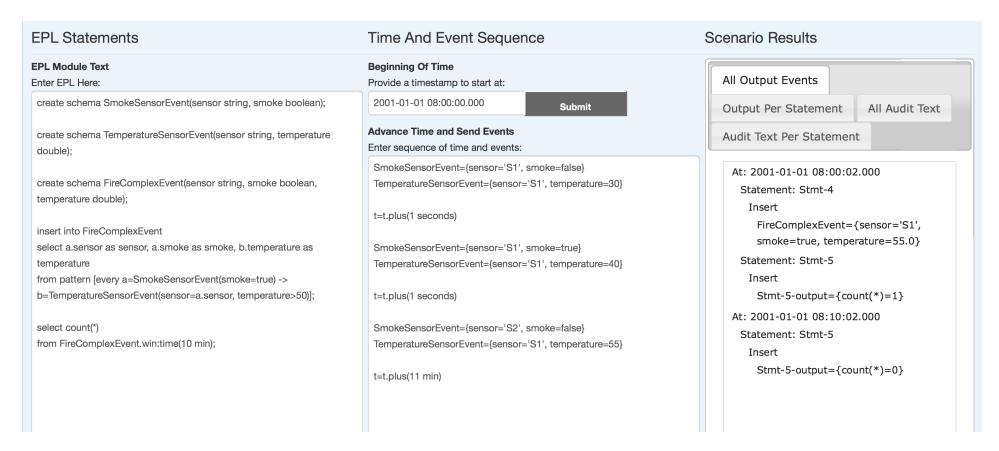
- EPL is similar to SQL
  - Select, where, ...
- Event streams and views instead of tables
  - Views define the data available for the query
  - Views can represent windows over streams
  - Views can also sort events, derive statistics from event attributes, group events, ...

#### EPL syntax

```
[insert into insert into def]
select select list
from stream def [as name]
[, stream def [as name]] [,...]
[where search conditions]
[group by grouping expression list]
[having grouping search conditions]
[output output specification]
[order by order by expression list]
[limit num rows]
```

#### Running example

#### http://esper-epl-tryout.appspot.com/epltryout/mainform.html



#### Time line

#### Filtering

```
to \frac{8:00:00}{T=30}

to \frac{S=False}{S:00:01}

to \frac{S=False}{S:00:01}

to \frac{S=True}{8:00:02}

to \frac{S=True}{S:00:03}
```

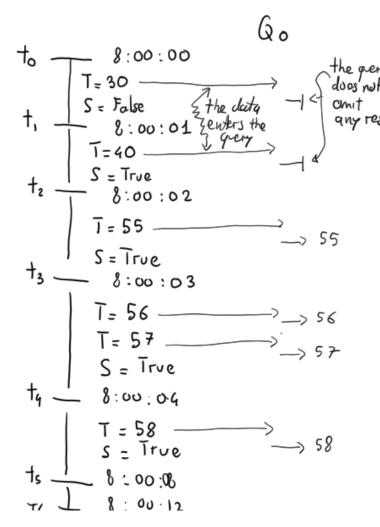
Q0 - the SQL-style

```
select *
from TemperatureSensorEvent
where temperature > 50;
```

Q0bis - the Event-Based System Style

```
select *
from TemperatureSensorEvent(temperature > 50);
```

#### Filtering the SQL style



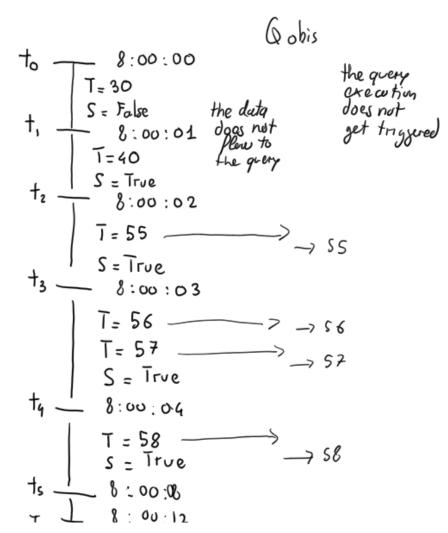
Q0 - the SQL-style

select \*
from TemperatureSensorEvent
where temperature > 50;

**Execution mode** 

Data points enter the query that filters them

#### Filtering the Event-Based System Style



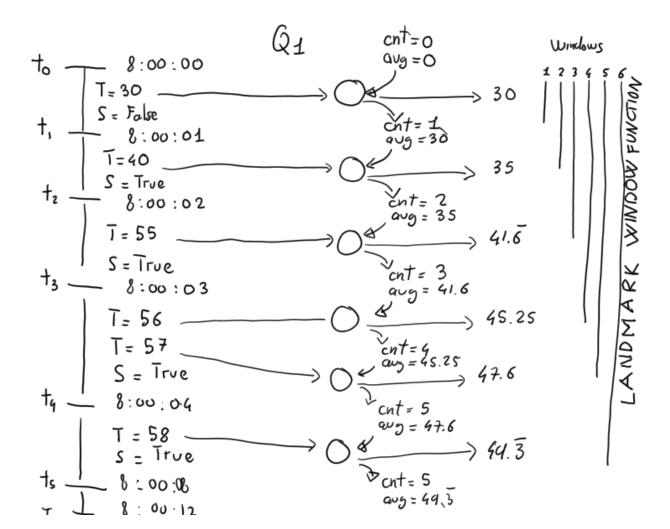
#### **Q0bis - the Event-Based System Style**

```
select *
from TemperatureSensorEvent(temperature > 50)
;
```

#### **Execution mode**

Data points are filtered before flowing into the query and the query execution does not get triggered

#### Basic Aggregation



select sensor, avg(temperature)
from TemperatureSensorEvent
group by sensor;

#### Execution mode

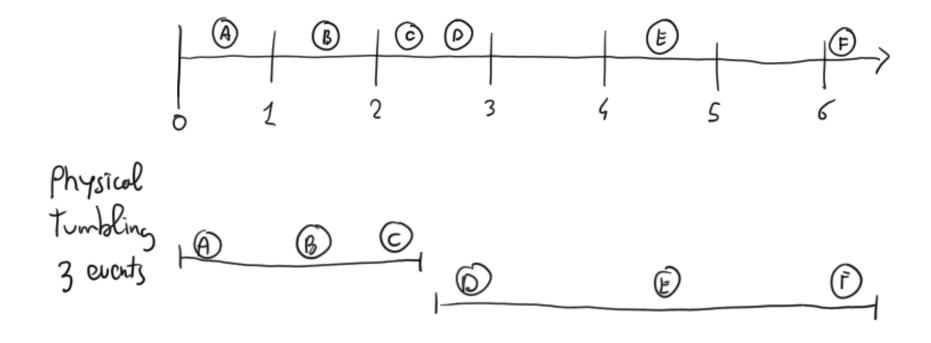
From a **logical** perspective, **data points cumulate in the landmark window,** and the query emits a new avg for every data point

From a physical perspective, the query is evaluated incrementally and maintains a state that captures the number of events (cnt) seen so far and the previous avg. The new average only depends on the state and the new data point entering the query. Old data can be forgotten.

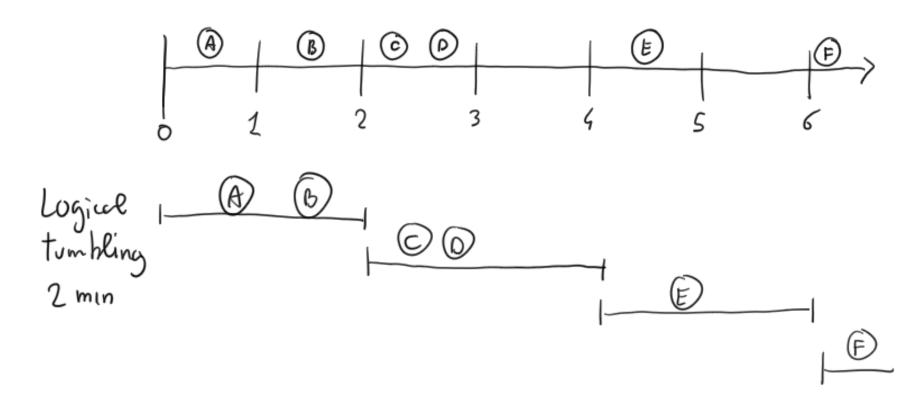
# Data Windowing

	Phy sical # events	Logical #Stime units	
sliding moves on overy event	native	hatice	
-			+ session window
tumbling	hutice	native	
ho pping	cutput claus	cutput clause	

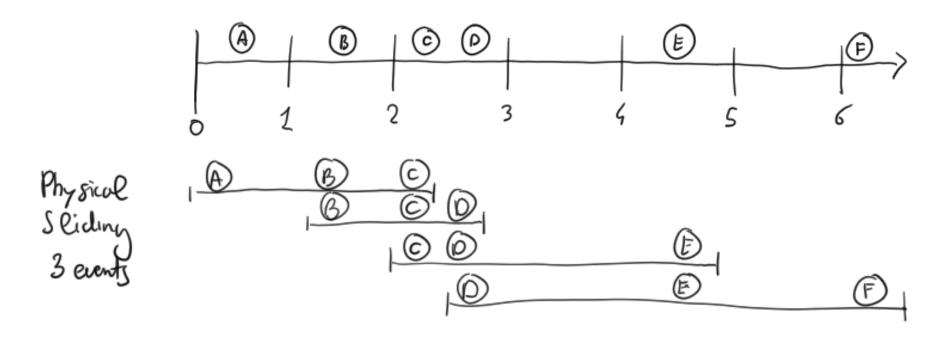
#### Physical tumbling windows



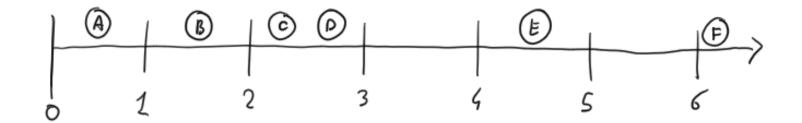
#### Logical tumbling windows

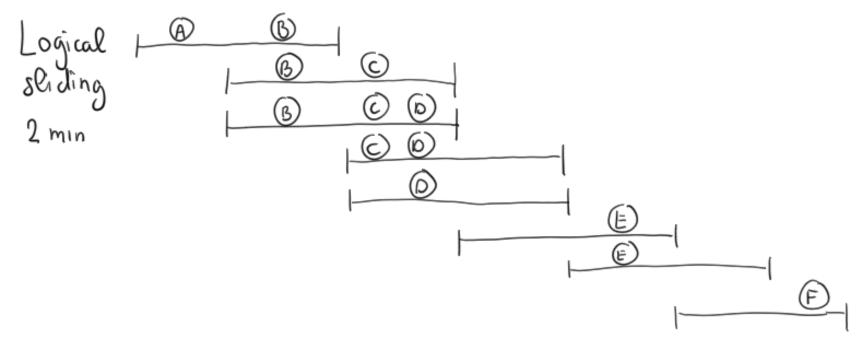


### Physical sliding windows



# Logical sliding windows

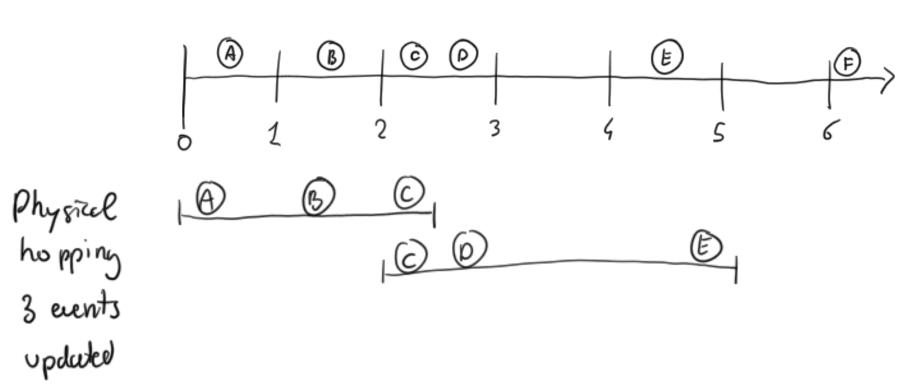




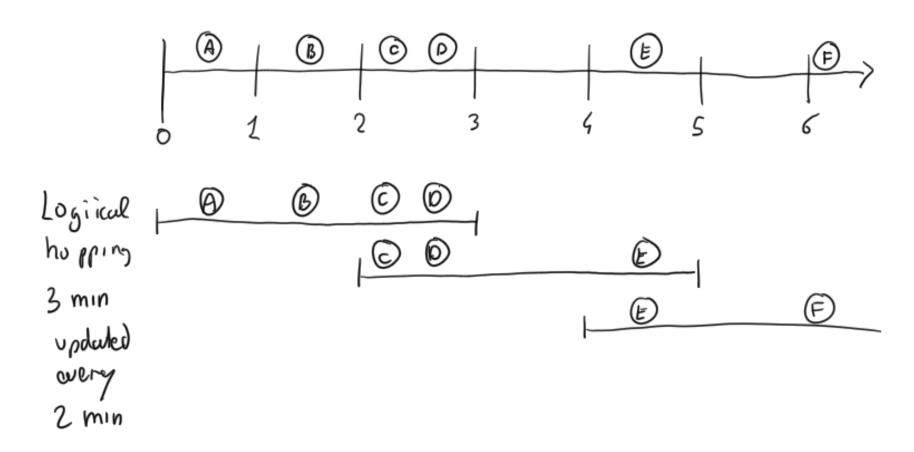
### EPL window syntax

Туре	Syntax	Description
Logical Sliding	win:time(time_period)	Sliding window that covers the specified time interval into the past
Logical Tumbling	win:time_batch(time_period [, reference point] [, flow control])	Tumbling window that batches events and releases them every specified time interval, with flow control options
Physical Sliding	win:length(size)	Sliding window that covers the specified number of elements into the past
Physical Tumbling	win:length_batch(size)	Tumbling window that batches events and releases them when a given minimum number of events has been collected

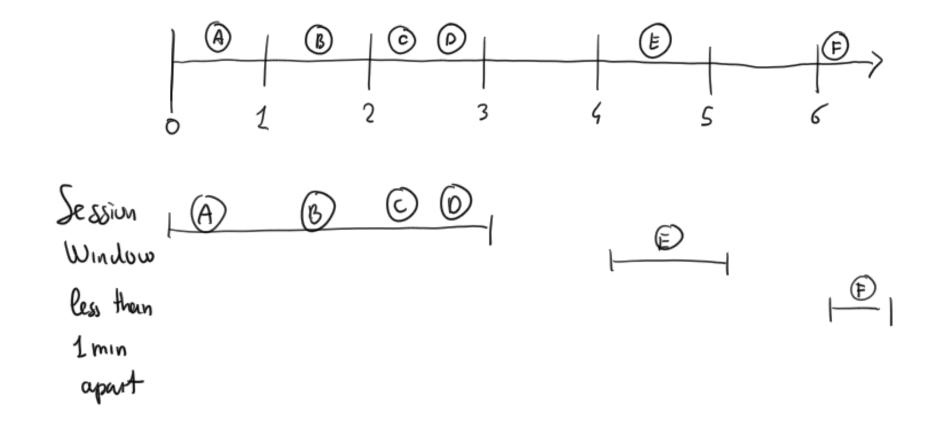
### Physical hopping windows



# Logical hopping windows



#### Session windows (not supported in EPL)



#### EPL Output control syntax

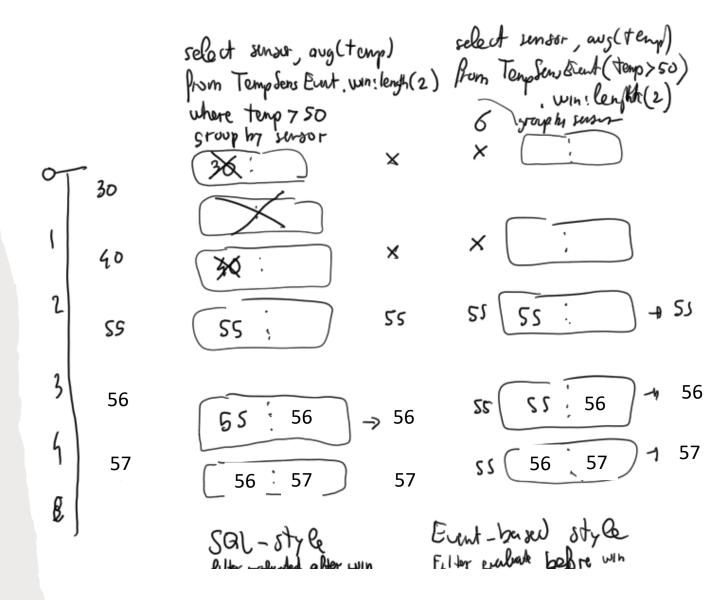
- The *output* clause is optional in Esper
- It is used to
  - Control the output rate
  - Suppress output events

```
output [all | first | last | snapshot]
everyoutput_rate [seconds | events]
```

#### Output control and Hopping Windows

Control advancement of sliding windows

A digression on data windowing and SQL-style vs. Eventbase style filtering



# The Pattern Matching clause of the Event Processing Language

#### Pattern matching

An event pattern emits when one or more event occurrences match the pattern definition, which can include

- Constraints on the content of events
- Constraints on the time of occurrence of events
- Conditions for pattern creation / termination

Content-based event selection

```
TempStream(sensor="S0", val>50)
```

Time-based event observers specify time intervals or time schedules

```
timer:interval(10 seconds)

timer:at(5, *, *, *, *)

Every 5 minutes
Syntax: minutes, hours, days of month, months, days of week
```

#### Pattern matching operators

- Logical operators
  - and, or, not
- Temporal operators that operate on event order
  - -> (*followed-by*)
- Creation/termination control
  - every, every-distinct, [num] and until
- Guards filter out events and cause termination
  - timer:within, timer:withinmax and while-expression

```
select a.sensor from pattern
[every (
   a = SmokeEvent(smoke=true)
   ->
   TempEvent(val>50, sensor=a.sensor)
   where timer:within(2 min)
)]
```

- every expr
  - When *expr* evaluates to true or false ...
  - ... the pattern matching for *expr* should re-start
- Without the every operator the pattern matching process does not re-start

This pattern fires when encountering an A event and then stops

A

 This pattern keeps firing when encountering A events, and does not stop

every A

A1 B1 B2 A2 A3 B3 A4 B4

every  $(A \rightarrow B)$ 

Detect an event A followed by an event B: at the time when B occurs, the pattern matches and restarts looking for the next A event

B1	{A1, B1}
В3	{A2, B3}
B4	{A4, B4}

A1 B1 B2 A2 A3 B3 A4 B4

every A -> B

The pattern fires for every A followed by a B event

B1	{A1, B1}
В3	{A2, B3}, {A3, B3}
B4	{A4, B4}

A1 B1 B2 A2 A3 B3 A4 B4

A -> every B

The pattern fires for an A event followed by every B event

B1	{A1, B1}
B2	{A1, B2}
В3	{A1, B3}
B4	{A1, B4}

A1 B1 B2 A2 A3 B3 A4 B4

every A -> every B The pattern fires for every A event followed by every B event

B1	{A1, B1}
B2	{A1, B2}
В3	{A1, B3}, {A2, B3}, {A3, B3}
B4	{A1, B4}, {A2, B4}, {A3, B4}, {A4, B4}

- With the every operator
  - Multiple (partial) instances of the same pattern can be active at the same time
  - Each instance can consume some resources when events enter the engine
- End pending instances whenever possible
  - With the *timer:within* construct
  - With the and not construct
- Note: the data windows on a pattern do not always limit pattern subexpression lifetime

A1 A2 B1

Pattern	Results
every A -> B	{A1, B1}, {A2, B1}
every A -> (B and not A)	{A2, B1}

The *and not* operator causes the sub-expression looking for {A1, B?} to end when A2 arrives

A1@1 A2@3 B1@4

Pattern	Results
every A -> B	{A1, B1}, {A2, B1}
every A -> (B where timer:within(2 sec))	{A2, B1}

The *timer:within* operator causes the sub-expression looking for {A1, B?} to end after 2 seconds

# Combine queries

 The insert into clause forwards events to other streams for further downstream processing

```
into FireComplexEvent
insert
select
            a.sensor as sensor,
            a.smoke as smoke,
            b.temperature as temperature
from
            pattern
            [every a=SmokeSensorEvent(smoke=true)
            ->
            b=TemperatureSensorEvent(
            sensor=a.sensor, temperature>50)];
select
            count(*)
            FireComplexEvent.win:time(10 min);
from
```

#### Reference

- Esper documentation online
  - https://esper.espertech.com/release-8.8.0/reference-esper/html\_single/
- Esper EPL online
  - <a href="http://esper-epl-tryout.appspot.com/epltryout/mainform.html">http://esper-epl-tryout.appspot.com/epltryout/mainform.html</a>

# License and Acknowledgment

 This work is licensed under the Creative Commons Attribution-ShareAlike International Public License

 The original slides where prepared by Alessandro Margara for the PhD course on "Stream and Complex Event Processing in the Big Data Era" offered in 2019

http://streamreasoning.org/events/scep2019

# Vertically Scalable solutions illustrated via EPL and Esper

Emanuele Della Valle

Politecnico di Milano