

Session typed Concurrent Programming (STCP)

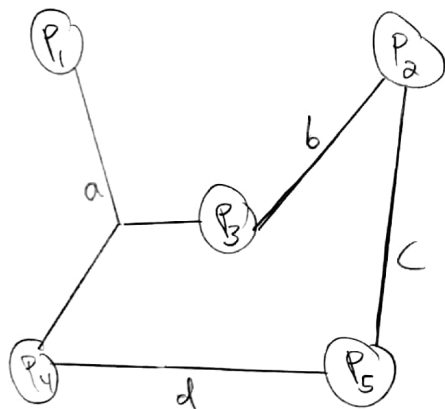
learning objectives:

- RoadMap:
- ① message-passing concurrent programming
 - ② session types
 - ③ Linear logic - Session types
 - ④ Manifest sharing

- ① how can we program using message-passing concurrent
- ② what session types are about
- ③ Benefits of linear logic to programming
- ④ How to accommodate sharing in a logically motivated way

Message-passing concurrency

processes that compute by exchanging messages along channels



channels are n -ary.

$$ar(a) = 3$$

$$ar(b) = ar(c) = ar(d) = 2$$

i.e., P_1 can send a message to either P_3 or P_4 or channel a .

thus, non-determinism.

a network of processes P_i connected by channels $\{a, b, c, d\}$.

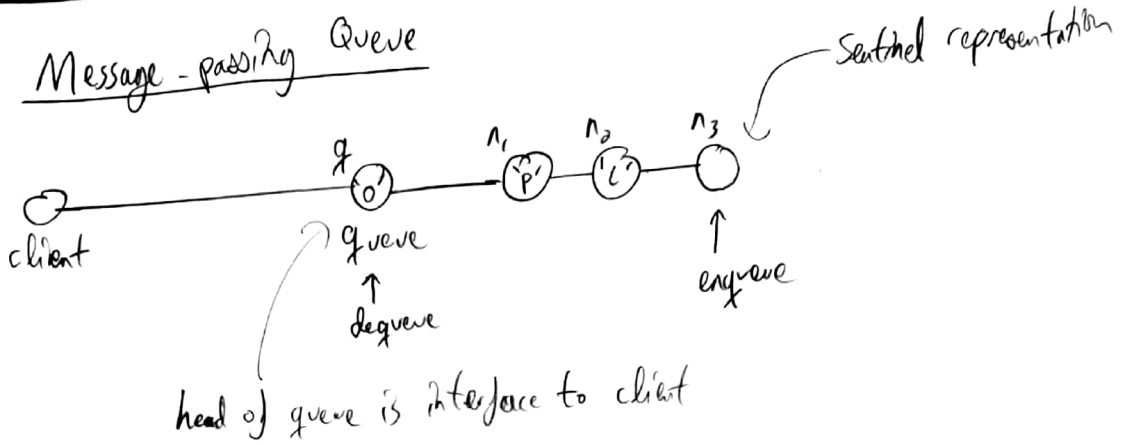
TT-calculus (Milner, 1992)

↳ "universal"

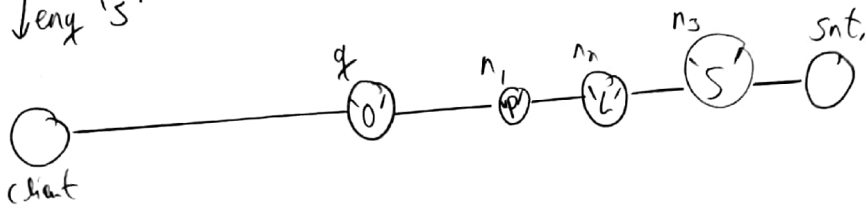
Functional Programming is to λ -calculus
as STCP are to process-calculus.

Message-passing Queue

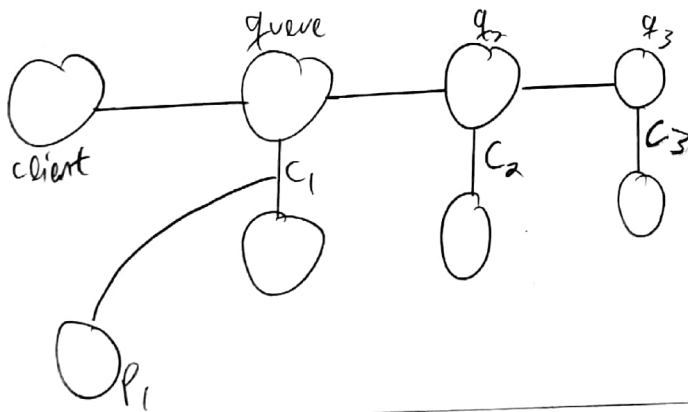
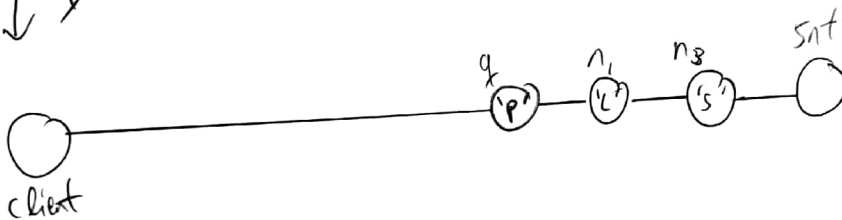
this is the process.



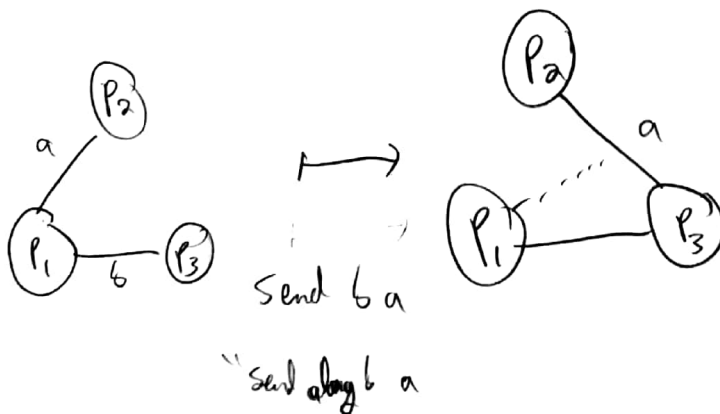
length '5'



this is the message



Channels passing other channels



"mobility"
"higher order channels"

How do you type message-passing protocols

- types for protocols of message-exchange

session types (Kohei Honda 1993)

$$A ::= ?[T].A' \mid ![T].A' \mid \overset{\text{"with"}}{\Delta} \{l_i:A_i, \dots, l_n:A_n\} \mid \overset{\text{"plus"}}{\oplus} \{l_i:A_i, \dots, l_n:A_n\} \mid \text{end} \mid x \mid Mx.A'$$

\uparrow \uparrow \uparrow \uparrow

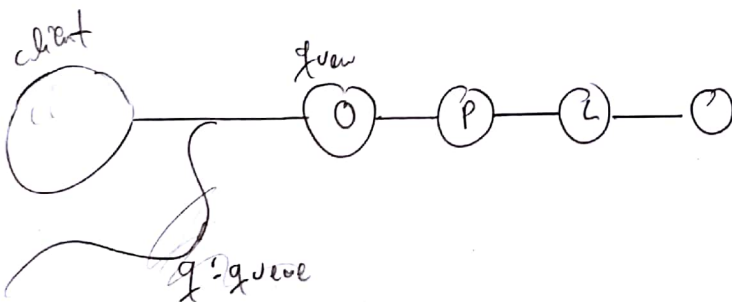
"I am willing to receive an input of type T, and after I will continue to be type A'"
 "I am willing to send ..."
 "choose between different types by label"
 give choice to process
 internal choice.

"give choice to client"
 external choice

$$T ::= A \mid \text{int} \mid \text{char} \mid \dots$$

a queue is an equi-recursive type

$$q_{\text{queue}} = \Delta \{ \text{eng} : ?[\text{char}]. q_{\text{queue}}, \text{dep} : \oplus \{ \text{none} : \text{end}, \text{some} : ![\text{char}]. q_{\text{queue}} \} \}$$



$q : \Delta \{ \text{eng} : \dots, \text{dep} : \dots \} \rightarrow$ "send 'eng' along q"
 $q : ?[\text{char}]. q_{\text{queue}} \rightarrow$ "send 's' along q"
 $q : q_{\text{queue}}$

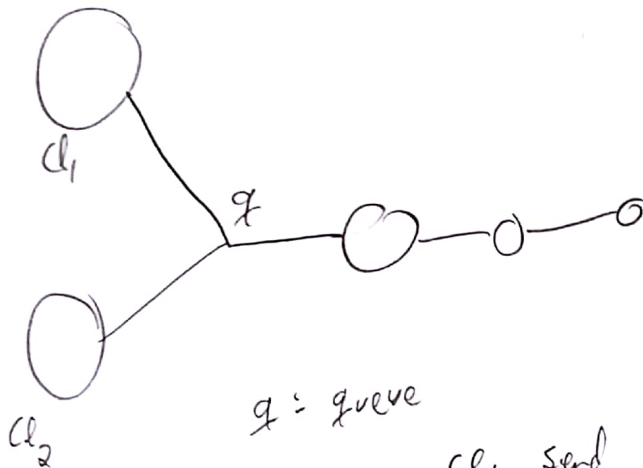
R stateful, types change over time

type of channel/process changes w message exchange.

Balzer 1.4

consider

two clients



$q = \text{queue}$

c_1 send msg along q

$q : ?[\text{char}], \text{queue}$

c_2 send msg along q

preservation in session types, called session fidelity, guarantees that expectation of client matches with the one of provider if they match initially

R channels are resources (from linear logic)

linear logic rejects weakening and contraction