

~~$\lambda x:\tau. \langle x, x \rangle : \tau \rightarrow \tau \otimes \tau$~~ \leftarrow rule set arbitrary copying

Linear λ -calculus

(Linear logic for PL)

Types

$\tau ::= \top \mid 0 \mid 1 \mid \tau_1 \rightarrow \tau_2 \mid !\tau \mid \tau_1 \& \tau_2 \mid \tau_1 \oplus \tau_2 \mid \tau_1 \otimes \tau_2$

$\Gamma; \Delta \vdash e : \tau$
 \nearrow nonlinear
 \uparrow linear

$$\frac{\Gamma, \Delta \vdash e : \tau \quad \Gamma; \Delta', x:\tau \vdash e' : \tau'}{\Gamma \Delta \Delta' \vdash \text{let } x = e \text{ in } e' : \tau'} \text{ (Cut)}$$

$$\frac{}{\Gamma; x:\tau \vdash x:\tau} \text{Id}^1 \quad \frac{}{\Gamma, x:\tau; \bullet \vdash x:\tau} \text{Id}^2$$

$$\frac{\Gamma; \Delta \vdash e : \tau_1}{\Gamma; \Delta \vdash l.e : \tau_1 \oplus \tau_2}$$

$$\frac{\Gamma; \Delta \vdash e : \tau_2}{\Gamma; \Delta \vdash r.e : \tau_1 \oplus \tau_2}$$

$$\frac{\Gamma; \Delta \vdash e : \tau_1 \oplus \tau_2 \quad \Gamma; \Delta', x:\tau_1 \vdash e_1 : \tau \quad \Gamma; \Delta', y:\tau_2 \vdash e_2 : \tau}{\Gamma; \Delta, \Delta' \vdash \text{case } e \text{ of } l.x \rightarrow e_1 \mid r.x \rightarrow e_2 : \tau}$$

$$\frac{\Gamma; \Delta_1 \vdash e_1 : \tau_1 \quad \Gamma; \Delta_2 \vdash e_2 : \tau_2}{\Gamma; \Delta_1, \Delta_2 \vdash \langle e_1, e_2 \rangle : \tau_1 \otimes \tau_2} (\otimes \text{I})$$

$$\frac{\Gamma; \Delta \vdash e : \tau_1 \otimes \tau_2 \quad \Gamma; \Delta', x : \tau_1, y : \tau_2 \vdash e' : \tau'}{\Gamma; \Delta, \Delta' \vdash \text{case } e \text{ of } \langle x, y \rangle \rightarrow e' : \tau'} (\otimes \text{E})$$

$$\frac{\Gamma; \Delta \vdash e : \tau_1 \& \tau_2}{\Gamma; \Delta \vdash e.l : \tau_1} (\& \text{E}_1) \quad \frac{\Gamma; \Delta \vdash e : \tau_1 \& \tau_2}{\Gamma; \Delta \vdash e.r : \tau_2} (\& \text{E}_2)$$

$$\frac{\Gamma; \Delta \vdash e_1 : \tau_1 \quad \Gamma; \Delta \vdash e_2 : \tau_2}{\Gamma; \Delta \vdash \{ l \rightarrow e_1, r \rightarrow e_2 \} : \tau_1 \& \tau_2} (\& \text{I})$$

(add to plus)

$$\frac{\Gamma; \Delta, x : \tau \vdash e : \tau'}{\Gamma; \Delta \vdash \lambda x : \tau. e : \tau \rightarrow \tau'} (\rightarrow \text{I})$$

$$\frac{\Gamma; \Delta \vdash e : \tau \rightarrow \tau' \quad \Gamma; \Delta' \vdash e' : \tau}{\Gamma; \Delta, \Delta' \vdash e e' : \tau'} (\rightarrow \text{E})$$

! "Of Course"

example: printer : ! (paper \otimes ink \rightarrow page)

$$\frac{\Gamma; \bullet \vdash e : \tau}{\Gamma; \bullet \vdash \text{many } e : !\tau} \text{ (many-I)}$$

$$\frac{\Gamma; \Delta \vdash e : !\tau \quad \Gamma, x : \tau; \Delta' \vdash e' : \tau'}{\Gamma; \Delta, \Delta' \vdash \text{case } e \text{ of many } x \rightarrow e' : \tau'} \text{ (}\leftarrow \text{oE)}$$

$$\frac{}{\Gamma; \Delta \vdash \{ \} : T}$$

$$\frac{}{\Gamma; \bullet \vdash \langle \rangle : 1}$$

$$\frac{\Gamma; \Delta \vdash e : 0}{\Gamma; \Delta \vdash \text{case } e \text{ of } \{ \} : \tau}$$

$$\frac{\Gamma; \Delta \vdash e : \perp \quad \Gamma; \Delta' \vdash e' : \tau}{\Gamma; \Delta, \Delta' \vdash \text{case } e \text{ of } \langle \rangle \Rightarrow e' : \tau}$$

Values

$$V ::= \langle \rangle \mid \langle v, v \rangle \mid l \cdot v \mid r \cdot v \mid \lambda x:\tau. e \mid \{l \Rightarrow e, r \Rightarrow e\} \\ \mid \xi \mid \text{many } v$$

Reduction (β) Rules

$$\text{let } x = v \text{ in } e \mapsto e[v/x]$$

$$(\lambda x. e) v \mapsto e[v/x]$$

$$\{l \Rightarrow e_1, r \Rightarrow e_2\}. l \mapsto e_1$$

$$\text{case } l \cdot v \text{ of } \{l \cdot x \Rightarrow e_1, r \cdot y \Rightarrow e_2\} \mapsto e_1[v/x]$$

$$\text{case } \{v_1, v_2\} \text{ of } \langle x, y \rangle \mapsto e[v_1/x, v_2/y]$$

$$\text{Case many } v \text{ of } x \Rightarrow e \mapsto e[v/x]$$

Lemma (Substitution)

◦ linear: If $\Gamma; \Delta \vdash e : \tau$ and $\Gamma; \Delta', x:\tau \vdash e' : \tau'$
then $\Gamma; \Delta, \Delta' \vdash e'[e/x] : \tau'$.

◦ nonlinear: If $\Gamma; \bullet \vdash e : \tau$ and $\Gamma, x:\tau; \Delta' \vdash e' : \tau'$
then $\Gamma; \Delta' \vdash e'[e/x] : \tau'$.