

Cost Model based on the  $\lambda$ -Calculus  
or  
The Church Calculus the Other Turing Machine

Guy Blelloch  
Carnegie Mellon University

July 9, 2018

## 1 Parallelism vs. Concurrency

### 1.1 Parallelism

Provides performance

Goal: Make parallel algorithms map sequential algorithms (no difference)

Don't change the semantics, just integrate a cost model.

## 2 $\lambda$ -Calculus with Cost Model

Church/Turing Hypothesis:  $(\lambda x.e_1)e_2 \Rightarrow_\beta e_1[e_2/x]$

### 2.1 Machine Models and Simulation

#### 2.1.1 Handbook of Theoretical Computer Science

Chapter 1: Machine Models and Simulations [Peter van Emde Boas]

In order to measure time and space, then need to specify which notions of time and space are meant.

Can use abstract machines to do this (i.e.: RAM Machine)  $\leftarrow$  General convection. We want to use a language-based cost model (okay because abstract).

However, the arbitrariness of model choice remains. Need a way to translate between the machines to compare **mutual simulations**.

Example: with a machine with unbounded registers can run exponential numbers of calculations simultaneously ( $\implies P = NP$ ). But this is not allowed, bits must be limited.

#### 2.1.2 Machine Models

Random Access Machines:

- SRAM - succ, pred
- RAM - add, sub
- MRAM - add, sub, mult
- LRAM - log length words
- RAM-L - cost of instruction is word length
- SMM, KUM, pure, impure

Parallel Machine Models:

- Circuit models
- PSPACE
- TM with alternation
- Vector models
- PRAM - EREW, CREW, CRCW (priority, arbitrary, ...)
- SIMDAG
- k-PRAM, MIND-RAM, PTM

### 2.1.3 The two parts

Part 1: The model

Well defined semantics

simple

close to programming paradigm

Part 2: Simulation

mapping of costs

good bounds when simulated on realistic machines

### 2.1.4 Language-Based Cost Model

A cost model base on a "cost semantics" instead of a machine

Why use the  $\lambda$ -calculus? Historically the first model, clean, well understood

What costs?

Advantages:

- naturally parallel (turing model not parallel)
- more elegant
- model is closer to code and algorithms

- closer - in terms of simulation costs - to "practical" machine models such as the RAM

Disadvantages:

- 50 years of history

### 2.1.5 CBV $\lambda$ -Calculus

$$e = x|(e_1, e_2)|\lambda x.e \quad (\text{expression})$$

Lambda Calculus Sequential without Cost:

$$e \Downarrow v \quad (\text{relation})$$

$$\lambda x.e \Downarrow \lambda x.e \quad (\lambda)$$

$$\frac{e_1 \Downarrow \lambda x.e \quad e_2 \Downarrow v \quad e[v/x] \Downarrow v'}{(e_1 e_2) \Downarrow v'} \quad (\text{App})$$

Lambda calculus parallel with cost:

$$e \Downarrow v; w, d \quad (\text{relation-par-cost})$$

$$\lambda x.e \Downarrow \lambda x.e; 1, 1 \quad (\lambda\text{-par-cost})$$

$$\frac{e_1 \Downarrow \lambda x.e; w_1, d_1 \quad e_2 \Downarrow v; w_2, d_2 \quad e[v/x] \Downarrow v'; w_3, d_3}{(e_1 e_2) \Downarrow v'; 1 + w_1 + w_2 + w_3, 1 + \max(d_1, d_2) + d_3} \quad (\text{App-par-cost})$$

### 2.1.6 Simulation (sequential)

CEK Machine: a state transition system

$$(C, E, K) \Rightarrow (C', E', K')$$

"control"  $C := e_1 e_2 | \lambda x.e | x$

"environment"  $E := x \rightarrow v$

"continuation"  $K :=$

Everything is constant time except perhaps look-up and insert.