

University of Bayreuth Institute for Computer Science

Bachelor Thesis

in Applied Computer Science

Topic: Randomly Generated CFGs, The CYK-Algorithm

And A GUI

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To my parents.

Abstract

The abstract of this thesis will be found here.

Zusammenfassung

 ${\it Hier steht die Zusammenfassung dieser Bachelorarbeit}.$

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1 Introduction

2 Technology Overview

3 Algorithms

Algorithm descriptions here.
RightHandSideElements (rhse) are terminals like "a, b, ..." and compound variables like "AB, AC, ...".

LeftHandSideElements (lhse) are variables like "A, B, ...".

setVAdvanced = Set<VariableKWrapper>[][] setV

setVSimple = Set<Variable>[][] setV

```
Algorithm: distributeDiceRollRightHandSideElements
2 Input: grammar, setRhse, minCount, maxCount, listVars;
3 Output: grammar;
  for (RightHandSideElement rhse : setRhse) {
          // countWillBeAdded is between [minCount, maxCount].
          countWillBeAdded = diceRoll();
          while (list Vars. size () > countWillBeAdded) {
                   Remove one variable of listVars via dice roll;
10
          for(Variable varLeft : listVars){
11
                   // An exception is thrown if the production
12
                   // already exists.
13
                   grammar.addProduction "varLeft --> rhse";
14
15
16
17 return grammar;
```

Listing 1: distributeDiceRollRightHandSideElements

```
Algorithm: CYK.calculateSetVAdvanced
Input: grammar, word;
Output: Set<VariableK > [][] cYKMatrix;

Set<VariableK > [][] cYKMatrix = new Set<VariableK > [wordSize][wordSize];
cYKMatrix = calculateCYKMatrix;
return cYKMatrix;
```

Listing 2: CYK.calculateSetVAdvanced

```
1 Algorithm: GeneratorGrammarDiceRollTopDownMartens
2 Input: word, settings;
3 Output: grammar;
4 Note: Keep in mind that the setV matrix is a upper right matrix. But the
5 description of how the algorithm works is done, as if the setV pyramid
6 points downwards (reflection on the diagonal + rotation to the left).
7 Regarding one cell, its upper left cell and its upper right cell
s | are looked at. setV[i][j] = down cell. setV[i + 1][j] = upper right cell
9 | \text{setV}[i][j-1] = \text{upper left cell.} With wordSize = 5, the visited indexes
  are as following: [01->12->23->34; 02->13->24; 03->14; 04;]
10
11
_{12}|\operatorname{Grammar} \operatorname{grammar} = \operatorname{new} \operatorname{Grammar}();
  // Part1: Distribute the terminals.
14 grammar = distributeDiceRollRightHandSideElements(
           grammar, settings Terminals, minCount Terminals,
15
           maxCountTerminals, settingsListVariables);
16
17 // Part2: Distribute the compound variables.
18 Set < Variable K Wrapper > [] [] set V Advanced;
19 // Fill the diagonal of the matrix with:
  setVAdvanced = CYK.calculateSetVAdvanced(grammar, word);
  for(Cell cellToBeVisited : setVAdvanced){
21
           setVariableCompound = calculate all the possible tupels of
22
                    ({varLeft}, {varRight});
23
           for(VariableCompound varComp : setVariableCompound){
24
                    // Because of dice rolling anyways and lots of grammars
25
                    // being generated, no varComp is added if the production
26
                    // already exists.
27
                    grammar = distributeDiceRollRightHandSideElements(
28
                            grammar, varComp, minCountVarComp,
29
                            maxCountVarComp, settingsListVariables);
30
                    setVAdvanced = CYK.calculateSetVAdvanced(grammar, word);
31
           }
32
  }
33
34 return grammar;
```

Listing 3: GeneratorGrammarDiceRollTopDownMartens

```
Algorithm: GeneratorGrammarDiceRollOnly
2 Input: settings;
3 Output: grammar;
4 Note: A lot of productions are generated, that later on are not needed
5 for parsing the specific word.
7 Grammar grammar = new Grammar();
8 // Part1: Distribute the terminals.
  grammar = distributeDiceRollRightHandSideElements(
          grammar, settingsTerminals, minCountTerminals,
10
          maxCountTerminals, settingsListVariables);
11
12 // Part2: Distribute the compound variables.
13 Set < Variables > vars = settings.get Variables();
14 Set < Variables Compound > set Var Comp;
setVarComp = calculate all the possible tupels of ({vars}, {vars});
16 grammar = distributeDiceRollRightHandSideElements(
          grammar, settingsTerminals, minCountVariableCompound,
17
          maxCountVariableCompound, setVarComp);
18
19 return grammar
```

Listing 4: GeneratorGrammarDiceRollOnly

```
1 Algorithm: GeneratorGrammarDiceRollOnlyBias
2 Input: settings;
3 Output: grammar;
4 Note: A lot of productions are generated, that later on are not needed
5 for parsing the specific word.
 Grammar grammar = \mathbf{new} Grammar ();
8 // Distribute the terminals.
  grammar = distributeDiceRollRightHandSideElementsBias(
          grammar, settingsTerminals, settingsMinCountTerminals,
10
          settingsMCountTerminals, settingsListVars, settingsFavouritism);
11
12
  // Distribute the compound variables.
14 Set < Variables Compound > set Var Comp;
  setVarComp = calculate all the possible tupels of ({vars}, {vars});
16 grammar = distributeDiceRollRightHandSideElementsBias (
          grammar, varComp, settingsMinCountVars,
17
          settingsMaxCountVars, settingsListVars, settingsFavouritism);
18
19 return grammar;
```

Listing 5: GeneratorGrammarDiceRollOnlyBias

```
1 Algorithm: distributeDiceRollRightHandSideElementsBias
2 Input: grammar, setRhse, minCount, maxCount, listVars, favouritismList;
3 Output: grammar;
4 Note: Because of dice rolling anyways and lots of grammars being
  generated, no rhse is added if the production already exists.
7 // Calculate the bloated varSet.
  Set < Variable > vars Bloated;
9 for (Variables varTemp : settings.getVariables()) {
          tempFavour = randomly pick favouritism[i];
10
          varsBloated.add({tempFavour times varTemp});
11
          favouritism.remove(tempFavour);
12
13 }
  // Because of dice rolling anyways and lot of grammars being generated,
15 // just no rhse is added if the production already exists.
grammar = distributeDiceRollRightHandSideElements ( grammar,
          varsBloated , minCount , maxCount , listVars );
17
18 return grammar;
```

Listing 6: distributeDiceRollRightHandSideElementsBias

Description of the checks here.

All test of the GrammarValidityChecker class are based on the simple setV matrix.

isValid = isWordProducible && isExamConstraints && isGrammarRestrictions

isWordProducible = CYK.algorithmAdvanced()

is Exam Constraints = is Right Cell Combinations Forced && is Max Sum Of Productions Count && is Max Sum Of Vars In Pyramid Count && count Right Cell Combinations Forced

 $is Grammar Restrictions = is Size Of Word Count \ \&\& \ is Max Number Of Vars Per Cell Count$

```
Algorithm: checksumOfProductions
Input: grammar, maxSumOfProduction;
Output: isSumOfProductions;

return grammar.getProductionsAsList().size() <= maxSumOfProductions;
```

Listing 7: checksumOfProductions

```
Algorithm: checkMaxNumberOfVarsPerCell
2 Input: setVSimple, maxNumberOfVarsPerCell;
3 Output: isMaxNumberOfVarsPerCell;
4 Note: Checking for maxNumberOfVarsPerCell <= zero isn't allowed;
6 int tempMaxNumberOfVarsPerCell = 0;
  int wordLength = tempSetV[0].length;
  for ( int i = 0; i < wordLength; i++ ) {</pre>
          for ( int j = 0; j < wordLength; j++ ) {
                  if ( tempSetV[i][j].size() > numberOfVarsPerCell ) {
                           numberOfVarsPerCell = tempSetV[i][j].size();
11
                  }
12
          }
13
15 return tempMaxNumberOfVarsPerCell <= maxNumberOfVarsPerCell;</p>
```

Listing 8: checkMaxNumberOfVarsPerCell

```
Algorithm: checkMaxSumOfVarsInPyramid
Input: setVSimple, maxSumOfVarsInPyramid;
Output: isMaxSumOfVarsInPyramid;

// put all vars of the matrix into one list and use its length.
List<Varaible> allVarsList = new ArrayList<>();
for ( int i = 0; i < setVSimple.length; i++ ) {
            for ( int j = 0; j < setVSimple.length; j++ ) {
                tempVars.addAll( setVSimple[i][j] );
            }

return allVarsList.size() <= maxSumOfVarsInPyramid;
```

Listing 9: checkMaxSumOfVarsInPyramid

```
Algorithm: rightCellCombinationsForced
2 Input: setVSimple, minCountForced, grammar;
3 Output: isForced, countForced, setVSimpleVarsThatForce;
4 Note: Keep in mind that the setV matrix is a upper right matrix. But the
5 description of how the algorithm works is done, as if the setV pyramid
6 points downwards (reflection on the diagonal + rotation to the left).
7 Regarding one cell, its upper left cell and its upper right cell
s | are looked at. setV[i][j] = down cell. setV[i + 1][j] = upper right cell
s | setV[i][j-1] = upper left cell.
10
11 int countForced = 0;
12 Set < Variable > [] [] set VMarked Vars That Force;
  for(Cell cell : setVSimple){
13
          // Trivial cases that would fulfil the restrictions each time.
14
          Ignore the upper two rows of the pyramid;
15
          isRightCellCombinationForced = true;
16
          if(!upperLeftCell.isEmpty() && !upperRightCell.isEmpty()) {
17
                   break;
18
19
          setVariableCompound = calculate all the possible tupels of
20
                   ({varLeft}, {varRight});
21
          for(Variable var : cellToBeVisited) {
22
                   varDownProdList = grammar.getProdList(varDown);
                   for(VariableCompound varComp : setVariableCompound) {
24
                           for(Production prod : varDownProdList){
25
                                    if(prod.getRhse() == varComp) {
26
                                            isForced = false;
27
                                    }
28
                           }
29
30
                   if(isRightCellCombinationForced) {
31
                           rightCellCombinationsForced++;
32
                           // Cell has index i and j.
33
                           setVMarkedVarsThatForce[i][j].add(var)
34
                   }
35
          }
36
37
  boolean isForced = countForced >= minCountRightCellCombinationsForced;
38
return is Forced, countForced, setVMarkedVarsThatForce;
```

Listing 10: rightCellCombinationsForced

```
Algorithm: Util.removeUselessProductions
Input: grammar, setVSimple, word
Output: grammar
Note: Very similar to the calculateSetVAdvanced algorithm. Additionally
to storing the k, it is also saved, which production have been used.
All productions that haven't been need are removed, from the grammar.

Set<Production> allProductions = grammar.getProductions();
Set<Production> onlyUsefulProductions;
onlyUsefulProductions = calculate useful productions with the input of
grammar, setVSimple and word;
grammar.remove(allProductions);
return grammar.add(onlyUsefulProductions);
```

Listing 11: Util.removeUselessProductions

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4 Course of Action

The informal goal is to find a suitable combination of a grammar and a word that meets the demands of an exam exercise. Also the CYK pyramid and one derivation tree of the word must be generated automatically as a "solution picture".

Firstly the exam exercise must have an upper limit of variables per cell while computing the CYK-pyramid.

Secondly the exam exercise must have one or more "special properties" so that it can be checked if the students have clearly understood the algorithm, e.g. "Excluding the possibility of luck."

The more formal goal is identify and determine parameters that in general can be used to define the properties of a grammar, so that the demanded restrictions are met. Also parameters could be identified for words, but "which is less likely to contribute, than the parameters of the grammar." [Second appointment with Martens]

Some introductory stuff:

Possible basic approaches for getting these parameters are the Rejection Sampling method and the "Tina+Wim" method.

Also backtracking plays some role, but right now I don't know where to put it. Backtracking is underapproach to Rejection Sampling.

Note: Starting with one half of a word and one half of a grammar.

Identify restrictions (=parameters) regarding the grammar.

Maybe find restrictions regarding the words, too.

Procedures for automated generation. Each generation procedure considers different restrictions and restriction combinations. The restrictions within one generation procedure can be optimized on its own. Up till now:

Generating grammars: DiceRolling, ...

Generating words: DiceRolling, ...

Parameter optimisation via theoretical and/or benchmarking approach.

Benchmarking = generate N grammars and test them, (N=100000).

Define a success rate and try to increase it.

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The overall strategy is as following:

1.) Identify theoretically a restriction/parameter for the grammar. Think about the influence it will have. Think also about correlations between the restrictions.

2.) Validate the theoretical conclusion with the benchmark. Test out the influence of this parameter upon the success rate. Try different parameter settings.

The ordering of step 1 and step 2 can be changed.

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Used software

Used Software 22

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Listings

Following are some interesting classes referenced in the thesis that were too long to fit into the text.

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Tables

This section contains all tables referenced in this thesis.

Tables 26

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References

[1] JSR 220: Enterprise Java Beans 3.0 https://jcp.org/en/jsr/detail?id=220, 09/09/2015