

Chapter 1: Exercises

Exercise 1: caches

Given: a L_1 cache of 32 Kbyte, 8-way set associative, cache line size of 64 Byte. Consider a matrix-vector multiplication $A * x$ (algorithm A) of a **256 x 256** matrix using double precision elements. The matrix is stored in column-major format. The matrix is not preloaded in cache memory.

What is the L_1 cache read miss rate + explain.

Exercise 2: caches

Given: a L_1 cache of 32 Kbyte, 8-way set associative, cache line size of 64 Byte. Consider a matrix-vector multiplication $A * x$ (algorithm A) of a **256 x 255** matrix using double precision elements. The matrix is stored in column-major format. The matrix is not preloaded in cache memory.

What is the L_1 cache read miss rate + explain.

Exercise 3: caches

Given: a L_1 cache of 32 Kbyte, 8-way set associative, cache line size of 64 Byte. Consider a matrix-vector multiplication $A * x$ (algorithm A) of a **255 x 256** matrix using double precision elements. The matrix is stored in column-major format. The matrix is not preloaded in cache memory.

What is the L_1 cache read miss rate + explain.

Exercise 4: caches

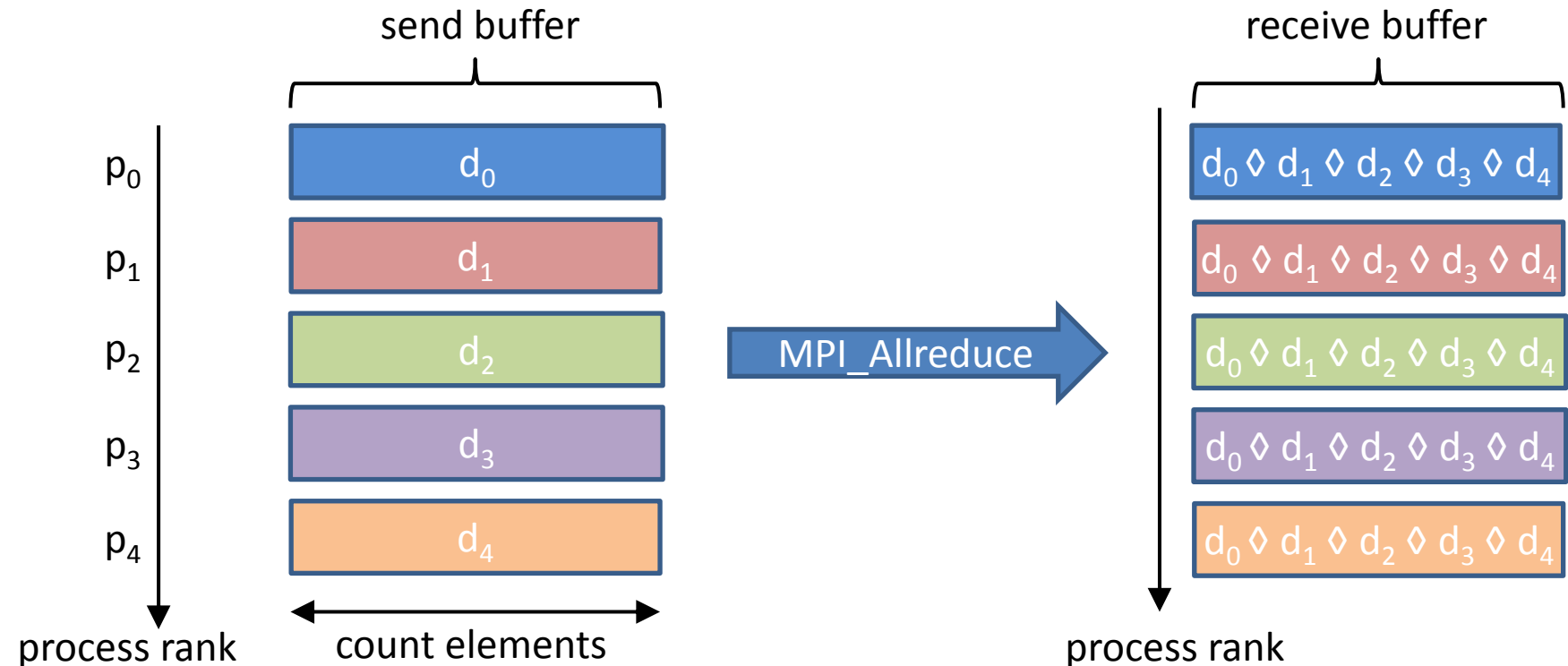
Given: a L_1 cache of 32 Kbyte, 8-way set associative, cache line size of 64 Byte. Consider a matrix-vector multiplication $A * x$ (algorithm A) of a **64 x 64** matrix using double precision elements. The matrix is stored in column-major format. The matrix is not preloaded in cache memory.

What is the L_1 cache read miss rate + explain.

Chapter 2: Exercises

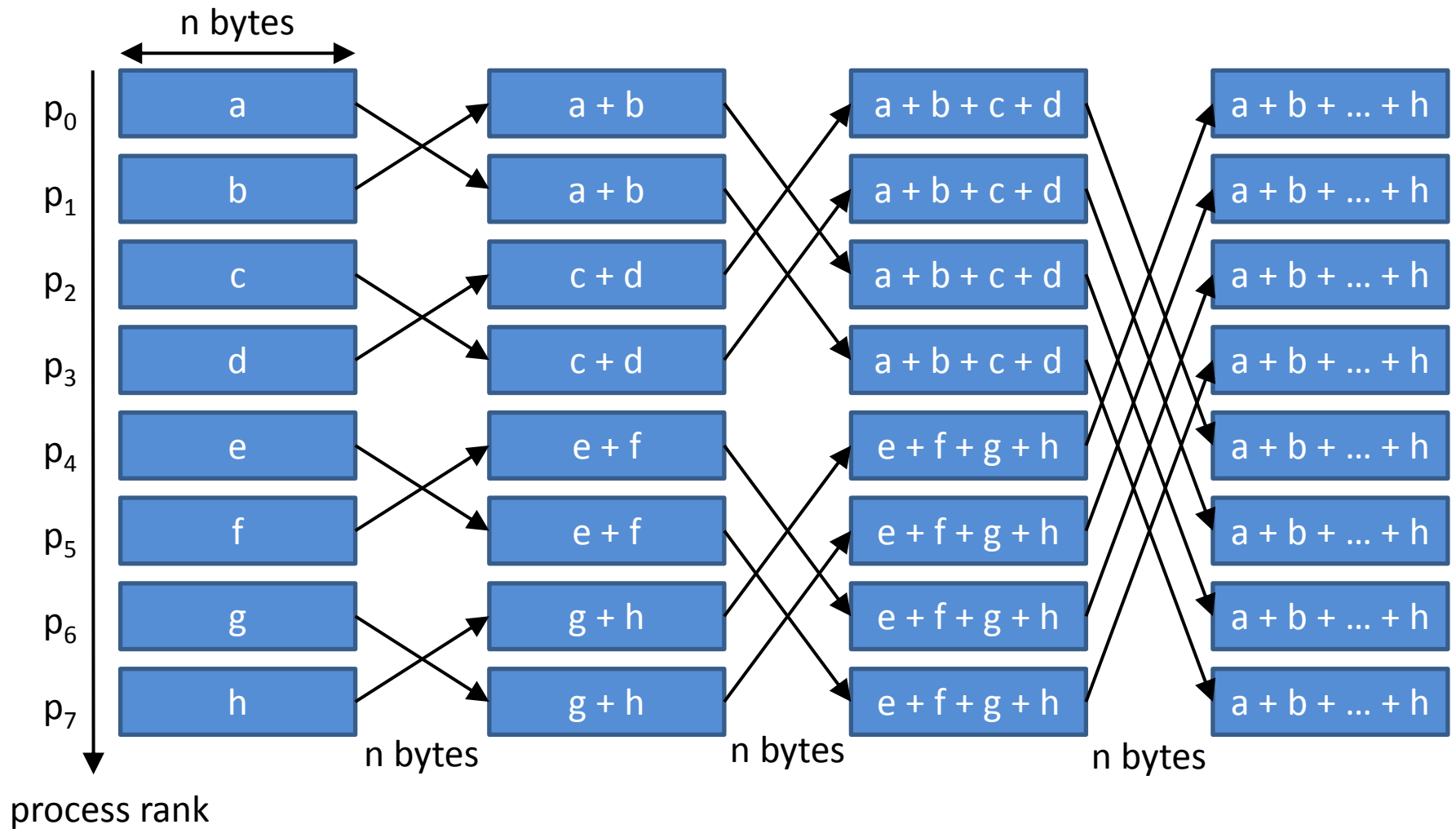
Allreduce algorithm

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MPI_Allreduce(void *sendbuf, void *recvbuf, int count,  
               MPI_Datatype dataType, MPI_Op op, MPI_Comm comm)
```



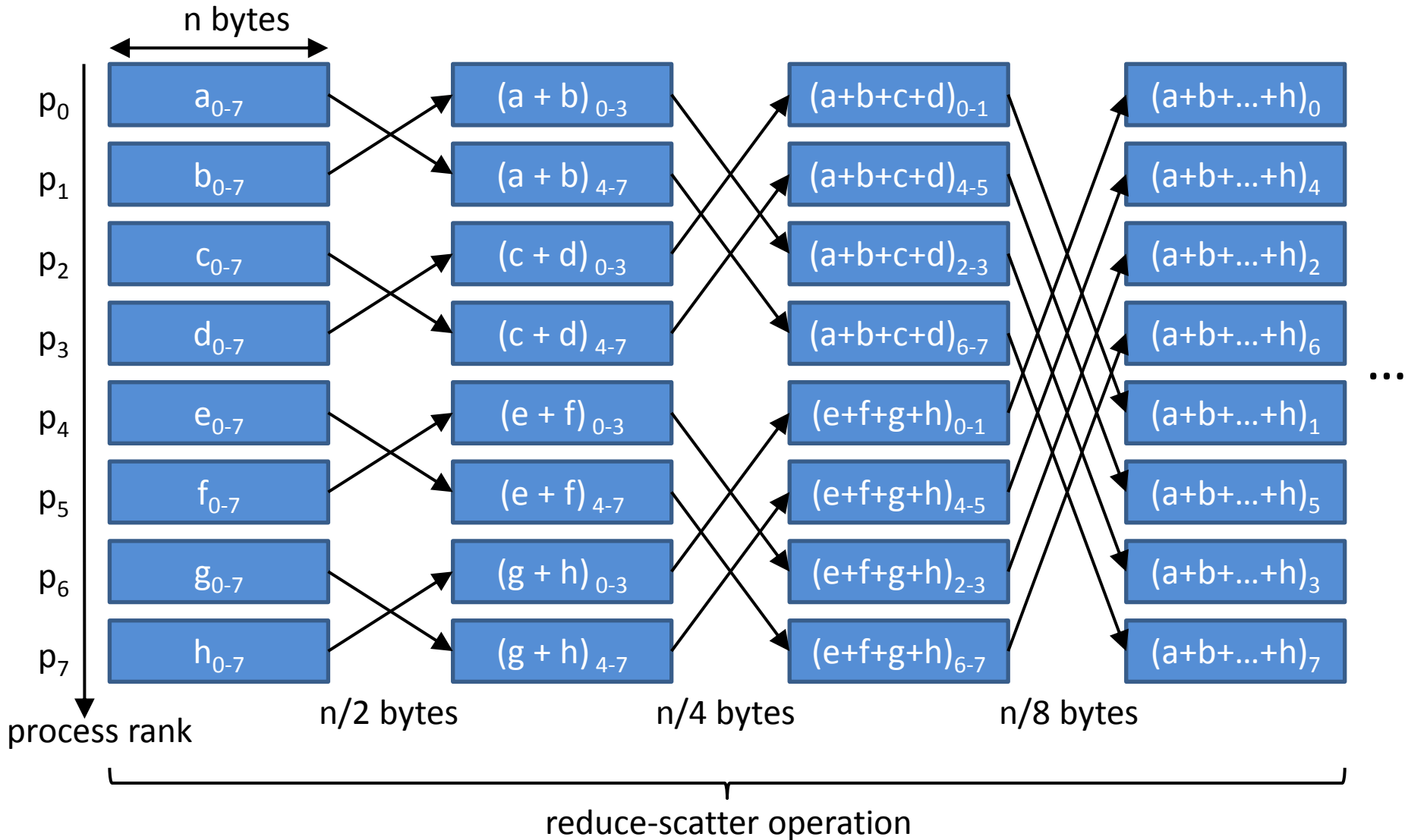
Q: Develop an algorithm for the allreduce routine. Derive the time complexity for that algorithm. Assume a fully non-blocking (duplex) network.

Solution 1: butterfly communication scheme



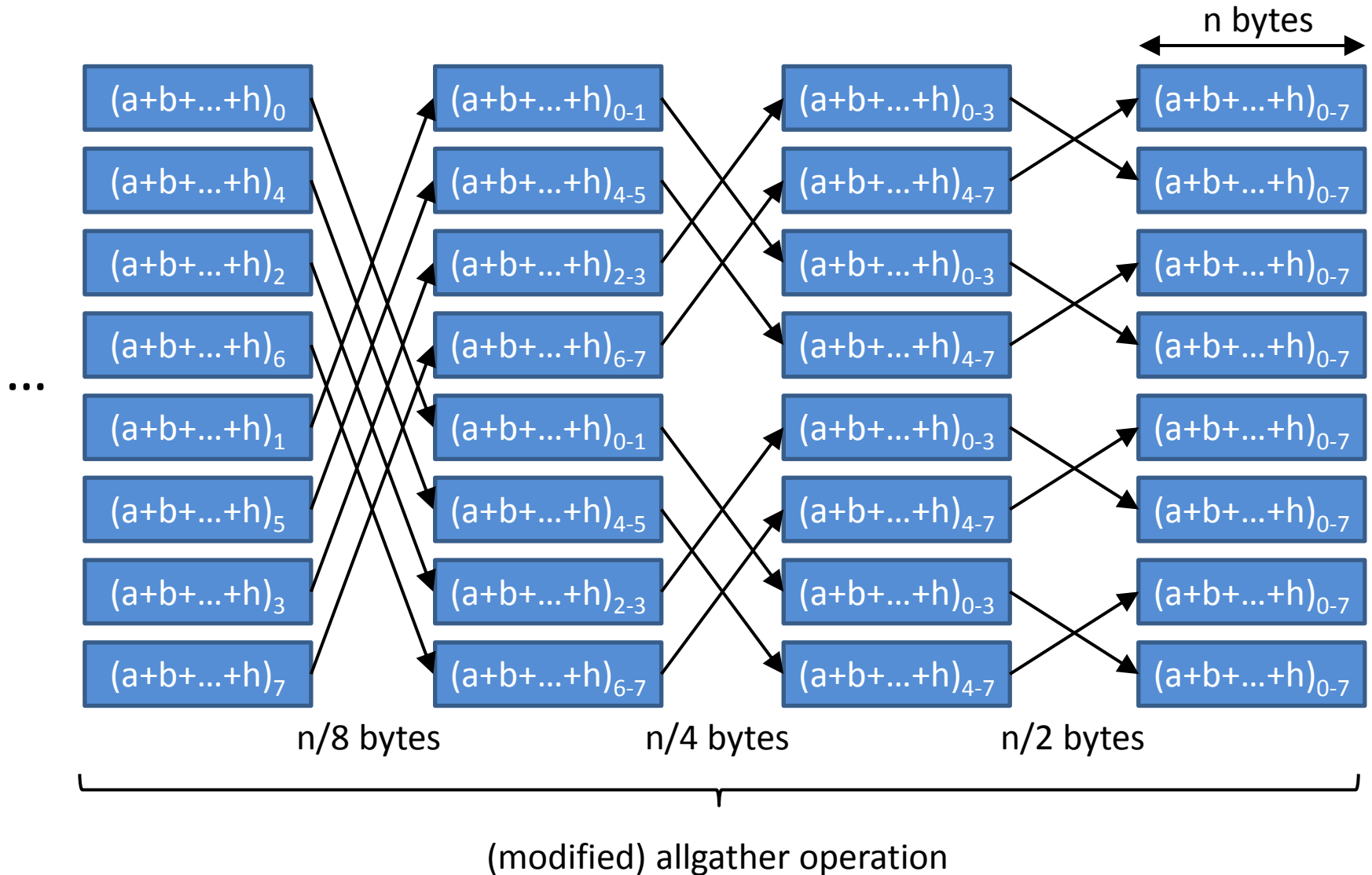
$$T_{\text{sol1}} = (\alpha + \beta n + \gamma n) \log_2 P \quad (\text{assume } P = 2^i)$$

Solution 2: reduce-scatter + allgather



$$T_{\text{sol2}} = \alpha \log_2 P + (\beta n + \gamma n)(1 - 1/P) + \dots$$

Solution 2: reduce-scatter + allgather (cont'd)



$$T_{\text{sol2}} = 2\alpha \log_2 P + (2\beta n + \gamma n)(1 - 1/P) \quad (\text{assume } P = 2^i)$$

Exercise 2: Sendrecv

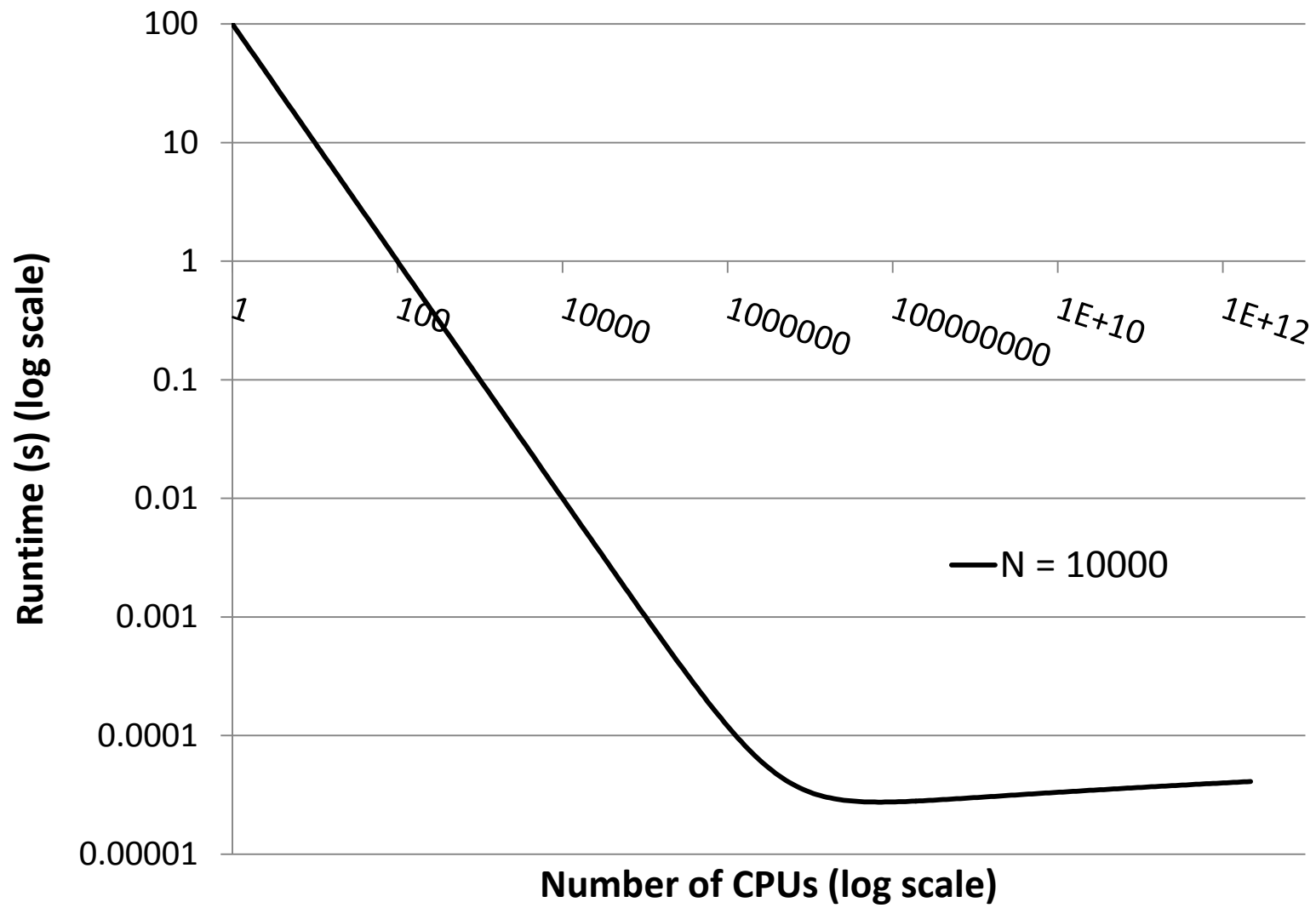
How would you implement **MPI_Sendrecv** so that it will never deadlock? What other benefits might you expect compared to using a blocking send / receive ?

Exercise 3: Speedup and efficiency

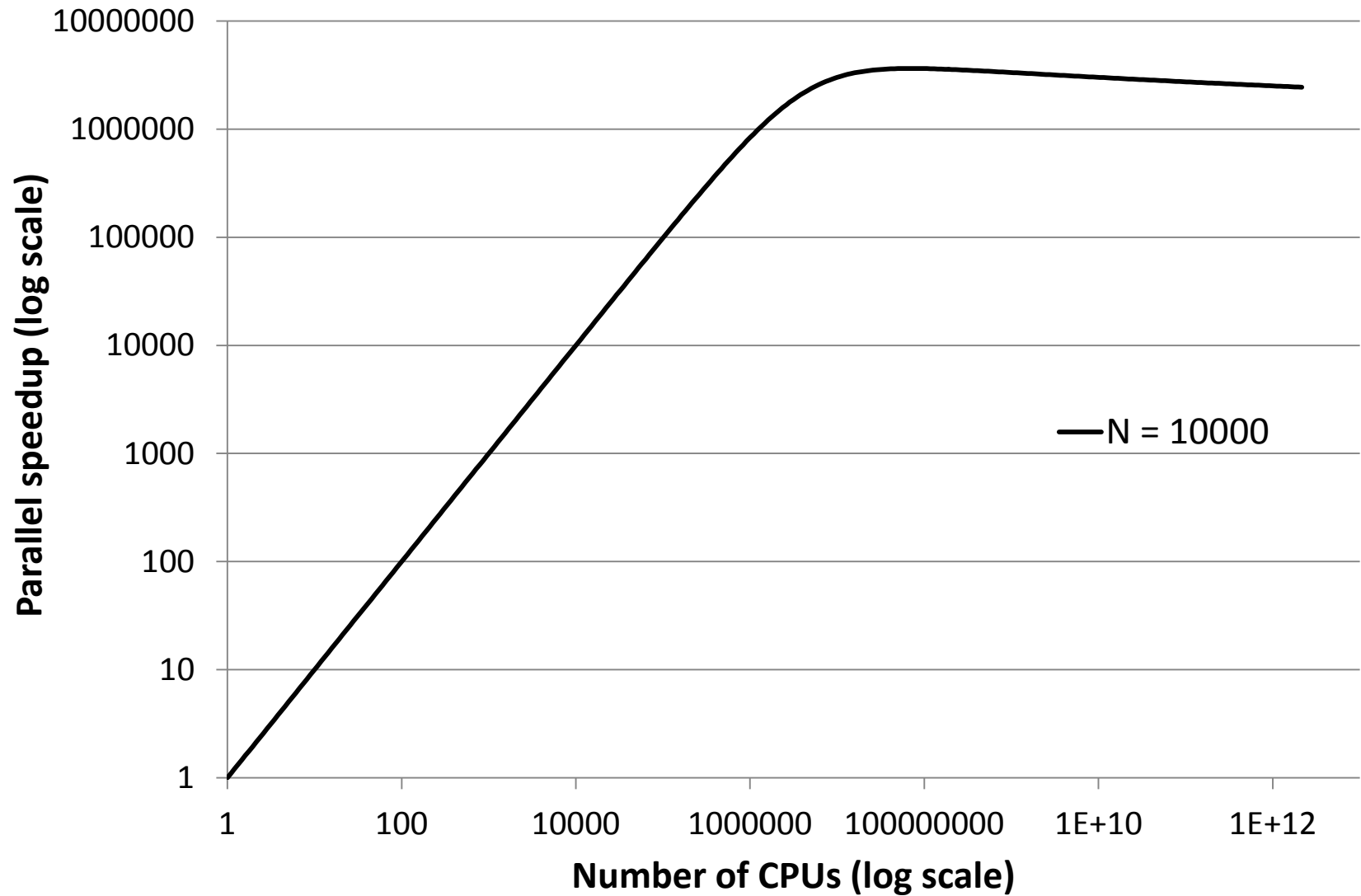
Given: $T_1(N) = N^2$; $T_p(N) = N^2/P + \log_2 P$

- What is the speedup ?
- What is the efficiency ?
- For fixed N , what is the limit for large P ?
- For fixed P , what is the limit for large N ?

Parallel runtime



Parallel speedup



Exercise 4: Isoefficiency

Given: $T_1(N) = N$; $T_p(N) = N/P + \log_2 P$

If we increase the P by a factor of k, by what factor do we need to increase N in order to maintain constant efficiency?

Exercise 5: Amdahl's law

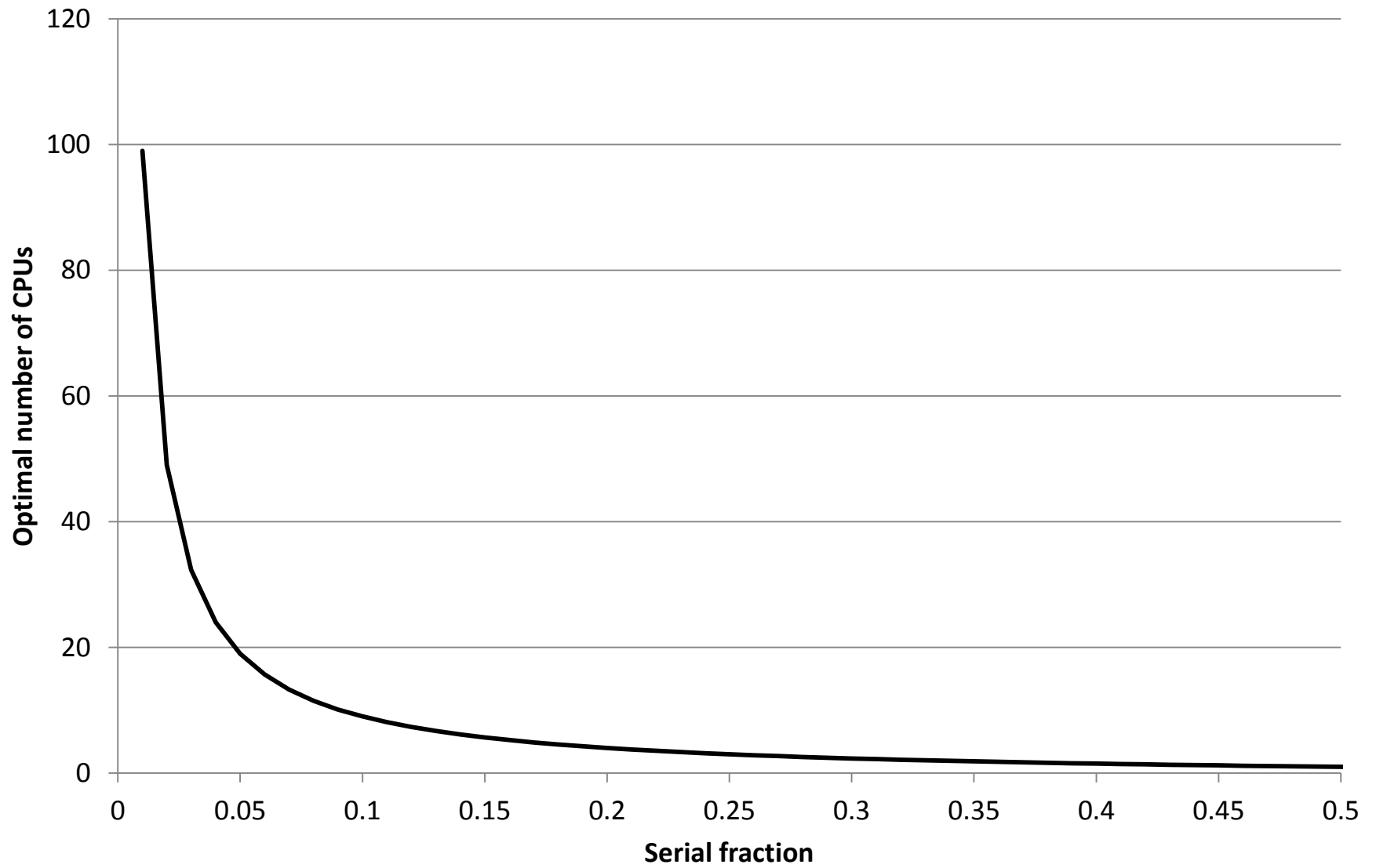
The financial cost C_f for renting compute power on a cluster is proportional to the product of the wall clock time T_p and the number of CPU cores allocated, i.e. $C_f = P \cdot T_p$

Derive an expression for the optimum number of CPUs (P_{opt}) that minimizes the cost function $C = C_f \cdot T_p$
Assume that T_p can be expressed by Amdahl's law.

What is the optimum P for $s = 1\%$ and $s = 5\%$?

What about efficiency?

Optimal number of CPUs

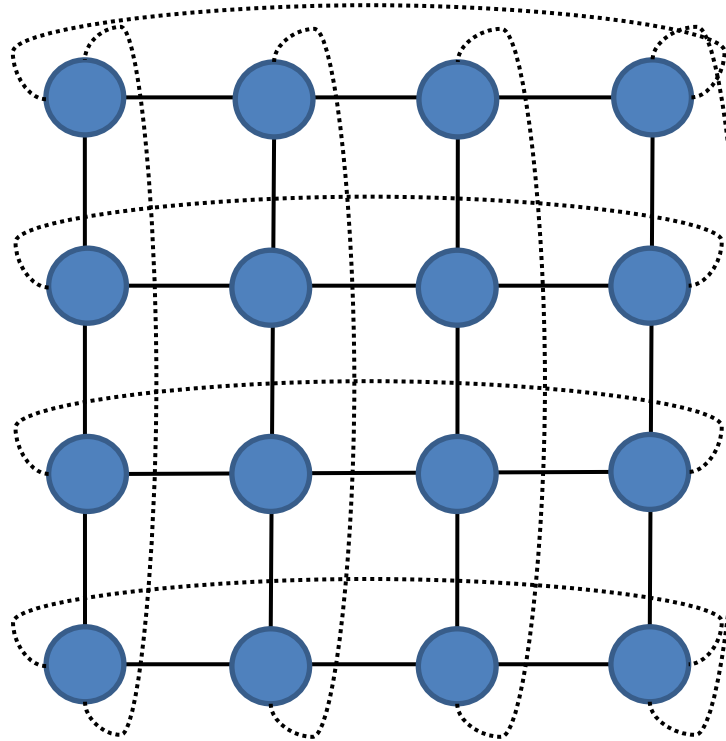


Exercise 6: Message Passing Interface

Write an MPI program that outputs data
(e.g. writing data to disc) in rank order (pseudocode)

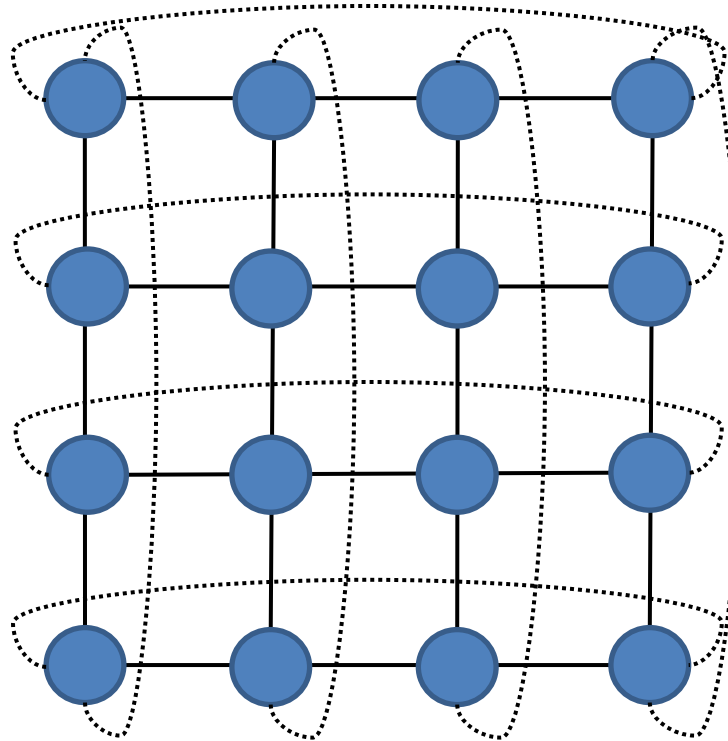
Exercise 7: Mesh networks

Given a square mesh network organized as $\sqrt{P} \times \sqrt{P}$ machines
Each machine can communicate to one of its four neighbors



What is the bisection bandwidth? What is the diameter?

Exercise 8: Mesh networks



Same mesh network. Derive a broadcast algorithm and derive time complexity.