

# 4. Propositional Logic



Language & Logic

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# Assignments

- Continuous assessment (20% of module)
  - 3 assignments in weeks 3, 6 and 10
    - worth 6%, 6% and 8%, respectively
- Assignment 1 out now, due Fri 13 Oct (5pm)
  - topic: grammars and truth tables
  - submission via Canvas
    - handwritten and scanned (via phone, tablet, scanner, etc.) or typeset (e.g. with Word, LibreOffice, Latex), PDF preferred
    - any problems – please ask
  - see solutions for Exercise Sheet 1 and practice quizzes

# Overview

- Last lecture: a first formal look at logic
  - atomic propositions and connectives
  - constructing truth tables, for...
  - 1. semantics of connectives and propositions
  - 2. checking validity of arguments
- Today
  - propositional logic as a formal language
  - semantics of propositions and arguments
  - proof: introduction to natural deduction

# Propositional logic

- Logical connectives:

- negation:  $\neg$  (not)
- conjunction:  $\wedge$  (and)
- disjunction:  $\vee$  (or)
- implication:  $\rightarrow$  (if then / implies)

Alternative notation  
(e.g., Tomassi book)

- negation:  $\sim\phi$
- conjunction:  $\phi \& \psi$
- disjunction:  $\phi \vee \psi$

- Formulas (or “sentences”) in propositional logic

- if  $P$  is an atomic proposition, then  $P$  is a formula
- if  $\phi$  and  $\psi$  are formula, then these are all formulas:  
 $\neg \phi$ ,  $\phi \wedge \psi$ ,  $\phi \vee \psi$  and  $\phi \rightarrow \psi$
- these are **atomic** and **compound** formulas, respectively

- To be more precise, these are **well-formed formulas**

# Propositional logic – syntax

- The **syntax** of propositional logic is specified by the **grammar**
  - $F \rightarrow Ap \mid [\neg] F \mid F [\wedge] F \mid F [\vee] F \mid F [\rightarrow] F$
  - $Ap \rightarrow [P] \mid [Q] \mid [R] \mid \dots$
- Question
  - $P \wedge Q \vee R$
  - is this a well-formed formula? what does it mean?
- Question
  - how do we write this in propositional logic?
  - If **a**rachnids have eight legs then **c**rabs spin webs and **s**corpions live underwater

# Avoiding ambiguity

- Use parentheses to avoid ambiguity
  - $(P \wedge Q) \vee R$  vs.  $P \wedge (Q \vee R)$
  - $A \rightarrow (C \wedge S)$  vs.  $(A \rightarrow C) \wedge S$
- General rule:
  - enclose all compound sub-formulas in parentheses
- Precedence
  - negation applies to the smallest formula following
  - e.g.  $\neg P \vee Q$  means  $(\neg P) \vee Q$
  - but: we won't make assumptions about precedence of  $\wedge$  and  $\vee$

# Some terminology

- The **scope** of a **connective**
  - the connective itself, plus what it connects
- The **main connective** of a formula
  - the connective whose scope is the whole formula
- **Example**
  - $(\neg(P \wedge Q)) \rightarrow (\neg P \vee \neg Q)$
- In terms of **parse trees**
  - scope = sub-tree
  - main connective = root node

# Semantics of propositions

- We have used truth tables to illustrate the **semantics** (meaning) of propositions
  - i.e., the **truth value** (true or false) of the proposition for each possible **truth assignment** (to atomic propositions)

P	Q	$P \rightarrow Q$
T	T	T
T	F	F
F	T	T
F	F	T

- Properties of propositions
  - **valid**: true for every possible truth assignment
  - **satisfiable**: true for some possible truth assignment
  - **unsatisfiable**: false for every possible truth assignment
  - **contingent**: true for some assignment and false for another
- Note
  - don't confuse validity of propositions and arguments



# Relationships between propositions

- Logical equivalence

- $\phi \equiv \psi$
- $\phi$  and  $\psi$  have the same truth value for every possible truth assignment
- e.g.  $\neg(P \wedge Q) \equiv \neg P \vee \neg Q$

P	Q	$\neg(P \wedge Q)$	$\neg P \vee \neg Q$
T	T	F	F
T	F	T	T
F	T	T	T
F	F	T	T

- Logical entailment

- $\phi \models \psi$
- for every truth assignment where  $\phi$  is true,  $\psi$  is also true
- e.g.  $P \wedge Q \models P$

P	Q	$P \wedge Q$	P
T	T	T	T
T	F	F	T
F	T	F	F
F	F	F	F

- Extends to arguments...

# Arguments in propositional logic

- Example
  - Premise 1:  $H \rightarrow O$
  - Premise 2:  $\neg O$
  - Conclusion:  $\neg H$
- Notation: written as a **sequent**
  - $H \rightarrow O, \neg O : \neg H$
  - i.e., comma-separated list of premises, colon, conclusion
  - (where premises/conclusion are formulas in propositional logic)
- If the argument is **valid** (and been proven to be), we write:
  - $H \rightarrow O, \neg O \vdash \neg H$
  - i.e., the  $:$  is replaced by the turnstile symbol  $\vdash$  (“entails”)

# Proofs in propositional logic

- For formal proofs, we need two things:
  1. A formal language
    - for representing propositions, arguments
    - for now, we are using propositional logic
  2. A proof theory
    - to prove (“infer”, “deduce”) whether an argument is valid
    - first option (last week): truth tables
    - another option (next few weeks): natural deduction
- Why natural deduction?
  - scales better when there are many atomic propositions
  - can also be adapted to predicate logic

# Natural deduction

- Natural deduction
  - “natural” style of constructing a proof (like a human would)
  - **syntactic** (rather than **semantic**) proof method
  - proofs are constructed by applying **inference rules**
  - showing that the conclusion can be inferred from the premises
- Some example inference rules (written as sequents):
  - Modus ponens:  $P \rightarrow Q, P \vdash Q$
  - Modus tollens:  $P \rightarrow Q, \neg Q \vdash \neg P$
  - $\wedge$ -Introduction:  $P, Q \vdash P \wedge Q$
  - $\wedge$ -Elimination:  $P \wedge Q \vdash P$
- Note:
  - the set of inference rules available for use may vary

# Inference rules

- Inference rules need to be **truth preserving**
  - as shown before, we can check this with a truth table
- E.g. for modus tollens:
  - $P \rightarrow Q, \neg Q \vdash \neg P$

		P1	P2	C
P	Q	$P \rightarrow Q$	$\neg Q$	$\neg P$
T	T	T	F	F
T	F	F	T	F
F	T	T	F	T
F	F	T	T	T

# Writing proofs

- A proof comprises a number of **lines**, each of which has
  - **line number** (consecutive)
  - **formula** (in propositional logic, for now)
  - **rule annotation** (which inference rule used?)
  - **dependency numbers** (depends on which other lines?)
- **Basic idea**
  - work from the premises towards the conclusion
- **Example**
  - proof for  $P, \neg Q \vdash P \wedge \neg Q$

1.	$P$	Premise	$\{1\}$
2.	$\neg Q$	Premise	$\{2\}$
3.	$P \wedge \neg Q$	$\wedge$ -Introduction <sub>1,2</sub>	$\{1,2\}$

# Summary

- Propositional logic
  - atomic propositions, combined with logical connectives
  - formal language for expressing propositions
  - ambiguity, precedence, scope, main connective
  - semantics, equivalence, entailment
- Arguments
  - written as sequents
  - can be proved using natural deduction
  - sequence of inference rule applications
- Exercise class: Tue 11am or Thur 10am
  - topic: natural deduction