

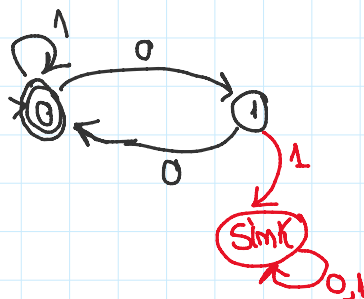
DFA / AFD (Automate Finite Deterministe)Def.

$$A = (K, \Sigma, \underline{\delta}, q_0, F)$$

$$\delta: K \times \Sigma \rightarrow K \quad (\text{funct. totală})$$

Cum funcționează?

- Primește un cuv.  $w$ .
- Inițial, în starea  $q_0$ .
- "Consumă" câte un caracter din  $w$  și prin  $\delta$  trece în altă stare.
- La final  $\rightarrow$  starea este finală ( $\in F$ )  $\rightarrow$  ACCEPT  
 $\rightarrow$  starea nu e finală ( $\notin F$ )  $\rightarrow$  REJECT

Ex.:

$$w = 11001 \Rightarrow \text{ACCEPT}$$

$$w = 1100011001 \rightarrow \text{REJECT}$$

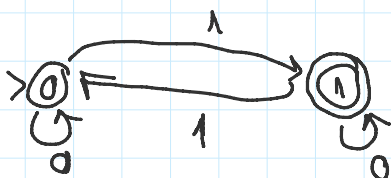
Configurație:  $(q, w) \in K \times \Sigma^*$  $\vdash_A$  1-par $\vdash_A^*$ 

0-sau-mai-mult-pari

Ex.

$$w = 11001$$

$$(0, 11001) \vdash_A (0, 1001) \vdash_A (0, 001) \vdash_A (1, 01) \vdash_A (0, 1) \vdash_A (0, \epsilon) \Big|_{0 \in F} \rightarrow \text{ACCEPT}$$

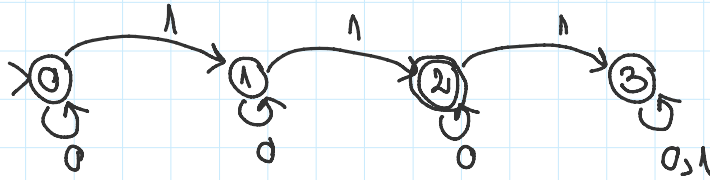
Def. DFA-ul  $A$  acceptă cuvântul  $w \Leftrightarrow (q_0, w) \vdash_A^* (q, \epsilon)$  și  $q \in F$ .Def.  $\mathcal{L}(A) = \{ w \in \Sigma^* \mid A \text{ acceptă } w \}$ Ex.  $\mathcal{L}(A) = \{ \text{cuv. pt. care orice } 0 \text{ apare însoțit de alt } 0 \}$ Exercițiu①  $\mathcal{L} = \{ w \in \{0,1\}^* \mid w \text{ conține un nr. impar de } 1 \}$ 

010 1100

$$(0, 0101100) \vdash_A (0, 101100) \vdash_A (1, 01100) \vdash_A (1, 1100)$$

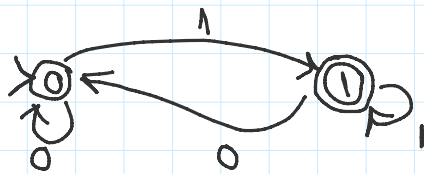
$$\vdash_A (0, 100) \vdash_A (1, 00) \vdash (1, \varepsilon) \Big|_{1 \in F} \Rightarrow \text{ACCEPT}$$

$$(2) \mathcal{L} = \{ w \in \{0,1\}^* \mid w \text{ contains exact 2 de 1} \}$$



$$(3) \mathcal{L} = \{ w \in \{0,1\}^* \mid w \text{ este un nr. impar } \}$$

$$w = \dots 1$$



$$q_0 = 0$$

$$(4) \mathcal{L} = \{ w \in \{0,1\}^* \mid w \% 3 = 0 \}$$

$$w = 3K$$

$$x \in \mathbb{N}, x \in \{3K, 3K+1, 3K+2\}$$

$$\begin{matrix} \downarrow & \downarrow & \downarrow \\ 0 & 1 & 2 \end{matrix}$$

$$\begin{matrix} 0011 \in \mathcal{L} \\ \varepsilon \in \mathcal{L} \end{matrix}$$

$$\begin{aligned} w &= 'v0' = 2v \\ w &= 'v1' = 2v+1 \end{aligned}$$

$$\text{I. } v = 3K \begin{cases} \xrightarrow{0} w = 2v = 6K = 3K' \rightarrow q_0 \\ \xrightarrow{1} w = 2v+1 = 6K+1 = 3K'+1 \rightarrow q_1 \end{cases}$$

$$\text{II. } v = 3K+1 \begin{cases} \xrightarrow{0} 2v = 6K+2 = 3K'+2 \rightarrow q_2 \\ \xrightarrow{1} 2v+1 = 6K+3 = 3K' \rightarrow q_0 \end{cases}$$

$$4v+1 = 6K+3 = 3n \rightarrow g_0$$

III.  $v = 3K+2$

$$\begin{array}{l} \xrightarrow{0} 2v = 3K'+1 \rightarrow g_1 \\ \xrightarrow{1} 2v+1 = 3K'+2 \rightarrow g_2 \end{array}$$

