

$S \rightarrow aAc$

$S \rightarrow aBd$

$S \rightarrow bBc$

$S \rightarrow bAd$

$A \rightarrow g$

$B \rightarrow g$

LR item sets

$S' \rightarrow S$  (0)

$S \rightarrow aAc$  (1)

$S \rightarrow aBd$  (2)

$S \rightarrow bBc$  (3)

$S \rightarrow bAd$  (4)

$A \rightarrow g$  (5)

$B \rightarrow g$  (6)

Pop Don Benild  
# 236

	First	Follow
$S'$	a, b	\$
$S$	a, b	\$
$A$	g	c, d
$B$	g	c, d

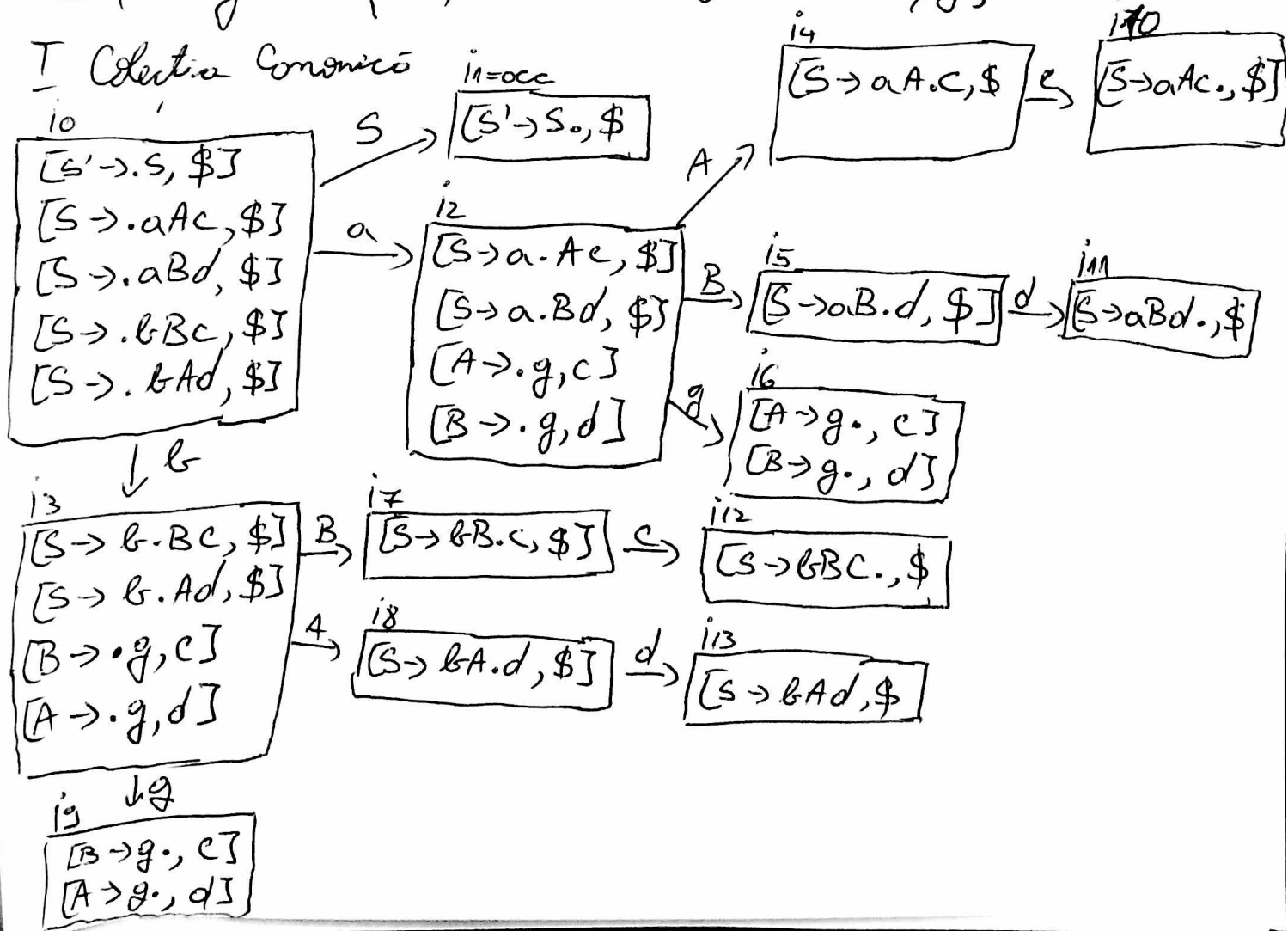
$\text{first}(S') = \text{first}(S)$

$\text{first}(S) = \{a, b\}$

$\text{first}(A) = \{g\}$

$\text{first}(B) = \{g\}$

LR Item Sets



# Tabelul de analiza LR(1)

Pop Ioan-Daniel  
08.236

	S	A	B	a	b	c	d	e	\$
i0	$\Delta_1$			$\Delta_2$	$\Delta_3$				
i1									acc
i2		$\Delta_4$	$\Delta_5$					$\Delta_6$	
i3		$\Delta_8$	$\Delta_7$					$\Delta_9$	
i4						$\Delta_{10}$			
i5							$\Delta_{11}$		
i6						$\Delta_5$	$\Delta_6$		
i7						$\Delta_{12}$			
i8							$\Delta_{13}$		
i9						$\Delta_6$	$\Delta_5$		
i10									$\Delta_1$
i11									$\Delta_2$
i12									$\Delta_3$
i13									$\Delta_4$

În tabelul de analiza LR(1) nu avem conflicte  $\Rightarrow$   
gramatica este de tip LR(1)

Conflict = într-o celulă facem mai mult de o reducere  
sau o reducere și o deplasare