### Lab 4

## Link for Demo:

5.1: https://youtu.be/j1jeYG0ZbL0

5.2: <a href="https://youtu.be/9zXEpe-KnGY">https://youtu.be/9zXEpe-KnGY</a>

# Group 4

Andrew Branum, Cade Garrett, Ankit Kanotra, Shawn Witt, Nathan Zabloudil

## **Abstract**

The purpose of this lab was to develop our knowledge and understanding of user processes and system calls within a raspberry Pi architecture. As per usual, we utilized the Raspberry Pi as well as our laptops to fulfill the obligations of the lab. Exercise 5.1 assigned us with the task of attempting to access system calls in user mode, and to return an exception that must be handled. Meanwhile, Exercise 5.2 tasked us with implementing a system call to create task priority, and then to demonstrate how the priority of the tasks was dynamically adjusted. Overall, the tasks of this lab were completed to their fullest.

### **Body**

Team 4 got off to an easy start on Lab 4 with none of the hiccups of the previous 3 labs. Having a fully functional Raspberry Pi and serial cable, they were easily able to begin work on Section 5. They started the process by reading chapter 5 within the provided GitHub, and then beginning work on the 2 provided questions. The only notable issue was a misplaced SD card, which returned no trouble as a replacement was immediately found.

Question 5.1 asked the team to return exception calls using a register after attempting to use system calls in user mode. To do this, a new register was created to handle system instructions that, once triggered, will cause a trap state to be formed and will thus throw an exception, in response.

Question 5.2 tasked the team to implement a system call for task priority and to showcase it dynamically altering task priority. To do this, a 5th system call was created within sys.h that would create a priority number. This priority number will be established to each process and increment based on the previous process. Using supervisor mode, these processes will be executed and the priority will dynamically change.

# Source Code (Software)

#### Exercise 5.1

```
From the sysregs.h file:
#ifndef SYSREGS H
#define SYSREGS H
// *************
// SCTLR EL1, System Control Register (EL1), Page 2654 of AArch64-Reference-Manual.
// **************
#define SCTLR RESERVED
                                (3 \le 28) \mid (3 \le 22) \mid (1 \le 20) \mid (1 \le 11)
#define SCTLR EE LITTLE ENDIAN
                                (0 << 25)
#define SCTLR EOE LITTLE ENDIAN
                                      (0 << 24)
#define SCTLR I CACHE DISABLED
                                      (0 << 12)
#define SCTLR D CACHE DISABLED
                                      (0 << 2)
#define SCTLR MMU DISABLED
                                (0 << 0)
#define SCTLR MMU ENABLED
                                (1 << 0)
#define SCTLR VALUE MMU DISABLED (SCTLR RESERVED)
SCTLR EE LITTLE ENDIAN | SCTLR I CACHE DISABLED |
SCTLR D CACHE DISABLED | SCTLR MMU DISABLED)
// *************
// HCR EL2, Hypervisor Configuration Register (EL2), Page 2487 of
AArch64-Reference-Manual.
#define HCR RW
                                 (1 << 31)
#define HCR VALUE
                           HCR RW
// *************
// SCR EL3, Secure Configuration Register (EL3), Page 2648 of AArch64-Reference-Manual.
// *************
#define SCR RESERVED
                           (3 << 4)
#define SCR RW
                            (1 << 10)
#define SCR NS
                           (1 << 0)
#define SCR VALUE
                                      (SCR RESERVED | SCR RW | SCR NS)
// *************
// SPSR EL3, Saved Program Status Register (EL3) Page 389 of AArch64-Reference-Manual.
// ***********************
```

```
#define SPSR MASK ALL
                                        (7 << 6)
#define SPSR EL1h
                                 (5 << 0)
#define SPSR VALUE
                                 (SPSR MASK ALL | SPSR EL1h)
// *************
// ESR EL1, Exception Syndrome Register (EL1). Page 2431 of AArch64-Reference-Manual.
// **************
#define ESR ELx EC SHIFT
                                 26
#define ESR ELx EC SVC64
                                 0x15
#define ESR ELx EC SYS INSTR 0x18
                                              //Establishing a register for instructions
#endif
From the entry.S file:
.globl vectors
vectors:
      ventry sync invalid el1t
                                              // Synchronous EL1t
      ventry irq invalid el1t
                                              // IRQ EL1t
      ventry fiq invalid el1t
                                              // FIQ EL1t
      ventry error invalid el1t
                                              // Error EL1t
      ventry sync invalid el1h
                                              // Synchronous EL1h
      ventry ell irq
                                              // IRQ EL1h
                                              // FIQ EL1h
      ventry fiq invalid el1h
      ventry error invalid el1h
                                              // Error EL1h
      ventry el0 sync
                                              // Synchronous 64-bit EL0
      ventry el0 irq
                                              // IRQ 64-bit EL0
      ventry fiq invalid el0 64
                                              // FIQ 64-bit EL0
      ventry error invalid el0 64
                                              // Error 64-bit EL0
      ventry sync invalid el0 32
                                              // Synchronous 32-bit EL0
      ventry irq invalid el0 32
                                              // IRQ 32-bit EL0
      ventry fiq invalid el0 32
                                              // FIQ 32-bit EL0
      ventry error invalid el0 32
                                              // Error 32-bit EL0
sync invalid ellt:
      handle invalid entry 1, SYNC INVALID EL1t
irq invalid ellt:
      handle invalid entry 1, IRQ INVALID EL1t
fiq_invalid el1t:
      handle invalid entry 1, FIQ INVALID EL1t
error invalid ellt:
      handle invalid entry 1, ERROR INVALID EL1t
sync invalid ellh:
      handle invalid entry 1, SYNC INVALID EL1h
fig invalid el1h:
      handle invalid entry 1, FIQ INVALID EL1h
```

```
error invalid el1h:
      handle invalid entry 1, ERROR_INVALID_EL1h
fig invalid el0 64:
       handle invalid entry 0, FIQ INVALID EL0 64
error invalid el0 64:
       handle invalid entry 0, ERROR INVALID EL0 64
sync invalid el0 32:
      handle invalid entry 0, SYNC INVALID EL0 32
irq_invalid el0 32:
       handle invalid entry 0, IRQ INVALID EL0 32
fig invalid el0 32:
      handle invalid entry 0, FIQ INVALID EL0 32
error invalid el0 32:
      handle invalid entry 0, ERROR INVALID EL0 32
ell irq:
      kernel entry 1
             handle irq
      bl
       kernel exit 1
el0 irq:
      kernel entry 0
             handle irq
      bl
      kernel exit 0
el0 sync:
       kernel entry 0
       mrs x25, esr el1
                                                // reads syndrome register
             x24, x25, #ESR ELx EC SHIFT
       lsr
                                                // exception class
             x24, #ESR ELx EC SVC64
                                                 // puts SVC in 64-bit state
       cmp
             el0 svc
       b.eq
       cmp x24, #ESR ELx EC SYS INSTR
       b.eq el0 sys instr
      handle invalid entry 0, SYNC ERROR
                                                // number of system calls
sc nr .req
             x25
                                                // syscall number
scno
      .req
             x26
                                                // syscall table pointer
stbl
             x27
      .req
el0 svc:
             stbl, sys call table
                                                // load syscall table
      adr
      uxtw
             scno, w8
                                                // syscall number in w8
             sc nr, #Num syscalls
      mov
       bl
             enable irq
                                                // check upper syscall limit
       cmp scno, sc nr
       b.hs
             ni sys
       ldr
             x16, [stbl, scno, lsl #3]
                                                // address in syscall table
      blr
                                                 // call sys * routine
             x16
      b
             ret from syscall
ni_sys:
      handle invalid entry 0, SYSCALL ERROR
```

```
ret from syscall:
      bl
             disable irq
             x0, [sp, #S X0]
                                                       // returned x0
      str
      kernel exit 0
el0 sys instr:
      bl show trapped sys instr
                                                       //branch and link to show trapped
             x22, [sp, #16 * 16]
                                                       // skip the trapped mrs instruction
       add x22, x22, #4
                                                       // add the register value
      str x22, [sp, #16 * 16]
      b ret to user
                                                       //return to user
.globl ret from fork
ret from fork:
      bl
             schedule tail
             x19, ret to user
                                                // not a kernel thread
      cbz
             x0, x20
      mov
             x19
      blr
ret to user:
      bl disable irq
      kernel exit 0
.globl err hang
err hang: b err hang
From the irg.c file:
#include "utils.h"
#include "printf.h"
#include "timer.h"
#include "entry.h"
#include "peripherals/irq.h"
const char *entry error messages[] = {
  "SYNC INVALID EL1t",
  "IRQ INVALID EL1t",
  "FIQ INVALID EL1t",
  "ERROR INVALID EL1T",
  "SYNC INVALID EL1h",
  "IRQ INVALID EL1h",
  "FIQ INVALID EL1h",
  "ERROR INVALID EL1h",
  "SYNC INVALID EL0 64",
  "IRQ INVALID ELO 64",
  "FIQ INVALID EL0 64",
  "ERROR INVALID EL0 64",
  "SYNC INVALID EL0 32",
```

```
"IRQ INVALID_EL0_32",
  "FIQ INVALID EL0 32",
  "ERROR INVALID EL0 32",
  "SYNC ERROR",
  "SYSCALL ERROR"
};
void enable interrupt controller()
  put32(ENABLE IRQS 1, SYSTEM TIMER IRQ 1);
void show invalid entry message(int type, unsigned long esr, unsigned long address)
  printf("%s, ESR: %x, address: %x\r\n", entry error messages[type], esr, address);
void show trapped sys instr() {
                                                //print if trapped
  printf("trapped\r\n");
void handle irq(void)
  unsigned int irq = get32(IRQ PENDING 1);
  switch (irq) {
       case (SYSTEM TIMER IRQ 1):
              handle timer irq();
              break;
       default:
              printf("Unknown pending irq: %x\r\n", irq);
From the utils.h file:
#ifndef UTILS H
#define UTILS H
extern void delay (unsigned long);
extern void put32 (unsigned long, unsigned int);
extern unsigned int get32 (unsigned long);
extern int get el (void);
extern unsigned int get ex();
                                                //Function to get an exception
#endif /* UTILS H */
```

#### Exercise 5.2

#### From the **irq.c** file:

```
#include "utils.h"
#include "printf.h"
#include "timer.h"
#include "entry.h"
#include "peripherals/irq.h"
const char *entry error messages[] = {
  "SYNC INVALID EL1t",
  "IRQ INVALID EL1t",
  "FIQ INVALID EL1t",
  "ERROR INVALID EL1T",
  "SYNC INVALID EL1h",
  "IRQ INVALID EL1h",
  "FIQ INVALID EL1h",
  "ERROR INVALID EL1h",
  "SYNC INVALID ELO 64",
  "IRQ INVALID EL0 64".
  "FIQ INVALID EL0 64",
  "ERROR INVALID EL0 64",
  "SYNC INVALID EL0 32",
  "IRQ INVALID EL0 32",
  "FIO INVALID EL0 32",
  "ERROR INVALID EL0 32",
};
void enable interrupt controller()
  put32(ENABLE IRQS 1, SYSTEM TIMER IRQ 1);
void show invalid entry message(int type, unsigned long esr, unsigned long address)
  printf("%s, ESR: %x, address: %x\r\n", entry error messages[type], esr, address);
void handle irq(void)
  unsigned int irq = get32(IRQ PENDING 1);
  switch (irq) {
```

```
case (SYSTEM_TIMER_IRQ_1):
                handle timer irq();
                break;
        default:
                printf("Unknown pending irq: %x\r\n", irq);
From the kernel.c file:
#include "printf.h"
#include "utils.h"
#include "timer.h"
#include "irg.h"
#include "sched.h"
#include "fork.h"
#include "mini uart.h"
#include "sys.h"
void user process1(char *array)
  char buf[2] = \{0\};
  long priority = 1;
                                              //intialize priority to 1
  while (1){
                                              //if initial priority == 1
        if(array[0] == '1'){}
                                                      //set next priority
                call sys priority(++priority);
        for(int j = 0; j < 4; j++){
                                              //write priority for each
               for (int i = 0; i < 5; i++)
                       buf[0] = array[i];
                       call sys write(buf);
                       delay(100000);
  }
void user process(){
  char buf[30] = \{0\};
  tfp sprintf(buf, "User process start\n\r");
  call sys write(buf);
  unsigned long stack = call sys malloc();
  if (stack < 0) {
        printf("Error while allocating stack for process 1\n\r");
        return;
  int err = call sys clone((unsigned long)&user process1, (unsigned long)"12345", stack);
```

```
if (err < 0)
        printf("Error while clonning process 1\n\r");
        return;
  stack = call sys malloc();
  if (stack < 0) {
        printf("Error while allocating stack for process 1\n\r");
        return;
  err = call sys clone((unsigned long)&user process1, (unsigned long)"abcd", stack);
                                                                         // err is pid of process2
  if (err < 0)
        printf("Error while clonning process 2\n\r");
        return;
  //call sys priority(err, 0xa);
  call_sys_exit();
void kernel process(){
  printf("Kernel process started. EL %d\r\n", get el());
  int err = move to user mode((unsigned long)&user process);
  if (err < 0)
        printf("Error while moving process to user mode\n\r");
void kernel main(void)
  uart init();
  init printf(0, putc);
  irq vector init();
  timer init();
  enable interrupt controller();
  enable irq();
  int res = copy process(PF KTHREAD, (unsigned long)&kernel process, 0, 0);
  if (res < 0) {
        printf("error while starting kernel process");
        return;
  }
  while (1)
        schedule();
```

```
}
```

#### From the sys.S file:

```
#include "sys.h"
.globl call sys write
call_sys_write:
  mov w8, #SYS WRITE NUMBER
  svc #0
  ret
.globl call sys malloc
call sys malloc:
  mov w8, #SYS MALLOC NUMBER
  svc #0
  ret
.globl call sys exit
call_sys_exit:
  mov w8, #SYS EXIT NUMBER
  svc #0
  ret
.globl call_sys_clone
call_sys_clone:
  /* Save args for the child. */
  mov x10, x0
                                         /*fn*/
                                         /*arg*/
  mov x11, x1
                                         /*stack*/
        x12, x2
  mov
  /* Do the system call. */
                                         /* stack */
  mov
         x0, x2
        x8, #SYS_CLONE_NUMBER
  mov
  svc
       0x0
  cmp x0, #0
  beq
       thread_start
  ret
globl call sys priority
                                                //makes global
call_sys_priority:
                                                //calls sys priority
                                                //move priority # to w8
  mov w8, #SYS PRIORITY NUMBER
                                         //trigger supervisor mode
  svc #0
  ret
                                         //return
```

```
thread start:
  mov x29, 0
  /* Pick the function arg and execute. */
  mov x0, x11
  blr x10
  /* We are done, pass the return value through x0. */
  mov x8, #SYS EXIT NUMBER
  svc 0x0
From the sys.c file:
#include "fork.h"
#include "printf.h"
#include "utils.h"
#include "sched.h"
#include "mm.h"
void sys write(char * buf){
  printf(buf);
}
int sys clone(unsigned long stack){
  return copy process(0, 0, 0, stack);
unsigned long sys malloc(){
  unsigned long addr = get free page();
  if (!addr) {
       return -1;
  return addr;
}
void sys exit(){
  exit process();
void sys_priority(long priority)
                                                   //schedules tasks based on priority
       current->priority = priority;
```

void \* const sys\_call\_table[] = {sys\_write, sys\_malloc, sys\_clone, sys\_exit, sys\_priority};

```
From the sys.h file:
#ifndef _SYS_H
#define _SYS_H
#define Num syscalls 5
                                        //Number of system calls is 5
#define SYS_WRITE_NUMBER
                                         // syscal numbers
#define SYS MALLOC NUMBER 1
#define SYS CLONE NUMBER
#define SYS EXIT NUMBER
#define SYS PRIORITY NUMBER 4
                                        //priority number is 4
#ifndef ASSEMBLER
void sys write(char * buf);
int sys fork();
void call sys write(char * buf);
int call sys clone(unsigned long fn, unsigned long arg, unsigned long stack);
unsigned long call sys malloc();
void call_sys_exit();
void call sys priority(long priority);
                                        //Calls System Priority
#endif
#endif /*_SYS_H */
```

# Schematics (Hardware) - None

### <u>Analysis</u>

All of the knowledge acquired in this lab was through tutorials and exercises created and hosted on Github by Sergei Matyukevich. The lessons in this lab pertain to user processes and system calls. Along with process management, process isolation should also be a present functionality in an operating system. One important technique that can be done is moving all user processes to exception level 0. Doing so ensures that a user cannot perform privileged processor operations. This is actually a necessary step to ensure proper process isolation. System calls are methods in an API provided by the OS that allow the OS to remain in exception level 0 while allowing interaction between a device and its users, thus maintaining process isolation. System calls are basically synchronous exceptions, and are generated by the svc instruction. Since system calls are synchronous exceptions, they are dealt with at exception level 1. The operating system then performs a series of of tasks, after which execution resumes at exception level 0. These tasks include verifying each argument, executes the required task, and initiates normal exception return. RPi OS contains four syscalls:

- 1. <u>Write</u> Uses the UART device to output something on a screen. The string to be printed is one its arguments.
- **2.** <u>Clone</u> Causes a new user thread to be created. The first argument for this system call is the location in the stack allocated for the new user thread.
- **3.** <u>Malloc</u> Allocates a memory page for a user process.
- **4.** Exit Is called by all processes after they are done executing.

These system cals are defined in the sys.c file. An array called "sys\_call\_table" contains pointers to each syscall handler. Each system call has a corresponding number, which acts as an index for each system call in sys\_call\_table. Before a system call can actually occur, a task must be running in user mode. There are two ways in which new user tasks can be created. Either a kernel thread can be moved to user mode, or a user task can fork itself to create a new user task.

### **Conclusion**

The purpose of this lab was to develop students' understanding of system calls and user processes within the Raspberry Pi architecture. Group 4 showcased this knowledge by returning exceptions to attempted system calls and creating task priority to deal with tasks as they dynamically come up. As computer scientists and engineers, it's important students understand low level processes such as system calls. Not only are they good practical knowledge, especially for anyone designing a system or kernel, but knowledge of such processes are often the difference between a computer science engineer, and someone who only took a programming bootcamp.

# <u>References</u>

S-Matyukevich. "S-Matyukevich/Raspberry-Pi-OS: Learning Operating System Development Using Linux Kernel and Raspberry Pi." *GitHub*, https://github.com/s-matyukevich/raspberry-pi-os.