



MS2 Project 1

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RISC-V-Single-Cycle-Microprocessor

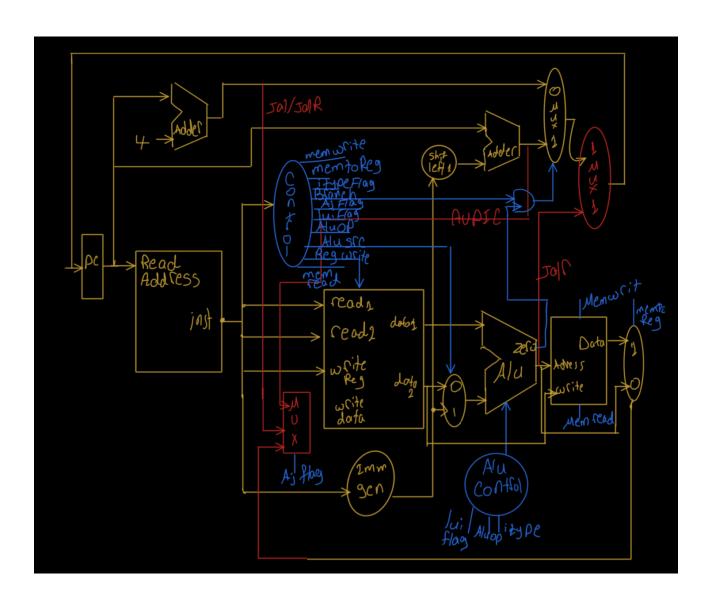
Introduction:

We are designing a RISC-V processor that supports the base integer instruction set according to the specifications found here: https://riscv.org/technical/specifications/ We implemented each instruction according to the specification found in page 130 of the RISC-V Instruction Set Manual –Volume I: Unprivileged ISA and explained in Chapter2 of the same manual.

	imm[31:12]			rd	0110111	LUI
	rd	0010111	AUIPC			
im	rd	1101111	JAL			
imm[11:	0]	rs1	000	rd	1100111	JALR
imm[12 10:5]	rs2	rs1	000	imm[4:1 11]	1100011	BEQ
imm[12 10:5]	rs2	rs1	001	imm[4:1 11]	1100011	BNE
imm[12 10:5]	rs2	rs1	100	imm[4:1 11]	1100011	BLT
imm[12 10:5]	rs2	rs1	101	imm[4:1 11]	1100011	BGE
imm[12 10:5]	rs2	rs1	110	imm[4:1 11]	1100011	BLTU
imm[12 10:5]	rs2	rs1	111	imm[4:1 11]	1100011	BGEU
imm[11:		rs1	000	rd	0000011	LB
imm[11:	- 1	rs1	001	rd	0000011	LH
imm[11:	- 1	rs1	010	rd	0000011	LW
imm[11:	- 1	rs1	100	rd	0000011	LBU
imm[11:		rs1	101	rd	0000011	LHU
imm[11:5]	rs2	rs1	000	imm[4:0]	0100011	SB
imm[11:5]	rs2	rs1	001	imm[4:0]	0100011	SH
imm[11:5]	rs2	rs1	010	imm[4:0]	0100011	SW
imm[11:		rs1	000	rd	0010011	ADDI
imm[11:		rs1	010	rd	0010011	SLTI
imm[11:	-	rs1	011	rd	0010011	SLTIU
imm[11:		rs1	100	rd	0010011	XORI
imm[11:	- 1	rs1	110	rd	0010011	ORI
imm[11:		rsl	111	rd	0010011	ANDI
0000000	shamt	rs1	001	rd	0010011	SLLI
0000000	shamt	rs1	101	rd	0010011	SRLI
0100000	shamt	rs1	101	rd	0010011	SRAI
0000000	rs2	rs1	000	rd	0110011	ADD
0100000	rs2	rs1	000	rd	0110011	SUB
0000000	rs2	rs1	001	rd	0110011	SLL
0000000	rs2	rs1	010	rd	0110011	SLT
0000000	rs2	rs1	011	rd	0110011	SLTU
0000000	rs2	rs1	100	rd	0110011	XOR
0000000	rs2	rs1	101	rd	0110011	SRL
0100000	rs2	rs1	101	rd	0110011	SRA
0000000	rs2	rs1	110	rd	0110011	OR
0000000	rs2	rsl	111	rd	0110011	AND

Design:

- To support JAL/JALR/AUIPC, we added a MUX that's colored in red in the below pic.
- To support JALR, we added an additional MUX that chooses between the output of the previous mux and the JALR signal.
- We added other singles in the control unit in order to differentiate between some new instructions
 - We added the Aj flag which distinguishes the Auipc and the jal instructions from each other where Aj flag selects from the jal/jalr, Auipc or from the alu
 - We added an i_type flag where some instructions such as addi can make conflict between sub and addi so we make this flag to differentiate between them.
 - We added lui flag which tells the ALU to shift left the immediate 12 times.



Algorithm:

In order to support the whole instructions, we did that by dividing them into groups (R-type, I-type,load instruction group, store instruction group, branches group, jal,jalr and AUIPC).

For supporting R- type instructions:

Based on the alu_selection line we choose the output of the ALU to be the result of the operation it does.

```
case(aluSel)
4'b0000: result = sum;
4'b0001: result= sub;
4'b0101: result= in1 & in2;
4'b0100: result= in1 | in2;
4'b0111:result=in1^in2;
4'b1000: result = in1>>in2; //SRL
4'b1010:result = in1>>in2; //SRA
4'b1001:result = in1<>>in2; //SRA
4'b1001:result={(in1<in2); //SLL
4'b1101:result={((in1<in2)?1:0); //SLT
4'b1111:result={(((in1<0?-in1:in1)<((in2<0?-in2:in2))?1:0); //SLTU
4'b0110:result={(in2,12'b0); ///LUI</pre>
```

For supporting I- type instructions:

Based on a flag we created called (I_type flag). As in the I-type instructions , the processor performs the same operations as R-type instructions and the only difference is the second operand where the ALU takes an immediate instead of a register operand. This flag is only activated in the control unit when the opcode equals the I-type opcode and deactivated for the rest of the instructions. Furthermore, in the ALU control unit we distinguished between the instructions using this flag.

```
else if (aluop == 2'b10 &&lui_flag==0 &&
i_type == 0&& instr2 == 0 && intr1 == 3'b000) // R- type
instruction (ADD) //instr2 VI
aluS = 4'b0000;
else if (aluop == 2'b10 &&lui_flag==0 && i_type == 1 &&
intr1 == 3'b000) // R- type instruction (ADDi) //same aluS
```

```
as add
aluS = 4'b0000;
```

For supporting B-type instructions:

For supporting the branching, we used the zero flag as an indication that the argument of any branch is true or false. Once it's true, the zero flag is set to 1, otherwise 0. For the unsigned branches, we settled on the sign first. If the number is negative, we multiply it by -1 and then compare. If it's positive, we compare directly.

```
4'b0010: // branching
begin
case(func3)
3'b000: if (in1 == in2) zero =1; else zero = 0; // beq zero =1 branch =1
3'b001: if (in1!= in2) zero = 1; else zero = 0; // bne branch =1 zero =0
3'b100: if (in1<in2) zero = 1; else zero = 0; // blt branch=1
3'b101: if (in1>= in2) zero = 1; else zero = 0; //bge branch=1
3'b110: if ({in1<0?-in1:in1}<{in2<0?-in2:in2}) zero = 1; else zero = 0; // bltu branch=1
3'b111: if ({in1<0?-in1:in1}>={in2<0?-in2:in2}) zero = 1; else zero = 0; // bgeu branch=1
```

For supporting Store-type instructions:

We modified the data memory module by passing function 3 for it so that it could distinguish between instructions. In order to support s-type, we need to activate the memwrite signal coming from the control unit and deactivate memread.

```
if (MemWrite==1 && func3 ==3'b010) //SW
mem[addr] <= data_in;
else if (MemWrite==1 && func3 ==3'b001)//sh
mem[addr][15:0] <= data_in;
else if (MemWrite==1 && func3 ==3'b000)//sb
mem[addr][7:0] <= data_in;
end</pre>
```

For supporting load-type instructions:

In order to support U-type, we need to activate the memread signal coming from the control unit and deactivate memwrite.

```
begin
if ( func3 == 3'b000) //LB
data_out = MemRead ? mem[addr][7:0]: 0;
else if ( func3 == 3'b001) //LH
data_out = MemRead ? mem[addr][15:0]: 0;
else if ( func3 == 3'b010) //LW
data_out = MemRead ? mem[addr]:0;
else if ( func3 == 3'b100) //LBU
data_out = MemRead ?{ 24'h0,mem[addr][7:0]}:0;
else if ( func3 == 3'b101) // LHU
data_out = MemRead ? {16'h0,mem[addr][15:0]}: 0;
```

For supporting JAL instruction:

We activated the branch signal in the control unit to allow the pc to branch. Also, we have added a new signal called AJ. it's a 2 bit signal. The first bit indicates the AUIPC signal and the second bit indicates the J signal related to jal/jalr.

```
5'b11011: //jal
begin
branch=1;
memread=0;
memtoreg=0;
aluop=11;
memwrite=0;
alusrc=1;
regwrite=0;
i_type=0;
lui_flag=0;
AJ_control=2'b01;
```

In the ALU control unit, we added this code that chooses the selection line.

```
if(aluop==2'b11) //jal
aluS=4'b0011;
```

We activate the zero flag to branch in the ALU module.

```
4'b0011:zero=1;
```

For supporting JALR instruction:

```
5'b11001: //jalr
begin
branch=0;
```

```
memread=0;
memtoreg=0;
aluop=10;
memwrite=0;
alusrc=1;
regwrite=1;
lui_flag=0;
i_type=0;
AJ_control=2'b01;
```

For supporting AUIPC instruction:

```
5'b00101: //auipc
begin
branch=0;
memread=0;
memtoreg=0;
aluop=00; //add as load
memwrite=0;
alusrc=1;
lui_flag=0;
regwrite=1;
i_type=0;
AJ_control=2'b11;
```

For supporting LUI instruction:

In the ALU module we added this code that loads the register with the upper 20 bits of the immediate and the rest of bits set them to zero.

```
4'b0110:result={in2,12'b0}; ///LUI
```

In the control unit we added this code.

```
5'b01101: //lui
begin
branch=0;
memread=0;
memtoreg=0;
aluop=10;
memwrite=0;
alusrc=1;
regwrite=1;
i_type=0;
lui_flag=1;
AJ_control=2'b00;
```

Testing:

We created 4 assembly files in order to test our single cycle processor

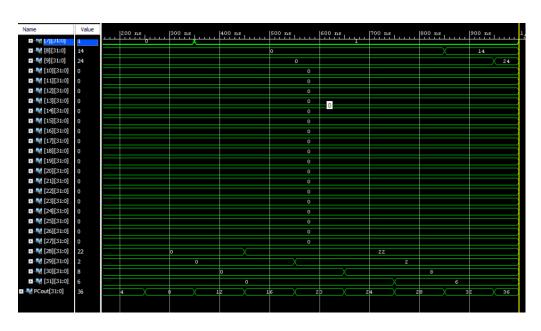
R type

- 1) we created R_typeTest.asm in order to test all the R type instructions,
- 2) then we converted them to binary to write them inside the instruction memory module

Assembly	Binary
lw t0, 0(zero) # 10 ->1010 lw t1, 4(zero) #12 ->1100 lw t2, 8(zero) #1 ->5 add t3,t1,t0 #22 sub t4,t1,t0 #2 and t5,t1,t0 #1000 ->8 xor t6,t1,t0 #0100 ->4 or s0,t1,t0 #1110 ->14 sll s1,t1,t2 #24 slt s2,t1,t0 #0 sltu s3,t1,t0 #0	000000000000000000010001010000011 000000
sra s4,t1,t2 #6 srl s5,t1,t0 #6	01000000011100110101101000110011 0000000

Screenshot

register

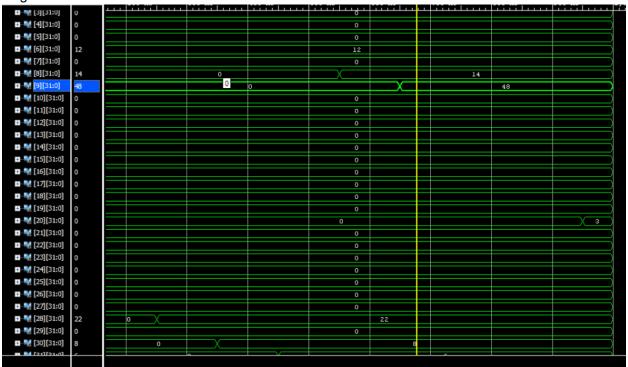


I type

Assembly	Binary
lw t1, 4(zero) #12 ->1100 addi t3,t1,10 #22 andi t5,t1,10 #1000 ->8 xori t6,t1,10 #0100 ->4 ori s0,t1,10 #1110 ->14 slli s1,t1,2 #24 slti s2,t1,10 #0 sltiu s3,t1,10 #0 srai s4,t1,2 #6 srli s5,t1,10 #6	0000000010000000010001100000011 00000000

Screenshot

Registers

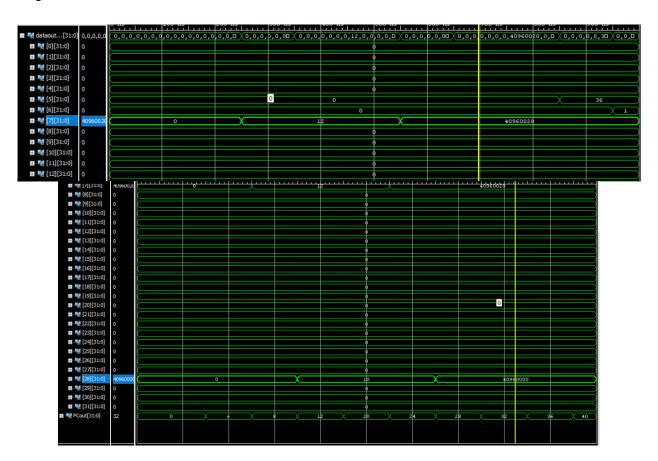


U type&jalr&jalr

Assembly	Binary	
lw t1,0(zero)#0 lw t2,4,(zero) #12 lw t3,8(zero) #10 jal jumb addi t1,t1,1 jumb: auipc t2,10000 lui t3,10000 addi t4,zero,0 jalr t0,t4,36 addi t1,t1,1 jumb2: sw t3,12(zero) #10 pc=36	00000000000000000010001100000011 0000000	

Screenshot

Registers and PC



S type

Assembly	Binary
Iw t1,0(zero)#0 Iw t2,4,(zero) #12 Iw t3,8(zero) #10 sw t1,12(zero)#0 sh t2,16,(zero) #12 sb t3,18(zero) #10	00000000000000000010001100000011 0000000

<u>Screenshot</u>

Memory:

