

# Hash Tables

...



<https://www.google.com/imgs?q=&url=https://i.pinimg.com/736x/57/a5/57a5f66950001NqjntDjCDN&size=AQ/VvndnF6950001NqjntDjCDN&src=60713103/688000&source=images&ed=rf&ved=0CwA1QJRoEwoTCG&5Glu40CE0AAAAAAAAAAAABAD>

# What are hash tables and why do we use them?

- Hash table: a data structure that can store and lookup elements efficiently (i.e. fast)
- Hash tables are comprised of two key parts:
  - Arrays
  - Hash Function
    - Takes in a key value and returns an index into the table
    - Should (*ideally*) provide a uniform distribution of keys across the hash table (array)
      - When keys *aren't* uniformly distributed, they are '*clustered*'.

# Collisions

Collisions occur when the index generated by the hash function for a certain key is already occupied.

- Example:  $\text{hash}(5) = 10 \rightarrow \text{table}[10] = 5$ 
  - $\text{hash}(40) = 10 \rightarrow \text{table}[10]$  is occupied!

Good hash functions will attempt to minimize collisions. *Great* ones will never have ‘em.

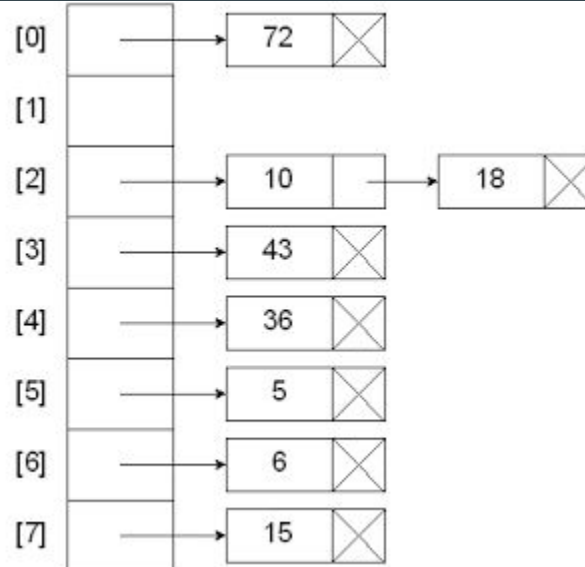
- Why not just always use great hash functions?
  - They’re hard to implement (\$\$\$)
  - They’re not always necessary

# Collision resolution techniques

- Chaining - Each bucket  $\sim$  a linked list.
  - Collision?  $\rightarrow$  Just link the new element. Simple as that!
- Probing
  - **Linear** - table[hash(x)] is occupied  $\rightarrow$  check table[hash(x) + 1]
    - Occupied?  $\rightarrow$  check table[hash(x)+2]
      - And so forth, until an empty bucket is found.
    - Drawbacks?
      - Clustering! If our function results in lots of collisions, values will be clustered around one another.
      - Takes a while to find the next empty bucket.
  - **Quadratic** - table[hash(x)] is occupied  $\rightarrow$  check table[hash(x) + 1<sup>2</sup>]
    - Occupied?  $\rightarrow$  check table[hash(x)+2<sup>2</sup>]
      - And so forth, until an empty bucket is found
    - Helps avoid clustering, but needs a big ol' table.

Hash key = key % table size

4 = 36 % 8  
2 = 18 % 8  
0 = 72 % 8  
3 = 43 % 8  
6 = 6 % 8  
2 = 10 % 8  
5 = 5 % 8  
7 = 15 % 8



# Real-world usage

University assigns an Identikey to each student.

Say they want to find the student 'Honda Rhoenigman'.

- Option 1: Start at 0000001. Is that her? No?
  - Check 0000002. Is that her? No?
- Option 2:  $\text{Hash}(\text{Honda Rhoenigman}) = 1294831$ .

# Inappropriate use cases

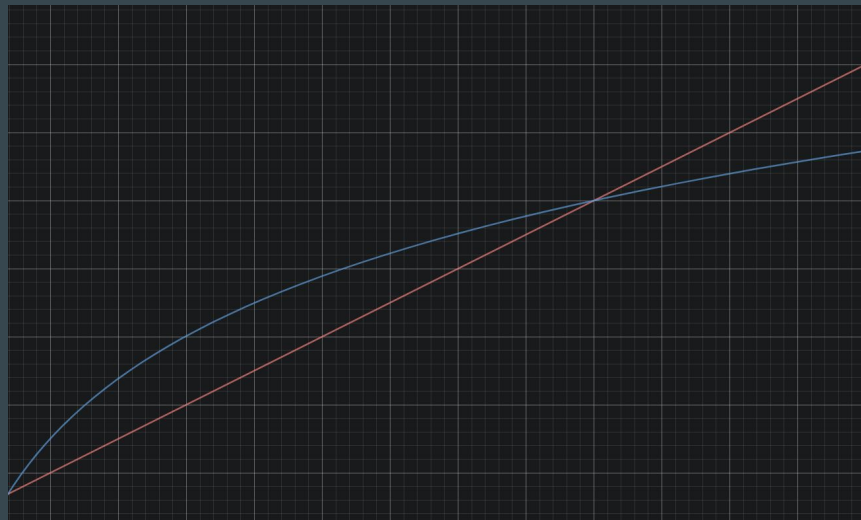
When the data we wish to store is sufficiently small, hashing may actually be more costly than it is worth.

Our knee-jerk is that ‘linear’ time is horrible.

- $\text{Value}(\text{linear search}) < \text{Value}(\text{hash}(\text{key}))$

But! For small datasets, it may be more sensible/cost-effective to just linear-search an array than to implement hashing.

Food for thought!





# Final Remarks

- Interview Grading Zoom links are under locations/description in your appointment slot!
  - Please show up with your code ready to run!
  - Email/Slack us if you're having trouble signing up or cannot make it to your appointment
    - Please make sure it's not a last minute notice unless an emergency arises

THANK YOU!