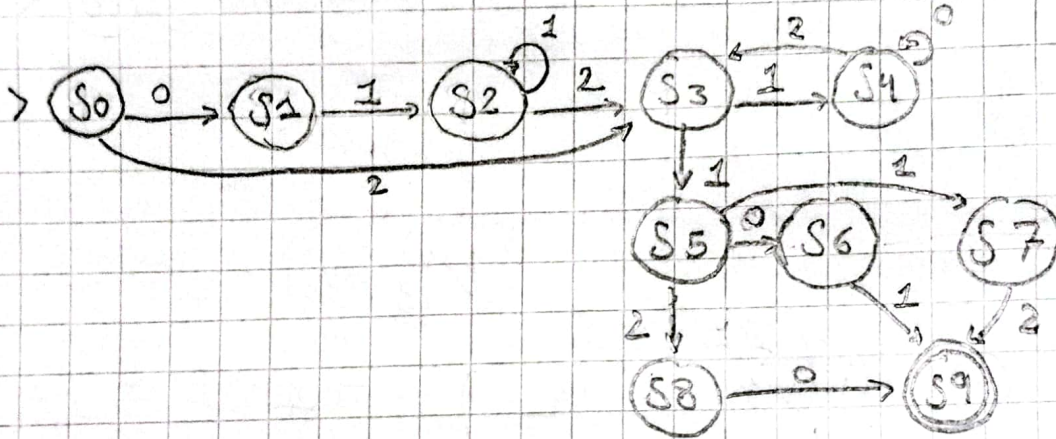
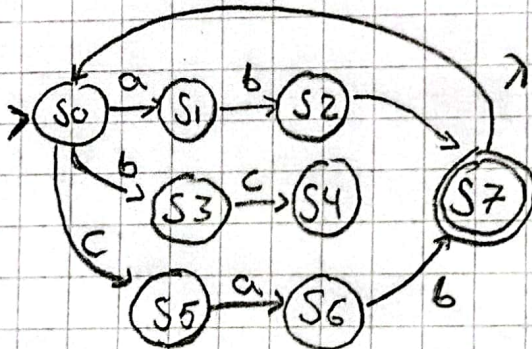


1. for each of the following languages, define the corresponding finite-state machine

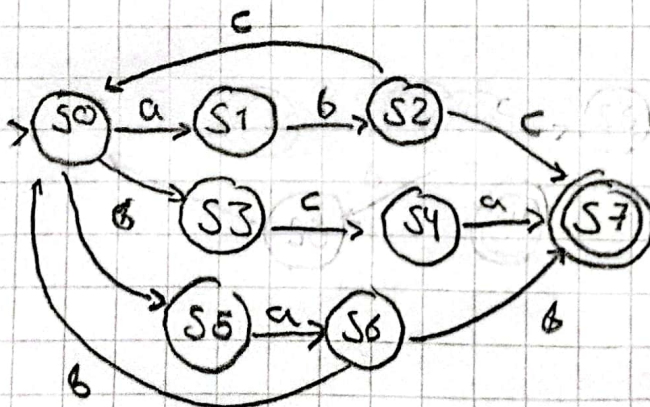
(i)  $\Sigma = \{0, 1, 2\}$ .  $L = (01^*2 \cup 2102)^* 101(01 \cup 12 \cup 20)^*$



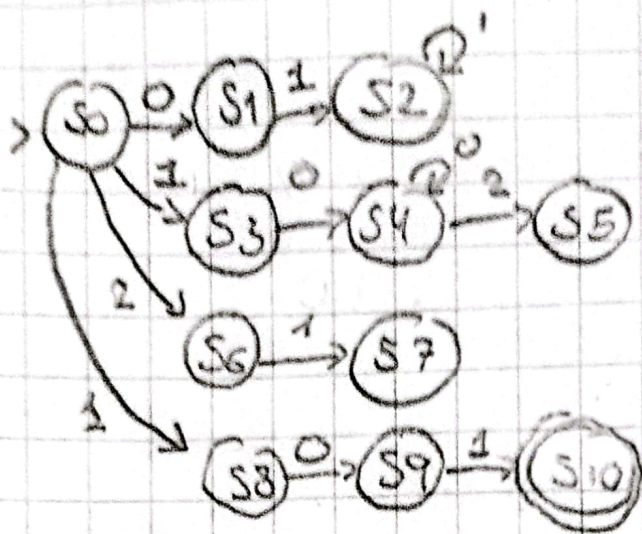
(ii)  $\Sigma = \{a, b, c\}$ .  $L = (abc \cup bca \cup cab)^* (abc \cup bca \cup cab)^*$



(iii)  $\Sigma = \{a, b, c\}$ .  $L = (abc \cup bca \cup cab)^* (abc \cup bca \cup cab)$



(iv)  $\Sigma = \{0, 1, 2\}$ .  $L = (01^*2 \cup 10^*2 \cup 21^*0)^* (01 \cup 12 \cup 20)^* 10$



2) For each one of the following finite-state machines, define the corresponding regular expression and generative grammar

i)  $\Sigma = \{0, 1\}$

$(0^* 10^* 1^*)^*$

generative grammar

$$G \begin{cases} S \rightarrow 0A \mid \lambda \\ A \rightarrow 0A \mid 1B \\ B \rightarrow 0B \mid 0C \\ C \rightarrow 1C \mid 1S \mid \lambda \end{cases}$$

ii)  $\Sigma = \{a, b, c\}$

$(a \cup b)^* c (\lambda \cup$



3) for each of the following regular expressions, define the corresponding generative grammar (all over the alphabet  $\Sigma = \{a, b, c, d\}$ ):

i)  $\{a^i b^j c^i d^j : i, j \geq 1\}$

$$G = \begin{cases} S \rightarrow aAd \\ A \rightarrow aAd \mid bBc \\ B \rightarrow bBc \mid \lambda \end{cases}$$

ii)  $\{a^i b^j c^i d^j : i, j \geq 1\}$

$$G = \begin{cases} S \rightarrow AB \\ A \rightarrow aAb \mid ab \\ B \rightarrow cBd \mid cd \end{cases}$$

iii)  $\{a^i b^j c^i d^j : i, j \geq 1\} \cup \{a^i b^j c^j d^i : i, j \geq 1\}$

iv)  $\{a^i b^j c^{i+j} : i \geq 0, j \geq 1\}$

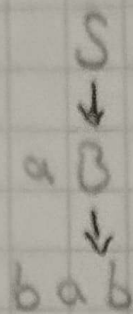
$$G = \begin{cases} S \rightarrow AD \\ A \rightarrow aAc \\ B \rightarrow bBc \mid bc \end{cases}$$

4) Be  $G$  a context-free grammar with the following productions

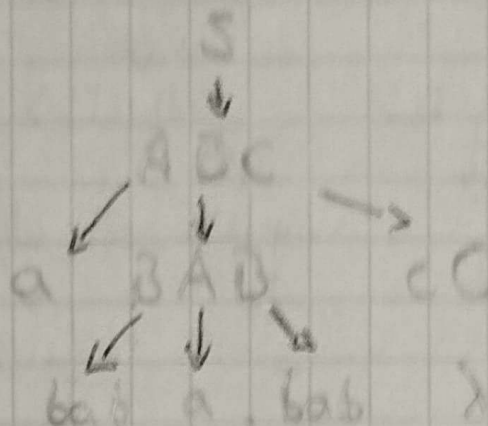
$$G = \begin{cases} S \rightarrow ABC \mid BAC \mid aD \\ A \rightarrow Fa \mid a \\ B \rightarrow BAB \mid bOb \\ C \rightarrow cC \mid \lambda \end{cases}$$

found derivation trees for the following strings

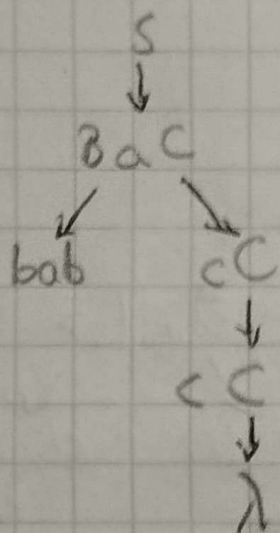
i)  $w_1 = abab$



iii)  $w_2 = abababab$



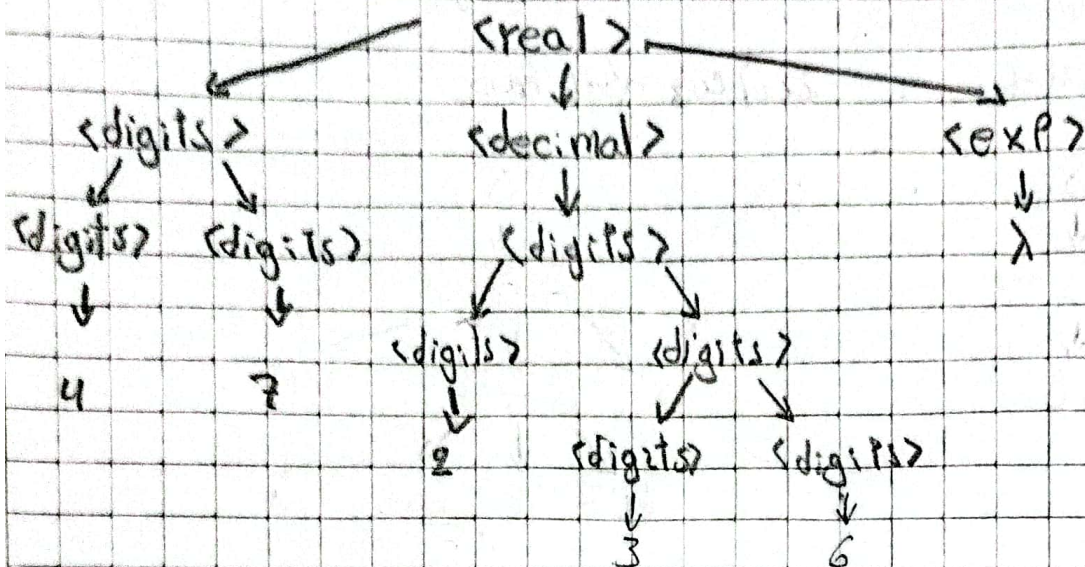
ii)  $w_2 = babab$



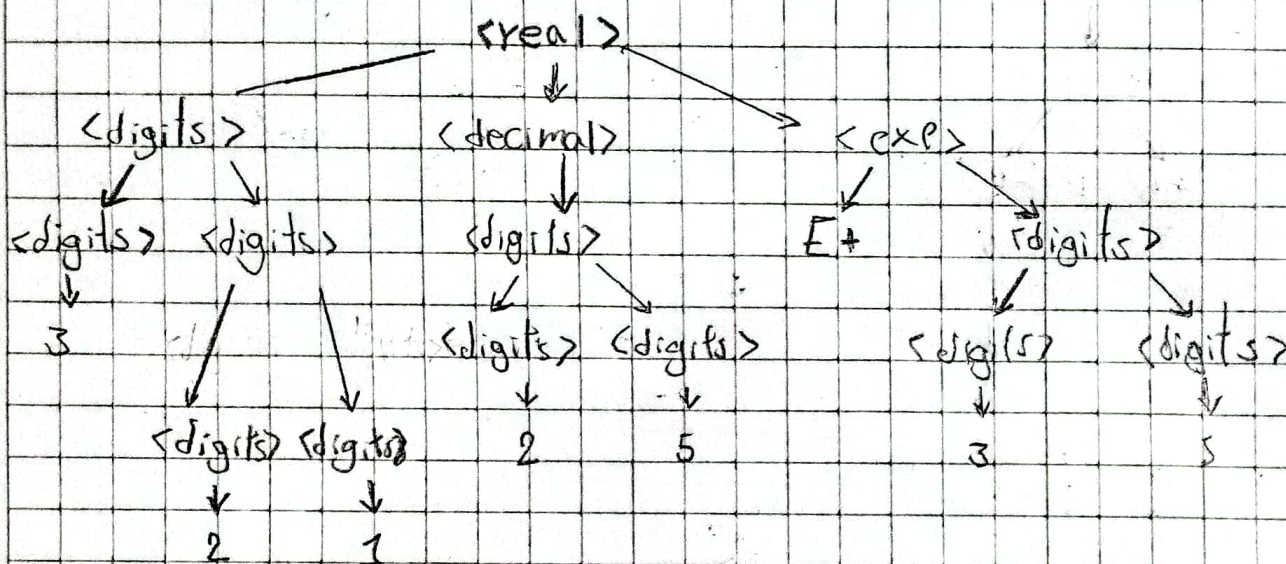
5) As follows there is a context-free grammar to generate real numbers without sign, the alphabet is  $\Sigma = \{0, 1, 2, 3, 4, 5, 6, 7, 8, \dots, +, -, \epsilon\}$ :

i)  $w_1 = 43, 236$

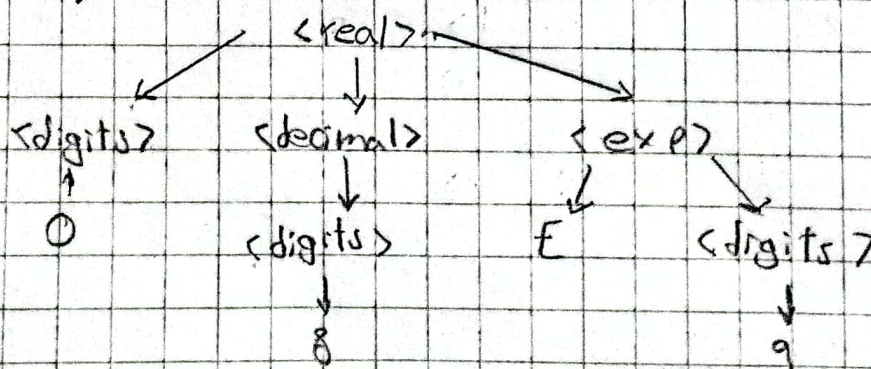




ii)  $w_2 = 321.25 E + 35$



iii)  $w_3 = 0.8 E 9$



iv)  $w_4 = 0.8E+9$

