



Department of Computer Science and Engineering

Data Structures and Object-Oriented Design

(CSE – 2050)

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Recap

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Quick Recap

- How to measure algorithm's performance in terms of timings?
 - Some issues with measuring timings
- We studied timings of an algorithm to find duplicates (duplicates_1) in input data
 - Discussed why duplicates_1 algorithm was not efficient
 - Improved and compared performance of both algos
- Asymptotic Analysis and Atomic operations
- Examples of operations in lists
- Activity # 5 for creating an algorithm to calculate the sum of K integers

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Activity # 5

Write an algorithm to sum first k integers and perform asymptotic analysis

Example:

k = 3

sum(k)

→ 6



Activity # 5 Solution

Write an algorithm to sum first k integers and perform asymptotic analysis

```

1  def sum_calc(k):
2      total = 0
3      for i in range(1, k+1):
4          total += i
5      return total
6
7  print(sum_calc(3))
8  print(sum_calc(6))
9  print(sum_calc(7))
10 print(sum_calc(100))

```

```

hasanbaig@HBMAC Prac % python3 sum_calc.py
6
21
28
5050

```

$$2k + 2$$



Activity # 5 Solution

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More clever approach of calculating the sum of k integers

```

12 def sum_clever(k):
13     total = k*(k+1)//2
14     return total
15
16 print(sum_clever(3))
17 print(sum_clever(6))
18 print(sum_clever(7))
19 print(sum_clever(100))

```

```

hasanbaig@HBMAC Prac % python3 sum_calc.py
6
21
28
5050

```

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Asymptotic Analysis

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Lets perform asymptotic analysis of duplicate_2 program

```

10 def duplicates_2(L):
11     n = len(L)
12     for i in range(1,n):
13         for j in range(i):
14             if L[i] == L[j]:
15                 return True
16     return False

```

$$n^2/2 - n/2 + 3$$

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Running Time Analysis

Asymptotic Analysis


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- Only higher order terms significantly affects a function
- In asymptotic analysis, we drop lower order terms and constants
 - Consider only higher order terms which grows faster with input size, n

$$5n^2 + 3n + 2$$

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Running Time Analysis


Big-O Notation

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- The formal mathematical definition which allows us to ignore lower order terms and constants is called **Big-O** notation

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Notes goes here

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Big-O Notation
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Example:

```

10 def duplicates_2(L):
11     n = len(L)
12     for i in range(1,n):
13         for j in range(i):
14             if L[i] == L[j]:
15                 return True
16     return False
          
```

$$n^2/2 - n/2 + 3$$

$$O(n^2)$$

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Activity # 6 (1/2)

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Calculate the time complexity (Big-O) for the following codes

```

1 #Practice question 1
2 def func(L):
3     x = 0
4     for i in L:
5         for j in L:
6             x += i * j
7     return x

```

$$1 + n(n \times 3) + 1 = 3n^2 + 2$$

$$= O(n^2)$$

```

1 # Practice question 2
2 a = 0
3 b = 0
4 for i in range(N):
5     a = a + random()
6     for j in range(M):
7         b = b + random()

```

$$N + M \Rightarrow O(N + M)$$

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Activity # 6 (2/2)

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Calculate the time complexity (Big-O) for the following codes

```

1 #Practice question 3
2 a = 0;
3 for i in range(N):
4     for j in reversed(range(i,N)):
5         a = a + i + j;

```

$$N = 10$$

$$9 \ 8 \ 7 \ 6 \ 5 \ 4 \ 3 \ 2 \ 1 \ 0$$

$$9 \ 8 \ 7 \ 6 \ 5 \ 4 \ 3 \ 2 \ 1$$

$$4 \ 3 \ 2$$

$$1 + 2 + 3 + \dots + N$$

$$= \frac{N(N+1)}{2} \Rightarrow O(N^2)$$

```

1 # Practice question 4
2 i = 1
3 while i <= n:
4     print(i)
5     i *= 2

```

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