HW 5 part 2

Part 2

T1	T2	T3
R(A)		
W(A)		
		R(A)
		W(A)
	R(A)	
R(B)		
		R(B)
W(B)		
		W(B)
	R(B)	
	commit	
commit		
		commit

→ Dirty Read
   $T_1 \rightarrow T_3$

→ Dirty Read
   $T_3 \rightarrow T_2$

→ Unrepeatable Read
   $T_3 \rightarrow T_1$

→ Lost Update
   $T_1 \rightarrow T_3$

→ Dirty Read
   $T_3 \rightarrow T_2$
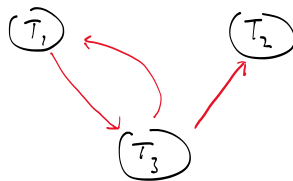
Precedence Graph:

There is a cycle between the T_1 & T_3 node, which means that the precedence graph is not a directed acyclic graph, which means that the schedule is not conflict serializable by the theorem we discussed in class.