



Network Infrastructures

Project:
Epidemic models for VANET message
dissemination

Presentation of the problem

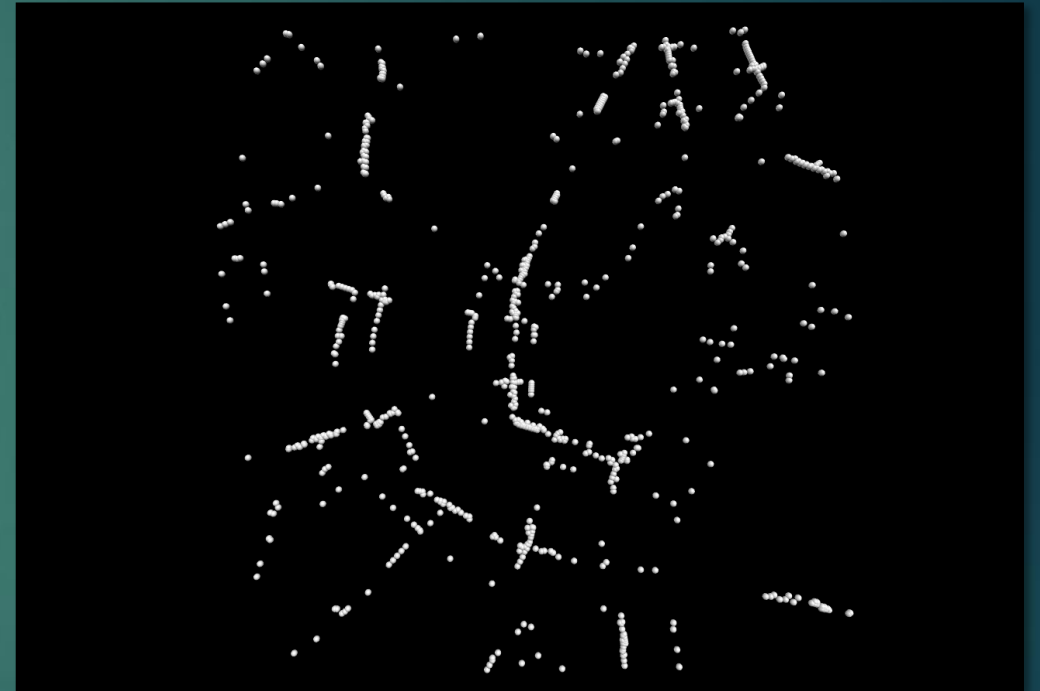
- ▶ **EPIDEMICS** are critical phenomena, not only from a biological point of view, such as infectious diseases, but also from a technological point of view, as malware propagation.
- ▶ In our specific case the propagation of a message in a vehicles urban network

Presentation of the problem

- ▶ Main purpose of the project is to develop an **algorithm** for the dissemination of a message in a vehicle network.
- ▶ Each vehicle knows the **position** of its neighbors
- ▶ Vehicles communicate only in **broadcast** with IEEE 802.11 standard wifi technology

Presentation of the problem

- ▶ Lets consider a graph where each node is a vehicle
- ▶ Vehicles can communicate only if they are neighbours
- ▶ The program provide an interface to simply interact with the graph and make the dissemination start
- ▶ We used VPython to represent the graph



[illegible]

Presentation of the problem

- ▶ We have created a dissemination algorithm of messages through our vehicles in order to reach the higher number of “infected” nodes, without overload the network.
- ▶ The program will perform the algorithm simply clicking the sphere from where we want to start.
- ▶ The result will be a graphic representation of the infected graph and the data results in the terminal.

Presentation of the problem

- ▶ We generated a series of graphs to show the variation of the outputs according to the changing of the cities and of the algorithm parameters like R_{Max} and T_{Max} .
- ▶ The utilized programming language is Python, the libraries are Matplotlib for the graphs and VPython for the user interface.

Algorithm's description

- ▶ Our algorithm runs for each vehicle and it returns whether or not a car must re-broadcast a message.
 - ▶ If all of this vehicles broadcast, then the network would be overloaded.
 - ▶ Else if too few cars broadcast, then few nodes would be reached.

Algorithm's description

- ▶ After receiving a message, each car starts a timer which becomes shorter if the distance from the last emitter becomes longer:

$$t = T_{\max}(1 - (dist / R_{\max}))$$

- ▶ Long distance, short timer.

During this timer, the car will wait for other messages arriving.

- ▶ The lists of these messages' emitters will be the central theme of our algorithm.

Algorithm's description

Algorithm 1 EvaluatePositions returns whether or not a car must re-broadcast a message

```
1: procedure EVALUATEPOSITIONS
2: input:
3:    $messages \leftarrow$  list of messages received during waiting time
4:    $neighbor\_positions \leftarrow$  list of positions of my neighbors
5:    $my\_pos \leftarrow$  my position
6:
7: initialization:
8:    $emitters = \emptyset$ 
9:   for  $msg$  in  $messages$  do
10:      $emitters = emitters \cup msg.emitters$ 
11:
12: main:
13:   for  $emitter$  in  $messages$  do
14:     if  $dist(my\_pos, emitter.pos) > 2RMIN$  then
15:       skip iteration
16:     for  $pos$  in  $neighbor\_positions$  do
17:       if  $dist(emitter, pos) < RMIN$  then
18:          $neighbor\_positions.remove(pos)$ 
19:   return  $neighbor\_positions.length > 0$ 
```

In the first part of the main function, for each emitter who relayed the message, if it is very far from me, then I will ignore it.

Algorithm's description



Algorithm 1 EvaluatePositions returns whether or not a car must re-broadcast a message

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```

Instead, about emitters not so far from me, if some of my neighbours are NOT contained in emitters' range of action, then I will broadcast.

Algorithm's description - Neighbours

Colors Legend

-  Actual Car
-  Neighbours



Algorithm's description - Previous Emitters

Colors Legend



Actual Car






Previous Emitters



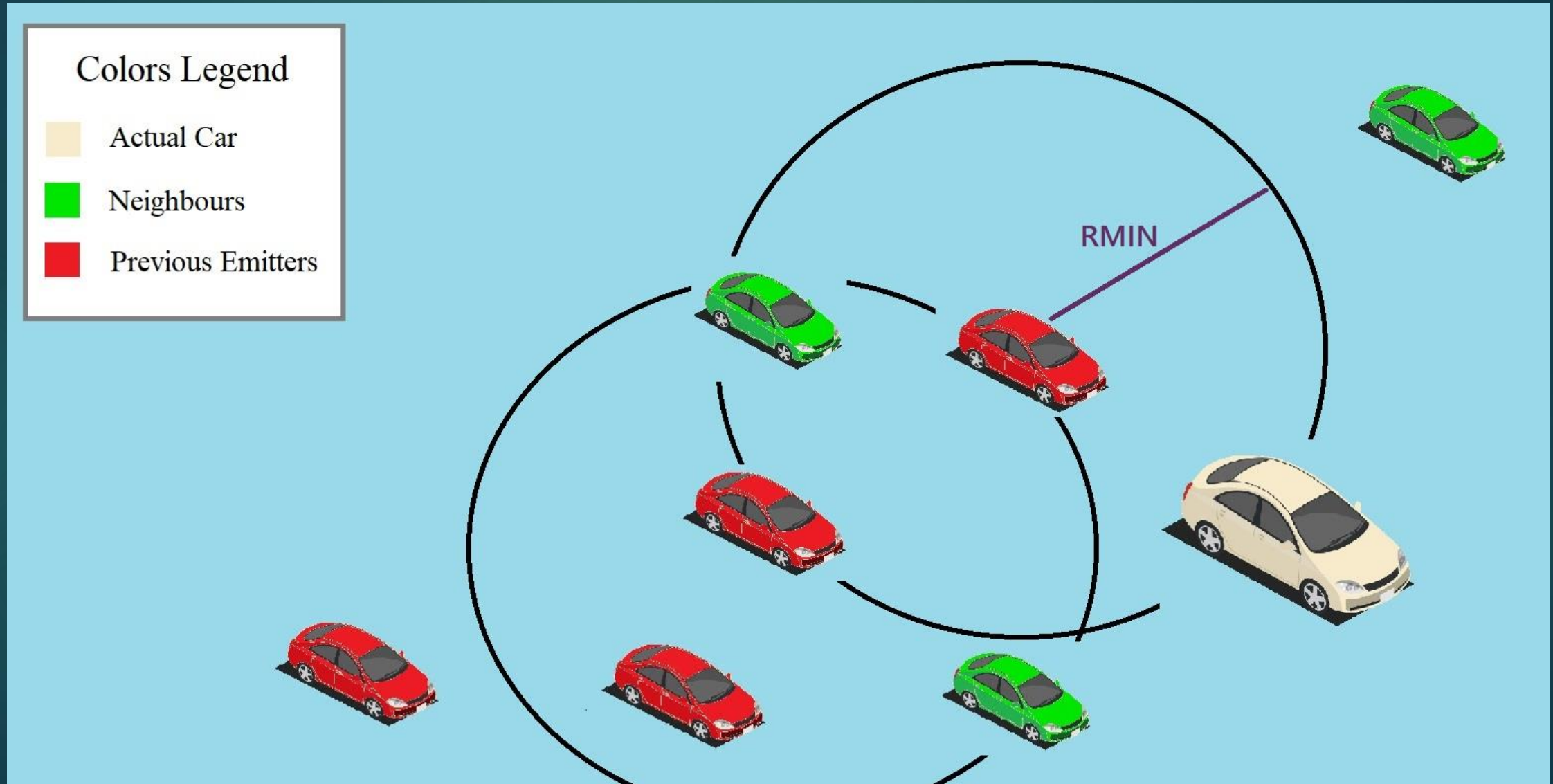
Algorithm's description - Previous Emitters & Neighbours

Colors Legend

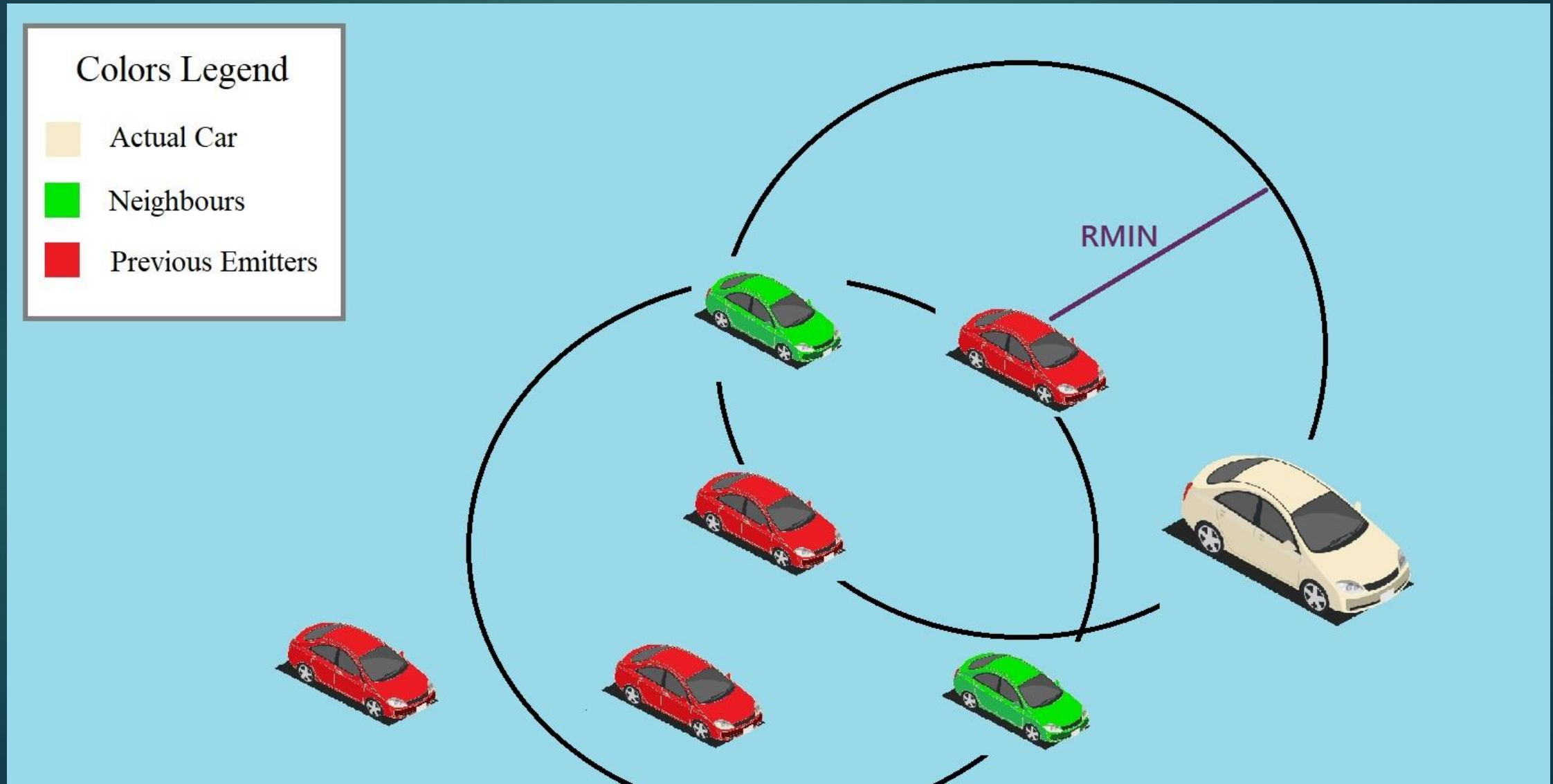
-  Actual Car
-  Neighbours
-  Previous Emitters



Algorithm's description - Re-Broadcast



Algorithm's description - NO Re-Broadcast



Results

data available

- ▶ We have real world data: position and adjacency matrix for Luxembourg, Cologne and New York, counting several hundred cars.
- ▶ Multiple graphs for a single city, more and less dense.
- ▶ Cities have different topology: impact on performances.

Results simulations

- ▶ We built a simulator capable of understanding those data and emulating the behavior of the cars.
- ▶ We can model the simulation with different parameters, both for the environment and the dissemination algorithm.
- ▶ With the obtained results, we can evaluate the performance of the algorithm and how these differ when we change some parameters.

Results evaluation

- ▶ We evaluated different metrics:
 - ▶ Cars reached;
 - ▶ Broadcasted messages;
 - ▶ Delay to reach last car;
 - ▶ Hops to reach last car.
- ▶ We ran multiple simulations for different cities for different graphs for different parameters: lots of data!

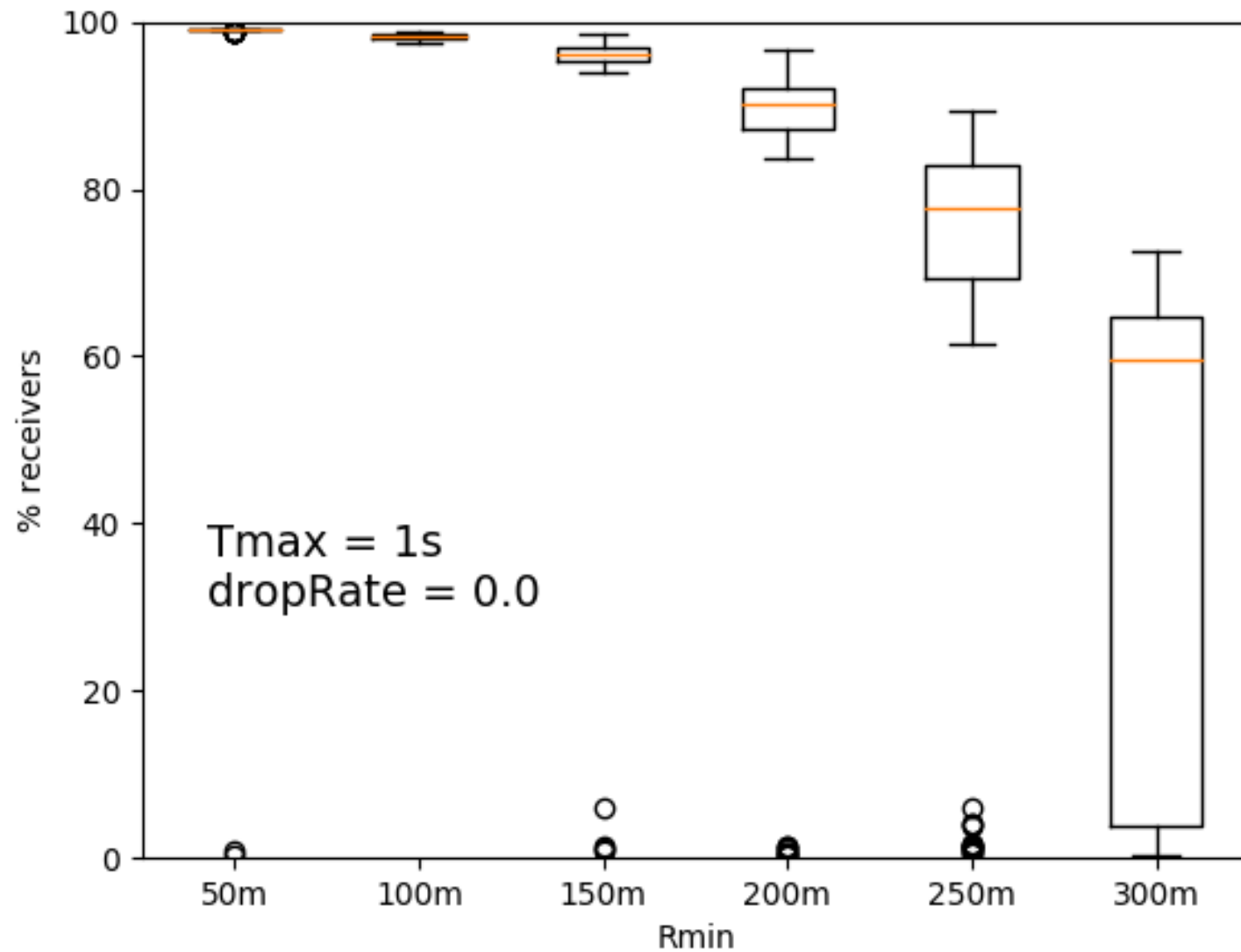
Results evaluation

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 - ▶ **Broadcasted messages;**
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Results

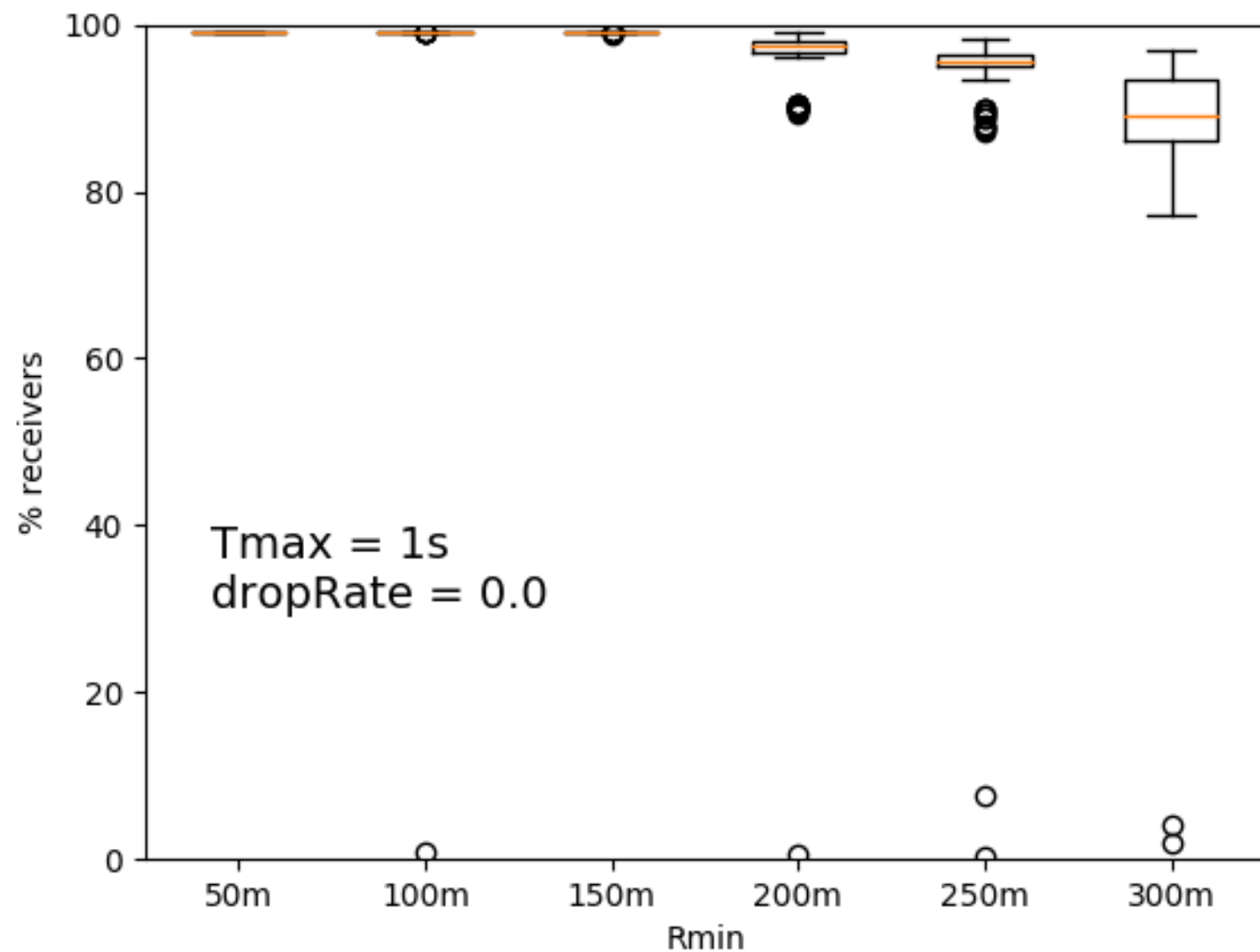
impact of RMIN

- ▶ Among all the parameters, RMIN impacted performance the most.
 - ▶ RMIN small: cars may decide to broadcast even if it is not necessary, impacting network traffic;
 - ▶ RMIN big: cars may decide to not broadcast even if it is necessary, the message won't reach some cars.
- ▶ Balance between performances and network congestion.



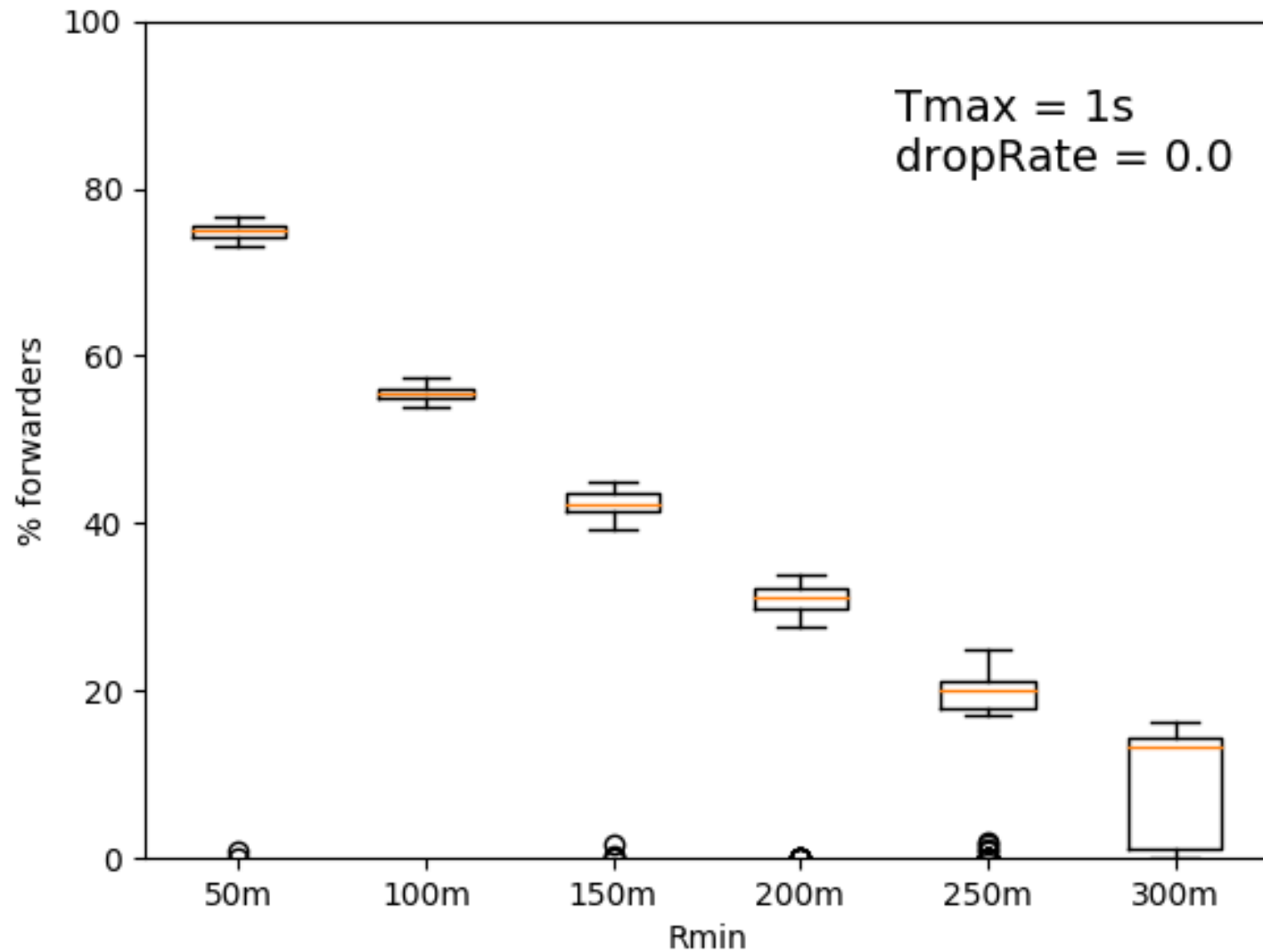
Luxemburg

- ▶ Percentage of receivers and RMIN.
- ▶ Least dense graph



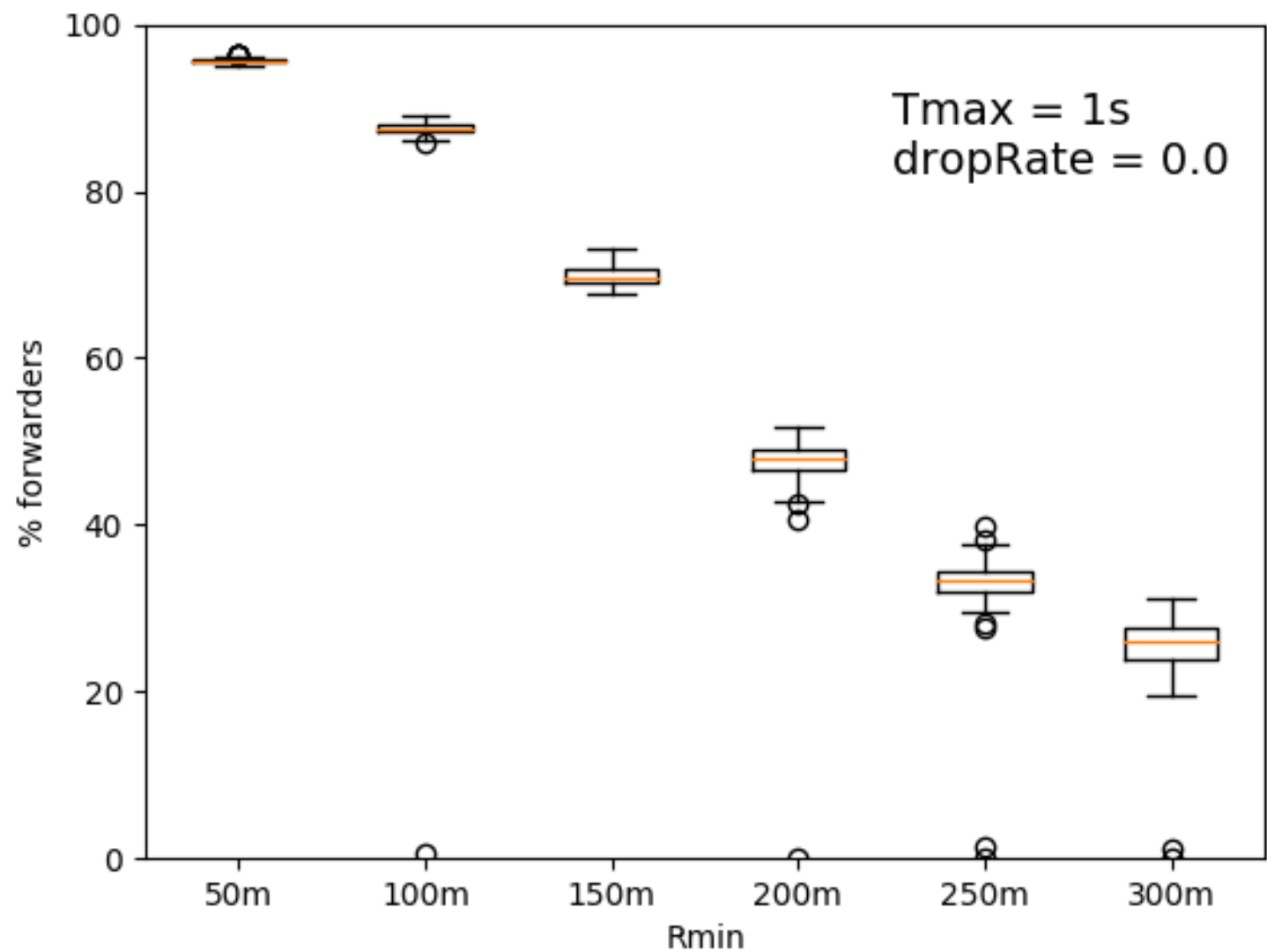
Luxemburg

- ▶ Percentage of receivers and R_{min} .
- ▶ Denser graph



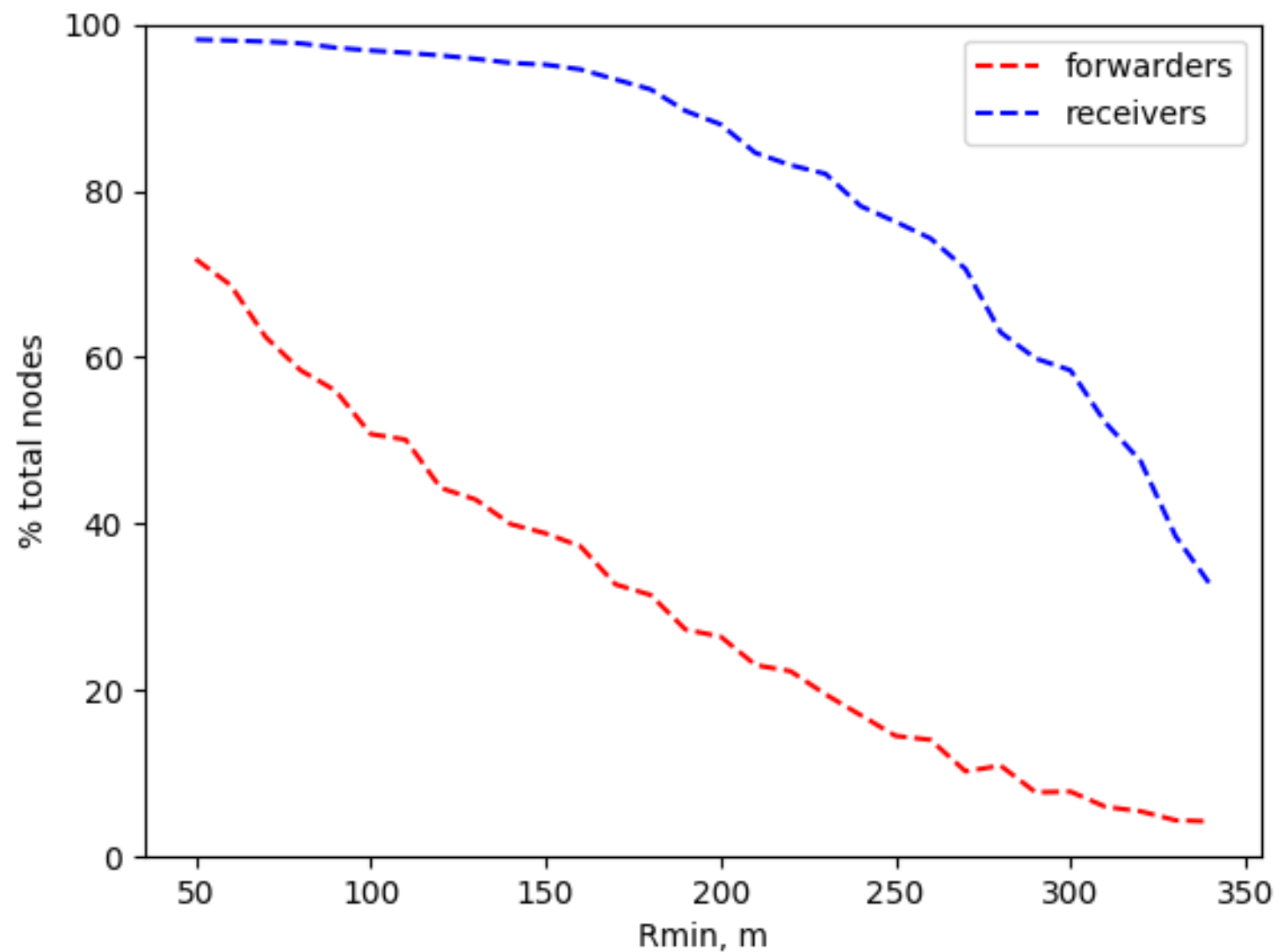
Luxemburg

- ▶ Percentage of forwarders and RMIN.
- ▶ Least dense graph



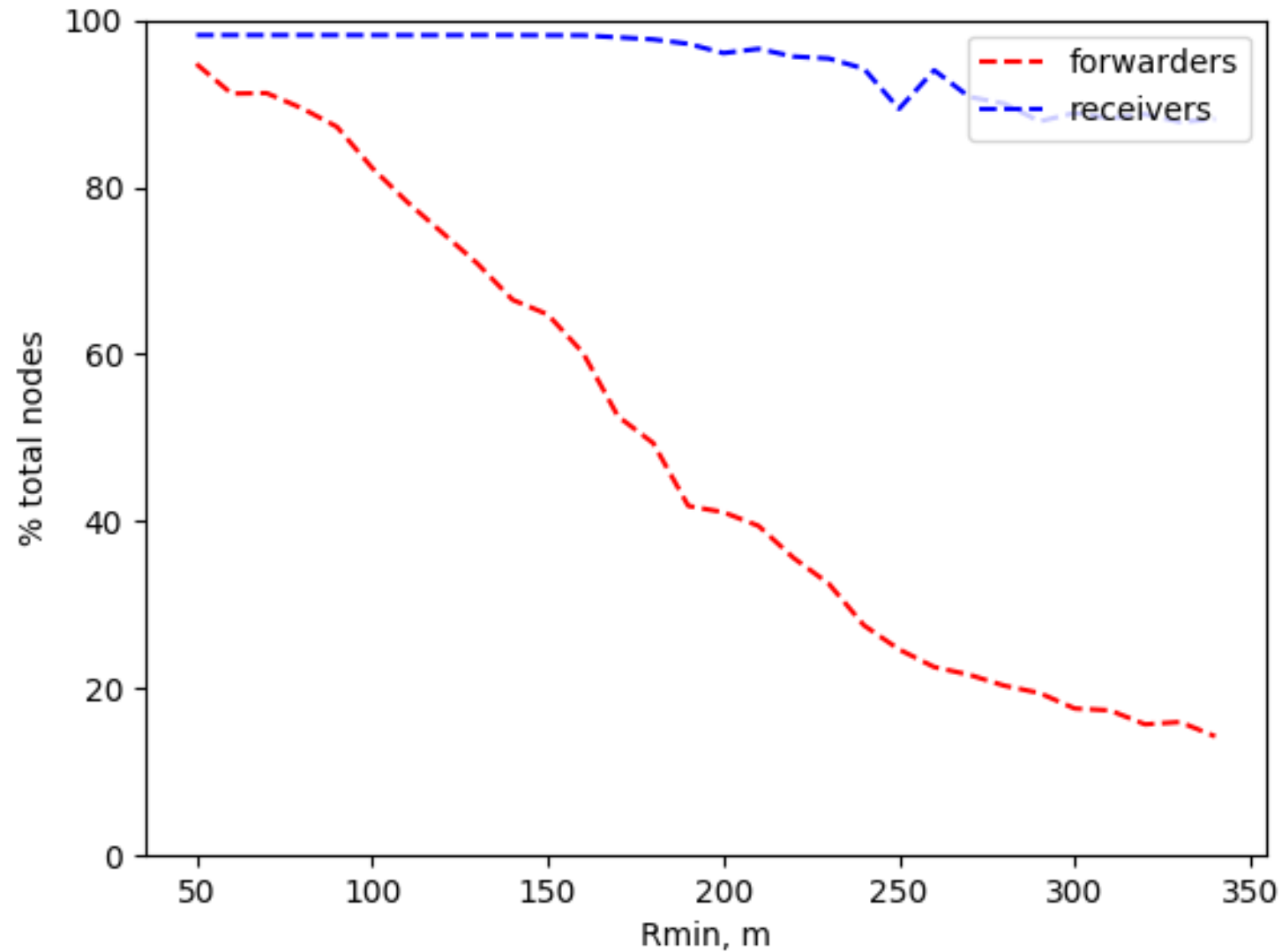
Luxemburg

- ▶ Percentage of forwarders and RMIN.
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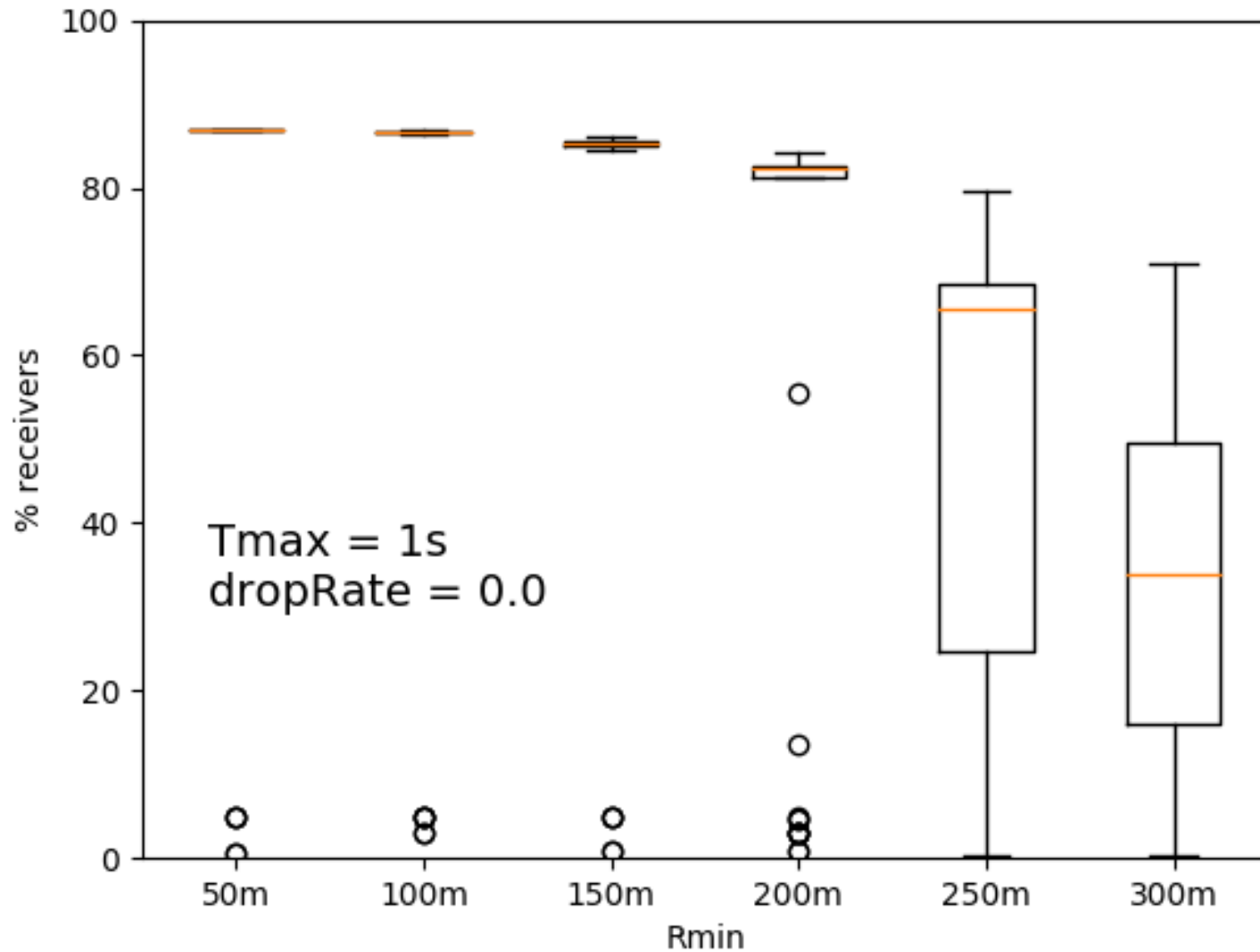
Luxemburg

- ▶ Forwarders and receivers compared
- ▶ Least dense graph



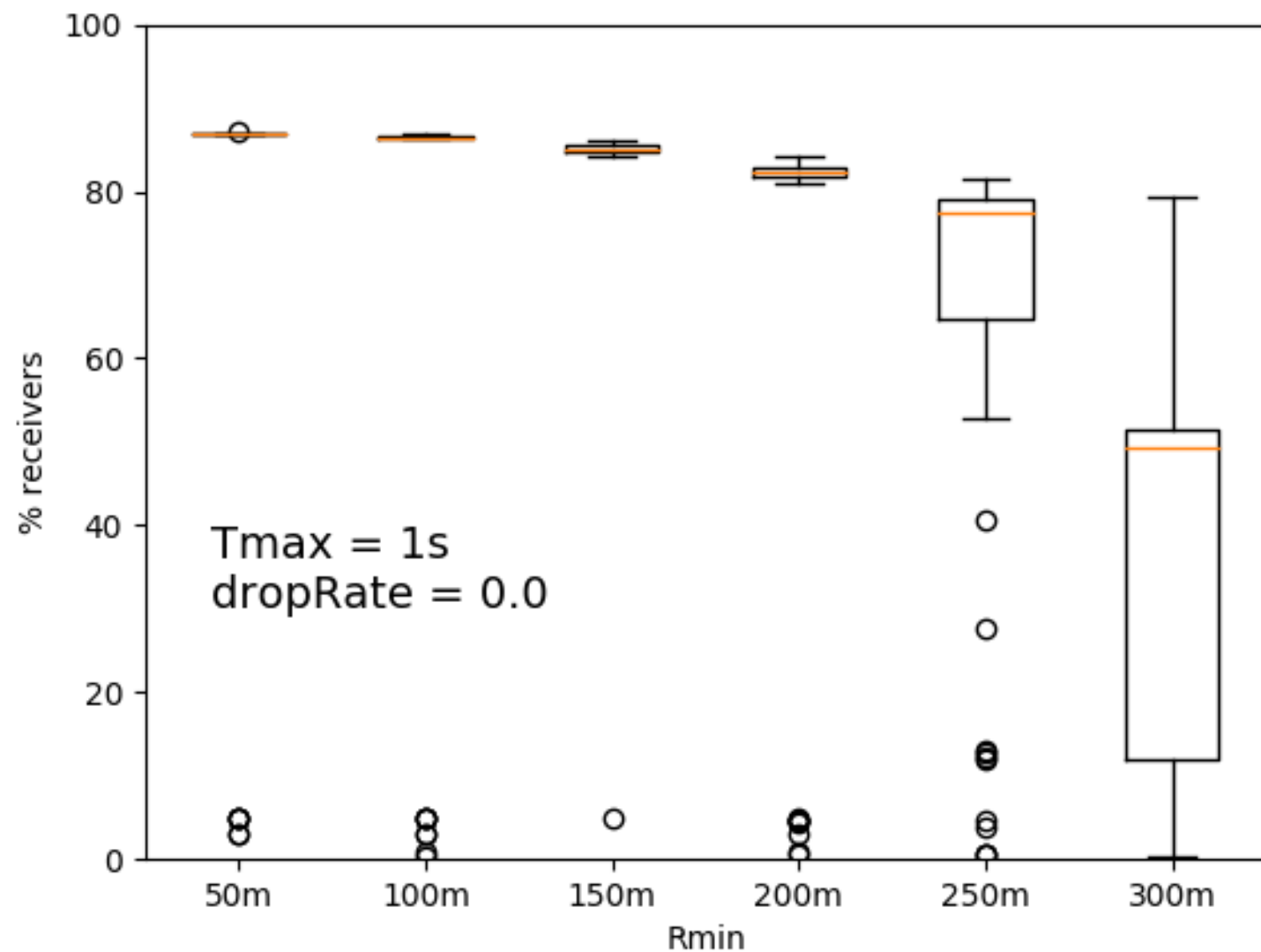
Luxemburg

- ▶ Forwarders and receivers compared
- ▶ Denser graph



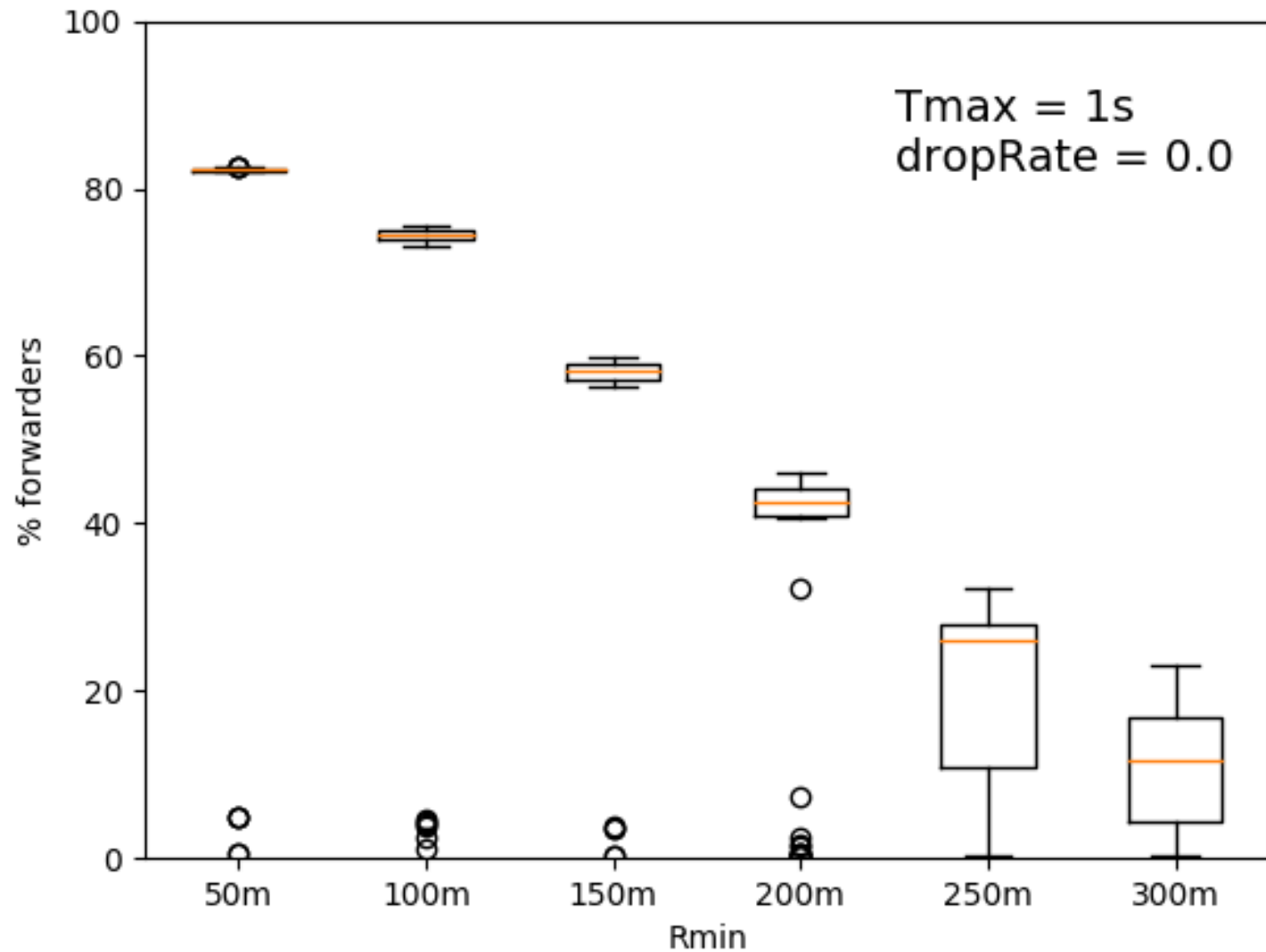
Cologne

- ▶ Percentage of receivers and RMIN.
- ▶ Least dense graph



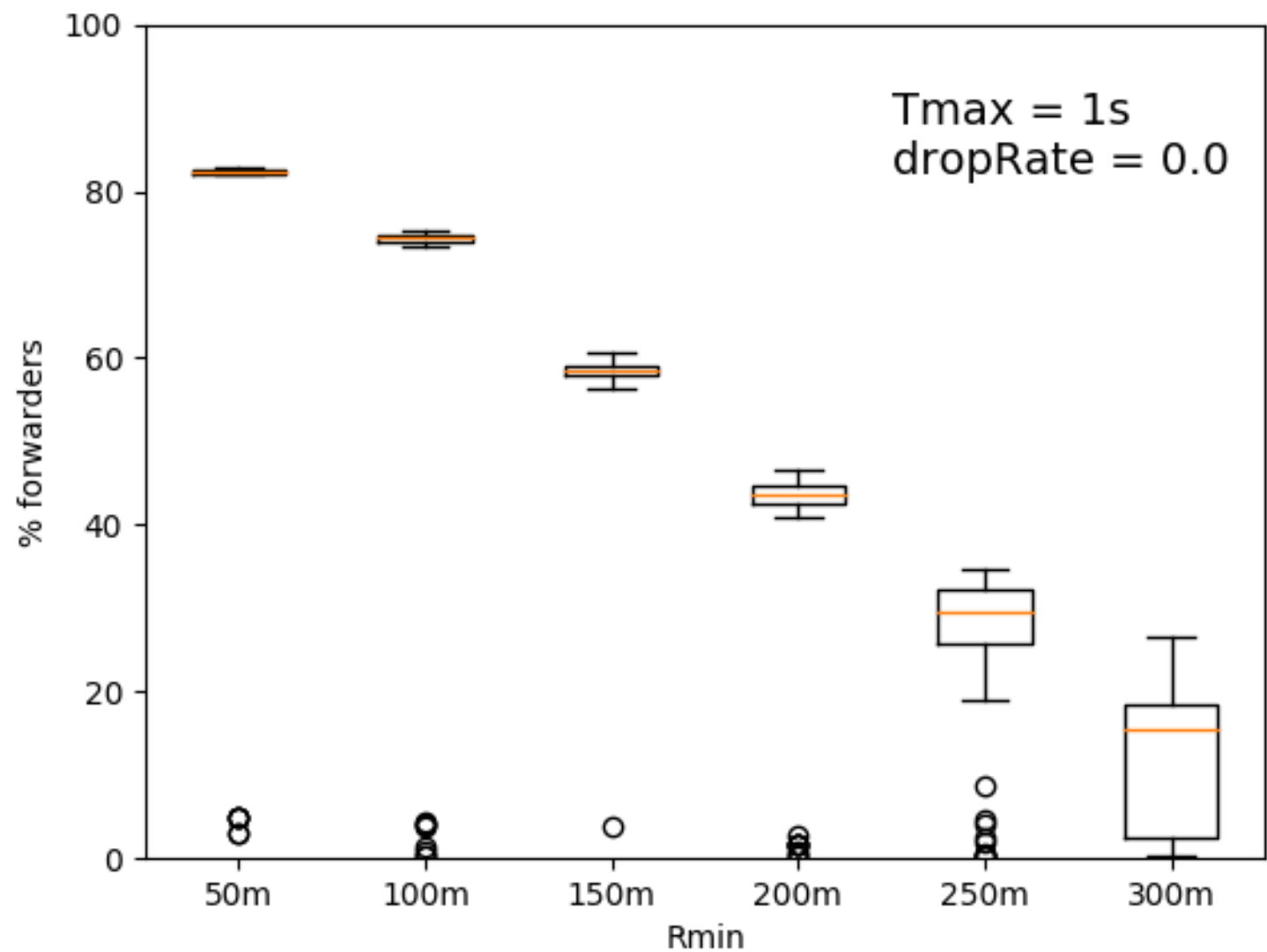
Cologne

- ▶ Percentage of receivers and RMIN.
- ▶ Denser graph



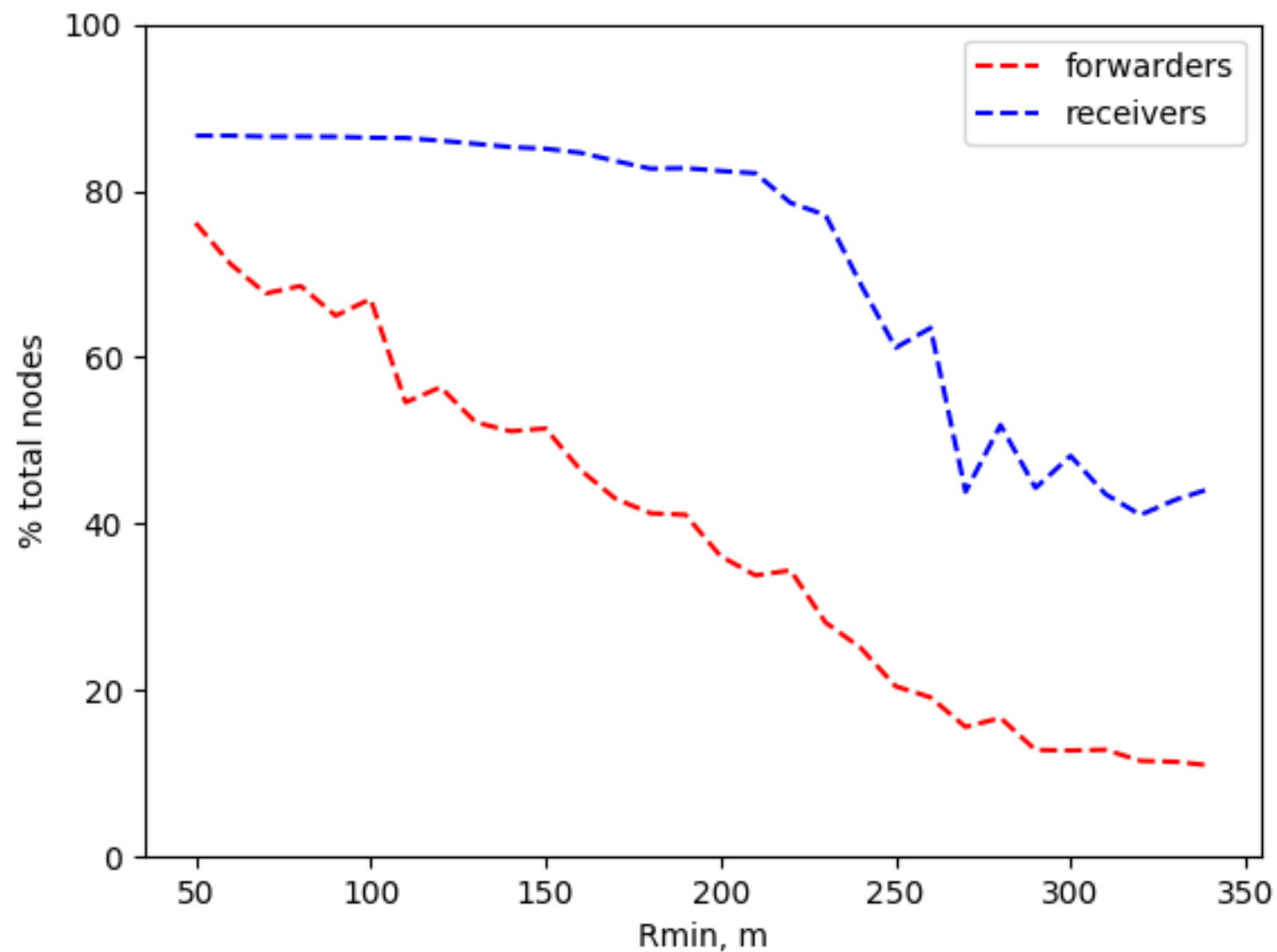
Cologne

- ▶ Percentage of forwarders and RMIN.
- ▶ Least dense graph



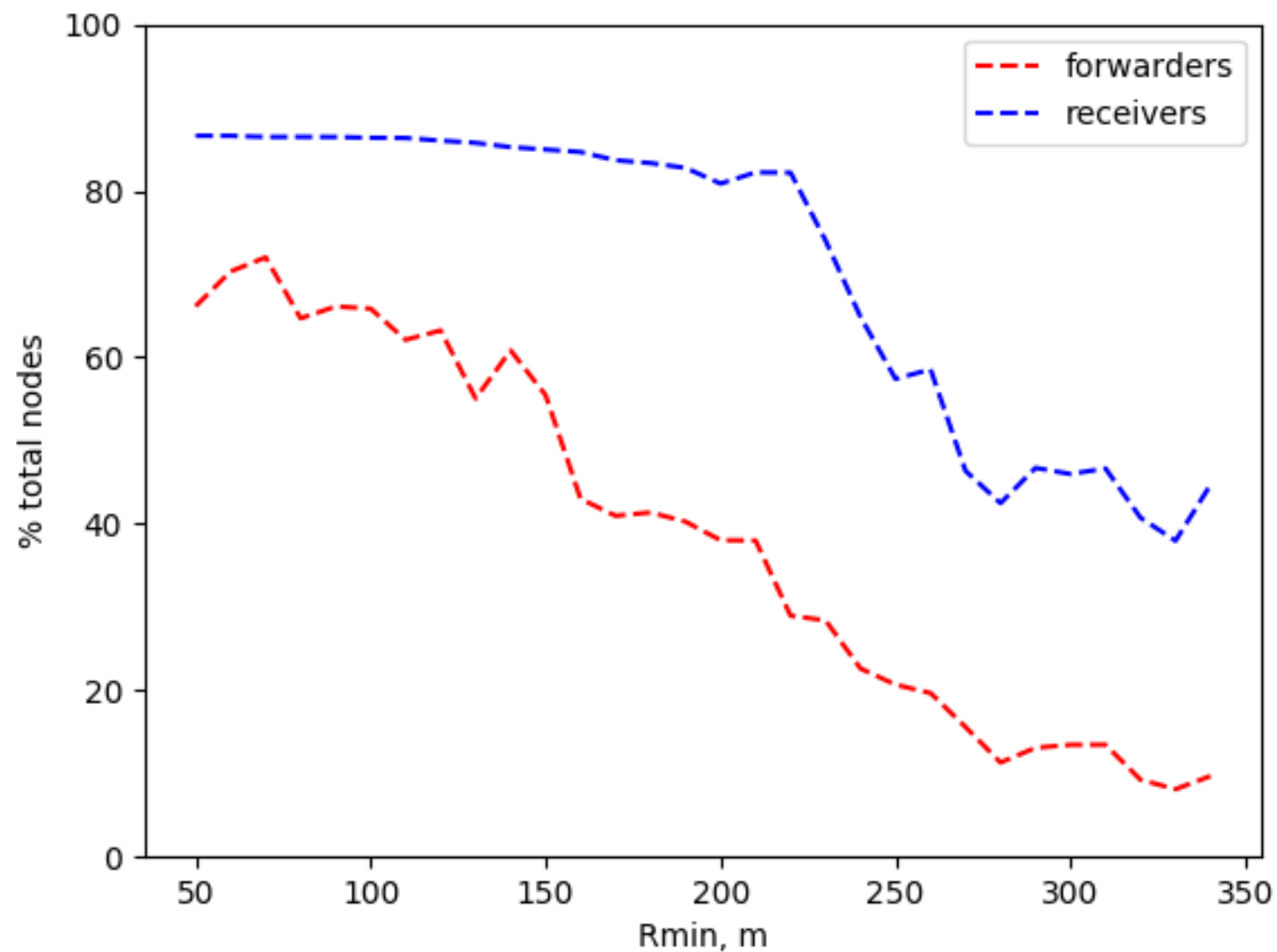
Cologne

- ▶ Percentage of forwarders and RMIN.
- ▶ Denser graph



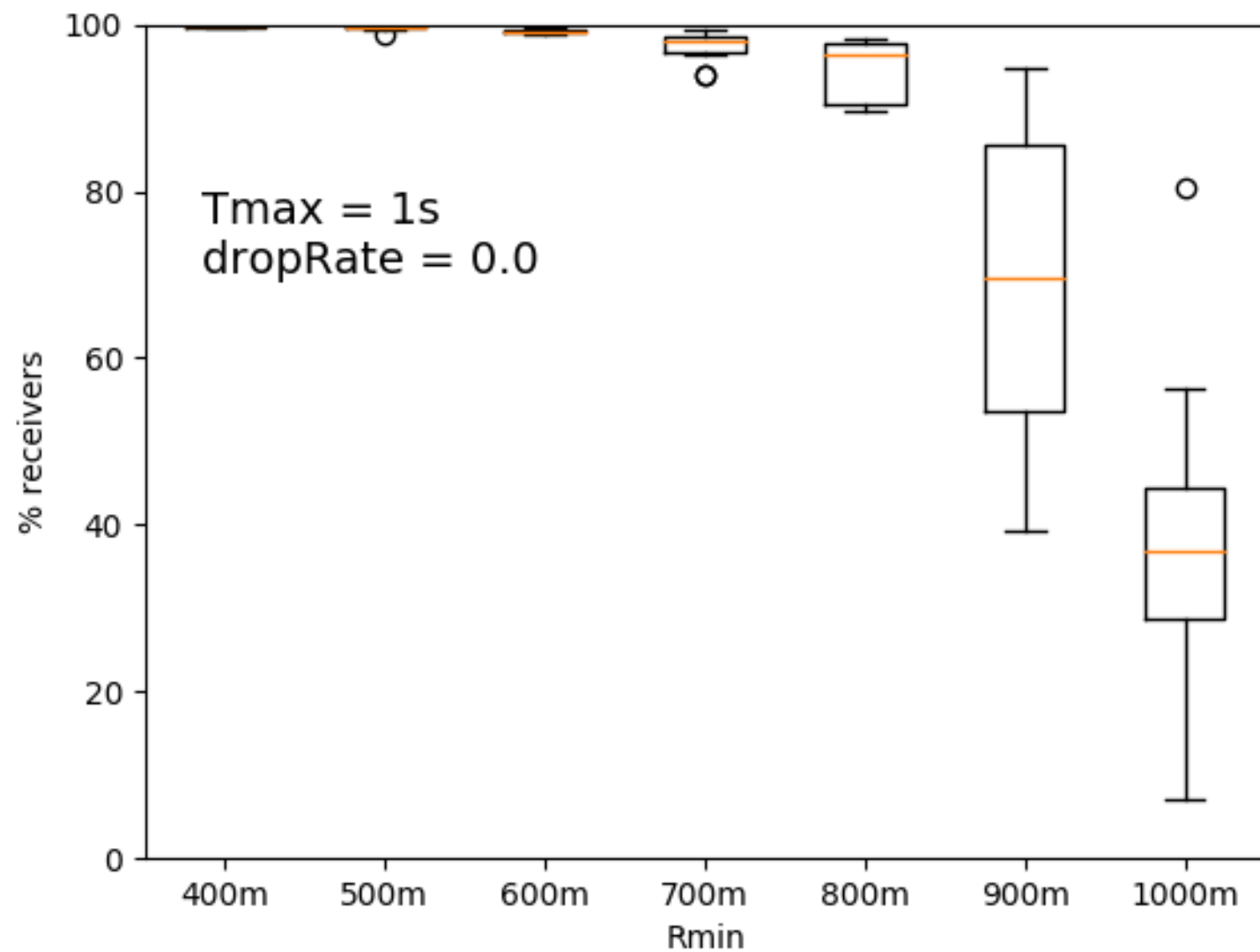
Cologne

- ▶ Forwarders and receivers compared
- ▶ Least dense graph



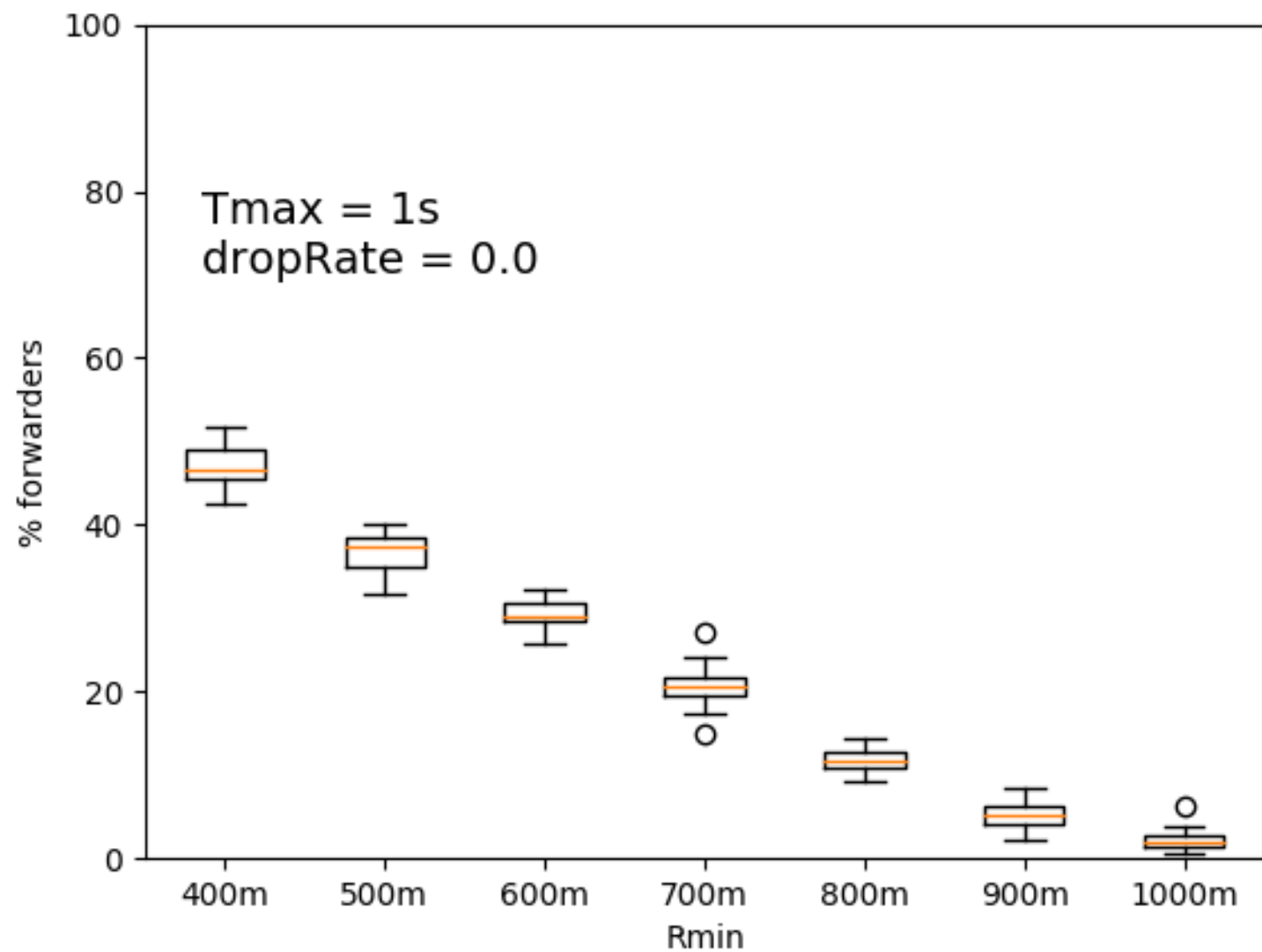
Cologne

- ▶ Forwarders and receivers compared
- ▶ Denser graph



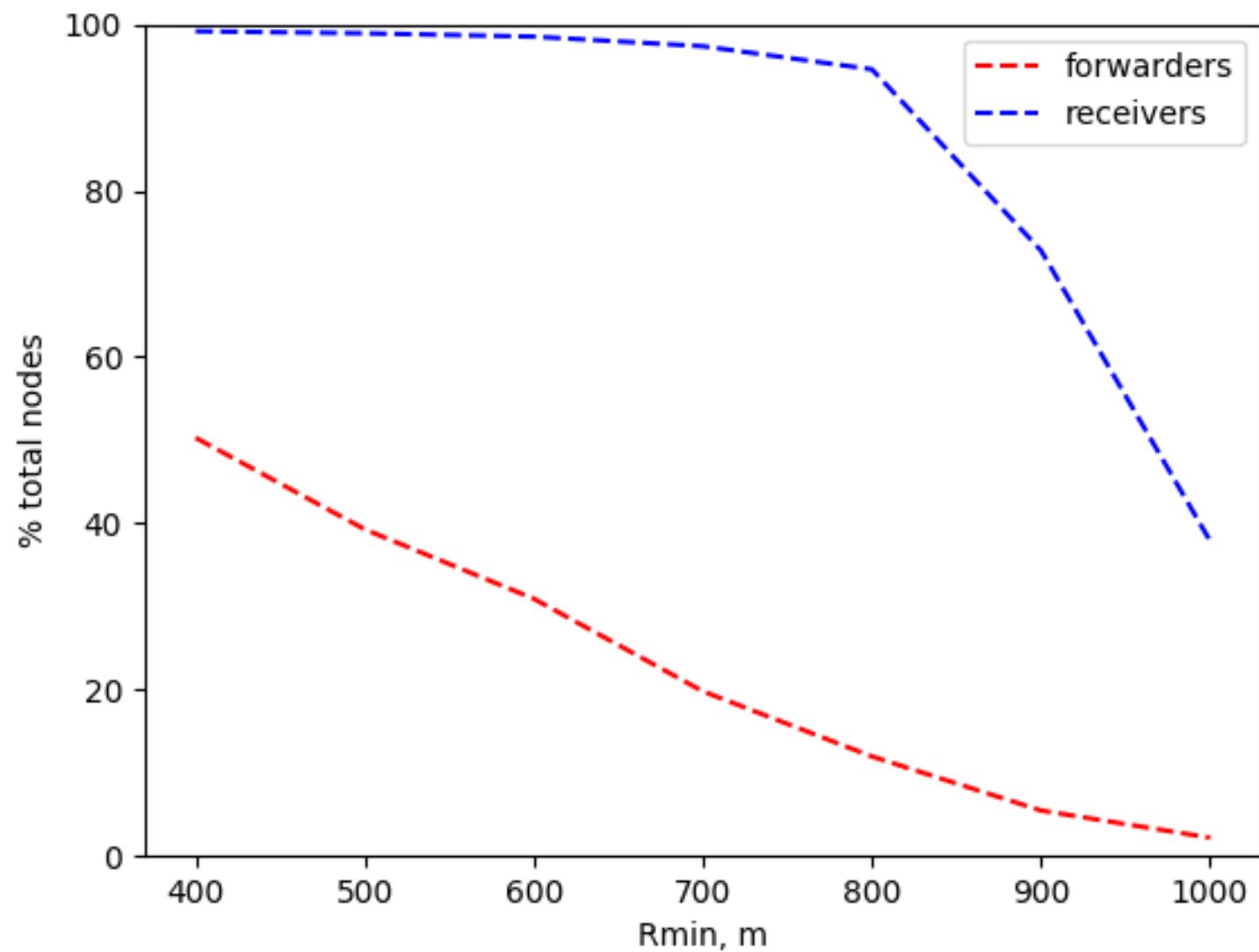
New York

- ▶ Percentage of receivers and RMIN.



New York

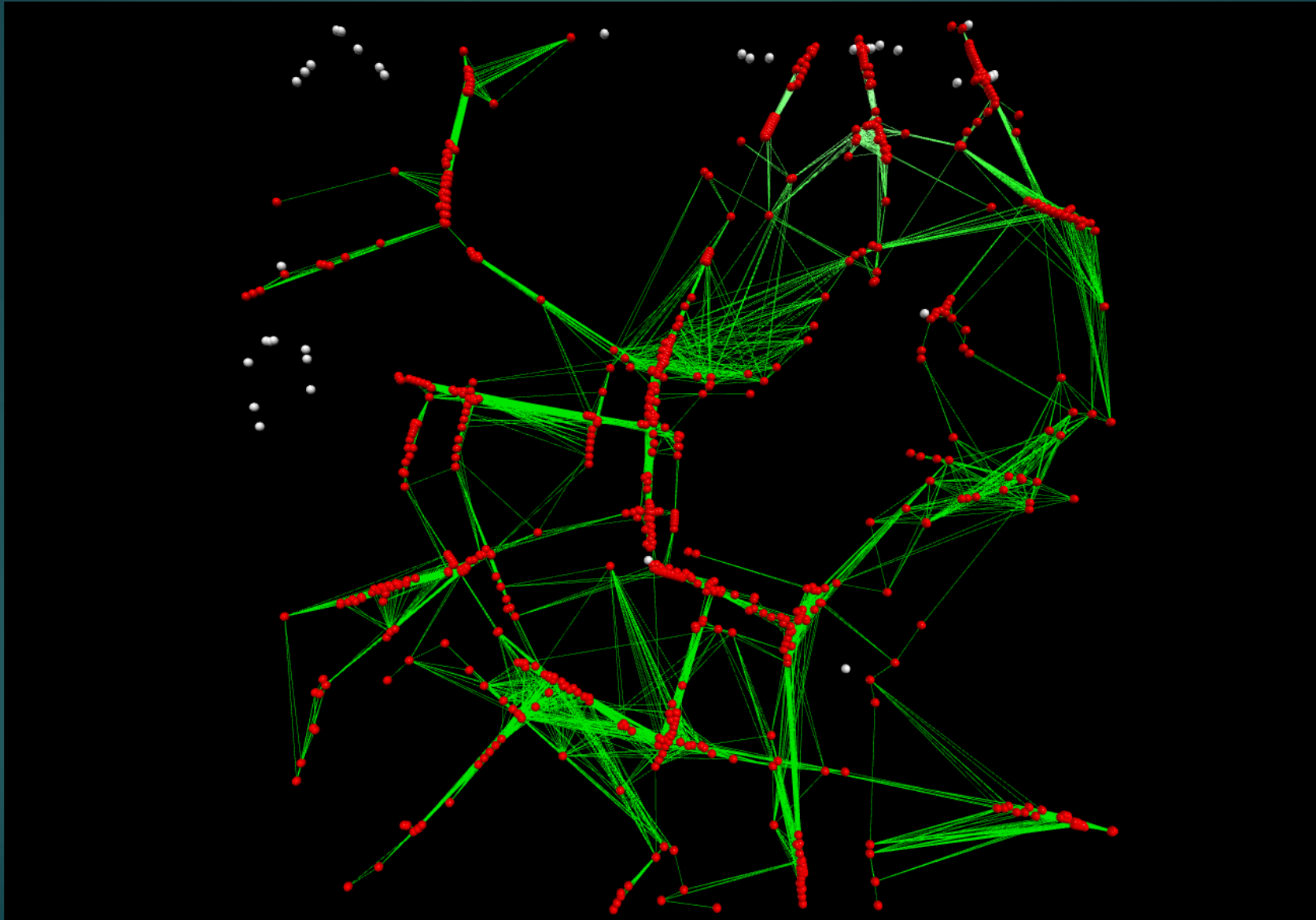
- ▶ Percentage of forwarders and RMIN.



New York

- Forwarders and receivers compared

Result's Presentation - Visual Algorithm



```
<7478.82, 5261.88, 0>
```

```
400
```

```
Simulation ended
```

```
Vulnerable: 36
```

```
Infected: 0
```

```
Recovered: 767
```

```
Average metrics with rmin = 200
```

```
#sent messages: 305.0
```

```
#received messages: 7864.0
```

```
time of last car infection: 7.3000000000000001
```

```
#hops to reach last infected car: 10.0
```

```
Cars infected ratio: 97.38%
```

DEMO