

R1 Compiler Project

Copy your `S1.java` to `R1.java`. Then replace every occurrence of "S1" in `R1.java` with "R1".

Be sure your name is in `R1.java` on line 10.

Modify `R1.java` so that it meets the specifications for R1 described below.

Compile your R1 compiler with

```
javac R1.java
```

Compile `S1.s` (which is in the J1 Software Package) with your R1 compiler with

```
java R1 S1
```

Assemble the output file `S1.a` created by your R1 compiler with

```
a S1.a
```

Finally run the executable program in `S1.e` created by the assembler with

```
e S1 /c
```

Submit `R1.java`, `S1.a` and `S1.<family name>.log` (the log file that the `e` program creates when it runs `s1.e`) in the form requested. If files requested, do not submit a ZIP FILE.

15.1 SPECIFICATIONS FOR R1

Your R1 compiler should compile the program in `S1.s`. Your Java code should correspond to the translation grammar in the file `R1.tg` (in the software package).

See specification of code generator methods on the following page.

The translation grammar for R1 differs from the translation grammar for S1 in two significant ways. First, it does not directly call the `emitInstruction()` method in the code generator. `emitInstruction()` is now a private method in the code generator. For example, to generate code for the assignment statement, the R1 parser calls the `assign()` method in the code generator (see line 32 in Fig. 19.6) which, in turn, calls `emitInstruction()`. Second, the `factor()` method does not generate any code. It simply returns a symbol table index. The code to load a factor into the `ac` register is emitted by either the `add()` method in `termList()` or the `mult()` method in `factorList()`.

The methods in the R1 code generator that are not in the S1 code generator are

```
public void directive(String d)
    Outputs the directive d to the assembly language file. The program()
    method in the parser calls the directive() method with

        codeGen.directive("!register");

    to output the !register directive (see line 5 in Fig. 19.6). This
    directive causes the a and e programs to reconfigure to the register
    instruction set.

public void assign(int left, int expVal)
    Outputs the ld-st sequence for an assignment statement by making calls to
    emitInstruction() (see line 26 in Fig. 19.6).

public void println(int expVal)
    Outputs the ld-dout-ldc-aout sequence for a println statement by
    making calls to emitInstruction() (see line 36 in Fig. 19.6).

public void add(int left, int right)
    Outputs the ld-add-st sequence for an add operation by making calls to
    emitInstruction() (see Fig. 19.5a and line 57 in Fig. 19.6)

public void mult(int left, int right)
    Outputs the ld-mult-st sequence for a multiplication operation by making
    calls to emitInstruction() (see line 57 in Fig. 19.6). mult() is
    similar to add().

private int getTemp()
    Generates the next temporary variable from the sequence "@t0", "@t1",
    "@t2", ..., enters it into the symbol table, and returns its symbol
    table index (see Fig. 19.5b). getTemp() is called by add() and mult()
    (see Fig 19.5a).

private void emitInstruction(String op, int opndIndex)
    The code in this method consists of a call to another emitInstruction()
    method (the one with two String parameters):

        emitInstruction(op, symTab.getSymbol(opndIndex));

    R1 also has a third emitInstruction() method (the one with a single
    parameter whose type is String).
```