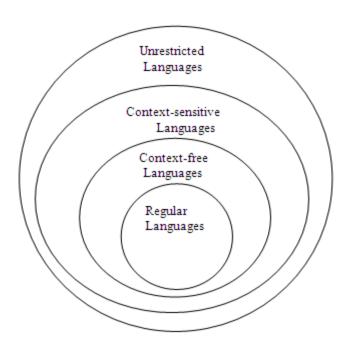
Chapter 5

Chomsky's Hierarchy

Chomsky's hierarchy



Context-sensitive produtions

```
S -> bSc
bS -> beA (replaces S with eA)
Ac -> Aec (replaces A with Ae)
Ae -> Aee (replaces A with Ae)
eAe -> ede (replaces A with d)
```

Also allowed productions like

BC -> CB

Context-sensitive grammar

Any essentially non-contracting grammar. Context-sensitive productions are allowed as long as they are non-contracting

Theorem

All context-free languages are context sensitive.

Proof:

Eliminate lambda productions.

Reinsert lambda if in original language.

Result is an equivalent CSG.

Example

Convert to CSG:

S -> bS

 $S \rightarrow \lambda$

Eliminate lambda production

S -> bS

S -> b

Reinsert null string. Get CSG:

S' -> S

 $S' \rightarrow \lambda$

S -> bS

S -> b

Unrestricted grammars

No restriction on productions except left side must be non-null.

Can define a language with an unrestricted grammar if and only if it can be defined by a Turing machine.