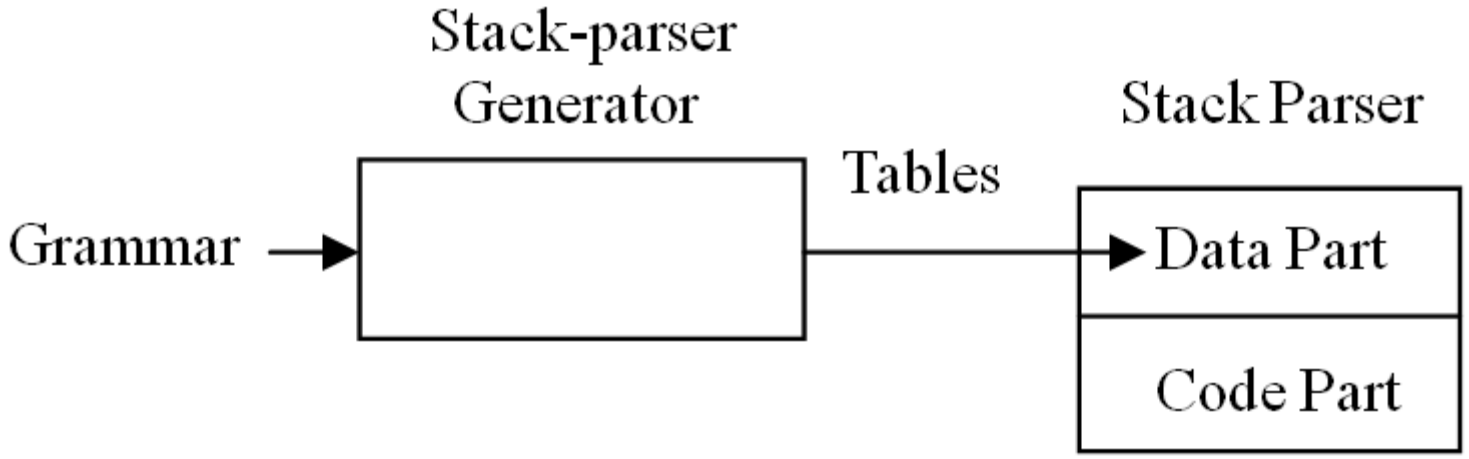
Chapter 8

Table-Driven Stack Parser

Parser generator



Uniform stack parser

Handle all productions in the same way:

1. pop
2. push entire right side

If the top of the stack is a terminal and it matches the current input, then pop and advance.

Construct uniform parser for

G8.1

		Selection Set
0)	S → fBC	{f}
1)	B → bb	{b}
2)	B → CD	{c, d, e}
3)	C → cC	{c}
4)	C → λ	{d, e, #}
5)	D → dD	{d}
6)	D → e	{e}

Parse table

Current token

Symbol on top of stack

	b	c	d	e	f	#
S					pop push(C) push(B) push(f)	
B	pop push(b) push(b)	pop push(D) push(C)	pop push(D) push(C)	pop push(D) push(C)		
C		pop push(C) push(c)	pop	pop		pop
D			pop push(D) push(d)	pop push(e)		
b	pop advance					
c		pop advance				
d			pop advance			
e				pop advance		
f					pop advance	
$						accept

Abbreviated parse table

	b	c	d	e	f	#
S					0	
B	1	2	2	2		
C		3	4	4		4
D			5	6		

Implementing uniform stack parser

```
int[ ][ ] parseTable  =
{
    {-1, -1, -1, -1,  0, -1},
    { 1,  2,  2,  2, -1, -1},
    {-1,  3,  4,  4, -1,  4},
    {-1, -1,  5,  6, -1, -1}
};
```

Representing productions

```
String [ ] pTab =
{
    "CBf",
    "bb" ,
    "DC" ,
    "Cc" ,
    "" ,
    "Dd",
    "e"
};
```

Mapping terminal/nonterminal to index

```
String nonTerms = "SBCD";

String tokens = "bcdef#";
```

Table-driven stack parser

Fig0804.txt

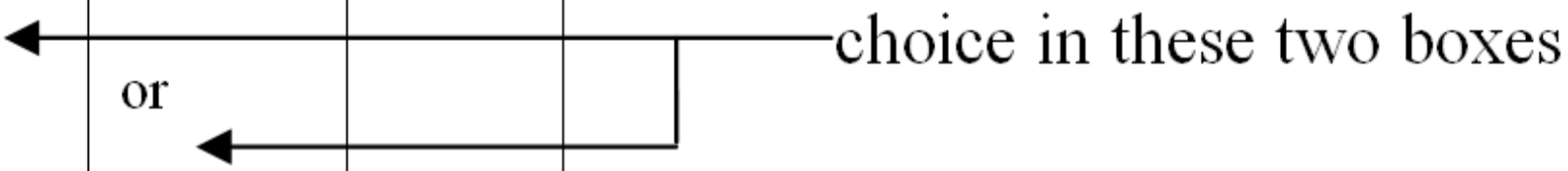
Non-deterministic parsers

G8.2

		Selection Set
1)	S → bSb	{b}
2)	S → cSc	{c}
3)	S → λ	{b, c, #}

Parse table has choices

	b	c	#
S	pop push(b) push(S) push(b) or pop	pop push(c) push(S) push(c) or pop	pop
b	pop advance		
c		pop advance	
$			accept



Interesting theoretical result

A language can be defined by a context-free grammar if and only if it can be defined by a stack parser.

In automata theory, the abstract model of a stack parser is called a pushdown automaton.