Course in Semantics · Ling 531 / 731 McKenzie · University of Kansas

## 1 Abstract functions

- **1.** Complete the following applications/ $\beta$ -reductions, until you run out of arguments.
  - 1. [  $\lambda f \in D_{\langle e, t \rangle}$ . f(x) ](Q) = Q(x)
  - 2.  $[\lambda g \in D_{\langle e, t \rangle}, g(x)](Q) = Q(x)$
  - 3.  $[\lambda f \in D_{\langle e, t \rangle}.\lambda x \in D_e. g(x)](Q) = \lambda x \in D_e. Q(x)$
  - 4.  $[\lambda f \in D_{\langle e, t \rangle}.\lambda x \in D_e. g(x)](Q)(x) = Q(x)$
  - 5. [  $\lambda f \in D_{(e,t)}.\lambda x \in D_e$ . g(x) ](x)(Q) = (watch out!) x(Q); the function is x and the argument is Q
  - 6.  $[\lambda f \in D_{(e, t)}.\lambda x \in D_e. g(f(x))](A)(b) = g(A(b))$

## 2 Similar but with natural language

- 1.  $[\lambda f \in D_{\langle e, t \rangle}, f(y)](\lambda x \in D_e, cat(x)) = [\lambda x \in D_e, cat(x)](y) = cat(y)$
- 2.  $[\lambda f \in D_{\langle e, t \rangle}. f(x)](\lambda y \in D_e. happy(y)) = [\lambda y \in D_e. happy(y)](x) = happy(x)$
- 3.  $[\lambda f \in D_{\langle e, t \rangle}. f(Marie)](\lambda x \in D_e. happy(x)) = [\lambda x \in D_e. happy(x)](Marie) = happy(Marie)$
- 4. [  $\lambda g \in D_{\langle e, t \rangle}$ . $\lambda x \in D_e$ . g(x) ]([ broad ])(the Mississippi River) = [ $\lambda x \in D_e$ . broad(x)](MR) = broad(MR)
- $$\begin{split} 5. & [ \ \lambda f \in D_{\langle e, \ t \rangle}. \lambda x \in D_e. \ Q(f(x))](\lambda y \in D_e. \ Greek(y))(Apollo) = \\ & [\lambda x \in D_e. \ Q([\lambda y \in D_e. \ Greek(y)](x))](Apollo) = \\ & Q([\lambda y \in D_e. \ Greek(y)](Apollo)) = \\ & Q(Greek(Apollo)) \end{split}$$