Announcements

- Homework 0 online, due F
- Course Discord linked on course webpage
- Remember to fill out FinAid survey

Today

- Levels of Algorithm Design
- Introduction to graph algorithms
 - Graph basics and representation
 - Explore/DFS

Naive Algorithms: Turn definition into algorithm easy to write, good first pass, often very slow

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Magic: Sometimes, an algorithm requires a surprising new insight

Graph and Connectivity (Ch 3)

- Graph basics and representation
- Depth First Search
- Connected components
- Pre- and Post- orderings
- DAGs / Topological Sort
- General directed graphs & strongly connected components

Graph Definition

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A collection V of vertices, or objects to be connected.

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Definition: A *graph* G = (V,E) consists of two things:

- A collection V of *vertices*, or objects to be connected.
- A collection E of edges, each of which connects a pair of vertices.

Question: Which are graphs?

Which of the following can't be modeled by a graph? (multiple correct answers)

- A) The internet, V = {websites}, E = {links}
- B) The internet, V = {computers},
 E = {physical connections}
- C) UCSD, V = {students}, E = {classes}
- D) Highway System, V = {intersections}, E = {roads}
- E) A book, $V = \{words\}$

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Examples of Graphs in CS

- The Internet (either webpages, or physical connections)
- Social Networks
- Transitions between states of a program
- Road maps

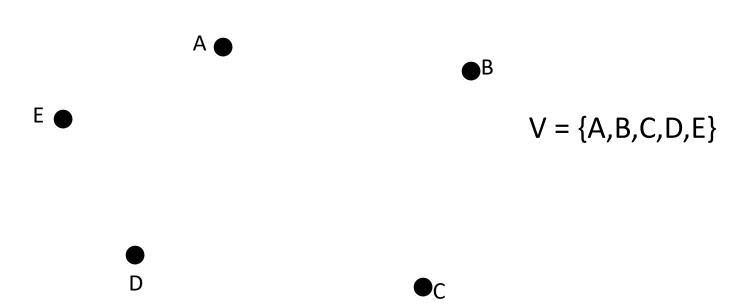
Draw vertices as points

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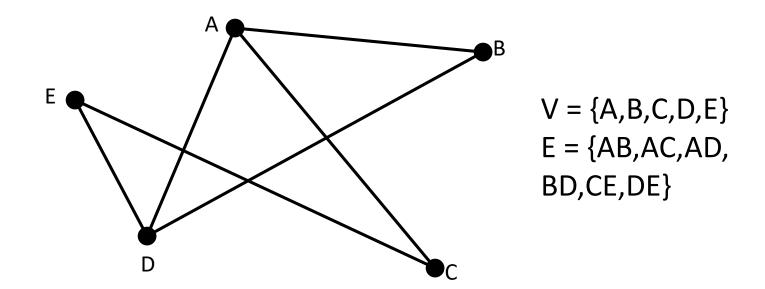
 \bullet^{B} $\mathsf{V} = \{\mathsf{A},\mathsf{B},\mathsf{C},\mathsf{D},\mathsf{E}\}$

OC

- Draw vertices as points
- Draw edges as line segments or curves connecting those points



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Exploring Graphs

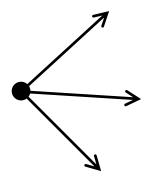
You're playing a video game and want to make sure that you've found all the areas in this level before moving on to the next one.

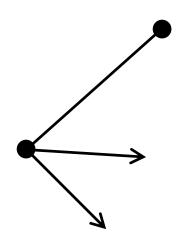
Exploring Graphs

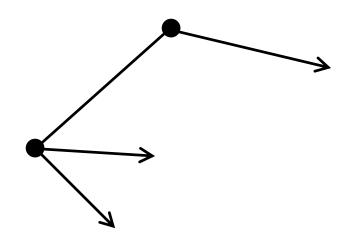
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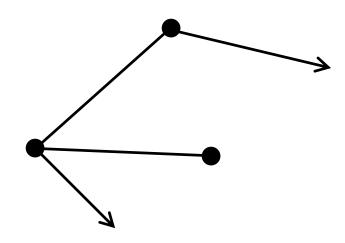
How do you ensure that you have found everything?

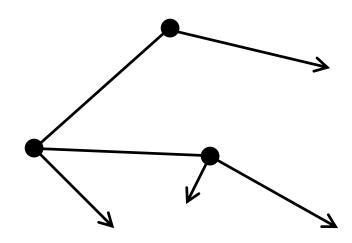
```
Keep track of all areas discovered
While there is an unexplored path,
  follow path
```

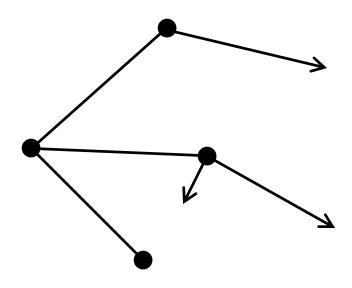


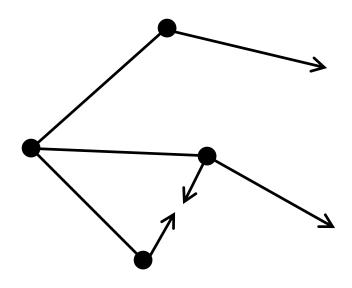


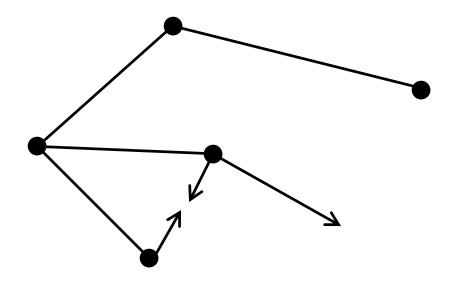


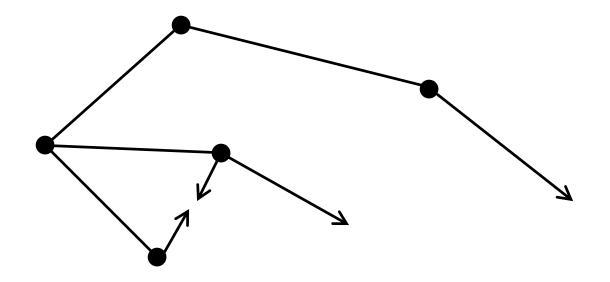


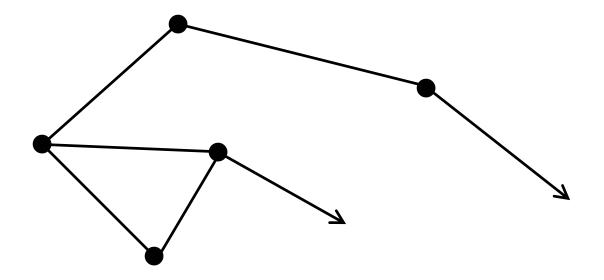


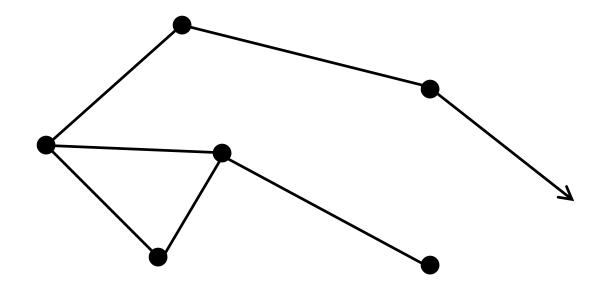


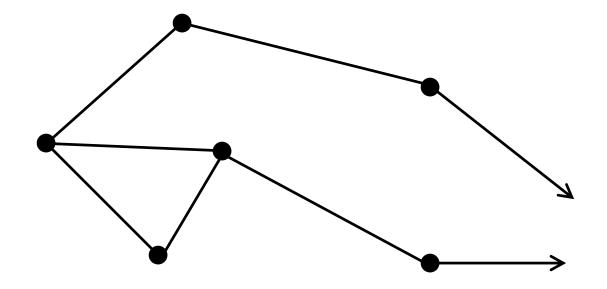


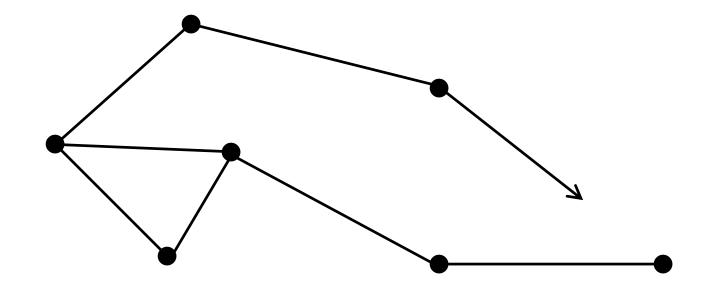


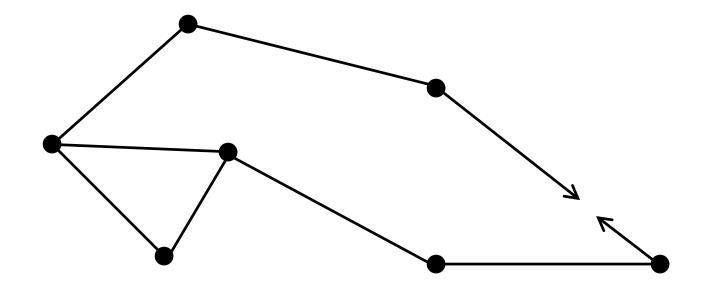






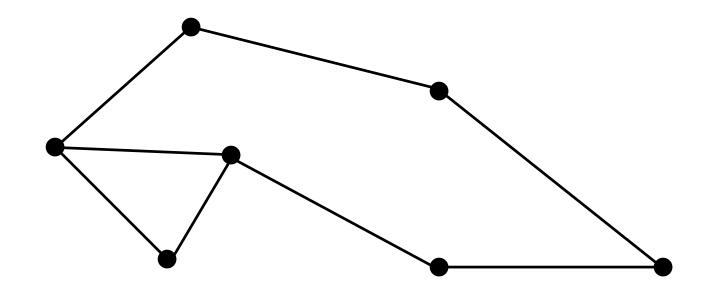






Basic Algorithm

Keep track of all areas discovered
While there is an unexplored path,
 follow path



Systematize

Need to keep track of:

- Which vertices discovered
- Which edges have yet to be explored

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explore Algorithm will:

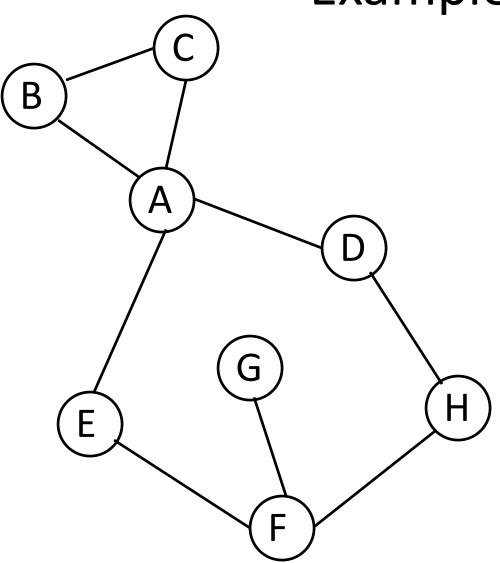
- Use a field v.visited to let us know which vertices we have seen.
- Store edges to be explored implicitly in the program stack.

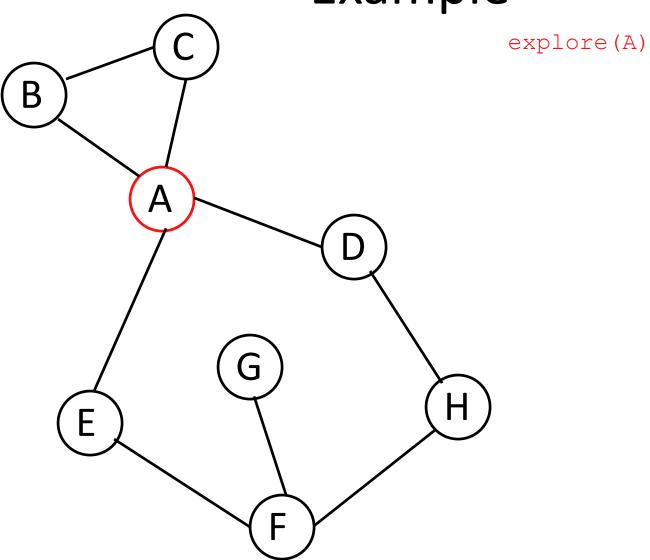
Explore

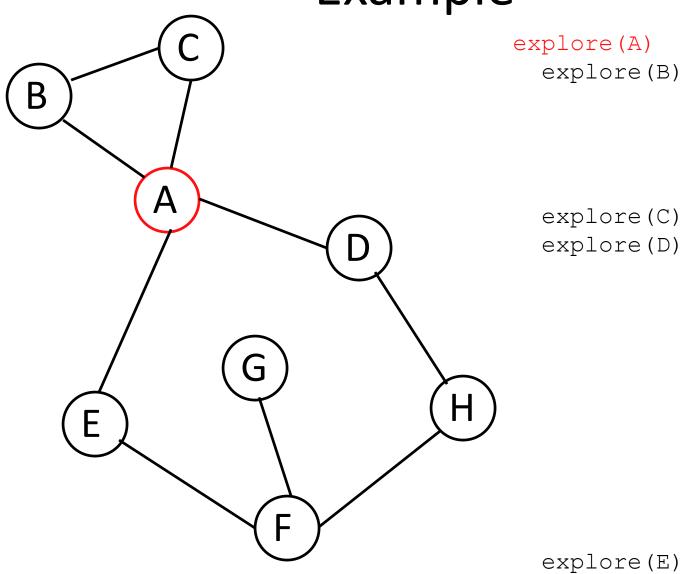
```
explore(v)
  v.visited ← true
  For each edge (v, w)
    If not w.visited
      explore(w)
```

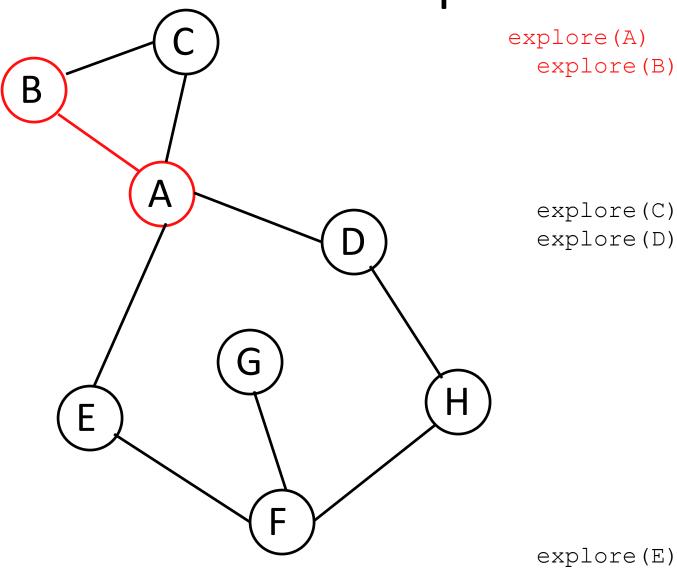
Explore

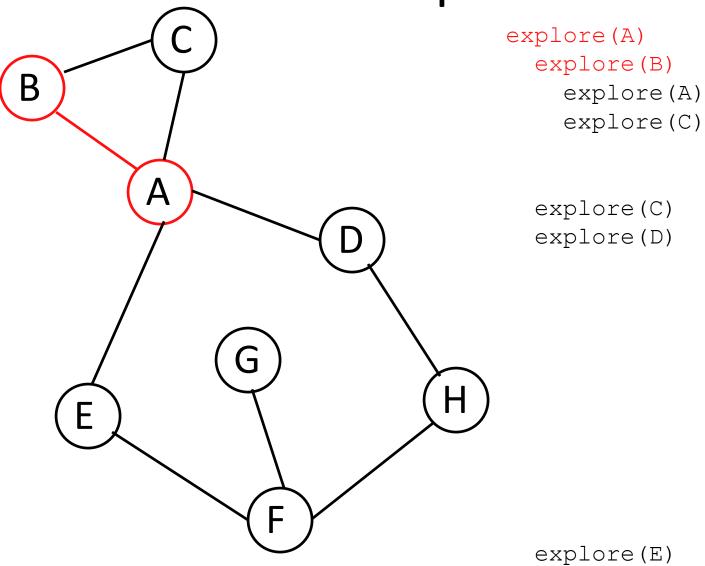
```
explore(v)
  v.visited ← true
  For each edge (v, w)
    If not w.visited
      explore(w)
      w.prev ← v
```

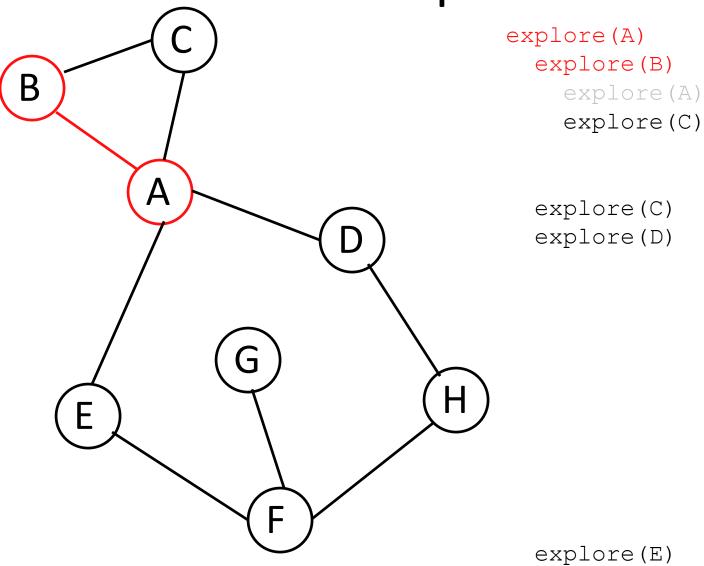


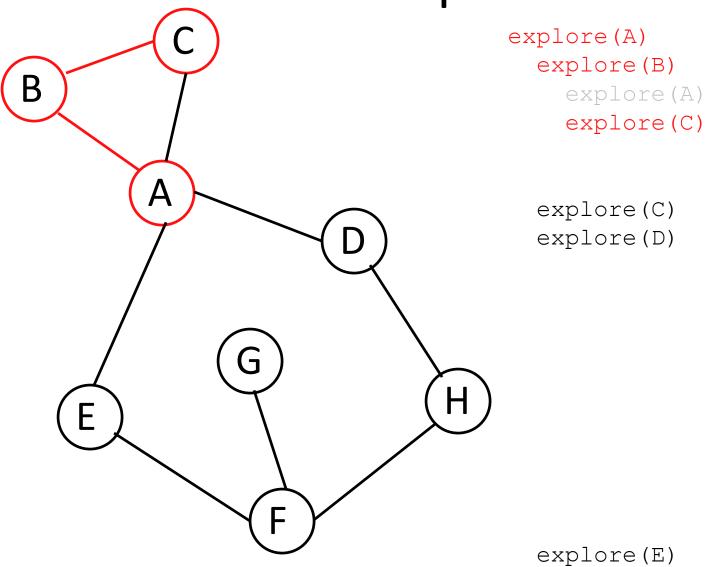


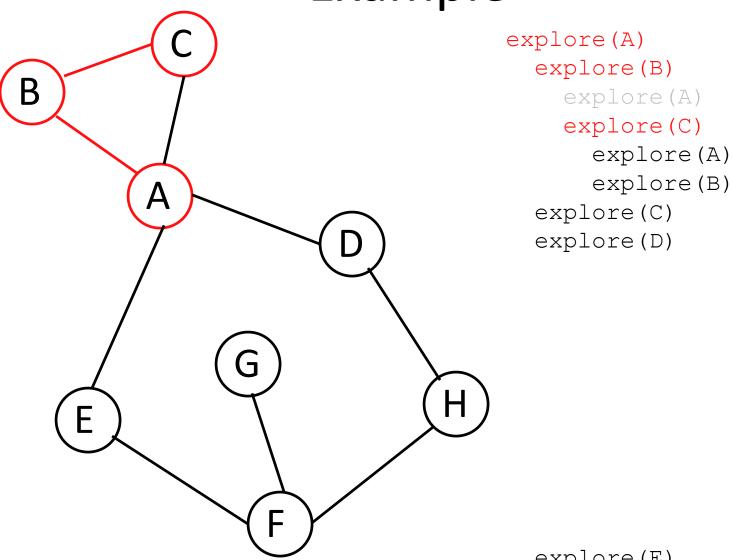


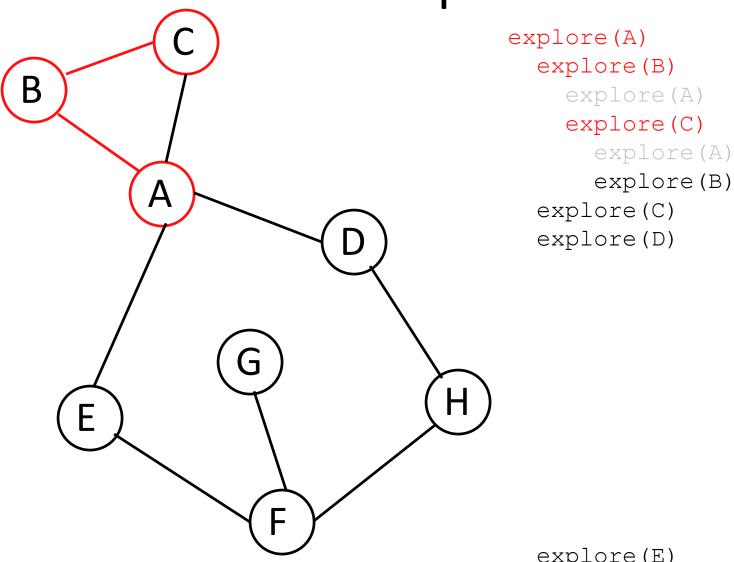


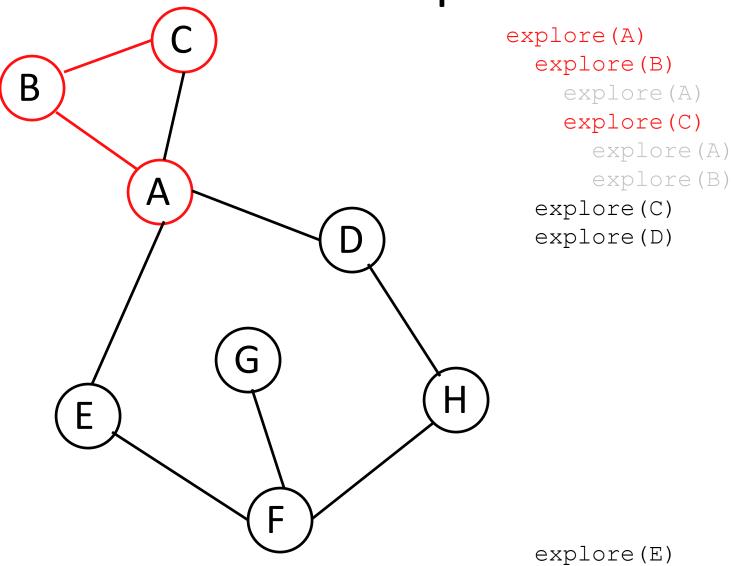


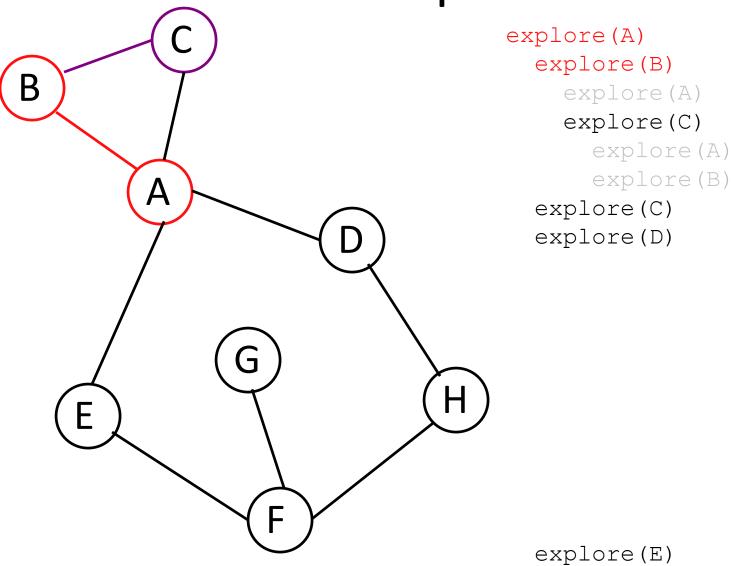


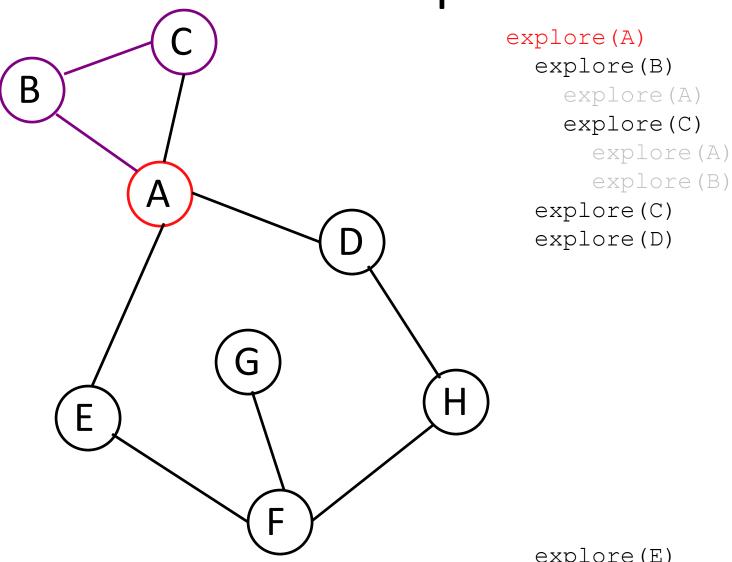


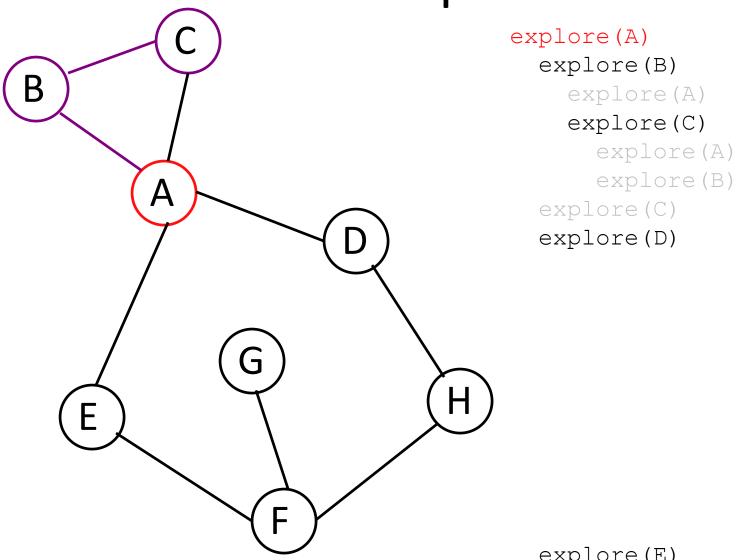


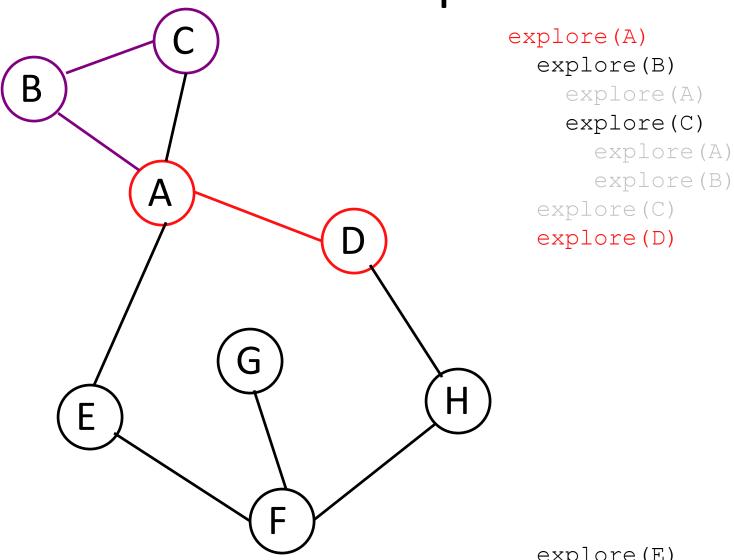


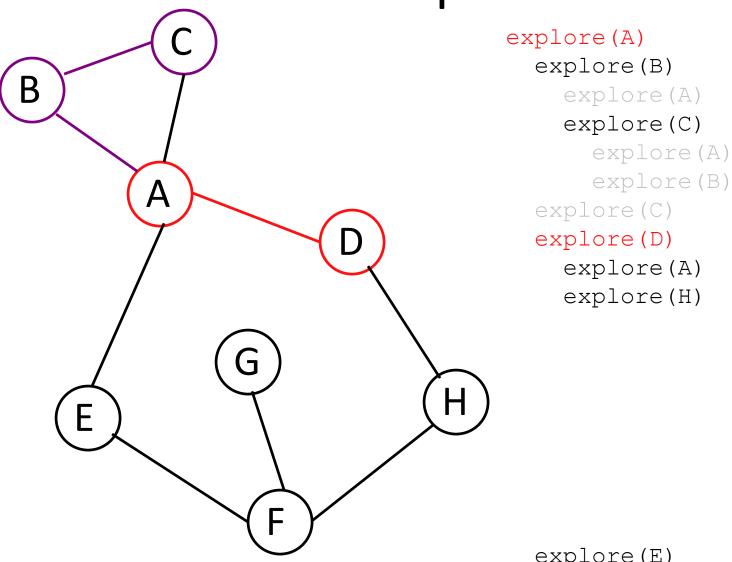


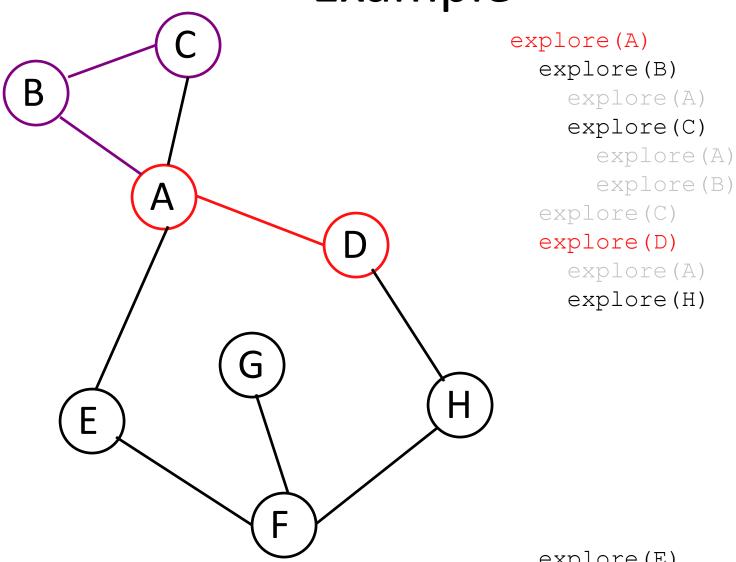


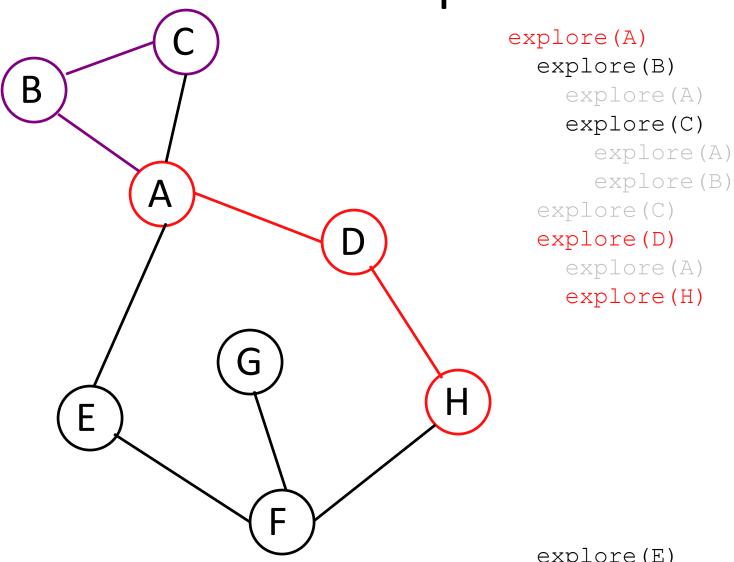


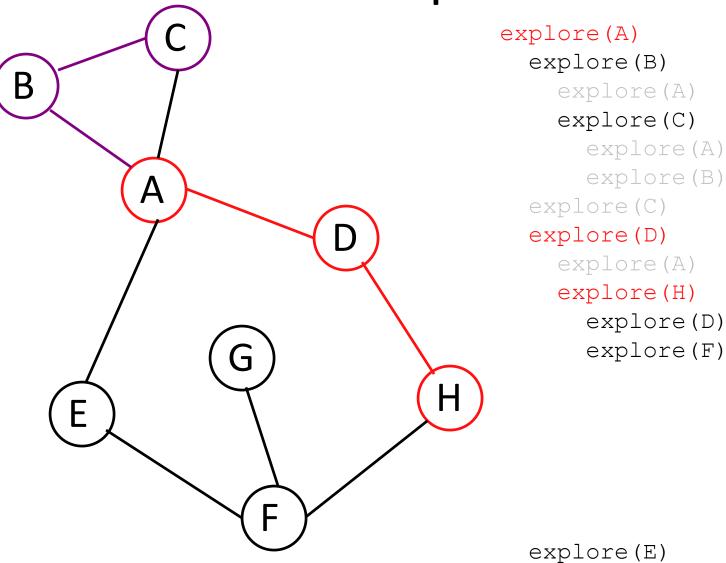


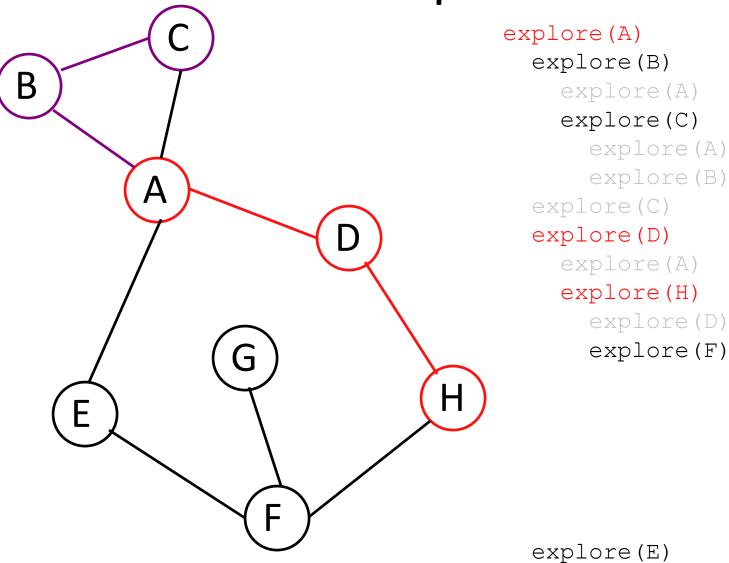


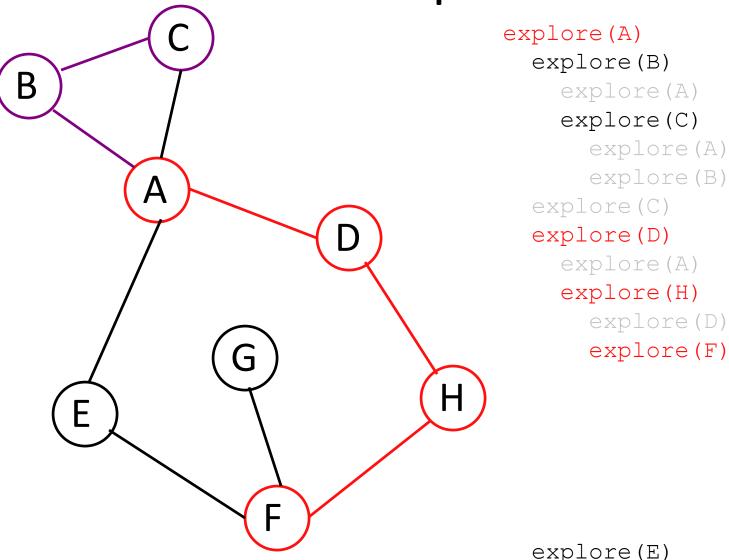


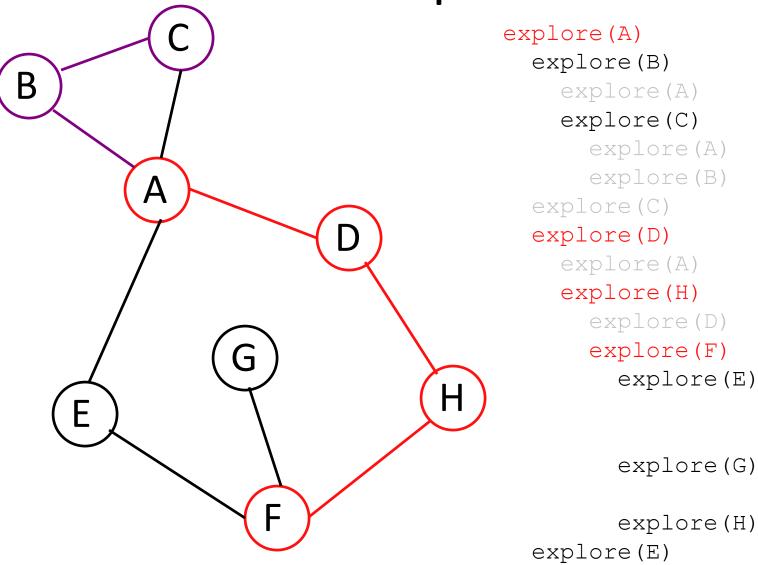


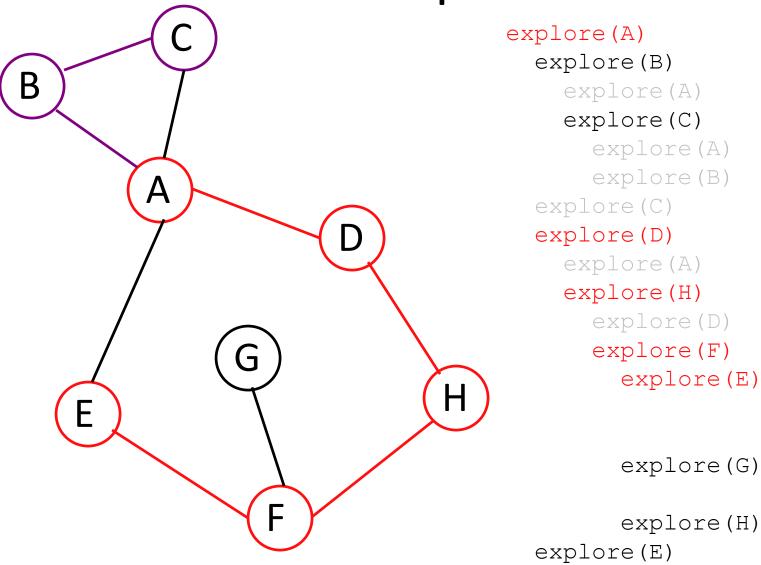


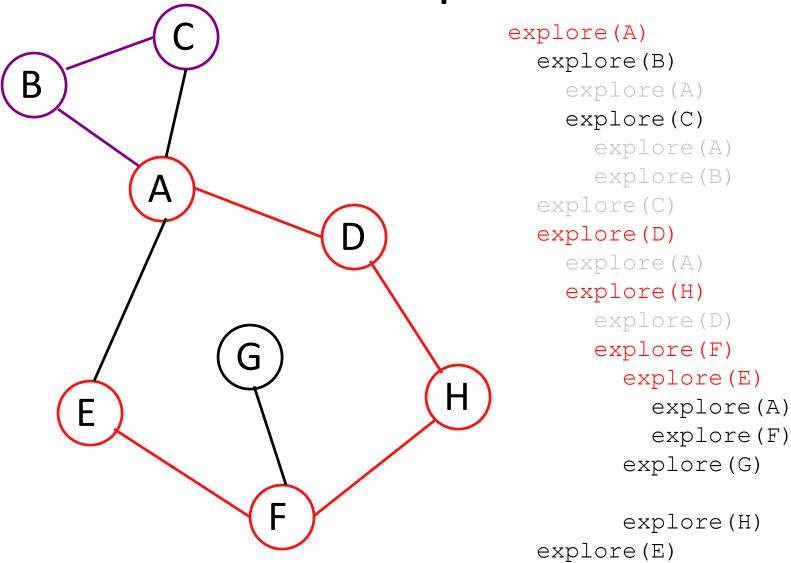


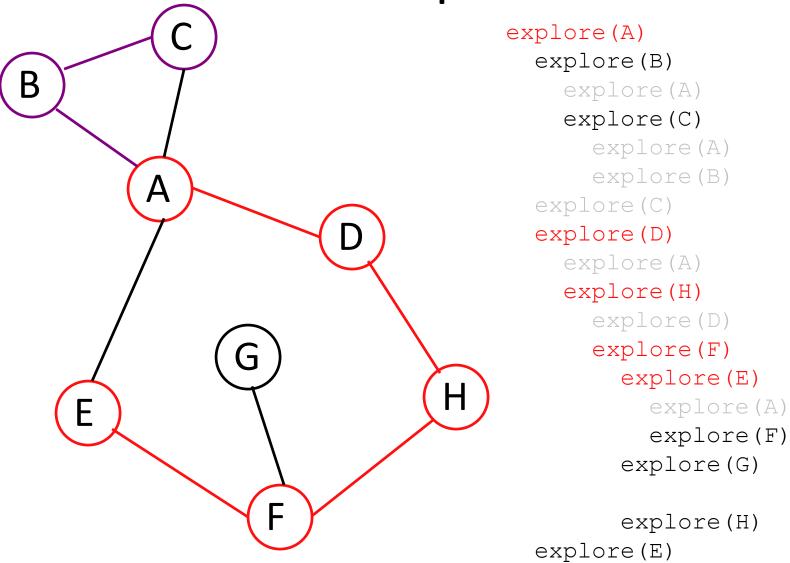


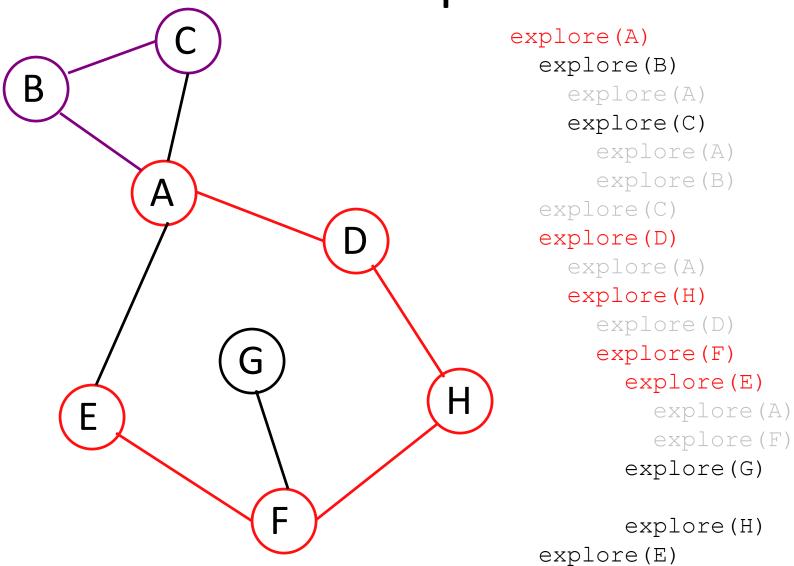


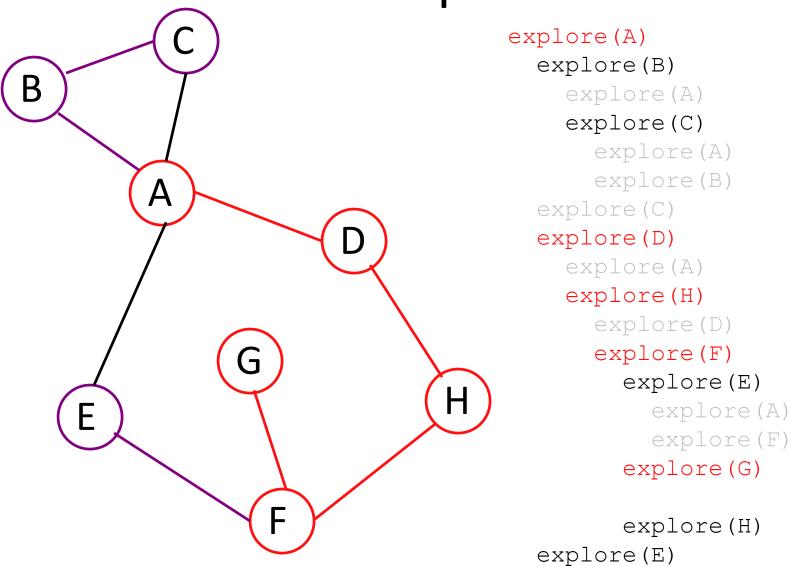


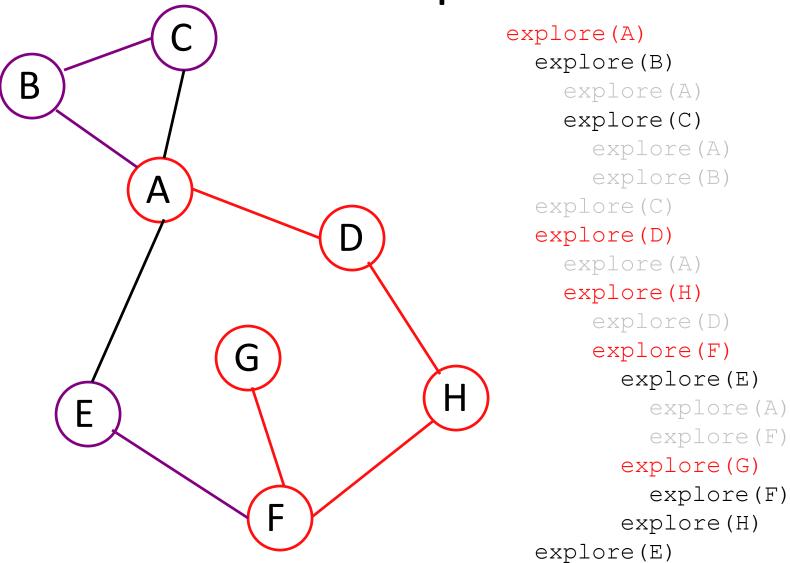


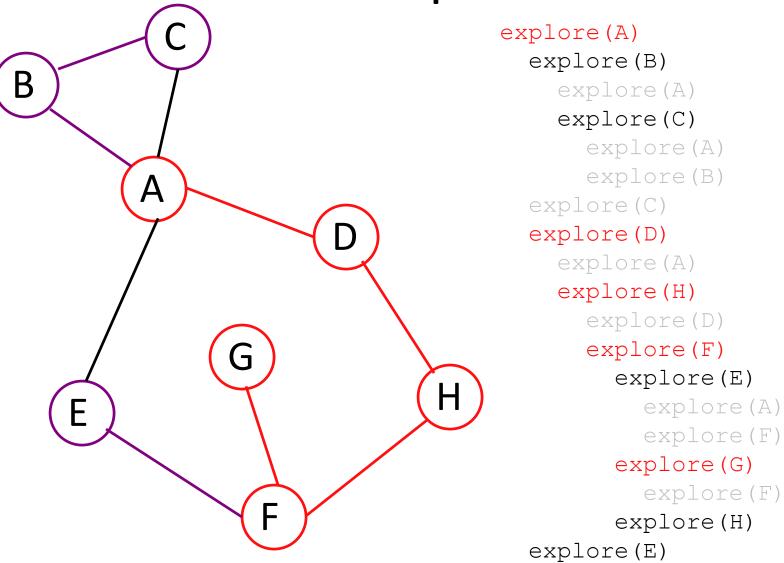


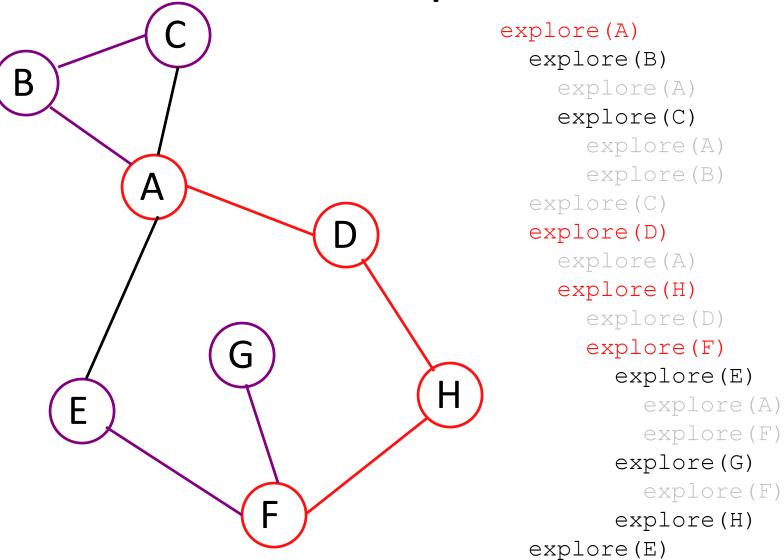


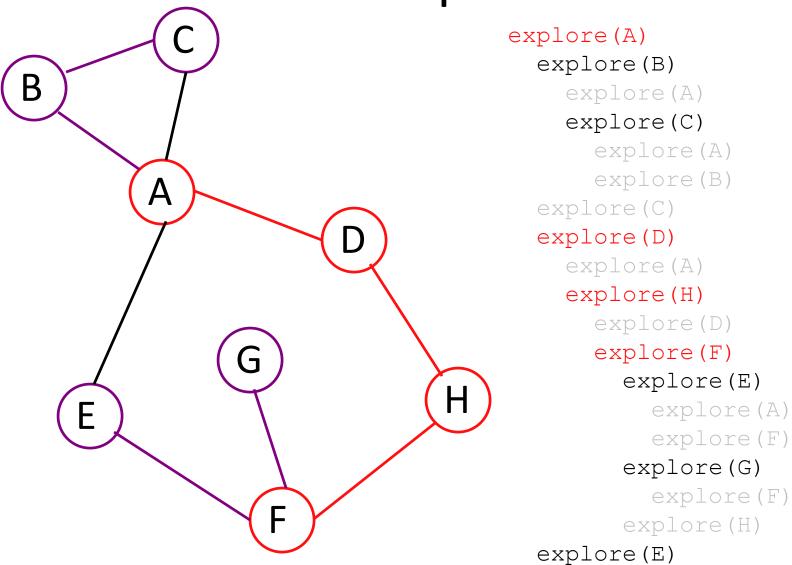


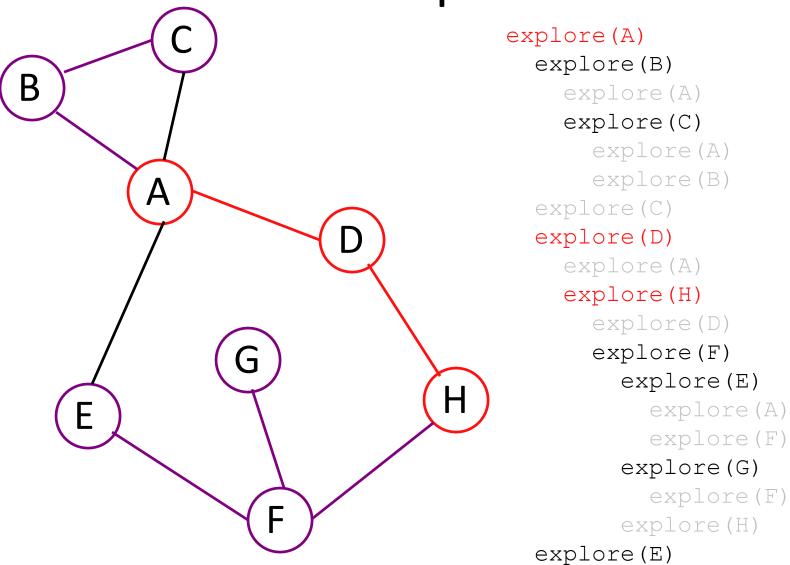


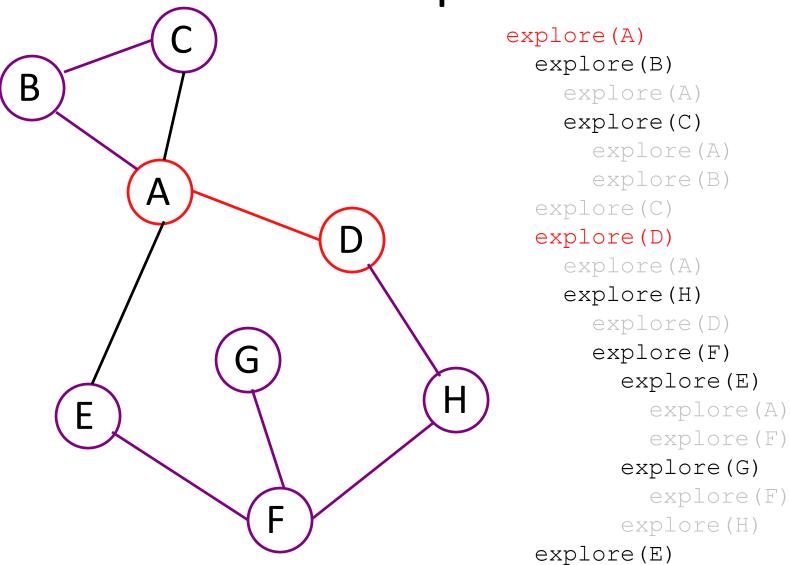


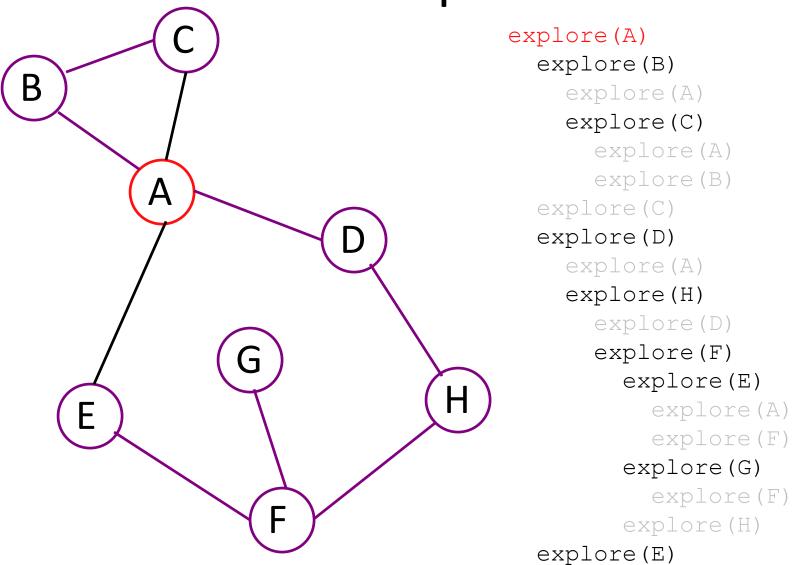


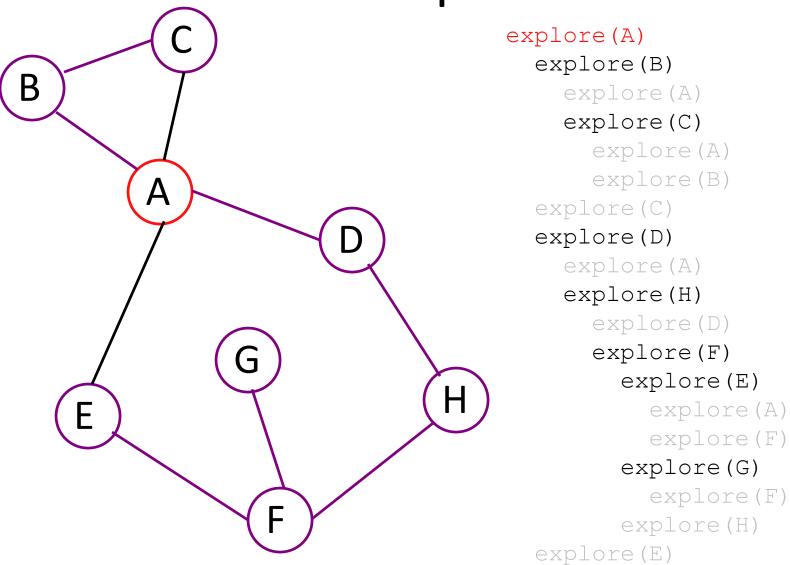


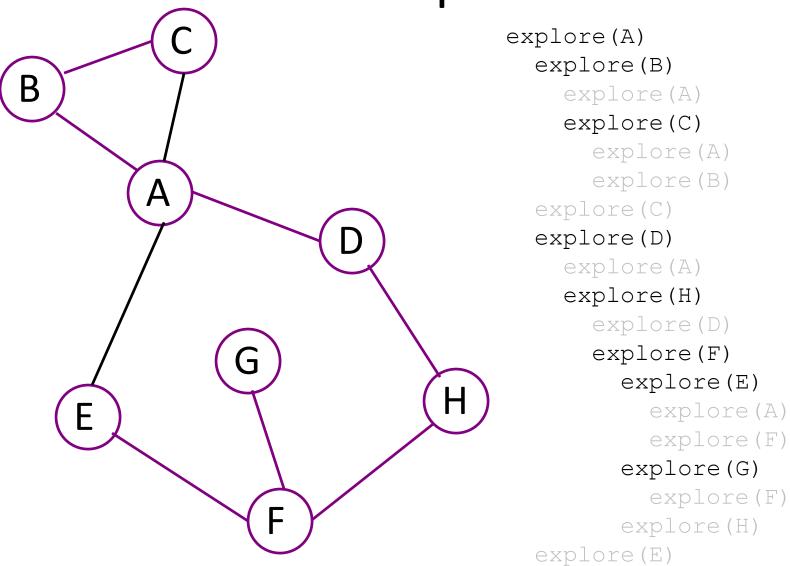




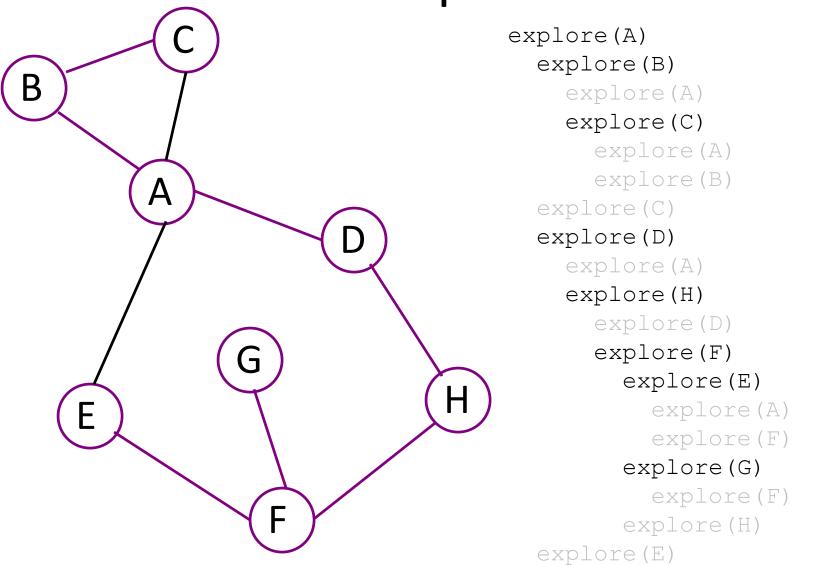








Note: edges used leave Example behind "DFS tree".



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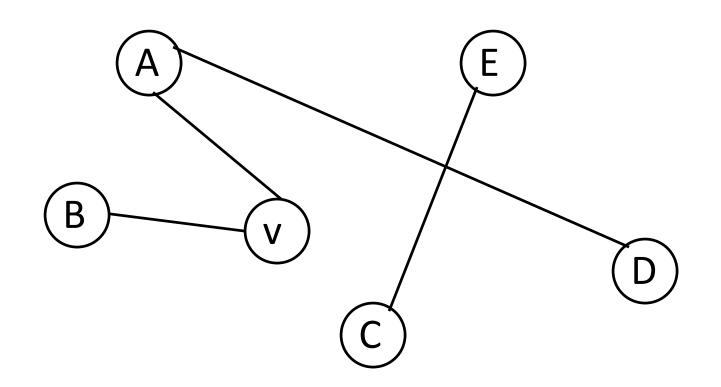
Theorem: If all vertices start unvisited, explore (v) marks as visited exactly the vertices reachable from v.

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 - If there is a chain of vertices $v \rightarrow u_1 \rightarrow u_2 \rightarrow ... \rightarrow w$ eventually visit all of them

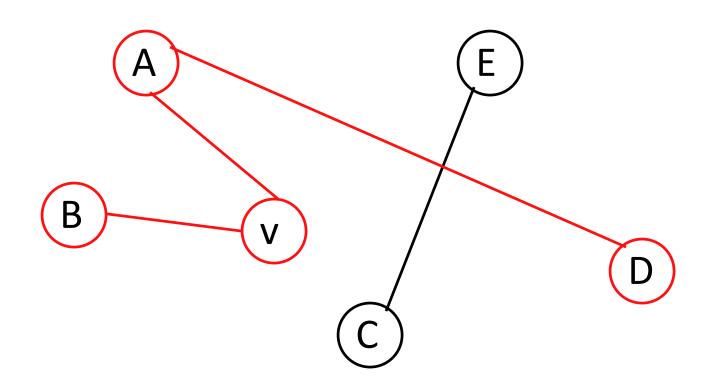
Question: explore

Which vertices does explore (v) mark as visited?



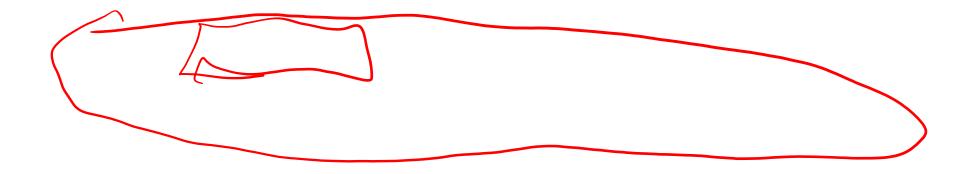
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Depth First Search

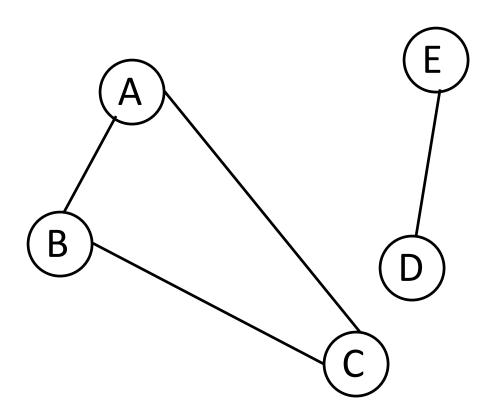
explore only finds the part of the graph reachable from a single vertex. If you want to discover the entire graph, you may need to run it multiple times.

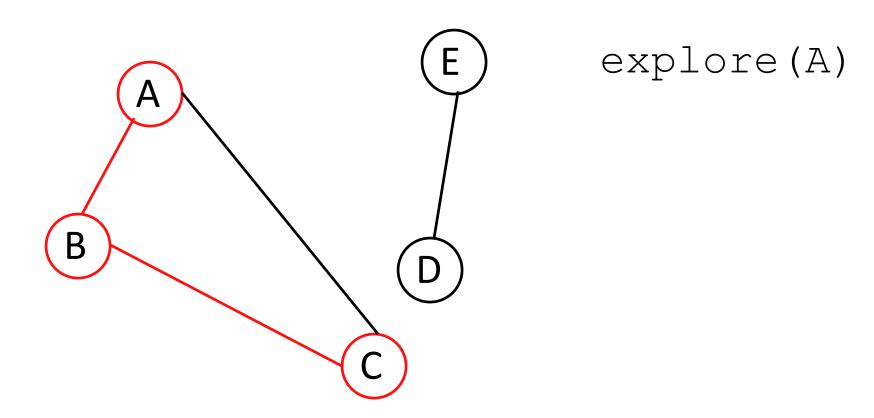


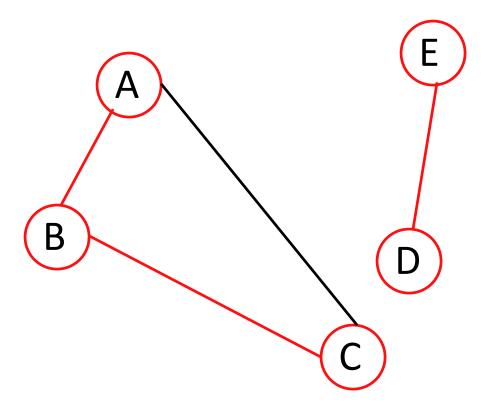
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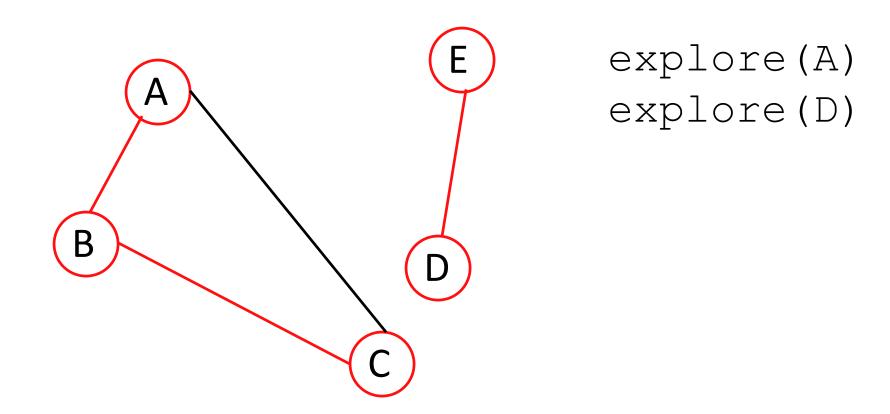
```
DepthFirstSearh(G)
   Mark all v ∈ G as unvisited
   For v ∈ G
   If not v.visited, explore(v)
```







explore(A)
explore(D)



DFS(G) eventually discovers all vertices in G.

```
explore(v)

v.visited ← true

For each edge (v,w)

If not w.visited

explore(w)
```

```
DFS(G)
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```
explore(v)
                       Run once per
                                 O(|V|) total
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                         vertex
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                             vertex
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Final runtime: O(|V|+|E|)

Note on Graph Algorithm Runtimes

Graph algorithm runtimes depend on both |V| and |E|. (Note O(|V|+|E|) is linear time)

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What algorithm is better may depend on relative sizes of these parameters.

Sparse Graphs:

 $|E| \text{ small } (\approx V)$

Examples:

- Internet
- Road maps

Dense Graphs:

|E| large ($\approx V^2$)

Examples:

- Flight maps
- Wireless networks

Question: Graph Runtimes

Suppose that you have two graph algorithms for the same problem

Alg1 has runtime $O(|V|^{3/2})$

Alg2 has runtime O(|E|)

Which of the following is likely true?

- A) Alg1 is faster on most graphs
- B) Alg2 is faster on most graphs
- C) Alg1 is faster on sparse graphs, slower on dense graphs
- D) Alg2 is faster on sparse graphs, slower on dense graphs

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 - Small space for dense graphs.
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 - Small space for dense graphs.
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- Edge list: List of all vertices, list of all edges
 - Hard to determine edges out of single vertex.
- Adjacency list: For each vertex store list of neighbors.
 - Needed for DFS to be efficient
 - We will usually assume this representation

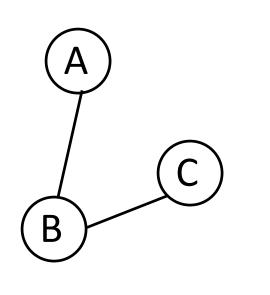
Connected Components

- Want to understand which vertices are reachable from which others in graph G.
- explore(v) finds which vertices are reachable from a given vertex.

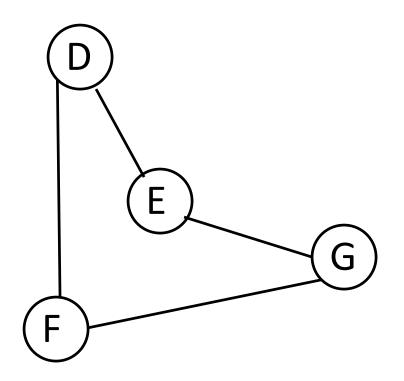
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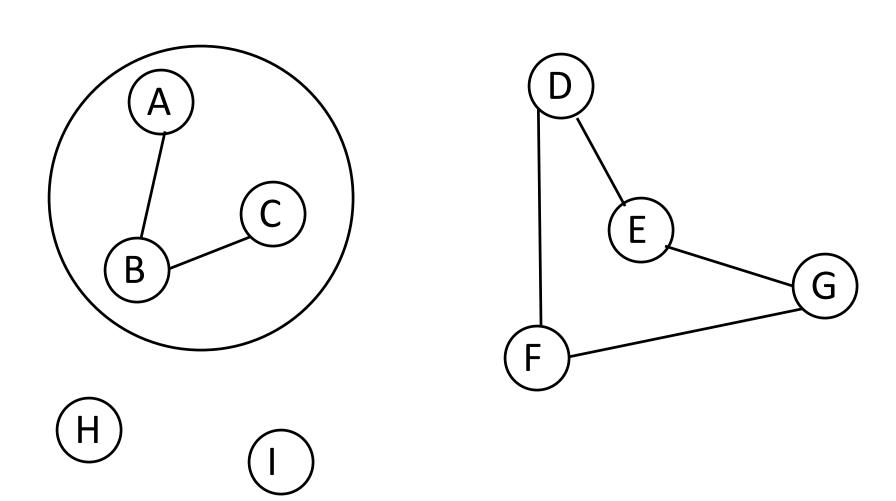
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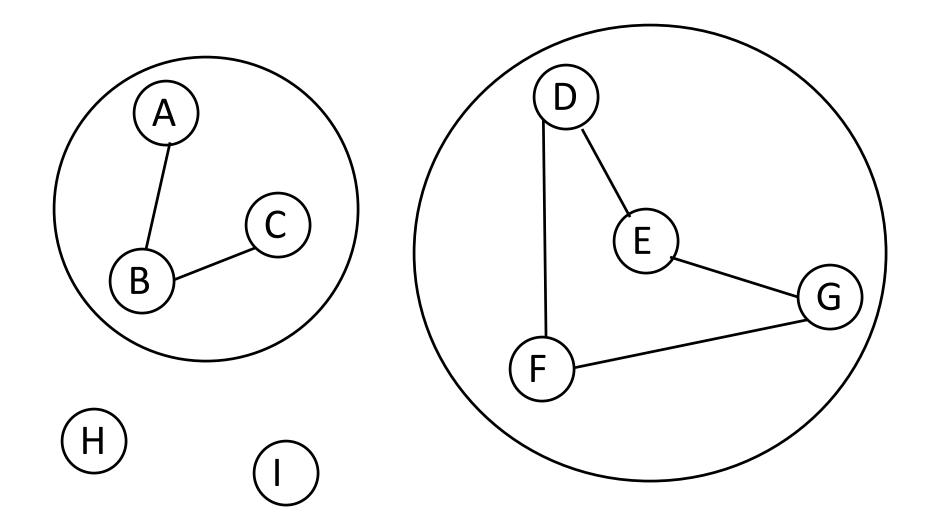
Theorem: The vertices of a graph G can be partitioned into connected components so that v is reachable from w if and only if they are in the same connected component.

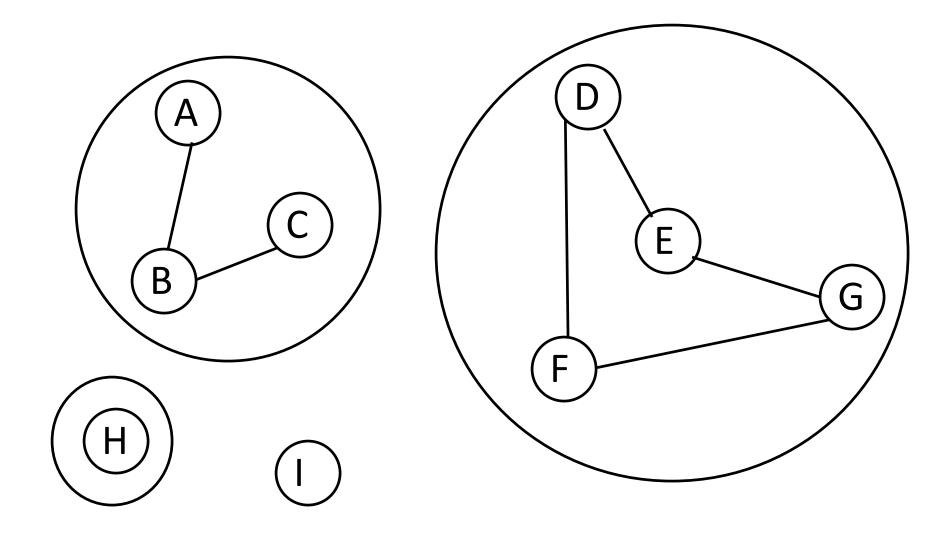


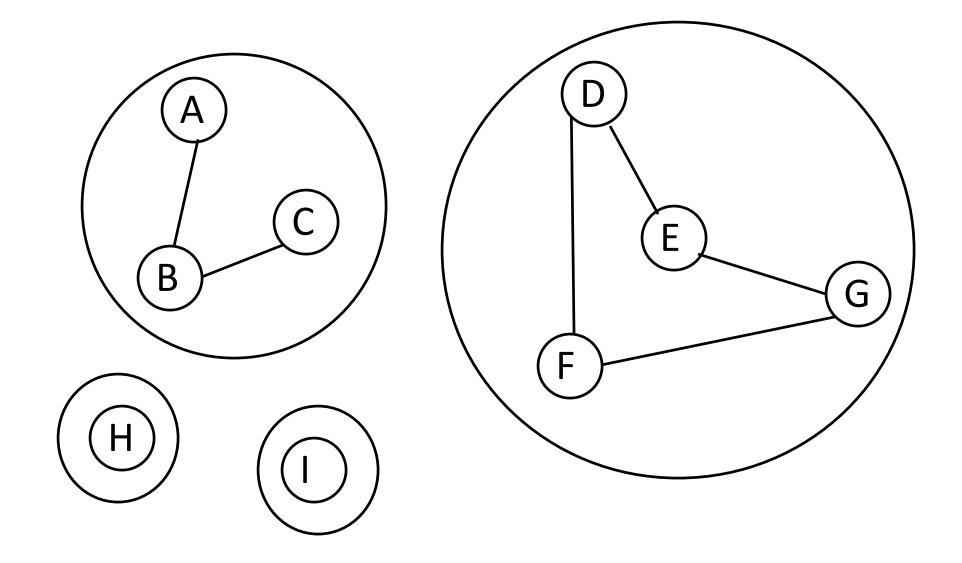








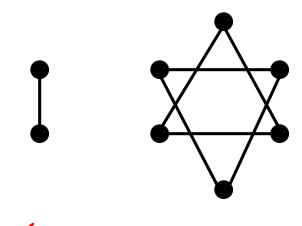




Question: Connected Components

How many connected components does the graph below have?

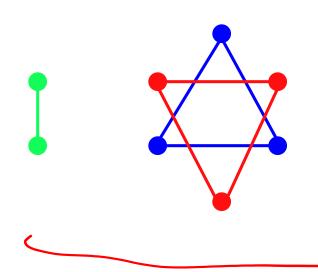
- A)1
- B)2
- C)3
- D)4
- E)5



Question: Connected Components

How many connected components does the graph below have?

- A)1
- B)2
- **C)3**
- D)4
- E)5



Problem: Computing Connected Components

Problem: Given a graph G, compute its connected components.