

Correlated predictor

- Prediction with correlation
- (m,n) predictor
 - m bits of correlation
 - n -bit predictor for branch
 - last m branches ($2m$) each with an n -bit predictor
- • Implementation: Global history with selected address bits (so called “gselect”)
 - m -bit shig register holds outcome of last m branches
 - BHT indexed by $m:low(PC)$
 - BHT can also be indexed just by m (global history prediction)

Correlated Predictor

- **Correlated (two-level) predictor** [Patt 1991]
 - Exploits observation that branch outcomes are correlated
 - Maintains separate prediction per (PC, BHR) pairs
 - **Branch history register (BHR)**: recent branch outcomes
 - Simple working example: assume program has one branch
 - BHT: one 1-bit DIRP entry
 - BHT + **2BHR**: $2^2 = \text{four}$ 1-bit DIRP entries
- Why didn't we do better?
 - BHT not long enough to capture pattern

Time	"Pattern"	State				Prediction	Outcome		Result?
		NN	NT	TN	TT				
1	NN	N	N	N	N	N	T	Wrong	
2	NT	T	N	N	N	N	T	Wrong	
3	TT	T	T	N	N	N	T	Wrong	
4	TT	T	T	N	T	T	N	Wrong	
5	TN	T	T	N	N	N	T	Wrong	
6	NT	T	T	T	N	T	T	Correct	
7	TT	T	T	T	N	N	T	Wrong	
8	TT	T	T	T	T	T	N	Wrong	
9	TN	T	T	T	N	T	T	Correct	
10	NT	T	T	T	N	T	T	Correct	
11	TT	T	T	T	N	N	T	Wrong	
12	TT	T	T	T	T	T	N	Wrong	

Correlated Predictor – 3 Bit Pattern

- Try 3 bits of history
- 2^3 DIRP entries per pattern

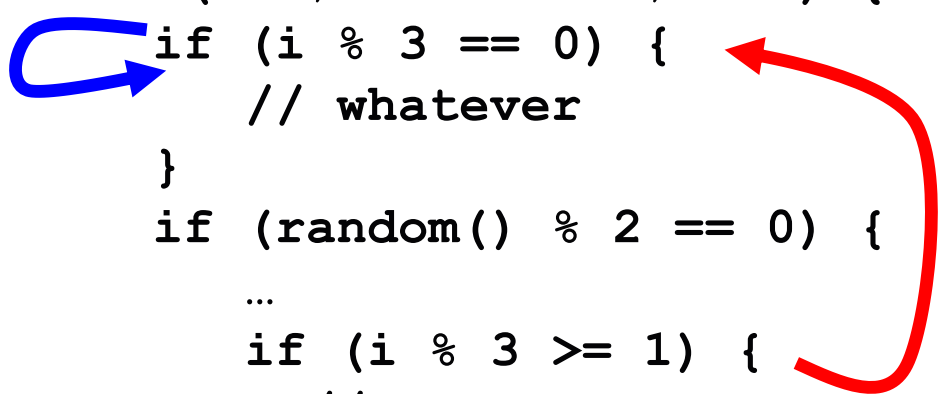
Time	"Pattern"	State								Prediction	Outcome	Result?
		NNN	NNT	NTN	NTT	TNN	TNT	TTN	TTT			
1	NNN	N	N	N	N	N	N	N	N	N	T	Wrong
2	NNT	T	N	N	N	N	N	N	N	N	T	Wrong
3	NTT	T	T	N	N	N	N	N	N	N	T	Wrong
4	TTT	T	T	N	T	N	N	N	N	N	N	Correct
5	TTN	T	T	N	T	N	N	N	N	N	T	Wrong
6	TNT	T	T	N	T	N	N	T	N	N	T	Wrong
7	NTT	T	T	N	T	N	T	T	N	T	T	Correct
8	TTT	T	T	N	T	N	T	T	N	N	N	Correct
9	TTN	T	T	N	T	N	T	T	N	T	T	Correct
10	TNT	T	T	N	T	N	T	T	N	T	T	Correct
11	NTT	T	T	N	T	N	T	T	N	T	T	Correct
12	TTT	T	T	N	T	N	T	T	N	N	N	Correct

+ No mis-predictions after predictor learns all the relevant patterns!

Correlated Predictor Design

- Design choice I: one **global** BHR or one per PC (**local**)?
 - Each one captures different kinds of patterns
 - Global history captures relationship among different branches
 - Local history captures "self" correlation
 - Local history requires another table to store the per-PC history
- Consider:

```
for (i=0; i<1000000; i++) { // Highly biased
    if (i % 3 == 0) { // "Local" correlated
        // whatever
    }
    if (random() % 2 == 0) { // Unpredictable
        ...
        if (i % 3 >= 1) {
            // whatever
        }
    }
}
```

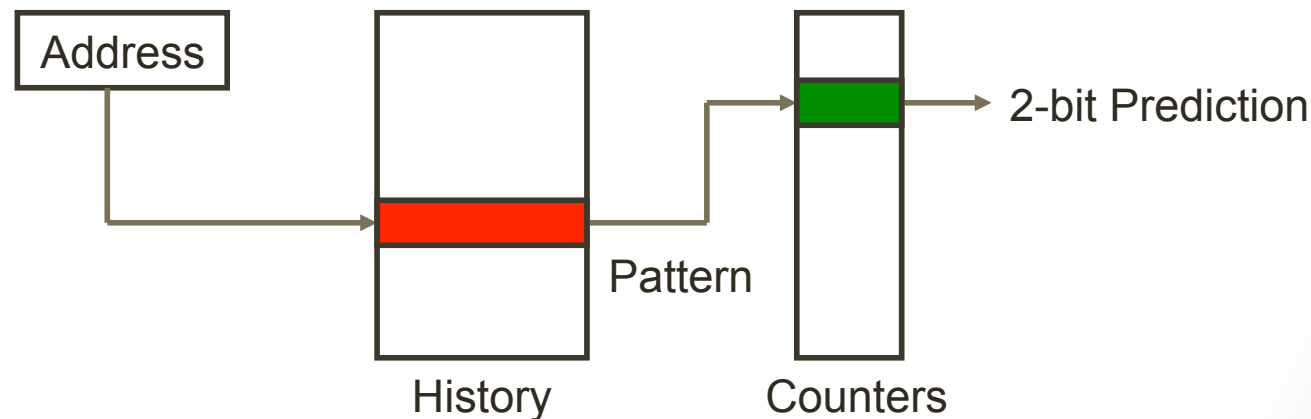


Correlated Predictor Design

- Design choice II: how many history bits (BHR size)?
 - Tricky one
 - + Given unlimited resources, longer BHRs are better, but...
 - BHT utilization decreases
 - Many history patterns are never seen
 - Many branches are history independent (don't care)
 - PC xor BHR allows multiple PCs to dynamically share BHT
 - BHR length $< \log_2(\text{BHT size})$
 - Predictor takes longer to train
 - Typical length: 8–12

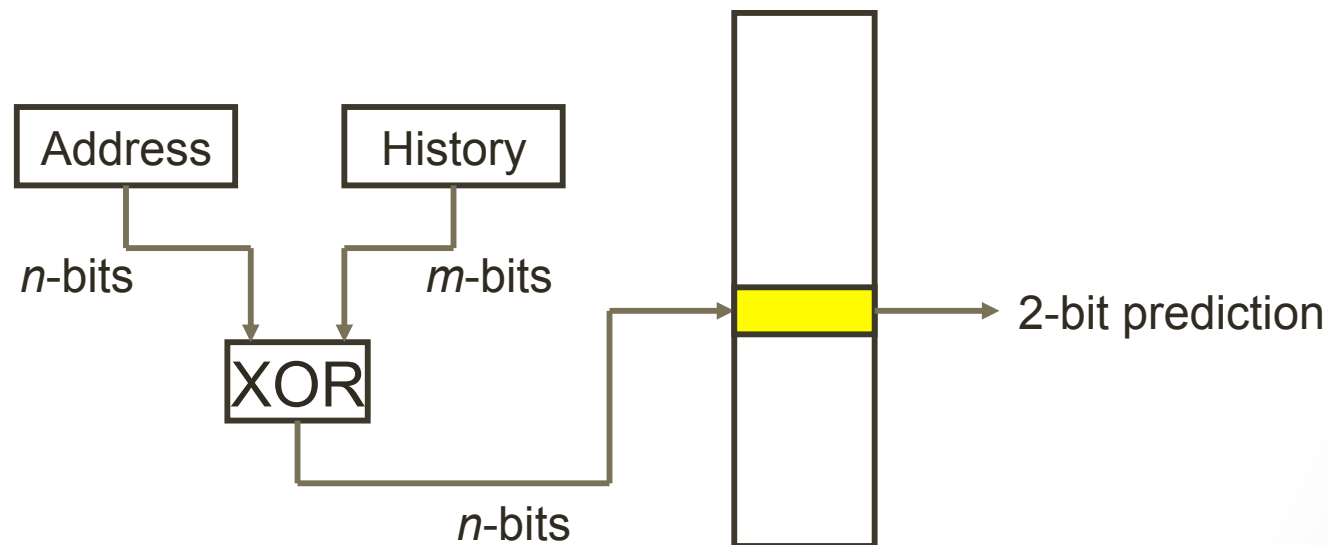
Local branch prediction

- A two-level history table
- Table 1: history of recent branches indexed by the low address bits of branch instruction PC
- Table 2: two-bit branch predictors indexed by the history from table 1

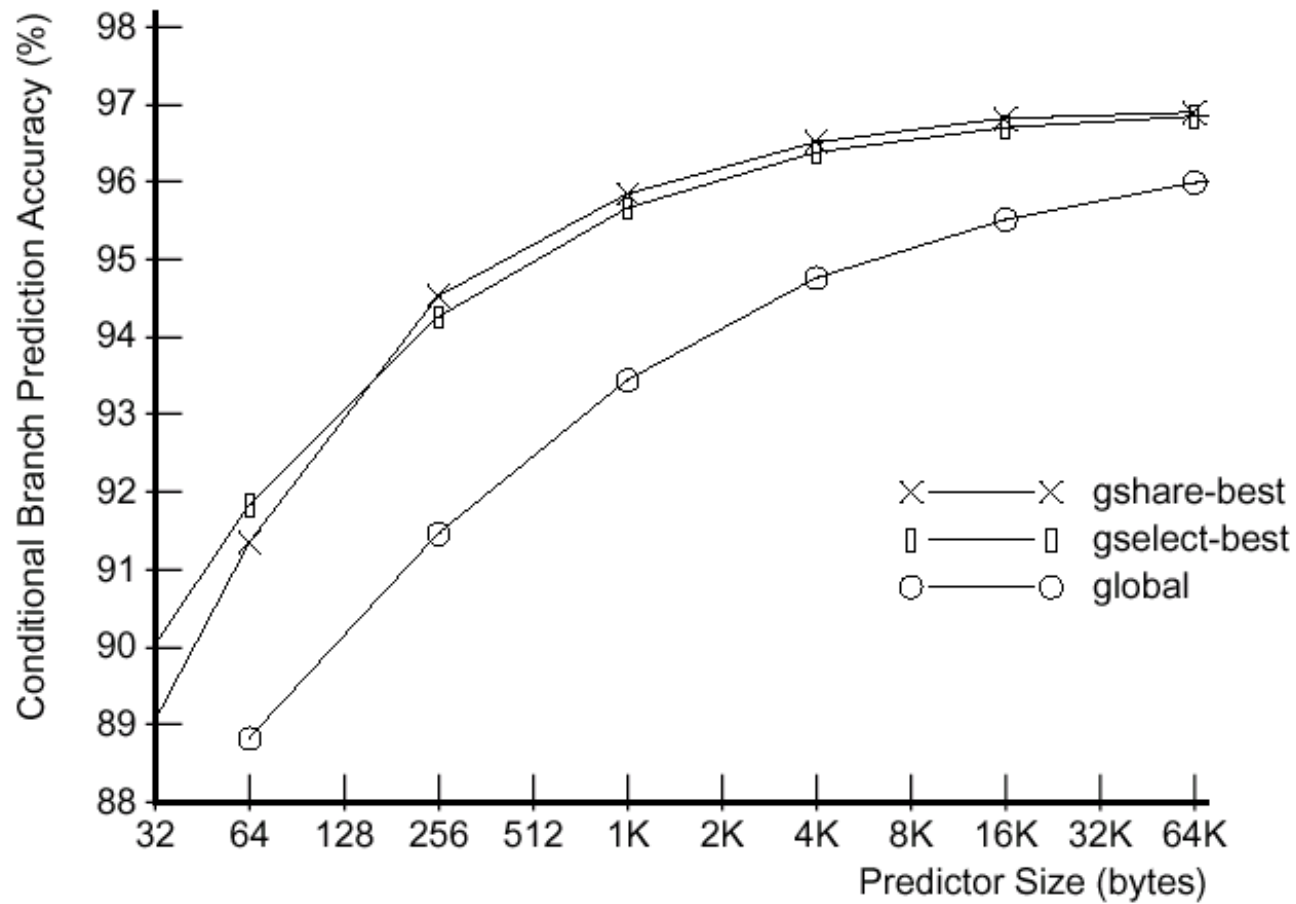


Gshare – global with index sharing

- Similar to gselect predictor, except the branch address and global history are combined by an XOR

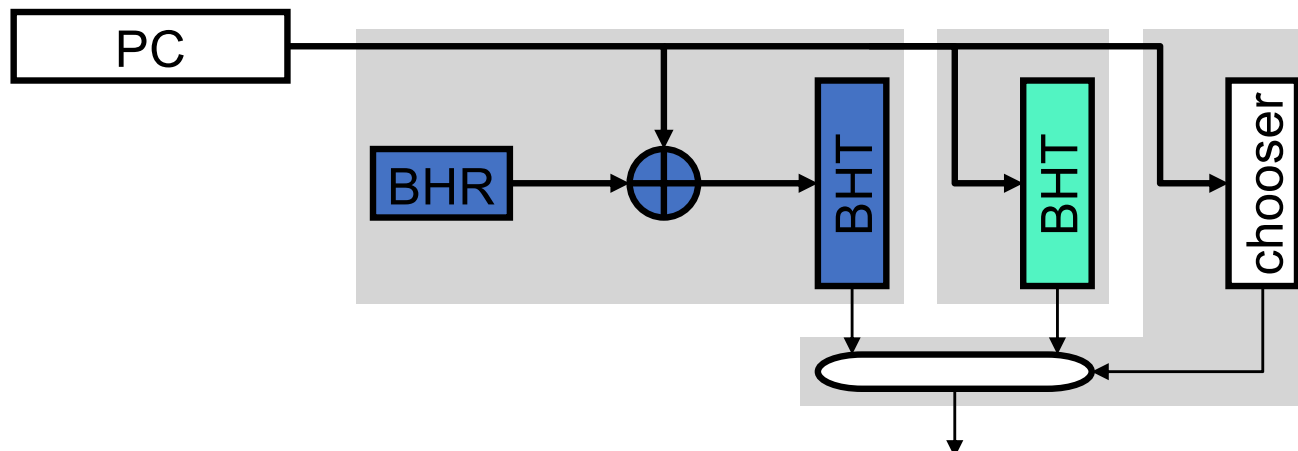


Gshare vs. Gselect

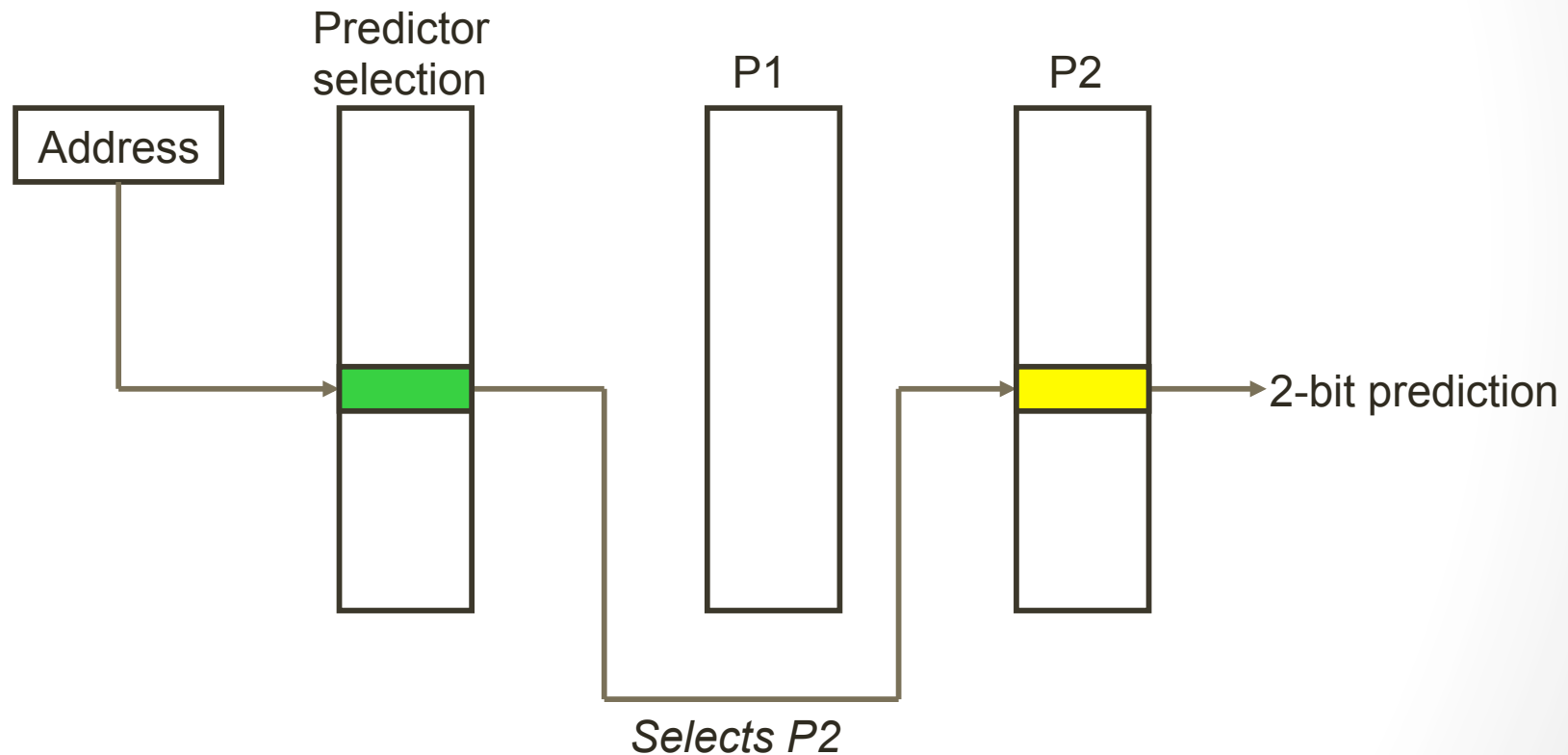


Tournament branch predictor

- **Hybrid (tournament) predictor** [McFarling 1993]
 - Also known as “combining predictors”
 - Attacks correlated predictor BHT capacity problem
 - Idea: combine two predictors
 - **Simple BHT** predicts history independent branches
 - **Correlated predictor** predicts only branches that need history
 - **Chooser** assigns branches to one predictor or the other
 - Branches start in simple BHT, move mis-prediction threshold
- + Correlated predictor can be made **smaller**, handles fewer branches
- + 90–95% accuracy



Tournament predictor implementation



Keeping track of predictor accuracy

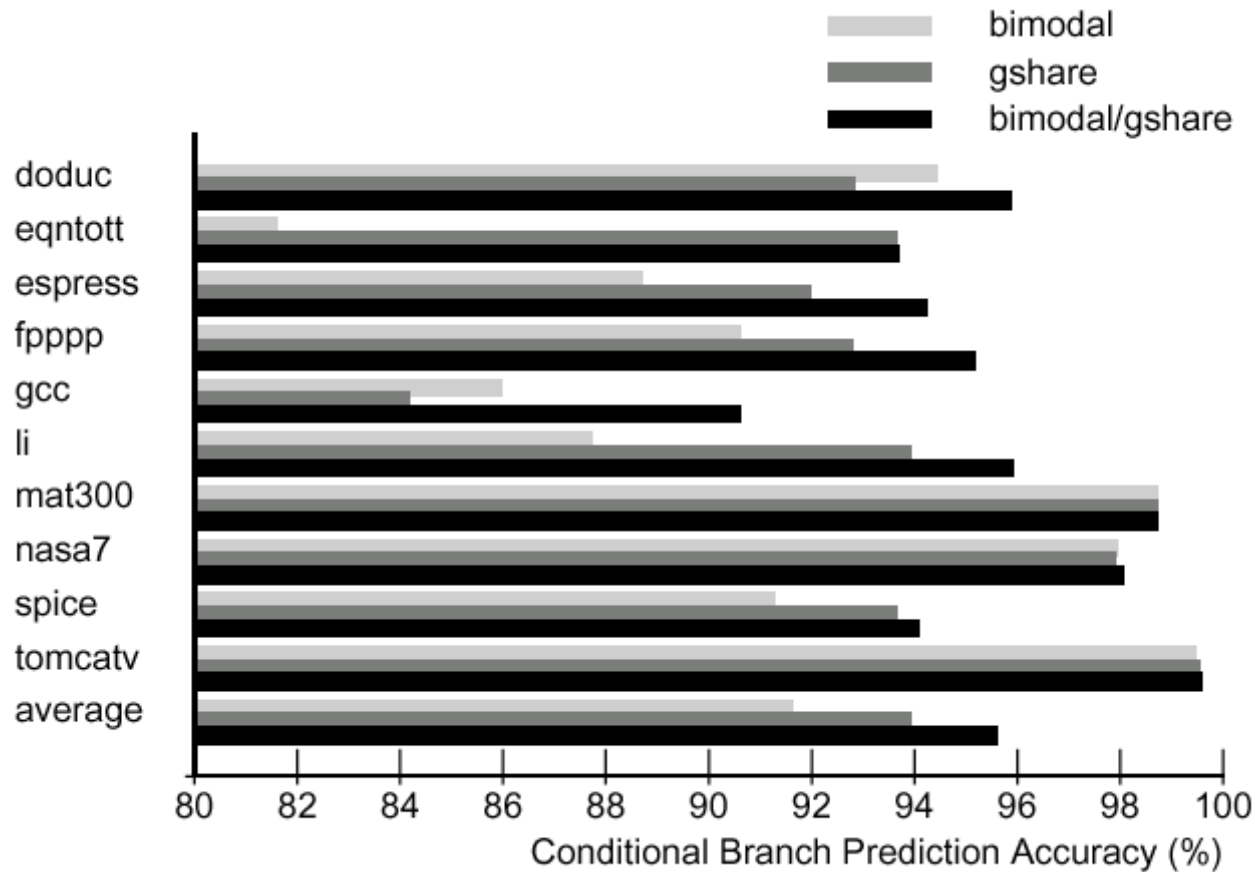
- 2-bit counter incremented/decremented

<u>P1-correct</u>	<u>P2-correct</u>	<u>P1-correct - P2-correct</u>	<u>Action</u>
0	0	$0-0 = 0$	None
0	1	$0-1 = -1$	Decrement
1	0	$1-0 = 1$	Increment
1	1	$1-1 = 0$	None

<u>Counter value</u>	<u>Use predictor</u>
00	P2
01	P2
10	P1
11	P1

} Selects which predictor table to use for the prediction

Tournament predictor performance



Tournament predictor always better than either predictor alone

Research: Perceptron Predictor

- **Perceptron predictor** [Jimenez]
 - Attacks predictor size problem using machine learning approach
 - History table replaced by table of function coefficients F_i (signed)
 - Predict taken if $\sum(BHR_i * F_i) > \text{threshold}$
- + Table size $\#PC * |BHR| * |F|$ (can use long BHR: ~ 60 bits)
 - Equivalent correlated predictor would be $\#PC * 2^{|BHR|}$
- How does it learn? Update F_i when branch is taken
 - $BHR_i == 1 ? F_i++ : F_i--;$
 - “don’t care” F_i bits stay near 0, important F_i bits saturate
- + Hybrid BHT/perceptron accuracy: 95–98%

