

### Measuring Runtime

Count how many times each line executes, then say which  $O()$  statement(s) is(are) true.

```
int maxDifference(int[] arr) {
    int max = 0;
    for (int i=0; i<arr.length; i++) {
        for (int j=0; j<arr.length; j++) {
            if (arr[i] - arr[j] > max) {
                max = arr[i] - arr[j];
            }
        }
    }
    return max;
}
```

Assume  $n = \text{arr.length}$

A  $f(n) = O(2^n)$   
 B  $f(n) = O(n^2)$   
 C  $f(n) = O(n)$   
 D  $f(n) = O(n^3)$   
 E Other/none/more

Big  $O$   $\leq$  Upper bound  
 $f(n) = O(g(n)), f(n) \leq c \cdot g(n)$   
 for all  $n \geq n_0$

Big  $\Omega$   $\geq$  Lower bound  
 $f(n) = \Omega(g(n)), f(n) \geq c \cdot g(n)$   
 for all  $n \geq n_0$

Big  $\Theta$  tight bound  
 $f(n) = \Theta(g(n)), f(n) = c \cdot g(n)$   
 for all  $n \geq n_0$

For each function in the list below, it is related to the function below it by  $O$ , and the reverse is not true. That is,  $n$  is  $O(n^2)$  but  $n^2$  is not  $O(n)$ .

- $f(n) = 1/(n^2)$
- $f(n) = 1/n$
- $f(n) = 1$
- $f(n) = \log(n)$
- $f(n) = \sqrt{n}$
- $f(n) = n$
- $f(n) = n^2$
- $f(n) = n^3$
- $f(n) = n^4$
- ... and so on for constant polynomials ...
- $f(n) = 2^n$
- $f(n) = n!$
- $f(n) = n^n$

Count how many times each line executes, then say which  $O()$  statement(s) is(are) true.

```
int sumTheMiddle(int[] arr) {
    int range = 100;
    int start = arr.length/2 - range/2;
    int sum = 0;
    for (int i=start; i<start+range; i++) {
        sum += arr[i];
    }
    return sum;
}
```

Assume  $n = \text{arr.length}$

A  $f(n) = O(2^n)$   
 B  $f(n) = O(n^2)$   
 C  $f(n) = O(n)$   
 D  $f(n) = O(1)$   
 E None of these

Count how many times each line executes, then say which  $O()$  statement(s) is(are) true.

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    int start = arr.length/2 - range/2;
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Assume  $n = \text{arr.length}$

A  $f(n) = O(2^n)$   
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 C  $f(n) = O(n)$   
 D  $f(n) = O(1)$   
 E None of these

```
void printAllItemsTwice(int arr[], int size) {
    for (int i = 0; i < size; i++) {
        printf("%d\n", arr[i]);
    }
    for (int i = 0; i < size; i++) {
        printf("%d\n", arr[i]);
    }
}
```

What is the tight bound?

$6n + 4 \rightarrow \Theta(n)$   
 $C = 6$   
 $N_0 = 4$

```
void printFirstItemThenFirstHalfThenSayHi100Times(int arr[], int size) {
    printf("First element of array = %d\n", arr[0]);
    for (int i = 0; i < size/2; i++) {
        printf("%d\n", arr[i]);
    }
    for (int i = 0; i < 100; i++) {
        printf("Hi\n");
    }
}
```

What is the tight bound?

$\frac{3}{2}n + 30 \rightarrow \Theta(n)$

```
void printAllNumbersThenAllPairSums(int arr[], int size) {
    for (int i = 0; i < size; i++) {
        printf("%d\n", arr[i]);
    }
    for (int i = 0; i < size; i++) {
        for (int j = 0; j < size; j++) {
            printf("%d\n", arr[i] + arr[j]);
        }
    }
}
```

What is the tight bound?

$3n^2 + 2n + 2 \rightarrow \Theta(n^2)$   
 $C = 3$   
 $N_0 = 2$

### Selection Sort

```
import java.util.Arrays;
public class Sort {
    public static void sortA(int[] arr) {
        for(int i = 0; i < arr.length; i++) {
            System.out.print(Arrays.toString(arr) + " -> ");
            int minIndex = i;
            for(int j = i; j < arr.length; j++) {
                if(arr[minIndex] > arr[j]) { minIndex = j; }
            }
            int temp = arr[i];
            arr[i] = arr[minIndex];
            arr[minIndex] = temp;
            System.out.println(Arrays.toString(arr));
        }
    }
}
```

Selection Sort – what does it print out?

Sort.sortA(new int[] { 53, 83, 15, 45, 49 });

[53, 83, 15, 45, 49] -> [15, 83, 53, 49, 45] → [15, 45, 83, 49, 53] → [15, 45, 49, 83, 53] → [15, 45, 49, 53, 83]

Worst case: reverse sorted array  
 83, 53, 49, 45, 15

Best case: sorted array  
 15, 45, 49, 53, 83

What is the runtime? Consider the shape of the input array.

Worse case:  $\Theta(n^2)$   
 Best case:  $\Theta(n)$

is Sorted ( )  
 $\Theta(n)$

### Insertion Sort

```
import java.util.Arrays;
public class Sort {
    public static void sortB(int[] arr) {
        for(int i = 0; i < arr.length; i++) {
            System.out.print(Arrays.toString(arr) + " -> ");
            for(int j = i; j > 0; j--) {
                if(arr[j] < arr[j-1]) {
                    int temp = arr[j-1];
                    arr[j-1] = arr[j];
                    arr[j] = temp;
                }
            }
            System.out.println(Arrays.toString(arr));
        }
    }
}
```

Insertion Sort – what does it print out?

Sort.sortB(new int[] { 53, 83, 15, 45, 49 });

[53, 83, 15, 45, 49] -> [53, 83, 15, 45, 49] → [53, 83, 15, 45, 49] → [15, 53, 83, 45, 49] → [15, 45, 53, 83, 49] → [15, 45, 49, 53, 83]

What is the runtime? Consider the shape of the input array.

Worse case:  $\Theta(n^2)$   
 Best case:  $\Theta(n)$