

Quick Sort

Sorting Quickly

```
public class SortQuickly {  
    O(1) public static void swap(String[] array, int i1, int i2) {  
        String temp = array[i1];  
        array[i1] = array[i2];  
        array[i2] = temp;  
    }  
  
    public static int partition(String[] array, int low, int high) {  
        → int pivotStartIndex = high - 1;  
        String pivot = array[pivotStartIndex];  
        int smallerBefore = low, largerAfter = high - 2;  
  
        while (smallerBefore <= largerAfter) {  
            if (array[smallerBefore].compareTo(pivot) < 0) {  
                smallerBefore += 1;  
            }  
            else {  
                → swap(array, smallerBefore, largerAfter);  
                → largerAfter -= 1;  
            }  
        }  
  
        swap(array, smallerBefore, pivotStartIndex);  
        return smallerBefore;  
    }  
  
    public static void qsort(String[] array, int low, int high) {  
        if (high - low <= 1) { return; }  
        → int splitAt = partition(array, low, high);  
        qsort(array, low, splitAt);  
        qsort(array, splitAt + 1, high);  
    }  
  
    public static void sortD(String[] array) {  
        qsort(array, 0, array.length);  
    }  
  
    public static void main(String[] args) {  
        String[] str = {"f", "b", "a", "e", "d", "c"};  
        int[] result = SortQuickly.sortD(str);  
        System.out.println(Arrays.deepToString(result));  
    }  
}
```

Quicksort: Another recursive algorithm
<https://www.youtube.com/watch?v=yWBy6L5g28>

14	4	9	12	15	8	19	2
----	---	---	----	----	---	----	---

Select a **pivot** element:

14	4	9	12	15	8	19	2
----	---	---	----	----	---	----	---

"Partition" the elements in the array (smaller or equal to pivot, larger or equal to pivot)

2	4	9	8	15	12	19	14
---	---	---	---	----	----	----	----

Magically sort the smaller elements and the larger elements (Quicksort)

2	4	8	9	12	15	19	21
---	---	---	---	----	----	----	----

Quick Sort: Using a "good" pivot

How many levels will there be if you choose a pivot that divides the list in half?

A. 1
B. $\log(n)$
C. n
D. $n \cdot \log(n)$
E. n^2

If the time to partition on each level takes N comparisons, how long does Quicksort take with a good partition?

A. $O(1)$
B. $O(\log(n))$
C. $O(n)$
D. $O(n \cdot \log(n))$
E. $O(n^2)$

Which of these choices would be the **worst** choice for the pivot?

A. The minimum element in the list
B. The last element in the list
C. The first element in the list
D. A random element in the list

Worst choice (log n)

best choice median value

Quick sort with a bad pivot

the partitions

n elements
n - 1 elements
n - 2 elements
...

Space complexity: $O(n)$ for the activation records

If the pivot always produces one empty partition and one with $n - 1$ elements, there will be n levels, each of which requires $O(n)$ comparisons: $O(n^2)$ time complexity

Which of these choices is a better choice for the pivot?

A. The first element in the list
B. A random element in the list
C. They are about the same

There are many ways to partition!

Quick sort - Middle Pivot

1. We always pick the middle location as pivot
2. The data we sort is (2, 3, 1, 5, 4, 6, 7)

After the first split, what is the order of elements in the list that was \leq pivot?

A. 1 2 3 4
B. 2 3 4
C. 4 3 2 1
D. 3 4 1 2
E. None of the above

low = 0, high = 6
*lowIndex = low (0)
highIndex = high - 1 (6)
pivotIndex = (high + low) / 2 = (8 - 0) / 2 = 4*

```
import java.util.Arrays;
public class Sort {
    static void selectionSort(int[] arr) {
        for(int i = 0; i < arr.length; i++) {
            int minIndex = i;
            for(int j = i + 1; j < arr.length; j++) {
                if(arr[minIndex] > arr[j]) minIndex = j;
            }
            int temp = arr[i];
            arr[i] = arr[minIndex];
            arr[minIndex] = temp;
        }
    }

    static void insertionSort(int[] arr) {
        for(int i = 0; i < arr.length; i++) {
            for(int j = i; j > 0; j--) {
                if(arr[j] < arr[j-1]) {
                    int temp = arr[j];
                    arr[j] = arr[j-1];
                    arr[j-1] = temp;
                }
            }
        }
    }

    static int[] combine(int[] p1, int[] p2) {
        int len = arr.length;
        if(len <= 1) return arr;
        else {
            int[] p1 = Arrays.copyOfRange(arr, 0, len / 2);
            int[] p2 = Arrays.copyOfRange(arr, len / 2, len);
            int[] sortedPart1 = mergeSort(p1);
            int[] sortedPart2 = mergeSort(p2);
            int[] sorted = combine(sortedPart1, sortedPart2);
            return sorted;
        }
    }

    static int partition(String[] array, int l, int h) {
        if(l >= h) return l;
        int splitAt = partition(array, l, h);
        quickSort(array, l, splitAt);
        quickSort(array, splitAt + 1, h);
    }

    public static void sort(String[] array) {
        quickSort(array, 0, array.length);
    }
}
```

Sorted Array $\rightarrow O(n)$
Reverse Sort array $\rightarrow O(n^2)$
1 2 3 4 5 6
6 5 4 3 2 1
4 5 6
3 4 5 6
2 3 4 5 6
1 2 3 4 5 6

low = 0, high = 6
*lowIndex = low (0)
highIndex = high - 1 (6)
pivotIndex = (high + low) / 2 = (8 - 0) / 2 = 4*

loop
*if value[lowIndex] < pivot value
lowIndex++*
*else swap (lowIndex, highIndex)
highIndex--*
*if value[highIndex] > pivot value
highIndex--*
*else swap (lowIndex, highIndex)
lowIndex++*
*if (low >= high)
done = true*
Swap chosen index?

low = 0, high = 6
*lowIndex = low (0)
highIndex = high - 1 (6)
pivotIndex = (high + low) / 2 = (8 - 0) / 2 = 4*

12 4 9 3 15 8 19 2
2 4 3 15
15
15
15

	Insertion	Selection	Merge	Quick
Best case time	Sorted Array $\Theta(n)$	$\Theta(n^2)$	$\Theta(N + \log(n))$	Median value $\Theta(N + \log(n))$
Worst case time	Reverse Sorted Array $\Theta(n^2)$	$\Theta(n^2)$	$\Theta(N + \log(n))$	$\Theta(n^2)$ Average case: $\Theta(N + \log(n))$
Key operations	swap(a, j, j-1) (until in the right place)	swap(a, i, indexOfMin) (after finding minimum value)	$l = \text{copy}(a, 0, \text{len}/2)$ $r = \text{copy}(a, \text{len}/2, \text{len})$ $rs = \text{sort}(r)$ merge(l, rs)	$p = \text{partition}(a, l, h)$ $\text{sort}(a, l, p)$ $\text{sort}(a, p + 1, h)$

Last note about sorting

Not only do we care about runtime, we also care about

- Space: do we need extra storage?
- Stable: if we have duplicates, do we maintain the same ordering?

Algorithm	Space	Stable
Bubble sort	$O(1)$	Yes
Selection sort	$O(1)$	No
Insertion sort	$O(1)$	Yes
Heap sort	$O(1)$	No
Merge sort	$O(n)$	Yes
Quick sort	$O(\log(n))$	No

key value {1, "6 reg"}, {1, "Dent"}
[{1, "6 reg"}, {1, "Dent"}]
[{1, "Dent"}, {1, "6 reg"}]
Unstable ordering

*Just use the Arrays.sort
to QuickSort primitives
to MergeSort objects*