

CS 445: Assignment 4

symatnics.cpp

semantics.h symbolTable.cpp

A simple bC file in assignment 3

```
#DRBC This program is as simple as I can get.
#DRBC This should compile without errors
main ()
{
  int x;
  x = 5;
}
```

Func: main returns type void [line: 3]

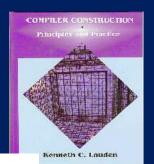
. Child: 1 Compound [line: 4]

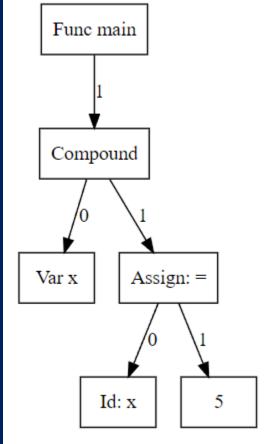
. Child: 0 Var: x of type int [line: 5]

. Child: 1 Assign: = [line: 6]

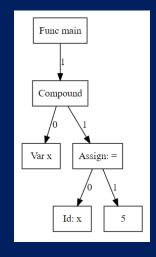
. Child: 0 Id: x [line: 6]

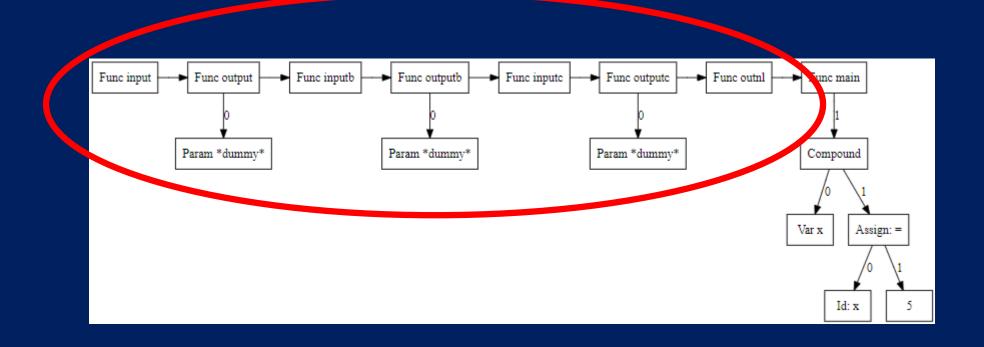
. Child: 1 Const 5 [line: 6]





A simple bC file in assignment 4: 10 new nodes for IOLib





A simple bC file in assignment 4

```
Func: input returns type int [line: -1]
Sibling: 1 Func: output returns type void [line: -1]
. Child: 0 Parm: *dummy* of type int [line: -1]
Sibling: 2 Func: inputb returns type bool [line: -1]
Sibling: 3 Func: outputb returns type void [line: -1]
. Child: 0 Parm: *dummy* of type bool [line: -1]
Sibling: 4 Func: inputc returns type char [line: -1]
Sibling: 5 Func: outputc returns type void [line: -1]
. Child: 0 Parm: *dummy* of type char [line: -1]
Sibling: 6 Func: outnl returns type void [line: -1]
Sibling: 7 Func: main returns type void [line: 3]
  Child: 1 Compound [line: 4]
    Child: 0 Var: x of type int [line: 5]
    Child: 1 Assign: = of type int [line: 6]
. . . Child: 0 ld: x of type int [line: 6]
. . . Child: 1 Const 5 of type int [line: 6]
```

```
Func: main returns type void [line: 3]

. Child: 1 Compound [line: 4]

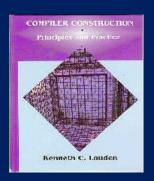
. Child: 0 Var: x of type int [line: 5]

. Child: 1 Assign: = [line: 6]

. Child: 0 Id: x [line: 6]

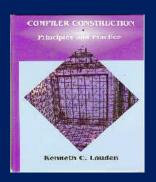
. Child: 1 Const 5 [line: 6]
```

Outline



- Changes to main
- 10 new IOLib nodes
- Types for everything
 - Symbol table to find types in scope
 - Dealing with type errors

main



```
SymbolTable *symtab;
symtab = new SymbolTable();
symtab->debug(false);
int globalOffset;
syntaxTree = semanticAnalysis(syntaxTree,
                    symtab,
                    globalOffset);
```

Set the default size to 1 in treeUtils and then update to something else if needed later

• Create this function. Some node kinds will call it:

```
void showAllocation(FILE *out, TreeNode *theNode){
  fprintf(out, " [mem: %s loc: %d size: %d]", varKindToStr(theNode->yarKind), theNode->offset, theNode->size);
}
```

• It will need this function:

```
char *varKindToStr(int kind){
  switch(kind){
    case None:
      return (char *)"None";
    case Local:
      return (char *)"Local";
    case Global:
      return (char *)"Global";
    case Parameter:
      return (char *)"Parameter";
    case LocalStatic:
      return (char *)"LocalStatic";
    default:
      return (char *)"unknownVarKind";
```

semanticAnalysis will fill in theses values

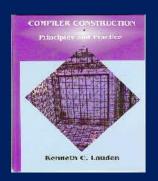
In treeUtils.cpp

semanticAnalysis(...)

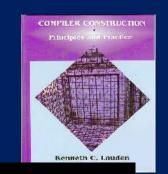
```
syntree = loadIOLib(syntree);
```

```
treeTraverse(syntree, symtab);
```

•••



Printing the tree printTree(stdout, syntaxTree);



```
In Assignment 3:
```

Func: main returns type void [line: 3]

. Child: 1 Compound [line: 4]

. . Child: 0 Var: x of type int [line: 5]

. . Child: 1 Assign: = [line: 6]

. . . Child: 0 ld: x [line: 6]

. . . Child: 1 Const 5 [line: 6]

Number of warnings: 0

Number of errors: 0

```
Func: input returns type int [line: -1]
```

Sibling: 1 Func: output returns type void [line: -1]

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Sibling: 5 Func: outputc returns type void [line: -1]

. Child: 0 Parm: *dummy* of type char [line: -1]

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Sibling: 7 Func: main returns type void [line: 3]

. Child: 1 Compound [line: 4]

. . Child: 0 Var: x of type int [line: 5]

. . Child: 1 Assign: = of type int [line: 6]

. . . Child: 0 Id: x of type int [line: 6]

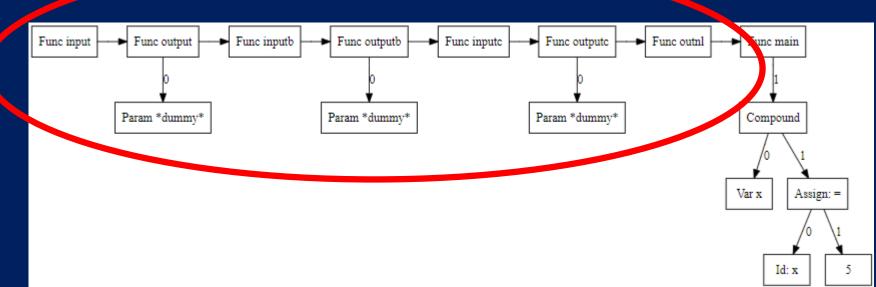
. . . Child: 1 Const 5 of type int [line: 6]

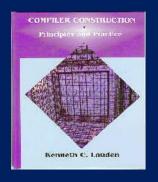
Number of warnings: 0

Number of errors: 0

Outline

- Changes to main
- 10 new IOLib nodes
- Types for everything
 - Symbol table to find types in scope
 - Dealing with type errors
 - Suggested Workflow





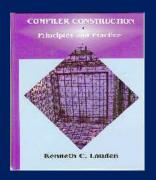
TreeNode *loadIOLib(TreeNode *syntree)

```
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Principles and Practice

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```

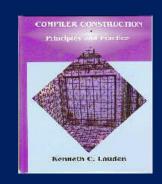
```
TreeNode *input, *output, *param_output;
TreeNode *inputb, *outputb, *param_outputb;
TreeNode *inputc, *outputc, *param outputc;
TreeNode *outnl;
/////// Stuff from next slides
// link them and prefix the tree we are interested in traversing.
// This will put the symbols in the symbol table.
input->sibling = output;
output->sibling = inputb;
inputb->sibling = outputb;
outputb->sibling = inputc;
inputc->sibling = outputc;
outputc->sibling = outnl;
outnl->sibling = syntree; // add in the tree we were given
return input;
```

TreeNode *loadIOLib(TreeNode *syntree)



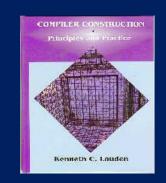
```
input = newDeclNode(FuncK, Integer);
        input->lineno = -1; // all are -1
        input->attr.name = strdup("input"); //We named the variables well
        input->type = Integer;
    inputb = newDeclNode(FuncK, Boolean);
    inputc = newDeclNode(FuncK, Boolean);
3.
    param_output = newDeclNode(ParamK, Void);
4.
    output = newDeclNode(FuncK, Void);
5.
6.
    param outputb = newDeclNode(ParamK, Void);
    outputb = newDeclNode(FuncK, Void);
7.
    param_outputc = newDeclNode(ParamK, Void);
8.
    outputc = newDeclNode(FuncK, Void);
9.
    outnl = newDeclNode(FuncK, Void);
```

TreeNode *loadIOLib(TreeNode *syntree) Exceptions



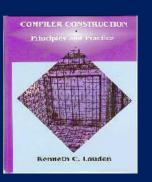
```
3. inputc = newDeclNode(FuncK, Boolean);
      inputc->type = Char;
4. param output = newDeclNode(ParamK, Void);
      param_output->attr.name = strdup("*dummy*");
      param output->type = Integer;
param outputb = newDeclNode(ParamK, Void);
      param outputb->attr.name = strdup("*dummy*");
      param outputb->type = Boolean;
8. param outputc = newDeclNode(ParamK, Void);
      param_outputc->attr.name = strdup("*dummy*");
      param outputc->type = Char;
```

TreeNode *loadIOLib(TreeNode *syntree) Linking



```
output = newDeclNode(FuncK, Void);
     output->child[0] = param output;
   param outputb = newDeclNode(ParamK, Void);
   outputb = newDeclNode(FuncK, Void);
     outputb->child[0] = param outputb;
   param outputc = newDeclNode(ParamK, Void);
   outputc = newDeclNode(FuncK, Void);
     outputc->child[0] = param outputc;
10. outnl = newDeclNode(FuncK, Void);
      outnl->child[0] = NULL;
```

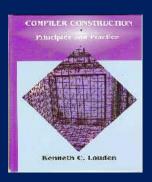
Outline



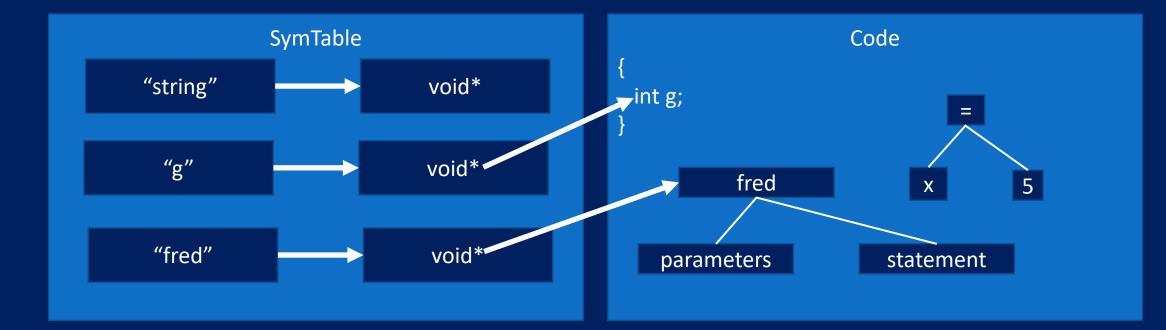
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```

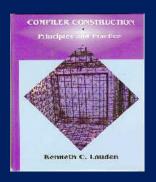
SymbolTable



- A mapping from a string to void*
 - The void* points to the declaration of the variable.
 - Anytime you have an identifier (like the name of a procedure or a variable) something that you need to tag with what its type is, then
 - we'll put that name into the as the string part and
 - the void pointer will point into the code for the declaration



We need to manage scope



- The symbol table is a stack of symbol tables.
 - when you look up the symbol table code, you'll find that in that class there are two types.
 - One is the symbol table type, which is a stack of scopes,
 - and these are listed as scopes.
- The global scope is always there, and the symbol table will yell at you if you try to get rid of it.
 - But you can push and pop scopes and those are the two next procedures you need, enter and leave.
- When I look things up, I start at the top of the stack and return the most recent occurrence.

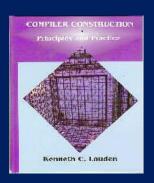
Global Scope

```
"g_xVar"
"g_yVar"
```

```
"local_x"
"local_y"
```

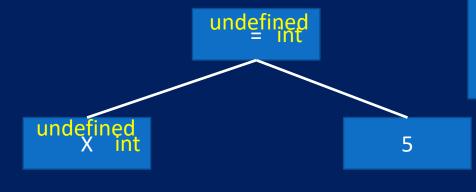
```
"compound_x"
"local_x"
"temp"
```

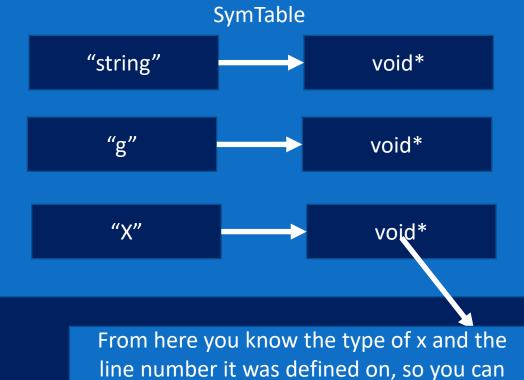
Symbol Table methods



- 4 routines
 - insert into symbol table (Any time you have a declaration)
 - Saves info like the type
 - Gives an error on multiple declarations of the same variable in the same scope
 - lookup (any time you get an id, like x=5, we need to look up x to see if we can put a 5 in it)
 - In x=5 we look up the type of the left and right side to see if they match. (We have no coercion)
 - If they don't match, we have an error for that.
 - X=5 the type of an assignment is the type of the left hand side of the assignment in bC
 - enter scope
 - leave scope

Assignment

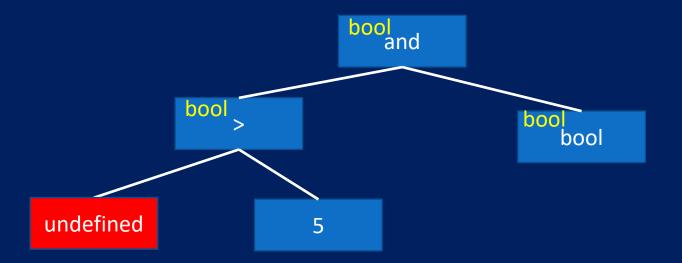




give an informative error message

- X=5 the type of an assignment is the type of the left-hand side of the assignment in bC
 - Look up X in the symbol table.
 - If X is an int, then = is int.
 - If X is undefined, then = is undefined.
- undefined = 5 sends undefined up the tree.
- int x = undefined sends int up the tree





- Because we do not want the error to cascade up the tree
 - Even if both values in the > are undefined the result is a bool

Useful global variables

```
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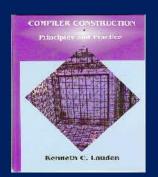
```
// memory offsets are GLOBAL
static int goffset;  // top of global space
static int foffset;  // top of local space
```

What YOU will write for the compound statement (roughly)

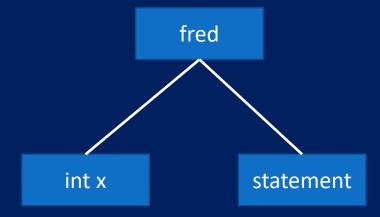
//because you can throw away the scope.

```
symtab->enter((char *)"compoundStmt");
  treeTraverse(current->child[0], symtab); // process declarations
  // More stuff
  current->size = foffset;
  treeTraverse(current->child[1], symtab);
  // More stuff
symtab->leave(); //end of processing of the compound statement
```

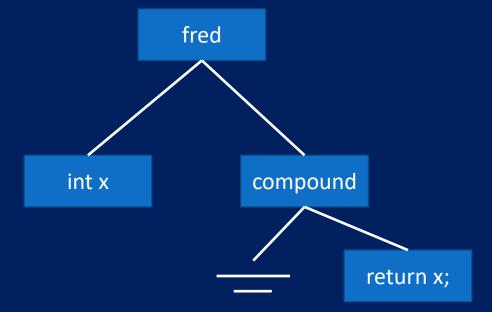
Remember: compound statements occurring right after a function are treated differently.



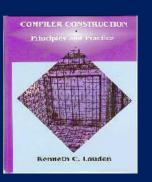
fred(int x) return x;



fred(int x) {return x;}



Outline

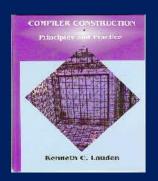


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```

. . . Child: 1 Const 5 of type int [line: 6]

Redefinition of variables



```
int x;
int x; // x is already defined
y = 73; //y is undefined
myFunction(int x) {
 int x; // Even in C this is not OK, It is a weird special case.
   int x; // This is OK
```

```
char dog(char x)
  char x; // SPECIAL CASE: params are in function compound statement
       // Symbol %s is already defined at line %d
  int b;
  return b; // Expecting return type of %s but got type %s
int dog(int x) // Symbol %s is already defined at line %d
    int x;
         // missing return warning
char cat(char x)
  int b;
  return; // Expecting a return type of %s but got none
```

basicAll

Kenneth C. Lauden

Outline

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Principles and Fractice

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- Changes to main
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Order of traversal

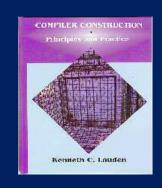
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- child[0]
- Current node
- child[1]
- child[2]
- sibling

New scope if:

- DeclK
 - FuncK
 - Also set foffset = -2;
- StmtK
 - CompoundK
 - ForK
 - IfK
 - WhileK



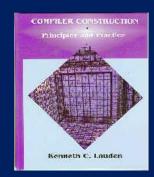
Scope

```
if(isCompound) {
    char *id = strdup("{");
    symtab->enter("NewScope from " + (std::string)id);
}
```

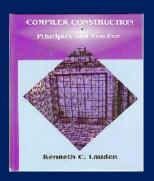
- child[0]
- Current node
- child[1]
- child[2]

```
if(isCompound) {
    symtab->leave();
}
```

• sibling



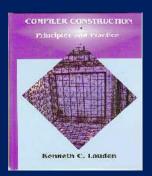
Types



- ExpK
 - AssignK & OpK
 - Boolean
 - AND
 - OR
 - EQ
 - NEQ
 - LEQ
 - <
 - GEQ
 - >
 - NOT

- Integer
 - Any value < LASTOP
- Whatever child[0]'s type is (and if it is an array)
 - =
 - [

Types

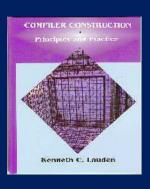


- ExpK
 - IdK

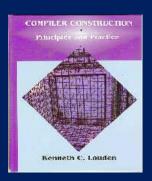
```
    Look up the type in the symtab.
        if ((tmp = (TreeNode *)(symtab->lookup(current->attr.name)))) {
                 current->type = tmp->type;
                  current->isStatic = tmp->isStatic;
                  current->isArray = tmp->isArray;
                  current->size = tmp->size;
                 current->varKind = tmp->varKind;
                  current->offset = tmp->offset;
                  }
```

- CallK
 - Look up like for IdK and set type and size

TMBC Offsets (See the tmDiscription)



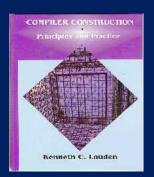
- goffset the global offset is the relative offset of the next available space in the global space.
- foffset the frame offset is the relative offset of the next available space in the frame being built
- toffset the temp offset is the offset from the frame offset next available temp variable.



- DeclK
 - VarK & ParamK

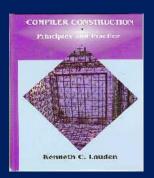
```
if (symtab->depth()==1) // This is a global variable since it is not in a function
     // Set the varKind to Global for VarK (Parameter for ParamK)
     current->offset = goffset;
     goffset -= current->size;
Otherwise, if current->isStatic // This is a static variable
     // Set varKind to LocalStatic for VarK (Parameter for ParamK)
     current->offset = goffset;
     goffset -= current->size;
     //symtab->insertGlobal with a unique name. (Keep a static int variable and append it to the end of the
     name)
Otherwise, treat it as normal
     // Set varKind to Local for VarK (Parameter for ParamK)
     current->offset = foffset;
                                                                         For VarK only
     foffset -= current->size;
```

if (current->isArray) current->offset--; // pt to array after size



- DeclK
 - FuncK

```
foffset = -2;
Traverse child[0]'s tree (where the parameters are)
current->size = foffset;
```



- StmtK
 - CompoundK
 - If newScope
 - Remember the current Offset
 - Traverse child[0]'s tree
 - current->size = foffset;
 - Traverse child[1]'s tree
 - Restore the current Offset
 - Otherwise
 - Traverse child[0]'s tree
 - current->size = foffset;
 - Traverse child[1]'s tree

StmtK

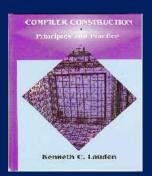
ForK

- symtab->enter((char *)"ForStmt");
- Remember the current Offset
- Traverse child[0]'s tree
- foffset-=2;
- current->size = foffset;
- Traverse child[1]'s tree
- Traverse child[2]'s tree
- Restore the current Offset
- symtab->leave();

What number will be output here? 10 or 37?

// Make space for the for loop var

```
main()
 int i;
 i = 37;
 int sum;
 sum = 0;
 for i=1 to 10 do sum += i;
 output(i);
                   You will get the
 output(sum);
                        error
                   Symbol 'i' is not
                      declared.
                   If you try to use
                  this i outside the
                       for loop
```



ExpK

ConstantK

• IdK

```
tmp = (TreeNode *)(symtab->lookup(current->attr.name) // Look up in the symbol table
current->offset = tmp->offset;
```