

Foundations of Computer Science – Exercise 6

$$1. x_{n+1} = x_n - \frac{x_n^2 - N}{(x_n^2 - N)'} = \frac{1}{2} \left(x_n + \frac{N}{x_n} \right)$$

2.

| Address | Command |
|---------|---------------|
| 101 | LOAD 1003 |
| 102 | STORE 1001 |
| 103 | LOAD 1005 |
| 104 | JUMPZERO 1007 |
| 105 | LOAD 1000 |
| 106 | DIVIDE 1001 |
| 107 | STORE 1004 |
| 108 | LOAD 1001 |
| 109 | SUBTRACT 1004 |
| 110 | DIVIDE 1006 |
| 111 | STORE 1001 |
| 112 | LOAD 1005 |
| 113 | SUBTRACT 1003 |
| 114 | STORE 1005 |
| 115 | LOAD 1001 |
| 116 | STORE 1002 |
| 117 | JUMP 103 |

| Address | Value |
|---------|---------|
| 1000 | n |
| 1001 | x_n |
| 1002 | result |
| 1003 | 1 |
| 1004 | N/x_n |
| 1005 | 5 |
| 1006 | 2 |

3. a)

| Compiler | Interpreter |
|---|--|
| Execute an entire program at a time | Execute a single line of code at a time |
| Create intermediate object code | No intermediate object code |
| Compilation takes place before execution | Compilation and execution are simultaneous |
| Faster | Slower |
| Require more memory | Require less memory |
| Difficult error detection as all errors are displayed after compilation | Easier error detection as errors of each line are displayed one by one |
| More efficient | Less efficient |

b)

Lexical Analysis:

- Identify the lexical units in a source code
- Classify lexical units into classes like constants, reserved words, and enter them in different tables. It will ignore comments in the source program
- Identify token which is not a part of the language

Syntax Analysis:

- Obtain tokens from the lexical analyzer
- Checks if the expression is syntactically correct or not
- Report all syntax errors
- Construct a hierarchical structure which is known as a parse tree

Semantic Analysis:

- Helps you to store type information gathered and save it in symbol table or syntax tree
- Allows you to perform type checking
- In the case of type mismatch, where there are no exact type correction rules which satisfy the desired operation a semantic error is shown
- Collects type information and checks for type compatibility
- Checks if the source language permits the operands or not

Intermediate Code Generation:

- It should be generated from the semantic representation of the source program
- Holds the values computed during the process of translation
- Helps you to translate the intermediate code into target language
- Allows you to maintain precedence ordering of the source language
- It holds the correct number of operands of the instruction

Code Optimization:

- It helps you to establish a trade-off between execution and compilation speed
- Improves the running time of the target program
- Generates streamlined code still in intermediate representation
- Removing unreachable code and getting rid of unused variables
- Removing statements which are not altered from the loop

Code Generation: Code generation is the last and final phase of a compiler. It gets inputs from code optimization phases and produces the page code or object code as a result. The objective of this phase is to allocate storage and generate relocatable machine code.

c) Lexical Analysis: Store in symbol table

Syntax Analysis: Store in parse tree.

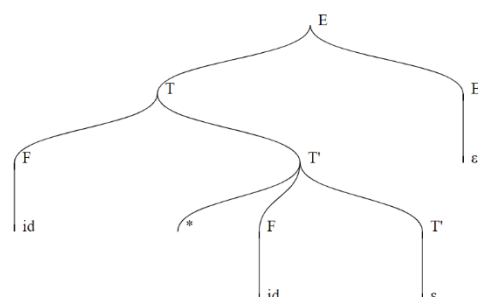
Code generation: Store in code.

4.

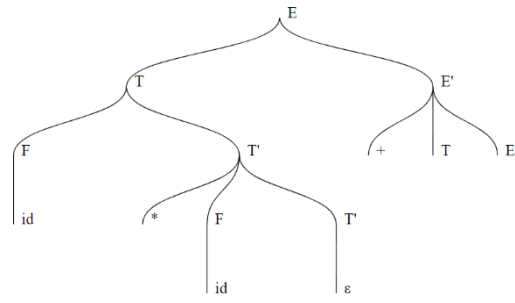
a) $E \rightarrow TE' \rightarrow FT'E' \rightarrow idT'E' \rightarrow idE' \rightarrow id+TE' \rightarrow id+FT' \rightarrow id+id*FT' \rightarrow id+id*id$

b) Yes as the sentence was parsed based on grammar provided by the parsing table.

c) **id * id**



d) Parsing failed:



5. a) Digit⁺ Other⁺
b) Tokens that do not follow the Digit – Other – Digit – Other ... pattern.