

# Building Calculus From Scratch

With Terence Tao's *Analysis I*

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## 1 Starting Axioms

1. 0 is a natural number.
2. if  $n$  is a natural number, then its successor  $n++$  is a natural number. We define the short-hand symbols

$$1 := 0++$$

$$2 := (0++)++$$

$$3 := ((0++)++)++$$

...

3. 0 is not the successor of any natural number. That is,  $0 \neq n++$  for any  $n$ .
4. Different natural numbers have different successors. That is,  $[n \neq m] \iff [n++ \neq m++]$  for any  $n, m$ . Equivalently, if we have that  $n++ = m++$ , then we have that  $n = m$ .
5. The principle of induction holds. That is, let  $P(n)$  be a property of the natural number  $n$ . Let  $P(0)$  be true, and let  $P(n) \implies P(n++)$ . Then  $P(n)$  is true for all natural numbers  $n$ .

## 2 Steps & Problems

1. Define the addition operation on two natural numbers,  $n$  and  $m$ .

We define incrementing  $m$  by zero as  $0 + m := m$ . Suppose inductively, that we know how to increment  $m$  by  $n$ . Then we can increment  $m$  by  $n++$  by defining

$$(n++) + m := (n + m)++$$

This recursive definition allows us to now add numbers (perform repeated incrementation). For example,  $2 + m = (1++) + m = (1 + m)++ = ((0++) + m)++ = ((0 + m)++)++ = ((m)++)++$ , which is exactly  $m$  incremented twice.

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2. Prove that  $n + 0 = n$  for any natural number  $n$

We will induct on  $n$ . We have that  $0 + 0 = 0$  by the fact that  $0 + m := m$  for any natural number  $m$ , including 0. Suppose, inductively, that  $k + 0 = k$  for some natural number  $k$ . Now consider the sum  $(k++) + 0 = (k + 0)++ \stackrel{\text{I.H.}}{=} (k)++$ . Hence for any natural number  $n$ , we have that  $n + 0 = n$ .

3. Prove that  $n + (m++) = (n + m)++$

We will induct on  $n$ . For  $n = 0$ , we have that  $0 + (m++) = m++ = (0 + m)++$ . Assume inductively that  $k + (m++) = (k + m)++$  for some natural number  $k$ . Now for the next number,  $k++$ , we have that  $(k++) + (m++) = (k + (m++))++$  (by the definition from problem 1)  $\stackrel{\text{I.H.}}{=} ((k + m)++)++ = ((k + m)++)++$ . Hence,  $n + (m++) = (n + m)++$  for all natural numbers  $n$ .

4. Prove that the natural numbers are closed under addition. That is, if  $a, b$  are natural numbers, then  $a + b$  is a natural number.

We will fix  $b$  and induct on  $a$ . Base case: 0 is a natural number by Axiom 1 and let  $b$  be a natural number. Then  $0 + b = b$  by our definition of incrementing by zero. But  $b$  is a natural number (by assumption). This closes the base case. Assume inductively that for some natural number  $a$ , if  $a$  and  $b$  are natural numbers then so is  $a + b$ . Consider the next natural number  $a++$ . By Axiom 2, since  $a$  is a natural number (I.H.), we must have that  $a++$  is also a natural number. Now consider the sum  $(a++) + b = (a + b)++$ . We have that  $a$  and  $b$  are natural numbers, so  $a + b$  is a natural number by the inductive hypothesis. By Axiom 2, we have that its successor  $(a + b)++$  must also be a natural number. This closes the induction. Hence the natural numbers are closed under addition.

5. Prove that the addition operation is commutative. That is, for any two natural numbers,  $a$  and  $b$ , we have that  $a + b = b + a$

We will fix  $b$  and induct on  $a$ . We have that  $0 + b = b$  from our first definition—incrementing a natural number by 0. Then  $b = b + 0$ , which we have from problem 2's proof. Hence we have our base case  $0 + b = b + 0$ . Suppose inductively that  $a + b = b + a$  for some natural number  $a$ . Now consider the next number,  $a++$ : we have  $(a++) + b = (a + b)++ \stackrel{\text{I.H.}}{=} (b + a)++ = b + (a++)$  by our previous proof. Hence for some natural number  $b$ , we have that  $a + b = b + a$  for all natural numbers  $a$ .

6. Prove that the additive cancellation law holds. That is, for any three natural numbers  $a$ ,  $b$ , and  $c$  such that  $a + b = a + c$ , we have  $b = c$ .

We will fix  $b, c$  and induct on  $a$ . For the base case  $a = 0$ , we will show that  $[0 + b = 0 + c] \implies [c = b]$ . Let us assume that  $0 + b = 0 + c$ . We have that  $0 + b = b$  and that  $0 + c = c$  from a previous proof. Then by transitive equality, we have  $b = c$ . This closes the base case. Assume inductively that for some natural number  $a$ , we have that  $[a + b = a + c] \implies [b = c]$ . Now consider the next natural number  $a++$ . Assume that  $(a++) + b = (a++) + c$ . Re-writing both sides, we have  $(a + b)++ = (a + c)++$ .  $a + b$  is

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a natural number and  $a + c$  is a natural number (by our additive closure proof). Recall by Axiom 4 that if the successors of two natural numbers are equal, then the natural numbers themselves must be equal. Hence we have that  $a + b = a + c$ . By the inductive hypothesis ( $[a + b = a + c] \implies [b = c]$ ), we must now have that  $b = c$ . This closes the inductive step. Thus, for any  $a, b, c$  such that  $a + b = a + c$ , we also have that  $b = c$ .

7. A natural number  $a$  is said to be positive iff  $a \neq 0$ . Prove that the sum of a natural number and a positive number must be positive.

Fix a positive number  $b$ . We will show that  $a + b$  is positive for all  $a$ . Base case:  $0 + b = b$  by our definition of incrementing by zero, and  $b$  is positive (by assumption). This closes the base case. Assume inductively that  $a + b$  is positive for some natural number  $a$ . Then consider the statement for the next natural number  $a++$ : we have  $(a++) + b = (a + b)++$ . By the inductive hypothesis, we have that  $a + b$  is positive. But this means that  $(a + b)++$ , the successor of  $a + b$ , must also be positive—it cannot be 0 because of Axiom 3: zero is not the successor of any natural number. This closes the induction.

8. Prove that, for natural numbers  $a$  and  $b$ , if  $a + b = 0$  then  $a = 0$  and  $b = 0$ .

We will prove the equivalent contraposition statement: if  $a \neq 0$  or  $b \neq 0$ , then  $a + b \neq 0$ . Assume that either  $a \neq 0$  or  $b \neq 0$ . In the case that  $a \neq 0$ ,  $a + b$  must be positive (according to the result of the previous proof) and hence non-zero. In the case that  $b \neq 0$ ,  $a + b$  must similarly be positive and hence non-zero.

9. (**EXERCISE 1**) Prove that the addition operation is associative. That is, for any three natural numbers,  $a, b, c$ , we have that  $(a + b) + c = a + (b + c)$

We will fix  $b, c \in \mathbb{N}$  and induct on  $a$ . Base case:

$$\begin{aligned}(0 + b) + c &= b + c \\ &= 0 + (b + c)\end{aligned}$$

Assume inductively that  $(a + b) + c = a + (b + c)$  for some natural number  $a$ . Then

$$\begin{aligned}((a++) + b) + c &= ((a + b)++) + c \\ &= ((a + b) + c)++ \\ &= (a + (b + c))++ \quad (\text{I.H}) \\ &= (a++) + (b + c)\end{aligned}$$

This closes the induction.

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10. **(EXERCISE 2)** Let  $a$  be a positive number. Prove that there exists exactly one natural number  $b$  such that  $b++ = a$

We effectively need to prove that each positive  $a$  has a unique predecessor  $b$ . Hence, it suffices to show that each  $a$  has both: 1) *at least* 1 predecessor and 2) *at most* 1 predecessor. If  $a$  is positive, then  $a \neq 0$ , which guarantees the existence of a  $b$  such that  $b++ = a$  (a predecessor of  $a$ ). Axiom 4 (different natural numbers have different successors) guarantees that no two distinct natural numbers have the same successor. Hence,  $a$  cannot have more than 1 predecessor  $\implies b$  must be the only predecessor of  $a$ .

Will will induct on  $a$ . Base case: for  $a = 1$ , we have  $b = 0 \implies b++ = a$ . This is the *only* such  $b$  for  $a = 1$  because a different  $b \neq 0$  would have a different successor  $b++ \neq 1$  according to Axiom 4: different natural numbers have different successors. Suppose inductively that for some natural number  $a$  there exists a  $b$  such that  $b++ = a$ . Now consider the next natural number  $a++$ : the number  $a++$ , by definition, is the successor of  $a$ . So we have found one  $b = a$  such that  $b++ = a++$ . This is the only such  $b$  because, similarly, by Axiom 4, any other  $b \neq a$  has a different successor  $b++ \neq a++$ . This closes the induction.

11. Define the ordering of the natural numbers.

Let  $n$  and  $m$  be natural numbers. We say that  $n$  is *greater than or equal to*  $m$ , and write  $n \geq m$  or  $m \leq n$ , iff we have  $n = m + a$  for some natural number  $a$ . We say that  $n$  is *strictly greater than*  $m$ , and write  $n > m$  or  $m < n$ , iff  $n \geq m$  and  $n \neq m$ .

12. **(EXERCISE 3)** Prove properties of order.

(a) (Reflexivity) Prove that  $a \geq a$  for any natural number  $a$ .

For any natural  $a$ , we can find an  $m$  such that  $a = a + m$  ( $a \geq a$ )  $\implies$  choose  $m = 0$

(b) (Transitivity) Prove that if  $a \geq b$  and  $b \geq c$ , then  $a \geq c$ .

$$\begin{aligned}
 a \geq b, b \geq c &\implies a = b + k_0, b = c + k_1 \text{ for some } k_0, k_1 \in \mathbb{N} \\
 &\implies a = (c + k_1) + k_0 \\
 &\implies a = c + (k_1 + k_0) \text{ (associativity)} \\
 &\implies a = c + m \text{ for some natural } m = k_0 + k_1 \text{ (additive closure)} \\
 &\implies a \geq c
 \end{aligned}$$

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(c) (Anti-symmetry) Prove that if  $a \geq b$  and  $b \geq a$ , then  $a = b$

$$\begin{aligned}
a \geq b, b \geq a &\implies a = b + k_0, b = a + k_1 \text{ for some } k_0, k_1 \in \mathbb{N} \\
&\implies a = (a + k_1) + k_0 \\
&\implies a = a + (k_1 + k_0) \text{ (associativity)} \\
&\implies a + 0 = a + (k_1 + k_0) \\
&\implies 0 = k_1 + k_0 \text{ (cancellation law)} \\
&\implies k_0 = 0 \text{ and } k_1 = 0 \text{ (zero-sum proof)} \\
&\implies a = b
\end{aligned}$$

(d) (Order preservation under addition) Prove that  $a \geq b$  if and only if  $a + c \geq b + c$ .

$$\begin{aligned}
a \geq b &\iff \exists k(a = b + k) \\
&\iff a + c = (b + k) + c \\
&\iff a + c = b + (k + c) \\
&\iff a + c = b + (c + k) \\
&\iff a + c = (b + c) + k \\
&\iff a + c \geq b + c
\end{aligned}$$

(e) Prove that  $a < b$  if and only if  $a++ \leq b$ .

$$\begin{aligned}
a < b &\iff \exists k(b = a + k \text{ and } a \neq b) \\
&\iff a + k \neq a \\
&\iff k \neq 0 \\
&\iff b = a + m++ \text{ with } k = m++ \\
&\iff b = a + (m + 1) \\
&\iff b = a + (1 + m) \\
&\iff b = (a + 1) + m \\
&\iff b = (a++) + m \\
&\iff a++ \leq b
\end{aligned}$$

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(f) Prove that  $a < b$  if and only if  $b = a + d$  for some positive number  $d$ .

$$\begin{aligned}a < b &\iff \exists k(b = a + k \text{ and } a \neq b) \\&\iff a + k \neq a \\&\iff k \neq 0 \\&\iff \text{choose } d = k \\&\iff b = a + d \text{ and } d \text{ is positive}\end{aligned}$$

**13. EXERCISE 4** Justify the following:

(a)  $0 \leq b$  for all  $b$ .

For any  $b$ , we can find an  $m$  such that  $b = 0 + m$  ( $0 \leq b$ )  $\implies$  choose  $m = b$

(b) If  $a > b$ , then  $a++ > b$

$$\begin{aligned}a > b &\implies a++ > a > b \text{ (definition of successor)} \\&\implies a++ > b \text{ (transitivity)}\end{aligned}$$

(c) If  $a = b$ , then  $a++ > b$

$$\begin{aligned}a = b &\implies a++ > a = b \text{ (definition of successor)} \\&\implies a++ > b\end{aligned}$$

**14. EXERCISE 5** Prove the principle of strong induction.

$$x \implies x$$