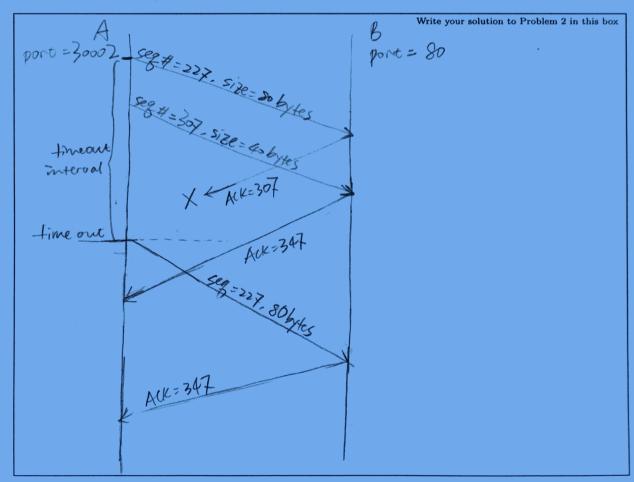
Answer True or False to the following questions and briefly justify your answer:

- (a) With the Selective Repeat protocol, it is possible for the sender to receive an ACK for a packet that falls outside of its current window.
- (b) With Go-Back-N, it is possible for the sender to receive an ACK for a packet that falls outside of its current window.
- (c) The Stop&Wait protocol is the same as the SR protocol with a sender and receiver window size of 1.
- (d) Selective Repeat can buffer out-of-order-delivered packets, while GBN cannot. Therefore, SR saves network communication cost (by transmitting less) at the cost of additional memory.

Write your solution to Problem 1 in this box A simple example would be a window of size 2 sender sends packet 1, 2, and receiver sends back Ack, ACK, The sender times one Ack1 & Ack2, and then retransmit Parket 1 & 2. Receiver receives the duplicate and resends Ack 1-1, Ack 1-2. After this the sender receives Ack 1 Ack 2 due to delay, and the window moves on to packet 324. Then Ack 1-1 & ACK 1-2 arrives, and parter 1& 2 are one of the window Same example as question (a) If the window site is 1, there would be no out-of-order packets (within the window). Cumulative Alks are just ordinary Alks. Thus, in this situation Stops Wort is the same as SR (d) True Instead of olis (arding the out-of-order packets, as GBN does, SR protocol holds them in a buffer (wost of additional memory), so that the send only needs to retromsmit the host packet, rather than everything after it.

Host A and B are communicating over a TCP connection, and Host B has already received from A all bytes up through byte 226. Suppose Host A then sends two segments to Host B back-to-back. The first and second segments contain 80 and 40 bytes of data, respectively. In the first segment, the sequence number is 227, the source port number is 30002, and the destination port number is 80. Host B sends an acknowledgment whenever it receives a segment from Host A. Fill in the blanks for questions (a) - (c) directly; work out the diagram in the box for question (d).

- (a) In the second segment sent from Host A to B, the sequence number is 307, source port number is 3002, and destination port number is 80.
- (b) If the first segment arrives before the second segment, in the acknowledgment of the first arriving segment, the ACK number is 207, the source port number is 20002.
- (c) If the second segment arrives before the first segment, in the acknowledgment of the first arriving segment, the ACK number is $22 \frac{1}{2}$.
- (d) Suppose the two segments sent by A arrive in order at B. The first acknowledgment is lost and the second acknowledgment arrives after the first timeout interval. Draw a timing diagram in the box below, showing these segments and all other segments and acknowledgment sent. Assume no additional packet loss. For each segment in your diagram, provide the sequence number and the number of bytes of data; for each acknowledgment that you add, provide the ACK number.



In Fast Retransmit algorithm, we saw TCP waits until it has received three duplicate ACKs before performing a fast retransmit. Why do you think the TCP designers chose not to perform a fast retransmit after the first or second duplicate ACKs for a segment received?

Whenever the receiver receives and of-order packet, it issues despicate Alks. Therefore, performing retrumsmittion after only the first or swand displicate. Alk would rause too many unlessary redundant packets in the natural. This would make network congestion even worse, while at the same time, the retromsmittions are mostly triggered but one-of-order packets, routher than actual packet loss.

Suppose that three measured SampleRTT values are 106 ms, 120 ms, and 140 ms. Compute the EstimatedRTT after each of these SampleRTT values is obtained, assuming that the value of EstimatedRTT was 100 ms just before the first of these three samples were obtained. Compute also the DevRTT after each sample is obtained, assuming the value of DevRTT was 5 ms just before the first of these three samples was obtained. Last, compute the TCP TimeoutInterval after each of these samples is obtained.

```
Write your solution to Problem 4 in this box
First:
   Fistimated RTT = 0.125 x Sample RTT + (1-0.125) x Ustimouted RTT
                   = 0.175 x 106 + 0.875 x 100
                   = 100.75 ms
  DevRTT = 0.25 x | SampleRTT- EstsmatedRTT | + 11-0.25) x DevRTT
           = 0.25 x 5,25 + 0.75 x 5
           = 5.0625ms
  TCP Timeout Interval = Estimated RTT + 4x DevRTT
                      = 100.75 + 4×5.0625
                      = 121 mc
Second
  Estimated RTT = 0.125 x 120 + 0.875 x 100.75
                = 103, 15625 ms
  DerRTT = 0.25 x 16.84 +0.75 x 5.0625
          = 8 mc
  TOTIMEOUT INTERVAL = 103, 15625 + 4×8
                      = 135.16 ms
Third:
  Extimated RTT = 0.125 x 140 + 0.875 x 103. 15625
                = 107 75 ms
  DevRTT = 0,25 x 32.25 + 0.75 x8
          = 14.06 mc
  TCD7 meout Interval = 107.75 + 4×14.06
                        = 164 ms
```

Compare Go-Back-N, Selective Repeat, and TCP (no delayed ACK). Assume that timeout values for all three protocols are sufficiently long, such that 5 consecutive data segments and their corresponding ACKs can be received (if not lost in the channel) by the receiving host (Host B) and the sending host (Host A), respectively. Suppose Host A sends 5 data segments to Host B, and the 2nd segment (sent from A) is lost. In the end, all 5 data segments have been correctly received by Host B.

- (a) How many segments has Host A sent in total and how many ACKs has Host B sent in total? What are their sequence numbers? Answer this question for all three protocols.
- (b) If the timeout values for all three protocols are much longer than 5RTT, then which protocol successfully delivers all five data segments in shortest time interval?

```
Write,

9 Segmences sent 1 from A to B); 123452345

8 ACK with seg # 1 from B to A): 1 1112345
                                                        Write your solution to Problem 5 in this box
 6 segments sent (from A to B): 123452
 5 Ack (from B to A) : 1 3452
 6 segments sent (from A to B): 123452

5 ACK (from B to A): 2 2226
   Seconse TCP uses fast retromsmit.
    i.e. TCP retransmits after three duplicate Acks
        TOP retransmits after three augusta while GBN & SR protocols don't retransmit until timeout.
```