

CSC265 Fall 2020 Homework Assignment 10

Solutions

The list of people with whom I discussed this homework assignment:

1. Suppose that, during the BFS of an undirected graph $G = (V, E)$, node a is first visited before node b , which is first visited before node c . Prove that if $\{a, c\} \in E$, but $\{a, b\} \notin E$, then there exists a neighbour d of b which is visited before node a .

Solution: We will use the following pseudocode from the text (page 595).

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BFS( $G, s$ )
1  for each vertex  $u \in V \setminus \{s\}$ 
2       $u.color = \text{WHITE}$ 
3       $u.d = \infty$ 
4       $u.\pi = \text{NIL}$ 
5   $s.color = \text{GRAY}$ 
6   $s.d = 0$ 
7   $s.\pi = \text{NIL}$ 
8   $Q = \emptyset$ 
9  ENQUEUE( $Q, s$ )
10 while  $Q$  is nonempty
11      $u = \text{DEQUEUE}(Q)$ 
12     for each  $v$  such that  $\{u, v\} \in E$ 
13         if  $v.color == \text{WHITE}$ 
14              $v.color = \text{GRAY}$ 
15              $v.d = u.d + 1$ 
16              $v.\pi = u$ 
17             ENQUEUE( $Q, v$ )
18      $u.color = \text{BLACK}$ 
```

Suppose $\{a, c\} \in E$, $\{a, b\} \notin E$. We know from lemma in class that if a is visited before b , which is visited before c , then $a.d \leq b.d \leq c.d$. However, since $\{a, c\} \in E$ and c is visited after a , we know when a is dequeued, $c.color = \text{WHITE}$, so $c.d = a.d + 1$ by line 15. This means we must have $a.d < b.d = c.d$ or $a.d = b.d < c.d$.

Case 1: $a.d < b.d = c.d$. We know that c is visited when a is dequeued but b cannot be visited when a is dequeued because $\{a, b\} \notin E$. This means b is visited before a is dequeued. However, notice that b is visited also when its parent $b' = b.\pi$ is dequeued. This means b' is enqueued before a , and thus b' is visited before a .

Case 2: $a.d = b.d < c.d$. We know from line 15 and 16 that $(b.\pi).d + 1 = b.d$, so $(b.\pi).d < b.d = a.d$. By lemma in class, we know this implies that $b.\pi$ is visited before a .

In either case, we know that $b.\pi$ (which is a neighbour of b) is visited before a , so we are done.

2. Suppose that, during the DFS of an undirected graph $G = (V, E)$, node a is first visited before node b , which is first visited before node c . Complete the following sentence and prove it is correct (I can think of two versions, so I will show both are true): If $\{a, c\} \in E$, but $\{a, b\} \notin E$, then there exists a neighbour d of b such that when a is first visited, there is a path of unvisited nodes from a to d (this path is of length ≥ 1 , since $\{a, b\} \notin E$) except for a (call this version 1). Another version (call this version 2) says that there exists a neighbour d of b such that d is visited before node c .

Solution: We will use the following pseudocode from the text (page 604).

DFS(G)

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1  for each vertex  $u \in V$ 
2       $u.color = \text{WHITE}$ 
3       $u.\pi = \text{NIL}$ 
4   $time = 0$ 
5  for each vertex  $u \in V$ 
6      if  $u.color == \text{WHITE}$ 
7          DFS-VISIT( $G, u$ )

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DFS-VISIT(G, u)

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1   $time = time + 1$ 
2   $u.d = time$ 
3   $u.color = \text{GRAY}$ 
4  for each  $v$  such that  $\{u, v\} \in E$ 
5      if  $v.color == \text{WHITE}$ 
6           $v.\pi = u$ 
7          DFS-VISIT( $G, v$ )
8   $u.color = \text{BLACK}$ 
9   $time = time + 1$ 
10  $u.f = time$ 

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Suppose $\{a, c\} \in E$, $\{a, b\} \notin E$. From the nested/disjoint interval theorem in class (and page 606, theorem 22.7 in CLRS), we know that for any two nodes u, v , $(v.d, v.f) \subset (u.d, u.f)$ if and only if v is a proper descendant of u in the DFS tree. Since $\{a, c\} \in E$ and $a.d < c.d$, we know that c must be visited before DFS-VISIT(G, a) is completed by line 5 of DFS-VISIT. So we know $(c.d, c.f) \subset (a.d, a.f)$. We are given $a.d < b.d < c.d$ in the question. Combining with the fact that $b.f < c.f$, we have $(b.d, b.f) \subset (a.d, a.f)$. Thus, b is a proper descendant of a . Let $a_0 = a, a_1, \dots, a_k = b$ be the simple path from a to b in the DFS tree. Since $\{a, b\} \notin E$, we know that $k \geq 2$. Let d be a_{k-1} . It remains to show that a_1, \dots, a_{k-1} are all unvisited. This is clearly true since $a_{i+1}.d = a_i.d + 1$ for $i = 0, \dots, k-1$ by definition of DFS tree. So $a_j.d \geq a_1.d > a.d$ for all $j \in \{1, \dots, k-1\}$, meaning that a_1 to $a_{k-1} = d$ is a path containing unvisited nodes. This proves version 1.

Moreover, this implies that a_{k-1} is visited before c because $(a_{k-1}).d < b.d < a.d$. This proves version 2.