

Clase 6 Expresiones Regulares

- AF \rightarrow Languages Regulares.

- ER \rightarrow forma declarativa
- LR

Expresión algebraica

$$\textcircled{0} (0+1)^* 1$$

Regex

○ percablos

1. Union $L \cup M$ L. Resulten.

$$L = \{001, 10, 111\}, M = \{\varepsilon, 001\}$$

$$L \cup M = \{\varepsilon, 10, 001, 111\}$$

2. concatenación

LM o' LM

$$[Y = \{ \underbrace{001, 10, 111}_{001}, \underbrace{001001, 10001, 111001}_{1001} \}]^*$$

$L \in$

$$LM \neq ML$$

3. Clausura de Kleene

L^*

$L = \{0, 1\} \Rightarrow L^*$ son todos
los cadenas
con 0s y 1s

$L = \{0, 11\} \Rightarrow L^*$ todos cadenas
que tiene 11
1 en pares
011, 11110, ...

$$L^* = \bigcup_{i \geq 0} L^i$$

$$L^0 = \{\varepsilon\}, L^1 = L, \dots$$

$$L^i = L L L \dots L \text{ (i veces)}$$

Ej. $L = \{0, 11\}$

$$L^0 = \{\varepsilon\}$$

$$L^1 = \{\varepsilon, 0, 11\}$$

$$L^2 = \{00, 011, 110, 1111\}$$

...

$L^* \sim$ infinito

ϵ_1 : $L \sim$ conjunto de tiras \emptyset

$$L \rightarrow \infty$$

$$L^* \rightarrow \infty$$

$$L^1 = L$$

$$L^2 = L$$

$$L^3 = L$$

$$\epsilon_1: L = \emptyset \rightarrow \emptyset = L^*$$
$$\{\epsilon\} \neq \emptyset$$

Construcción de Expresiones Regulares.

$L(E) \sim$ language descrito por $E \in ER$

$L(A) \sim$ language descrito por $A = AFD$

Base: 1. ϵ y \emptyset son $\in ER$

$$L(\epsilon) = \{\epsilon\}$$

$$L(\emptyset) = \emptyset$$

2. si a es un símbolo

a es un $\in ER$

\nearrow negrita $L(a) = \{a\}$

3. $L \sim$ Representa un lenguaje.

Paso Inductivo:

1. $E, F \in R$

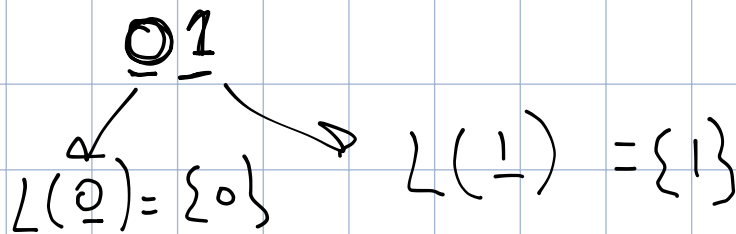
$$L(E + F) = L(E) \cup L(F)$$

$$2. L(EF) = L(E)L(F)$$

$$3. L(E^*) = (L(E))^*$$

$$4. L((E)) = L(E)$$

Ej.



01

- cadenas 01?

$01^* \neq (01)^* \leftarrow$ 1 o más veces

\downarrow	\downarrow
01	01 ✓
011	0101 ✓
01111	010101 ✓
⋮	⋮

no: acepta

- 0 y 1 alternados

$$(\epsilon + 1)(01)^*(\epsilon + 0)$$

\swarrow ó \nwarrow 1, mas veces \swarrow agrupar

101
010

$$(\epsilon + 0)^* = \{\epsilon, 0, 00, 000, 0000, \dots\}$$

$$\{101, 010, 1010, 0101, \dots\}$$

$$L((\epsilon + 1)(01)^*(\epsilon + 0)) =$$

$$\{\epsilon\} \rightarrow L^0 \cup L^1 \cup L^2 \cup \dots$$

Precedencia

$$* > \cdot > +$$

\nearrow
concatenación

\nearrow
concatenación

\searrow " "

$$(01)^*$$

$$xy^2 \neq (xy)^2$$

$$0.1^* \leftarrow 1$$

$$0.1^* = 0(1^*) = 0(1)^*$$

des pures

$$1^* = \{\epsilon, 1, 11, 111\}$$

$$11^* = \epsilon \notin 11^*$$

$$1^* = \{\epsilon, 1, 11, 111, \dots\}$$

$$11^* = \{1, 11, 111, 1111, \dots\}$$

$$\overset{a}{(\epsilon + a + b + c)^*} \overset{b}{(a + b)(\epsilon + a + b + c)^*} (a + b)(\epsilon + a + b + c)^*$$

$$\begin{array}{c} (\epsilon + 11)(01)^* \\ \uparrow \\ 0^* (10)^* (01)^* \end{array}$$

0 1 0 1 0 1 0 1 0 1

$$(0^* 1)^*$$

$$(0^* (01)^*)^*$$

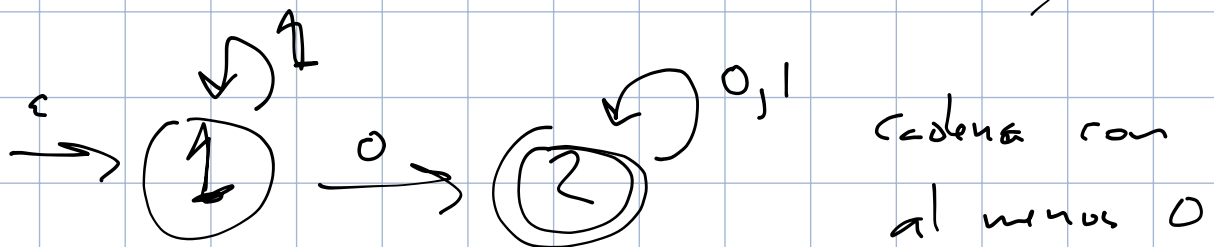
$$0^* (01)^*$$

$$\begin{array}{c} 11 \\ 0110 \\ 010110101 \end{array}$$

010110001
 0
 11000010100000 $0101, \dots$
 001001

$AFD \leftrightarrow AFN \leftrightarrow AFN-\epsilon \leftrightarrow ER$
 Language Regular

$AFD \rightarrow ER$ (3 métodos)



$R_{ij}^k \sim ER$ de i a j que tiene
 k estados intermedios

R_{ij}^0

$\gamma, R_{ij}^0 \sim i a j$

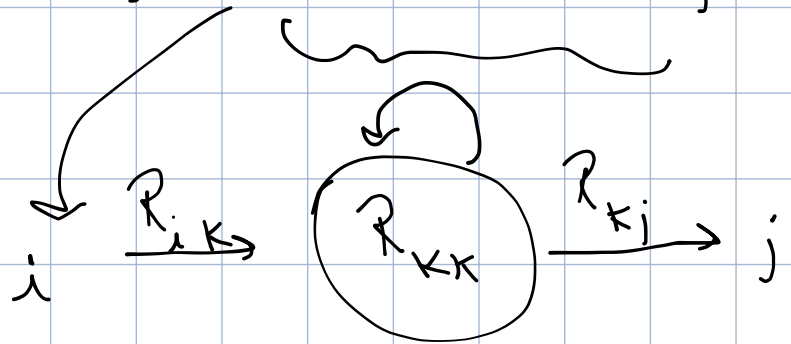
$$R_{11}^0 = \varepsilon + 1$$

$$R_{12}^0 = 0$$

$$R_{21}^0 = \emptyset$$

$$R_{22}^0 = \varepsilon + 0 + 1$$

2. Inductive $R_{ij}^k = R_{ij}^{k-1} + R_{ik}^{k-1} (R_{kk}^{k-1})^* R_{kj}^{k-1}$



$$R_{ij}^1 = R_{ij}^0 + R_{ik}^0 (R_{kk}^0)^* R_{kj}^0$$

R_{11}^1	$R_{11}^0 + R_{11}^0 (R_{11}^0)^* R_{11}^0 = \varepsilon + 1 + (\varepsilon + 1)(\varepsilon + 1)^*(\varepsilon + 1) = \varepsilon + 1 + (\varepsilon + 1)^+$
R_{12}^1	$R_{12}^0 + R_{11}^0 (R_{11}^0)^* R_{12}^0 = 0 + (\varepsilon + 1)(\varepsilon + 1)^* 0 = 0 + 1^+ 0 = 1^+$
R_{21}^1	$\emptyset + \emptyset (\varepsilon + 1)^* (\varepsilon + 1) = \emptyset$
R_{22}^1	$\emptyset + \emptyset (\varepsilon + 1)^* (\varepsilon + 1) = \emptyset$

$$R_{22} = \epsilon + 0 + 1 + \emptyset(\epsilon + 1)^* 0 = \epsilon + 0 + 1$$

$$(\epsilon + R)^* = R^*, \quad \emptyset R = R \emptyset = \emptyset$$

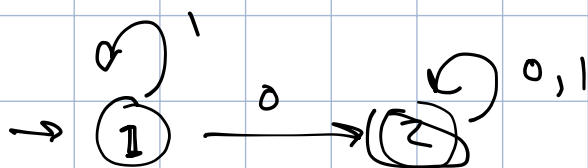
$$0 + 1^* 0 \Rightarrow \{0, 10, 110, 1110, \dots\}$$

$$R_{ij}^2 = R_{ij}^1 + R_{i2}^1 (R_{22}^1)^* R_{2j}^1$$

	R_{ii}^2	1^*
\rightarrow	R_{i2}^2	$1^* 0 (\epsilon + 1)^*$
	R_{2i}^2	\emptyset
	R_{22}^2	$(\epsilon + 1)^*$

$$L = 1^* 0 (\epsilon + 1)^*$$

Eliminación de estados



$$0, 1 = 0 + 1$$





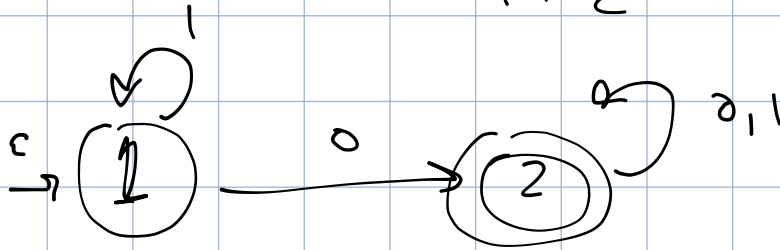
$$1^*0(0+1)^* \quad \checkmark$$

Método Arden

Teorema de Arden $R, Q, P \in ER$

$$R = Q + RP \Rightarrow R = QP^*$$

$$P \neq \epsilon$$



$$q_1 = \epsilon + q_1 \xrightarrow{1} q_1 = \underline{\epsilon(1)^*}$$

$$q_2 = q_1 0 + q_2 0 + q_2 1 = q_1 0 + q_2 (0+1)$$

$$\underbrace{Q + RP = QP^*}$$

$$q_2 = \underline{q_1 0 (0+1)^*}$$

$$q_2 = \underline{1^*0(0+1)^*} \quad \checkmark$$

ER