#### Workload Runtime And Placement (WRAP)

NGRM High-Level Review March 4 2013

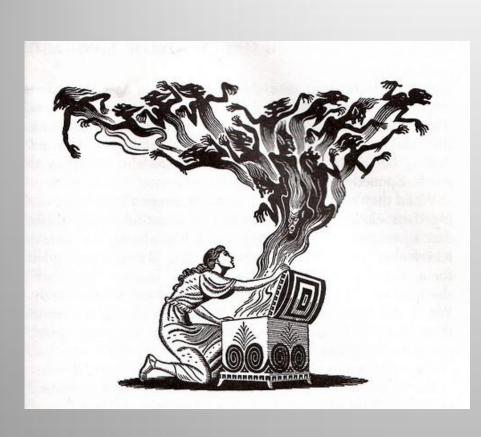
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## WRAP concerns all aspects of executing transactions of a job in the new paradigm.

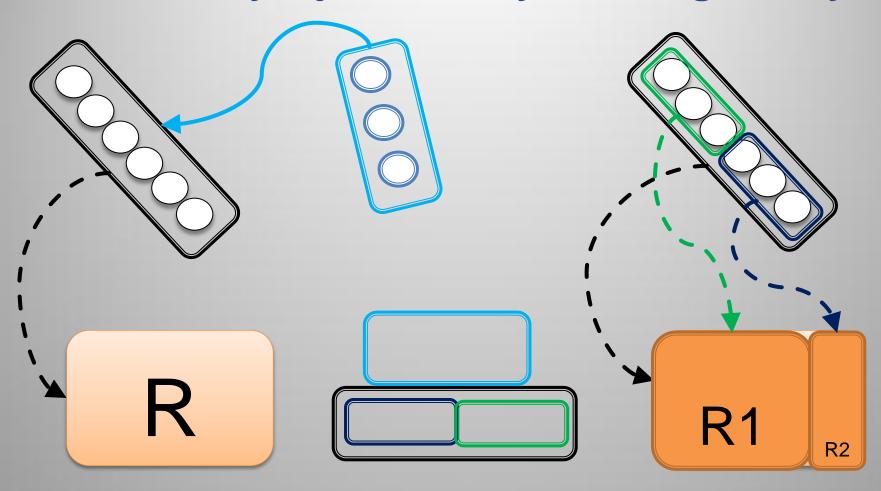


- Scheduler sets the overall bound for resources and the duration for a job, and now what?
- Opening the job will let loose various transactions on this bound to execute.
- WRAP must provide runtime services to execute them most effectively and efficiently under the new RM paradigm.

### The new paradigm needs new ways to organize/group processes a job executes.

- The traditional approach models transactions as a set of compute steps (e.g., job steps).
  - Limits the ability to seamlessly integrate various transactions beyond compute transactions (e.g., tools).
  - Limits the ability to relate one group of processes to another.
  - Limits the ability to enable resource elasticity.
- Designing the WRAP services after this traditional model would be an under-design.

## LWJ is our model to group processes executed by a job in many meaningful way.

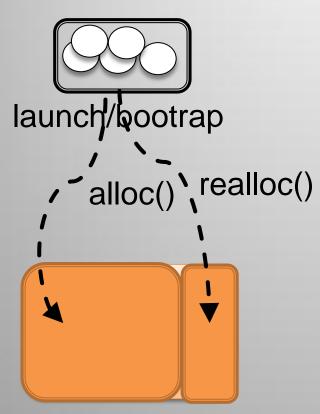


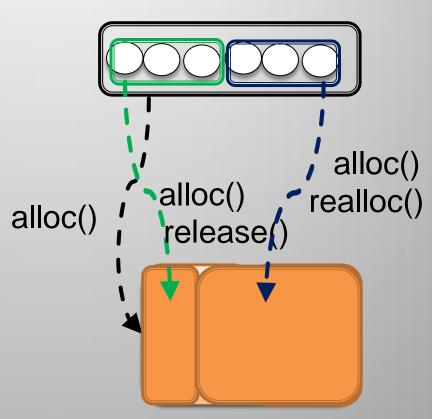
# WRAP uses LWJs to express run-time services concisely.

- LWJ serves as group identifiers for WRAP services to relate each LWJ to diverse resources and other LWJs.
- Resource allocation and elasticity: alloc(lwj, c), realloc(lwj,c), and release(lwj, r).
- Process management/confinement: launch(lwj), destroy(ljw), etc.
- Synchronization: sync(lwj(i), lwj(k)).
- Resource discovery and provenance: query(lwj) and record(lwj).

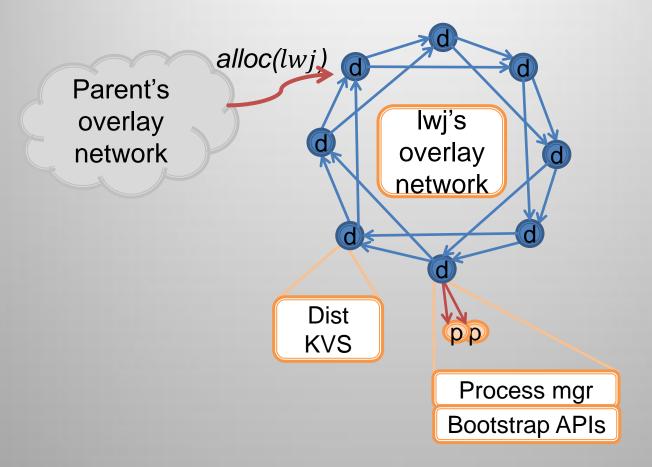
# We can compose the primitives to enable high-level services of the new paradigm.

grow and shrink for elasticity on any resource

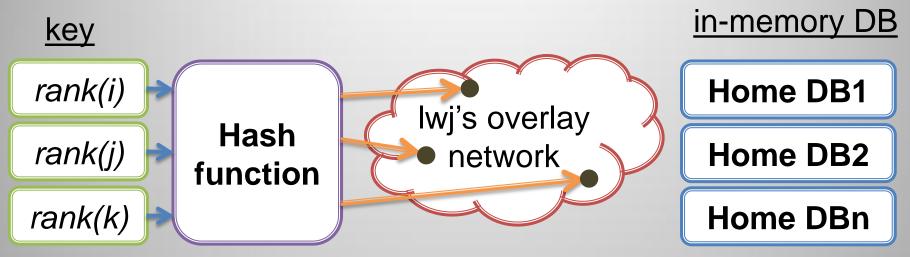




## The base architecture builds on comms. framework and distributed key-value store.



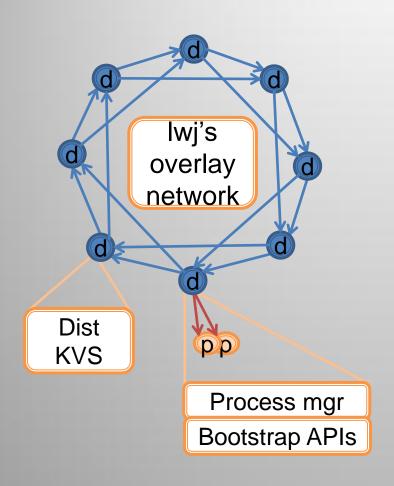
## **DKVS** is scalable distributed shared memory for an LWJ and its descendants.



- Get/put for data access
- Collective Fence for memory consistency

lwj(1)::resource	cores (128)	power(10KW)	lic (10 tokens)	
lwj(1)::rank(10)	host(1)	pid(345)	port(445)	
lwj(1)::record	info1	info2		
lwj(1)::lwj(2)::resource	cores(64)	power(4KW)	lic (2 tokens)	

#### All process management operations must be scalable.

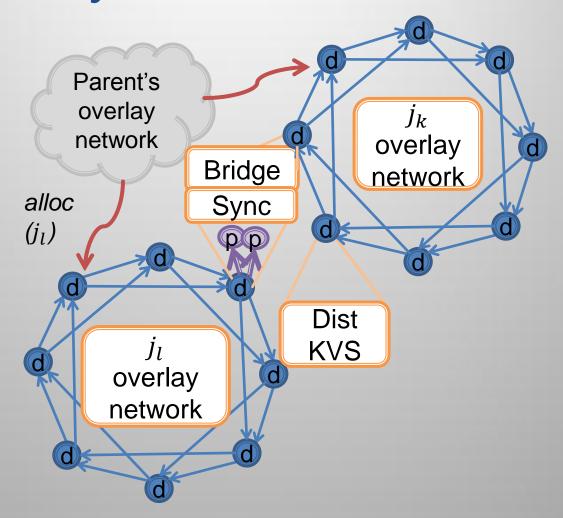


- Environment propagation: Bcast problem O(log(N))
- stdin/signal: Bcast problem, O(log(N))
- stderr/stdout handling:
  Reduction problem, O(log(N))
  with an ability to reduce the
  outputs on the fly
- Process termination detection/analysis: Reduction problem: O(log(N)) with an ability to reduce the outputs on the fly

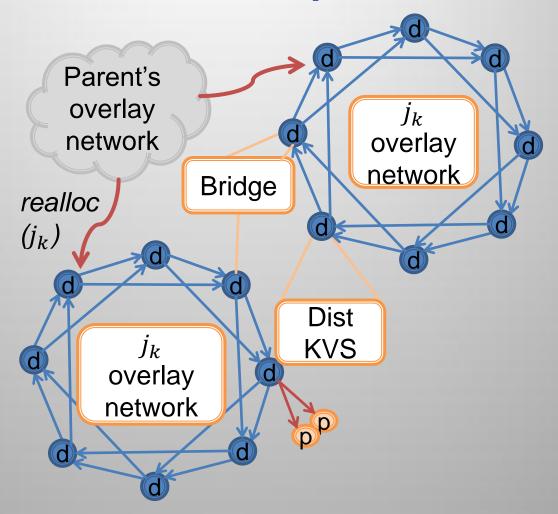
### **DKVS** allows ease integration with various types of LWJs beyond MPI

- PMI 1, 2 will be a very thin layer on top of DKVS.
- PMGR, PMGR Collective, COBO, LaunchMON, and LIBI use essentially the same bootstrapping technique: launched processes discover open TCP ports of their children processes or parent process in the binomial tree.
- Our plug-ins for these well-known bootstrappers will serve as the reference implementation.
- Other types of LWJs can write their own plug-ins for ease integration into WRAP.

#### We can easily extend the base architecture to implement *sync*.



### We can easily extend the base architecture to implement *realloc* w/ respect to CNs.



#### We devised a WRAP architecture that will allow us to create an effective development plan.

- WRAP has analyzed the requirements and use cases to conceive novel models as well as run-time services.
- WRAP devised a high-level architecture that can the services scalably.
- WRAP's architecture allowed us to identify the major components and their subsystems to develop (Unsolicited section).
- WRAP plans to use a phased approach to mitigate the risk of developing these components (Unsolicited section).
- WRAP work items will need to be considered in the context of the global Work Breakdown Structure (WBS).

