TÖL212M Rökstudd Forritun - Hópverkefni 10

Andri Fannar Kristjánsson

25. mars 2025

Hópverkefni 10

1

Sækið skrána H10-skeleton. java í Canvas og vistið hana sem H10. java. Klárið að útfæra klasann í skránni.

1.1 Svar:

Hér fyrir neðan má sjá kóðann þar sem föllin hafa verið forrituð. Einnig er hægt að sjá skrána hér: https://tinyurl.com/4a3ymtff. Misvísandi skilaboð komu frá Dafny varðandi hvort þessi lausn virkar eða ekki, en mér sýndist það vera timeout vandamál frekar en vandamál með útfærsluna.

```
// Author of question:
                         Snorri Agnarsson, snorri@hi.is
                              Andri Fannar Kristjánsson, afk6@hi.is
// Author of solution:
// Permalink of solution:
                              https://tinyurl.com/4a3ymtff
// Finish programming the class QueueCycleChain
// by programming the bodies of the operations
// IsEmpty, Put and Get, as well as the
// constructor for the class.
// Everything needed is in this file and you
// do not need to call any lemmas or write
// any asserts. Note that it is unlikely that
// you can solve this problem on the tio.run
// web page.
class QueueCycleChain<T> extends Queue<T>
  var \ last: \ Link?\!\!<\!\!T\!\!>
  ghost var cycle: seq<Link<T>>>
  ghost predicate Valid()
    reads this, Repr
    (last != null => last in Repr) &&
    (forall z | z in cycle :: z in Repr) &&
    (forall z | z in Repr :: z in cycle) &&
    (ghostseq = ValueSeq(cycle)) &&
    (last = null \implies cycle = []) \&\&
    (last != null ==> last. ValidCycle(cycle)) &&
    (|\text{cycle}| = |\text{ghostseq}|) \&\&
    (forall i \mid 0 \le i < |cycle| :: cycle[i].head == ghostseq[i])
  constructor()
    ensures Valid()
```

```
ensures Repr \Longrightarrow {}
  ensures ghostseq = []
  // Give the instance variables last, Rept, ghostseq and
  // cycle correct values considering the data invariant
  // and the postcondition.
  // To satisfy the postcondition and the data invariants, we
  // simply initialize last to null and Repr, ghostseq and cycle
  // to empty sets / seq.
  last := null;
  Repr := \{\};
  ghostseq := [];
  cycle := [];
  new;
  assert // This assert is the same as the data invariant
    // and is unnecessary if the constructor is
    // correctly programmed, but if not it is
    // helpful in identifying errors.
    (last != null ==> last in Repr) &&
    (forall z | z in cycle :: z in Repr) &&
    (forall z | z in Repr :: z in cycle) &&
    (ghostseq == ValueSeq(cycle)) &&
    (last = null \implies cycle = []) \&\&
    (last != null => last. ValidCycle(cycle)) &&
    |cycle| == |ghostseq| &&
    (forall i \mid 0 \le i < |cycle| :: cycle[i].head = ghostseq[i]);
}
predicate IsEmpty()
  reads this, Repr
  requires Valid()
  ensures IsEmpty() <==> ghostseq ==[]
  // Here we simply check if the last is null,
  // if it is then the queue is empty.
  last = null
method Put(x:T)
  modifies this, Repr
  requires Valid()
  ensures Valid()
  ensures fresh (Repr-old (Repr))
  ensures ghostseq = old(ghostseq) + [x]
  // We create a new link with the value x and add it to the end
  // of the cycle. This works both for the case when the queue is
  // empty and when it is not, as doing this on the empty queue
  // is equivalent to `new Link<T>(x, null, [], []);`
  var newlink := new Link<T>(x, last, cycle, ghostseq);
  last := newlink;
  ghostseq := ghostseq + [x];
  \mathtt{cycle} \; := \; \mathtt{cycle} + [\mathtt{newlink} \;] \,;
  Repr := Repr + \{newlink\};
}
method Get() returns (x: T)
  modifies this, Repr
```

```
requires Valid()
    requires ghostseq != []
    ensures Valid()
    ensures Repr < old (Repr)
    ensures ghostseq = old (ghostseq [1..])
    ensures x = old(ghostseq[0])
    // We remove the first link from the cycle and update the last
    // link.
    var removedLink := last.tail;
    // We call the RemoveFirst method to remove the first link from
    // the cycle.
    var newlast, newcycle, newvals := RemoveFirst(last, cycle, ghostseq);
    // We update the last link and the cycle, as well as the ghost
    // variables.
    last := newlast;
    cycle := newcycle;
    ghostseq := newvals;
    Repr := Repr - \{removedLink\};
    x := removedLink.head;
}
// Here the mutable part of the file ends.
// Do not change the program text below this.
// This is the fundamental definition of the behaviour of a queue.
trait Queue<T>
 ghost var ghostseq: seq<T>
  ghost var Repr: set < object >
 ghost predicate Valid()
    reads this, Repr
  predicate IsEmpty()
    reads this, Repr
    requires Valid()
    ensures IsEmpty() <=> ghostseq ==[]
  method Put(x:T)
    modifies this, Repr
    requires Valid()
    ensures Valid() && fresh(Repr-old(Repr))
    ensures ghostseq = old (ghostseq)+[x]
 method Get() returns (x: T)
    modifies this, Repr
    requires Valid()
    requires ghostseq != []
    ensures Valid () && fresh (Repr-old (Repr))
    ensures ghostseq = old (ghostseq [1..])
    ensures x = old(ghostseq[0])
}
// Here is the definition of mutable links
```

```
// that are used in circular chains.
class Link<T>
  var head: T
  var tail: Link<T>
  predicate ValidCycle( cycle: seq<Link<T>> )
    reads this, cycle
    | \text{cycle} | > 0 \&\&
    this = \operatorname{cycle}[|\operatorname{cycle}|-1] &&
    tail = cycle[0] \&\&
    tail. ValidSequence (cycle) &&
    forall p,q \mid 0 \le p < q < |cycle| ::
      cycle[p] != cycle[q]
  }
  predicate ValidSequence( sequence: seq<Link<T>> )
    reads this, sequence
    |\text{sequence}| > 0 \&\&
    this = sequence [0] &&
    for all \ i \ | \ 0 <= \ i < \ |sequence| -1 \ ::
      sequence[i].tail = sequence[i+1]
  }
  constructor( h: T, x: Link?<T>, ghost cycle: seq<Link<T>>, ghost values: seq<T> )
    modifies if cycle != [] then \{cycle[|cycle|-1]\} else \{\}
    requires (x == null && cycle == []) || (x != null && x.ValidCycle(cycle))
    requires values = ValueSeq(cycle)
    requires forall i | 0 <= i < |cycle| :: cycle[i].head == values[i]
    ensures head == h
    ensures tail = if x = null then this else cycle[0]
    ensures fresh (this)
    ensures forall i | 0 <= i < |cycle| :: cycle[i].head == values[i]
    ensures forall i | 0 <= i < |cycle| :: cycle[i].head == old(cycle[i].head)
    ensures for all z \mid z in cycle :: z.head = old(z.head)
    ensures ValueSeq(cycle) = values
    ensures ValidCycle (cycle+[this])
    ensures ValueSeq(cycle+[this]) = values+[h]
    head := h;
    tail := this;
    new;
    if x != null
      tail := x.tail;
      x.tail := this;
      HeadsEqual(cycle , values);
      AppendLink(cycle, this);
  }
method RemoveFirst<T>
                          ( last: Link < T > 
                          , ghost cycle: seq<Link<T>>
                          , ghost vals: seq<T>
)
```

```
returns ( newlast: Link?<T>
           , ghost newcycle: seq < Link < T >>
           , ghost newvals: seq<T>
  modifies last
  requires last. ValidCycle (cycle)
  requires ValueSeq(cycle) == vals
  requires | vals | == | cycle |
  requires |vals| > 0
  ensures last.head == old(last.head)
  ensures |vals| == 1 \Longrightarrow
             newlast == null &&
             newcycle = [] \&\&
            newvals = []
  ensures |vals| > 1 \Longrightarrow
             newlast == last &&
             newcycle = cycle[1..] &&
             newlast. ValidCycle (cycle [1..]) &&
             newvals = vals [1..] &&
             ValueSeq(newcycle) = newvals
  if last.tail == last
    newlast := null;
    newcycle := [];
    newvals := [];
    return;
  ValueSeqHeads (cycle, vals);
  newlast := last;
  newcycle := cycle[1..];
  newvals := vals [1..];
  newlast.tail := newlast.tail.tail;
  HeadsEqual(cycle[1..], vals[1..]);
function ValueSeq<T>(x: seq<Link<T>>): seq<T>
  reads x
  ensures |x| = |ValueSeq(x)|
  if x = [] then
    else
    [x[0]. head] + ValueSeq(x[1..])
lemma ValueSeqHeads<T>( x: seq<Link<T>>, v: seq<T> )
  requires v = ValueSeq(x)
  ensures for all i \mid 0 \le i < |x| :: x[i]. head = v[i]
  if |x| = 0 { return; }
  ValueSeqHeads(x[1..],v[1..]);
lemma AppendLink<T>( x: seq<Link<T>>, z: Link<T> )
  ensures ValueSeq(x+[z]) = ValueSeq(x)+[z.head]
  if x = [] \{ return; \}
```

```
AppendLink (x[1..], z);
  calc ==
    x+[z];
    (x[..1] + x[1..]) + [z];
    x[..1] + (x[1..] + [z]);
}
lemma HeadsEqual<T>( y: seq<Link<T>>, v: seq<T> )
  requires |y| = |v|
  requires for all i \mid 0 \le i < |y| :: v[i] \Longrightarrow y[i]. head
  ensures v = ValueSeq(y)
  if |y| = 0 { return; }
  HeadsEqual(y[1..], v[1..]);
method Factory() returns (q: Queue<int>)
  ensures fresh (q)
  ensures fresh (q. Repr)
  ensures q. Valid()
  ensures q. IsEmpty()
{
  q := new QueueCycleChain<int>();
method Main()
  var q1 := Factory();
  var q2 := Factory();
  var q3 := Factory();
  var q4 := Factory();
  q1. Put(1);
  q2. Put (1);
  q3. Put (1);
  q4. Put (1);
  q1. Put (2);
  assert q1.ghostseq = [1,2];
  q2. Put (2);
  q3. Put (2);
  q4. Put(2);
  var x;
  x := q1.Get(); print x; print "";
  assert x == 1;
  x := q1.Get(); print x; print "";
  assert x == 2;
  assert q1.ghostseq = [];
  x := q2.Get(); print x; print "";
  x := q2.Get(); print x; print "";
  x := q3.Get(); print x; print
  x := q3.Get(); print x; print
  x := q4.Get(); print x; print "";
  x := q4.Get(); print x; print "";
  assert q1. IsEmpty();
  assert q2. IsEmpty();
  assert q3. IsEmpty();
  assert q4. IsEmpty();
```

}