# TÖL212M Rökstudd Forritun - Hópverkefni 8

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## Hópverkefni 8

### 1

Sækið skrána H8-skeleton.dfy og vistið hana hjá ykkur en breytið nafni hennar í H8.dfy. Klárið að forrita föllin í skránni. Þetta er afbrigði af quickselect sem C.A.R. Hoare fann upp um leið og hann þróaði quicksort.

#### 1.1 Svar:

Hér fyrir neðan má sjá kóðann þar sem föllin hafa verið forrituð. Dafny samþykkir þessa útgáfu. Einnig er hægt að skoða kóðann á þessari slóð: https://tinyurl.com/yntnzx4e.

```
// Author of question:
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// Author of solution:
// Permalink of solution:
                              https://tinyurl.com/yntnzx4e
// Finish programming the two methods
method Partition (m: multiset < int > )
  returns ( pre: multiset < int >, p: int , post: multiset < int > )
  decreases |m|
  requires |\mathbf{m}| > 0
  ensures p in m
  ensures m = pre+multiset\{p\}+post
  ensures for all z \mid z in pre :: z \le p
  ensures for all z \mid z in post :: z >= p
  // The body is missing.
  // You can use a loop or recursion.
  // Remove one value from m.
  var x : | x in m;
  var m' := m-multiset\{x\};
  // If m' is then empty, we're done and return x,
  // and the empty sets on both sides (m')
  if (m' = multiset \{\}) { return m', x, m'; }
  // If m' is not empty, then we recursively partition m'.
  pre, p, post := Partition(m');
  // Here we need to remind Dafny that m is the same as m' + multiset\{x\}.
  assert m == m' + multiset\{x\};
  // If x is less than or equal to the pivot from that partition,
  // we add x to the left partition and put p as the new pivot.
```

1.1 Svar:

```
if (x \le p) { return pre+multiset \{x\}, p, post; }
  // If x is greater than the pivot from the partition,
  // we add x to the right partition.
  else { return pre, p, post+multiset {x}; }
method QuickSelect (m: multiset < int >, k: int )
  returns ( pre: multiset < int > , kth: int , post: multiset < int > )
  decreases m
  requires 0 \ll k \ll |m|
  ensures kth in m
  ensures m = pre+multiset\{kth\}+post
  ensures | pre | == k
  ensures for all z | z in pre :: z <= kth
  ensures forall z | z in post :: z >= kth
  // The body is missing.
  // You can use a loop or recursion.
  // Use the Partition method as a helper method.
  // We partition the multiset m.
  var p, piv, r := Partition(m);
  // If the size of the left partition is equal to k, we're done.
  if (|p| = k) { return p, piv, r; }
  // If the size of the left partition is less than k,
  // we recursively partition the right partition.
  else if (|p| < k)
    \operatorname{var} p2, \operatorname{newpiv}, \operatorname{r2} := \operatorname{QuickSelect}(r, k-|p|-1);
    return p + multiset{piv} + p2, newpiv, r2;
  }
  // If the size of the left partition is greater than k,
  // we recursively partition the left partition.
  else {
    var p2, newpiv, r2 := QuickSelect(p, k);
    return p2, newpiv, r + multiset{piv} + r2;
  }
}
```