

The background of the slide features a soft, out-of-focus image of a sunset or sunrise. A semi-circular sun is visible on the horizon, casting a warm glow. In the distance, a range of mountains is silhouetted against the sky. The overall color palette is dominated by light blues, purples, and whites, creating a calm and professional atmosphere.

Day 3

Graphs for Discrete Maths

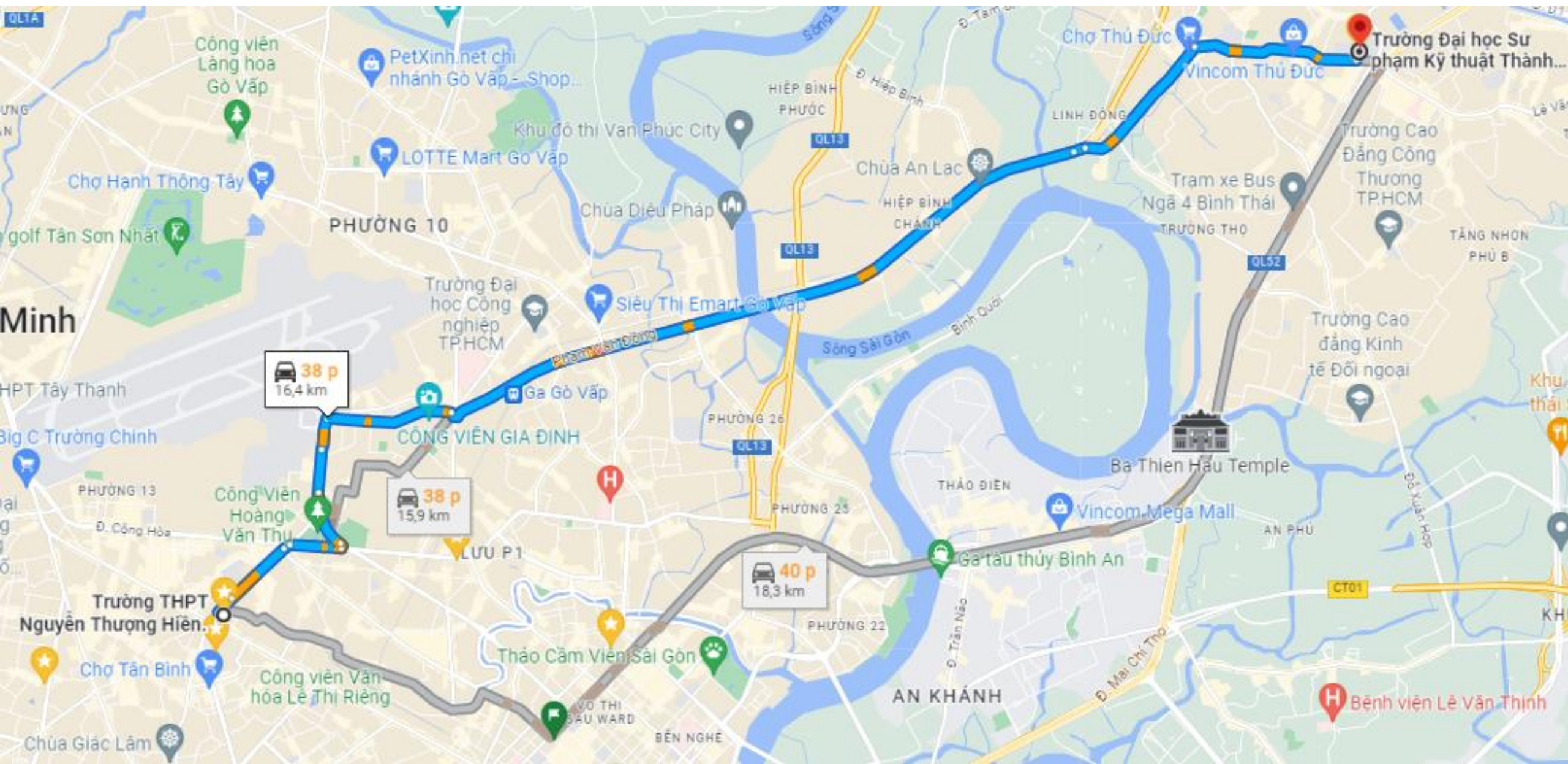
Lecturer: Msc. Minh Tan Le

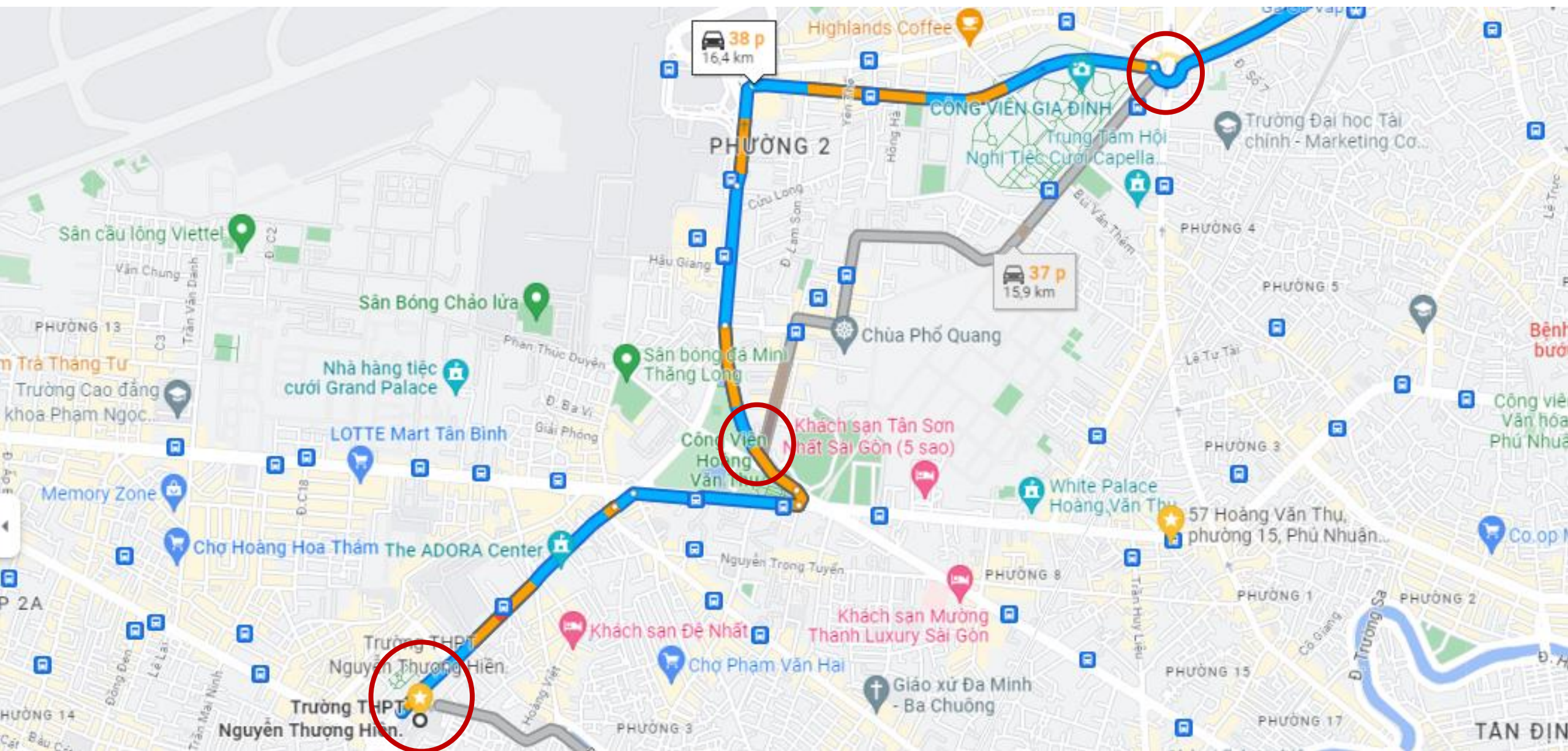
Overview of day 2

- Proposition calculation
- Relations of sets
 - Binary relation (2 sets)
 - N-ary relation (multiple sets)
- Type of relations
 - Reflexive
 - Symmetric
 - Antisymmetric
 - Transitive
 - Equivalent
- Operators: Union, intersection, minus.

Outline

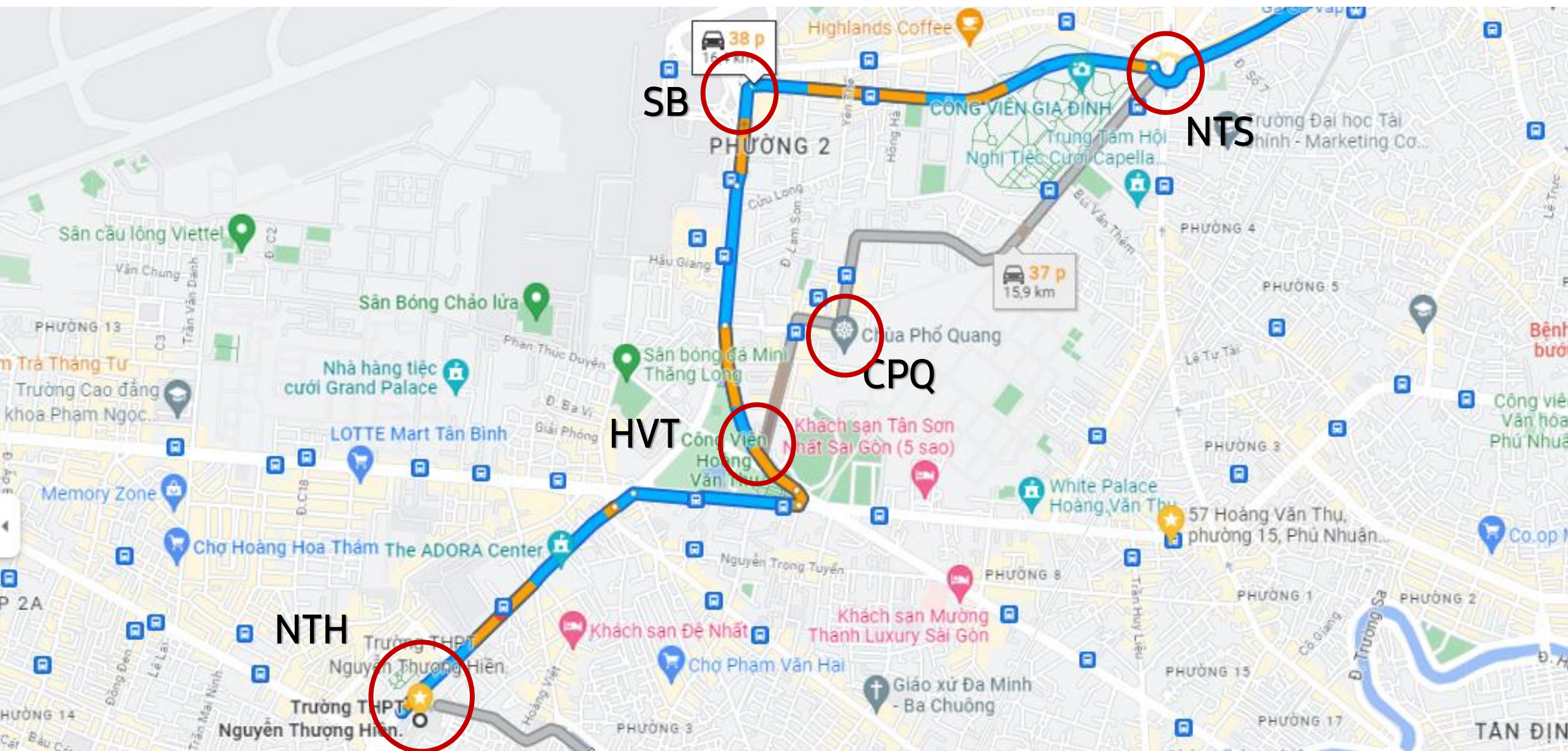
- I. Graph
- II. Digraph & Hasse diagram
- III. Boolean algebra: The essence
- IV. Graph theory: The essence





Definition

- A graph consists of a nonempty set of vertices (**nodes**) and a set of edges (**arcs**, links, lines, connections).
- Expression: $G = (V, E)$



Digraph

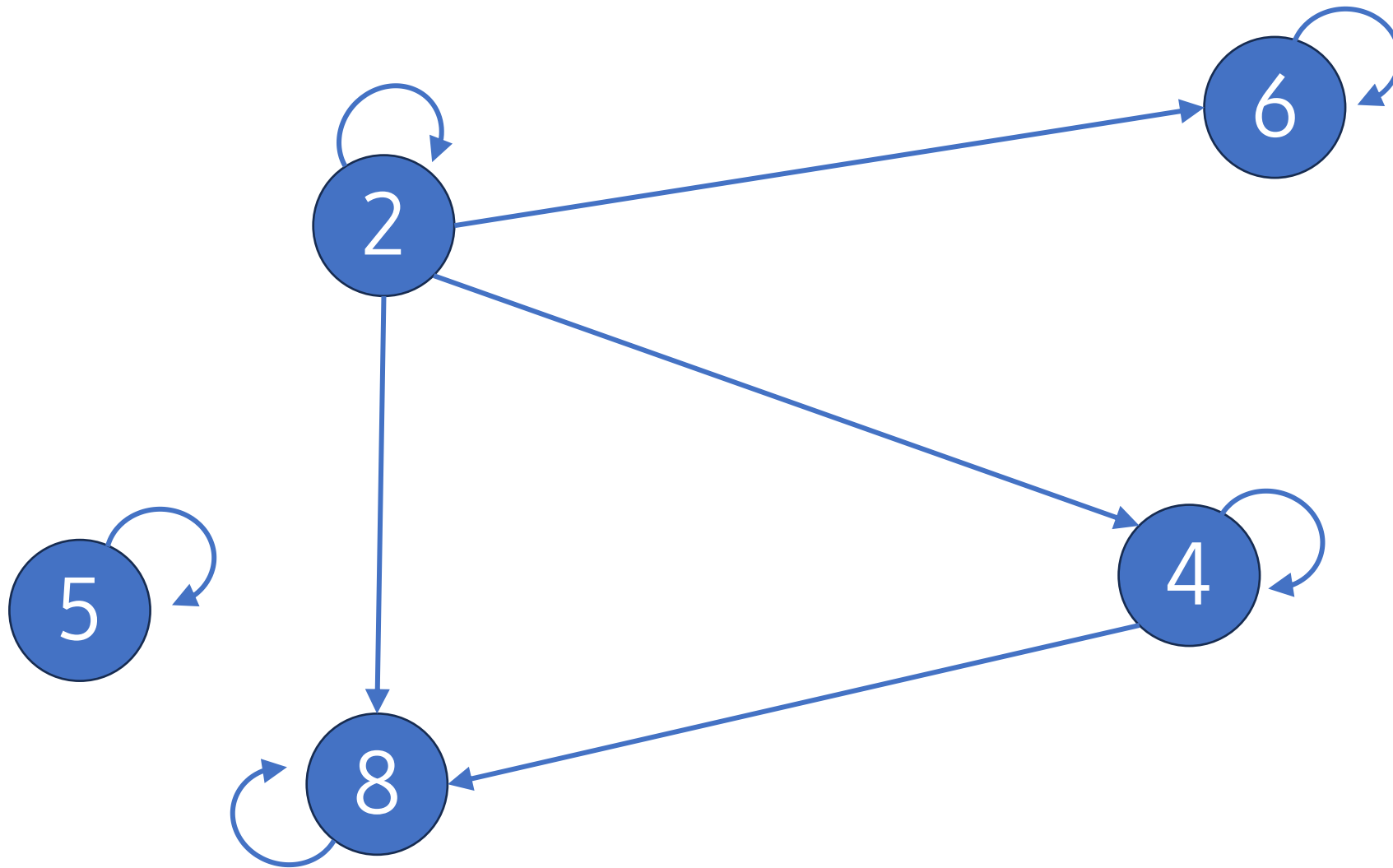
- A digraph is a graph, but with directional edges.
 - Ex: $v = (\text{NTH}, \text{HVT})$ represents by a line with arrow towards HVT from NTH.



$$A = \{2, 4, 5, 6, 8\}$$

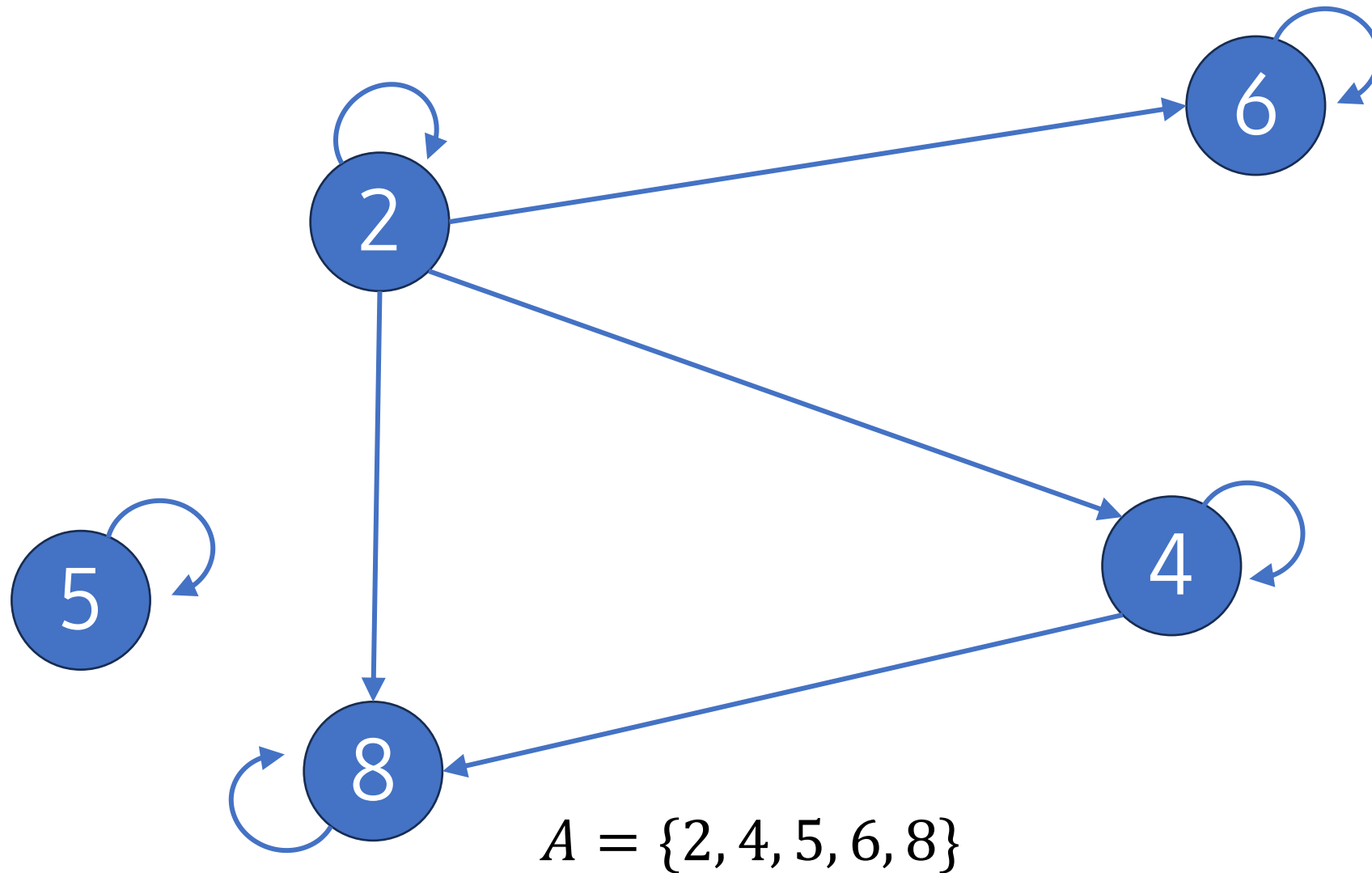
$$R = \{(2,2), (2, 4), (2, 6), (2, 8), (4, 4), (4, 8), (5, 5), (6, 6), (8, 8)\}$$

- Each $a \in A$ is a node/**vertex (vertices)**/point.
- Each $r \in R$ is a edge/**arc**.
- The number of edges connected to/from a is the degree of a , or $\deg(a)$.



$$A = \{2, 4, 5, 6, 8\}$$

$$R = \{(2,2), (2,4), (2,6), (2,8), (4,4), (4,8), (5,5), (6,6), (8,8)\}$$



$$R = \{(2,2), (2,4), (2,6), (2,8), (4,4), (4,8), (5,5), (6,6), (8,8)\}$$

R is partial ordering on set A , or A is partially order set (POSET)

Partially order is built on top of algebra

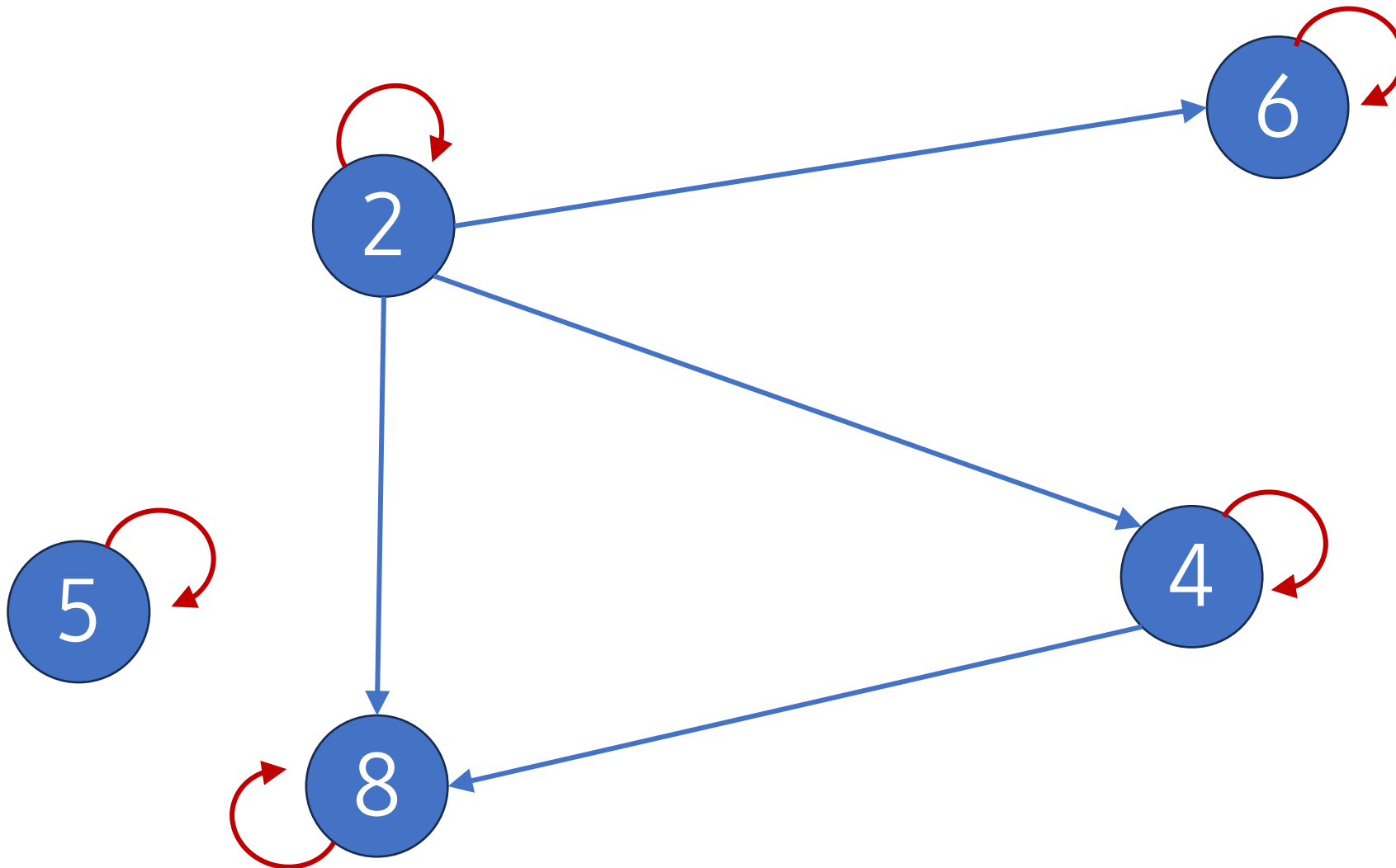
- If there are at least 2 **distinctive** elements, the sequence has **partially order**.
- If a, b, c ($a \neq b$, ex: $a < b$) are real numbers, then:
 - $a = a, b = b, c = c$ is always true (reflexive).
 - $b < a$ cannot happen (antisymmetric).
 - If $b < c$, then $a < c$ (transitive).
 - You **can sort** the sequence a, b, c .

How about equivalence?

- If all elements are **equal**, the order is **meaningless** in algebra.
- If a, b, c are equal real numbers, then:
 - $a = a, b = b, c = c$ is always true (reflexive).
 - $a = b$ and $b = a$ (symmetric).
 - $a = b, b = c, a = c$ (transitive).
 - You **cannot sort** a, b, c .

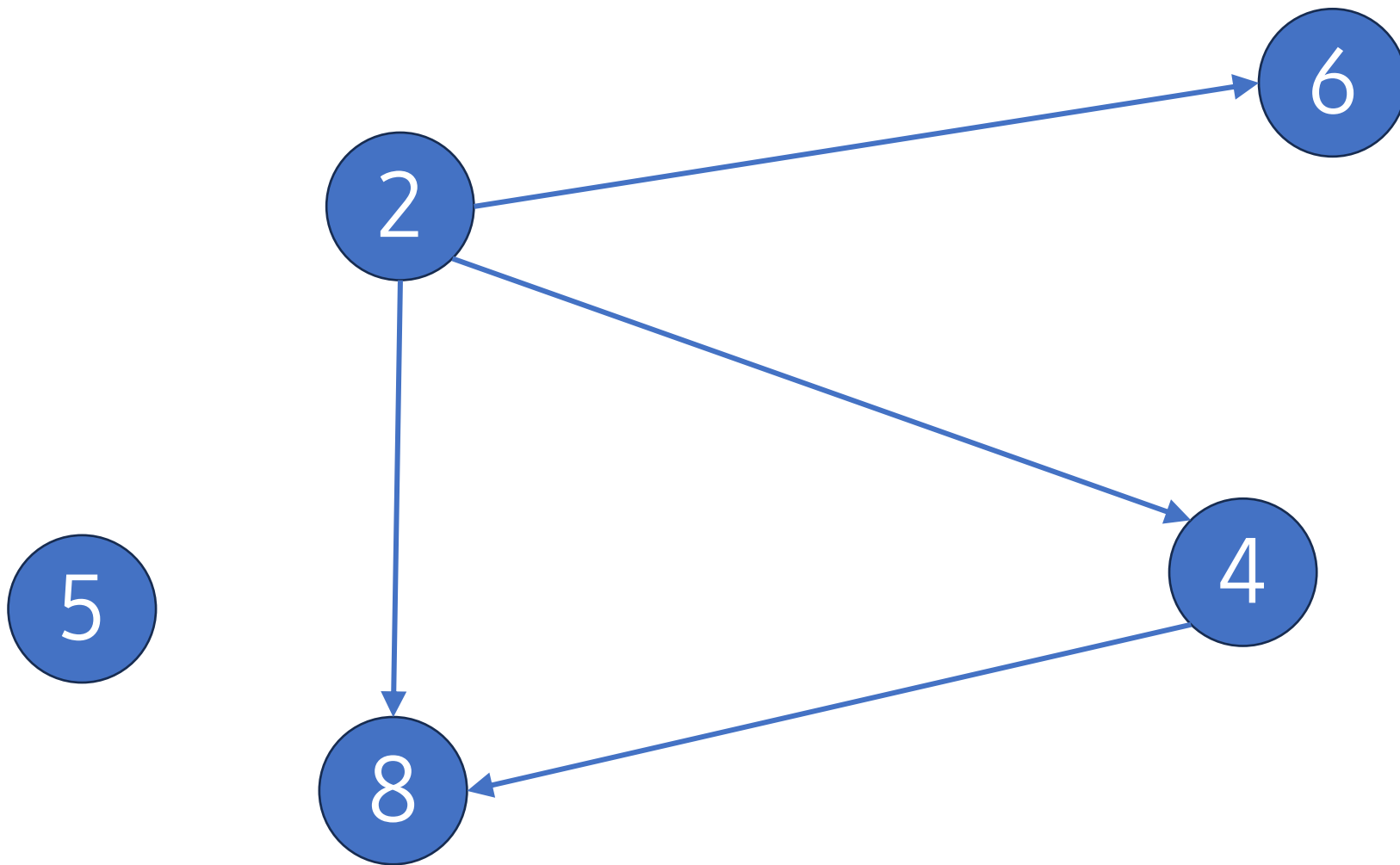
A set is partially order if...

- (By definition) there is at least a pair of element that has order.
- (By proving method) all these conditions are satisfied:
 - Reflexive
 - Antisymmetric
 - Transitive



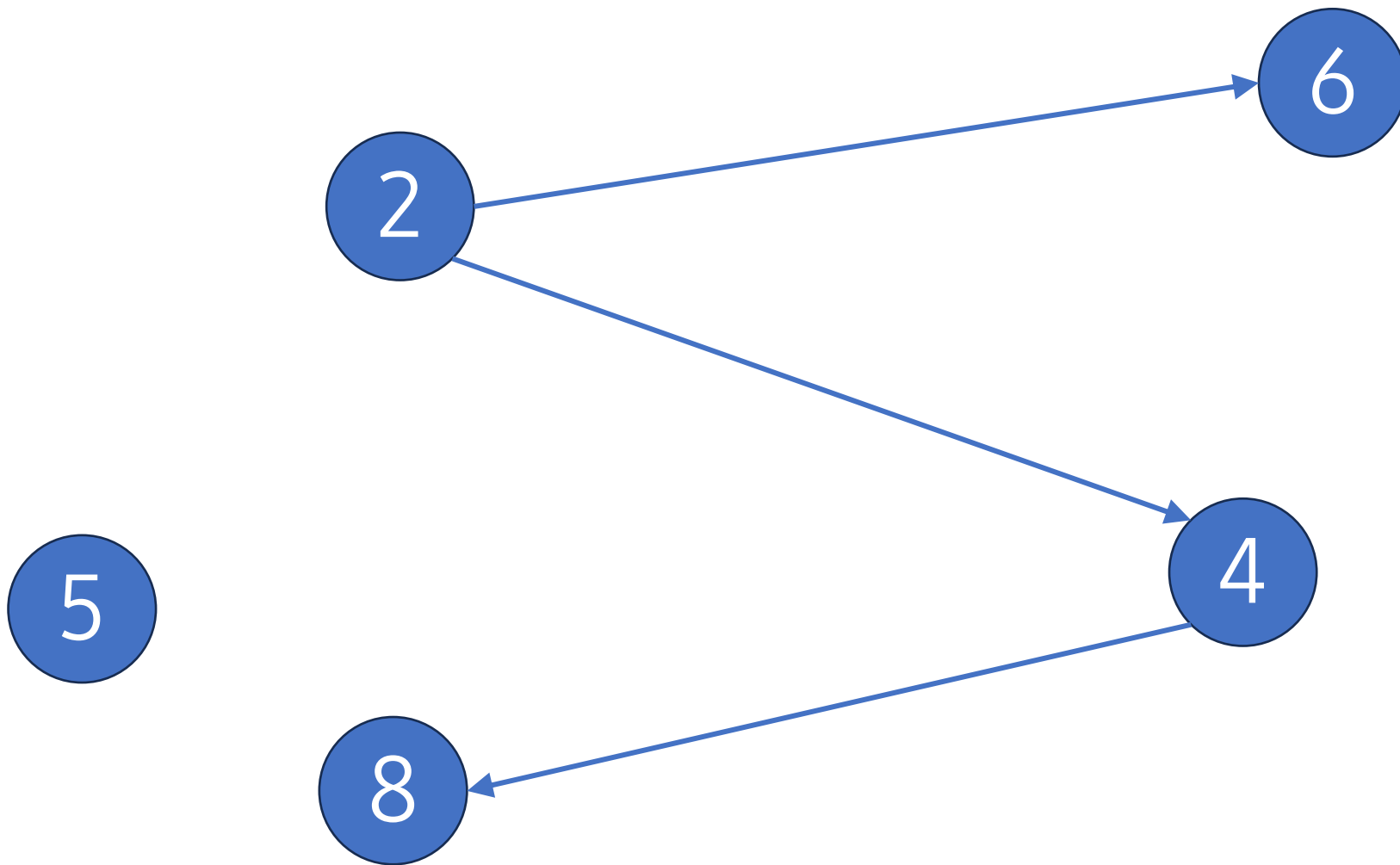
Graph of R

It seems the graph is quite complicated. Do you have any ideas how to simplify?



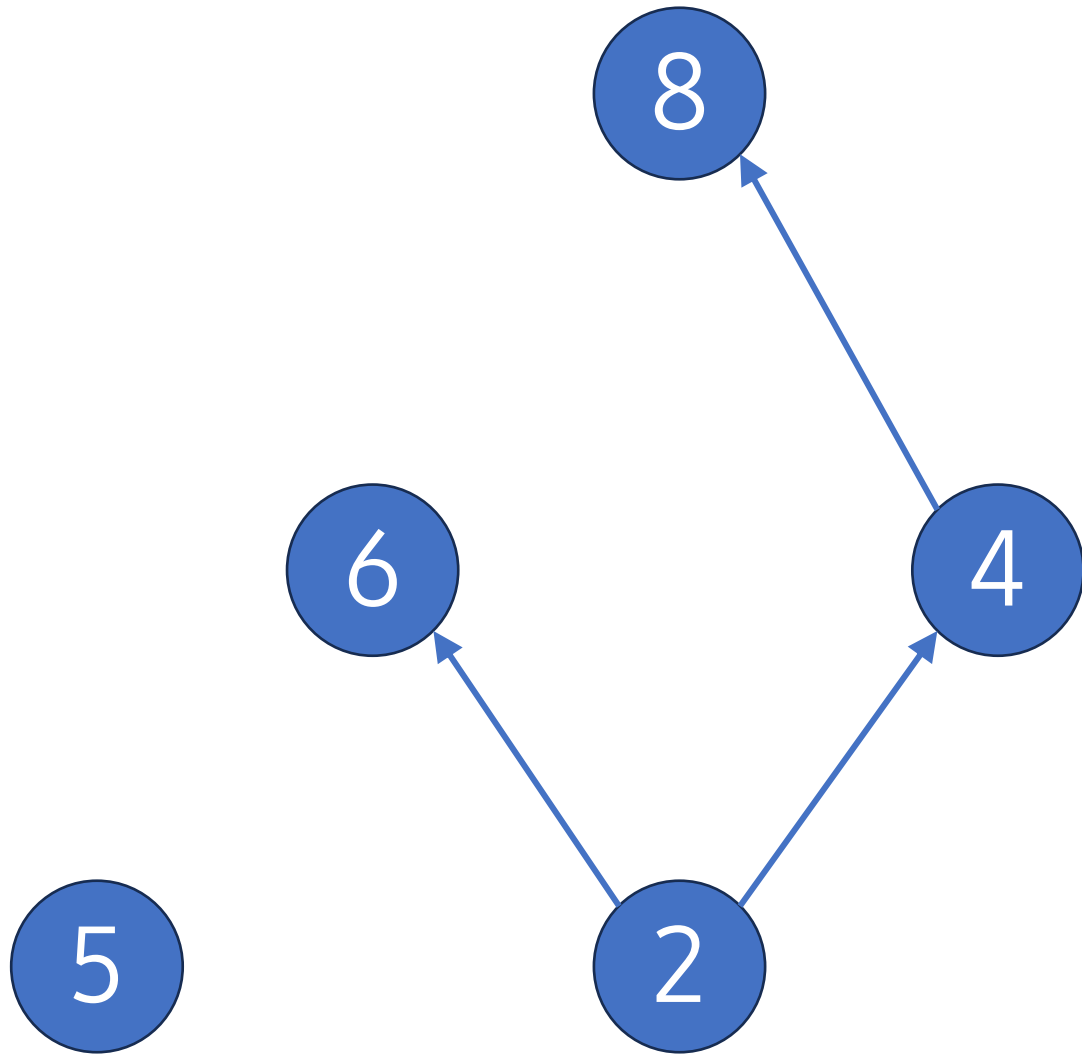
✓ Remove loops

Graph of partially order R



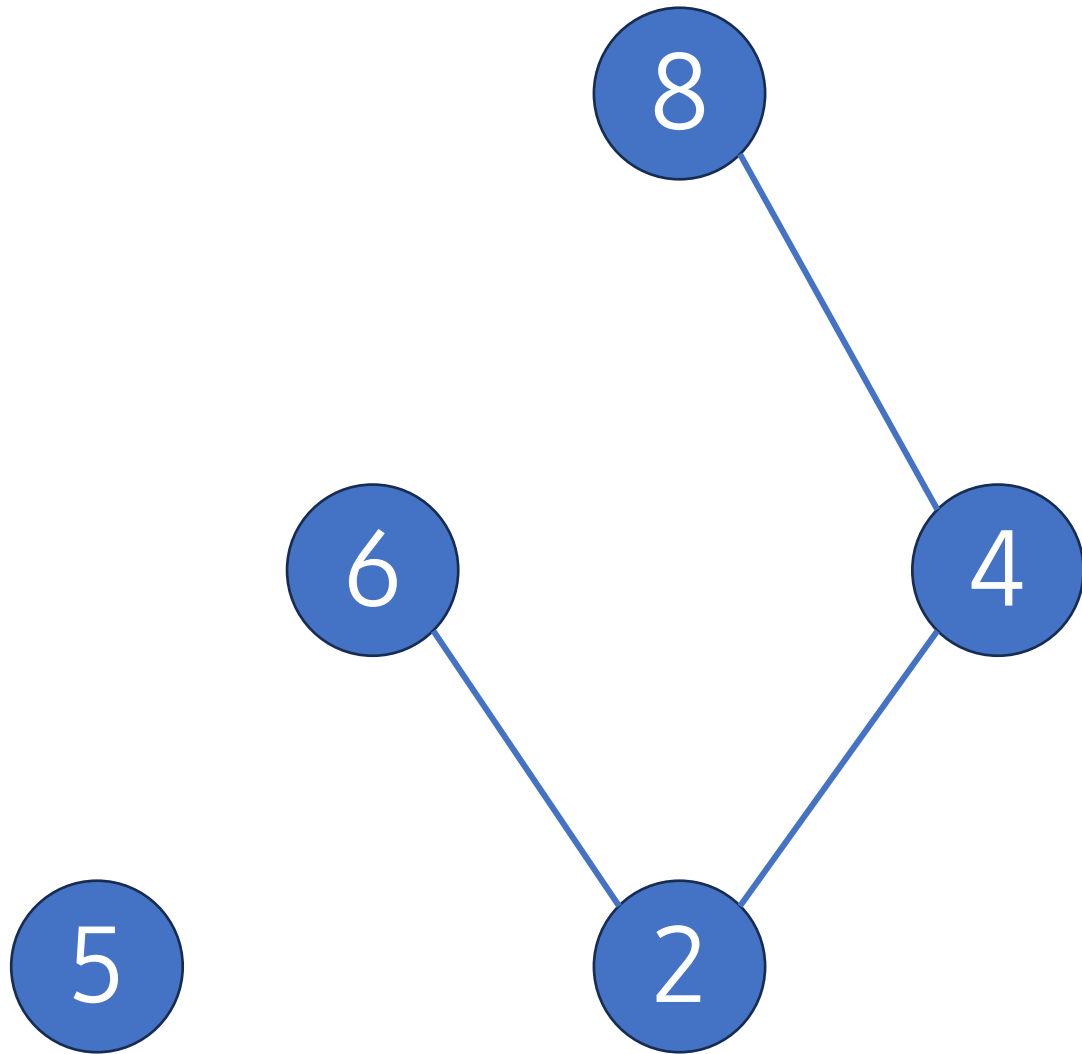
- ✓ Remove loops
- ✓ Remove transitive connection

Graph of partially order R



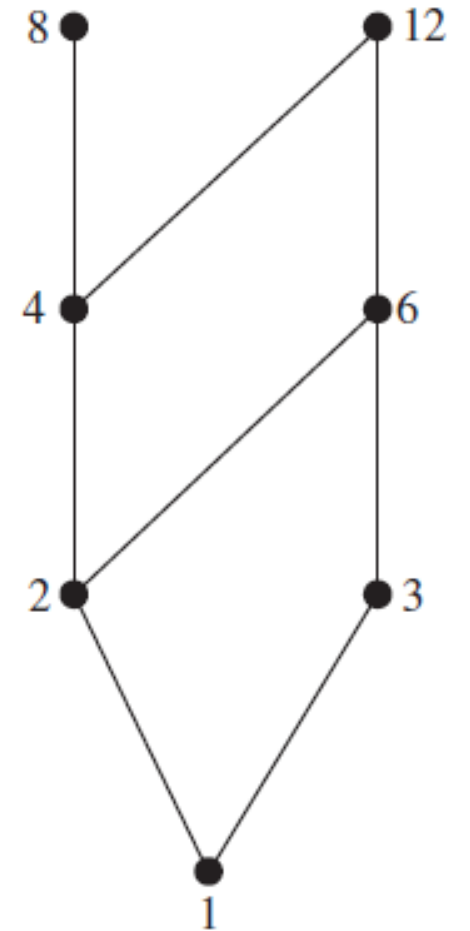
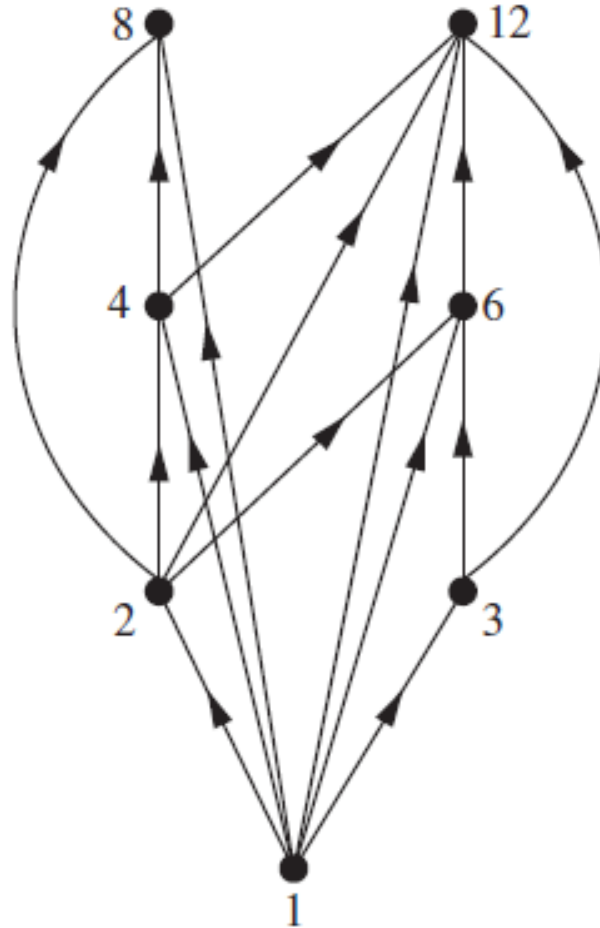
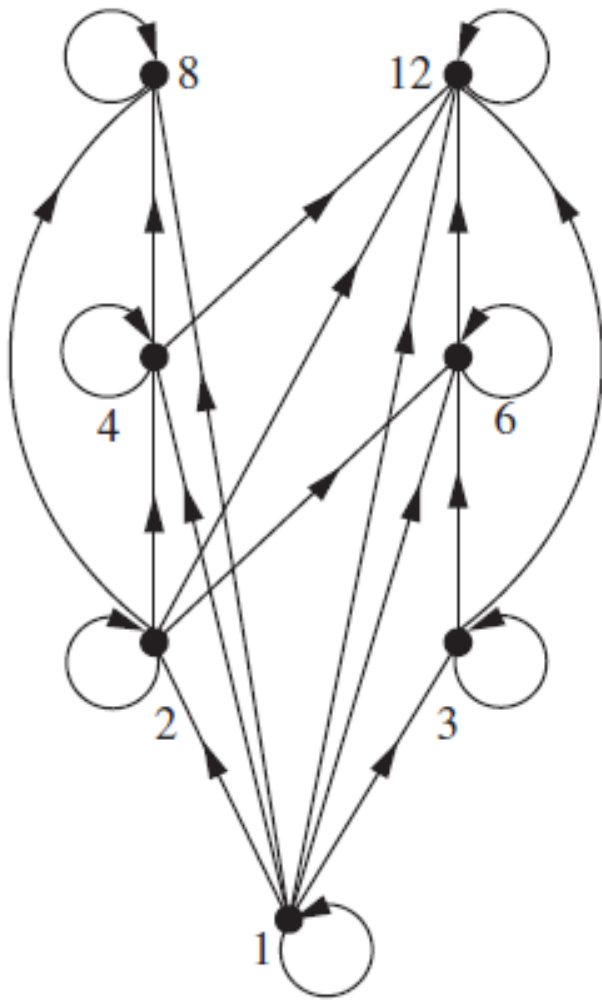
Graph of partially order R

- ✓ Remove loops
- ✓ Remove transitive connection
- ✓ Assuming arrows are pointing upward



Hasse diagram of partially order R

- ✓ Remove loops
- ✓ Remove transitive connection
- ✓ Assuming arrows are pointing upward and remove arrows



*An example of Hasse diagram creatation
p. 624, Discrete Mathematics & It's Application – Kenneth H. Rosen*

Hasse diagram: A simplified diagram

- Reflexive
 - Remove loops
- Transitive
 - If $a \rightarrow b, b \rightarrow c, a \rightarrow c$, remove $a \rightarrow c$.
- Arrow direction assumption
 - Move nodes, remove arrows by assuming they are pointing upward.

Warning!

- *You can only create Hasse diagram from partially ordering.*
- *You can restore back to the original diagram.*

III. Boolean algebra: The essence

- We don't have much thing to do with these sets:
 - {A, B, C}
 - {Tam, Tan, Trang}
 - {1, 2, 3}
- But what if there are only 2 discrete values?
 - {False, True}
 - {0, 1}

Boolean algebra

- Operator
- Expression
- Variable
- Function

Some operators in Boolean algebra

| Order | Propositional logic | Boolean algebra | Definition |
|-------|---------------------|-----------------------------|---------------------|
| 1 | $\neg, -$ | $-$ | NOT/Complementation |
| 2 | \wedge | \cdot (Can be omitted) | AND/Boolean product |
| 3 | \vee | $+$ | OR/Boolean sum |

Note: Calculate by order

Some operators in Boolean algebra

| Order | Propositional logic | Boolean algebra | Definition |
|-------|---------------------|-----------------------------|---------------------|
| 1 | $\neg, -$ | $-$ | NOT/Complementation |
| 2 | \wedge | \cdot (Can be omitted) | AND/Boolean product |
| 3 | \vee | $+$ | OR/Boolean sum |

$$1 \cdot 0 + \overline{0 + 1}$$

Expression vs. function

$$x \cdot 0 + \overline{0 + y}$$

An expression with variables

$$f(x, y) = x \cdot 0 + \overline{0 + y}$$

A function

Find all values of Boolean function

| x | y | $0 + y$ | $\overline{0 + y}$ | $x \cdot 0$ | $f(x, y)$ |
|-----|-----|---------|--------------------|-------------|-----------|
| 1 | 1 | 1 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 0 | 0 | 0 |
| 0 | 0 | 0 | 1 | 0 | 1 |

$$f(x, y) = x \cdot 0 + \overline{0 + y}$$

Part IV. Graph theory: The essence

- Degree, rank, path
- Types of graphs determined by visuals

Degree: The number of edges connected to/from a .

$\deg(v)$: Degree

$\deg^-(v)$: Number of edges that are connection to v .

$\deg^+(v)$: Number of edges that are connection from v .

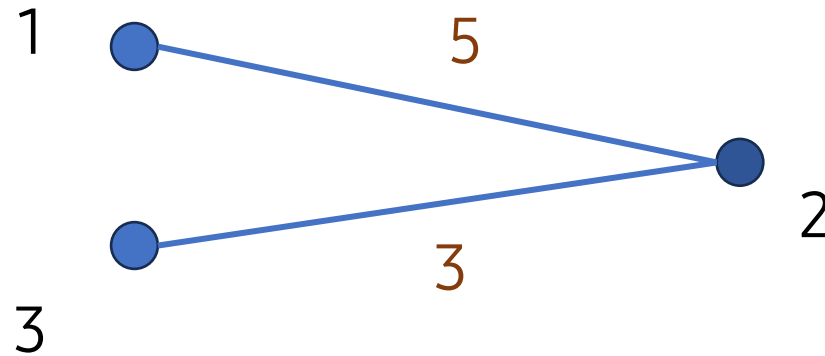
$$\sum_{v \in V} \deg^-(v)$$

$$\sum_{v \in V} \deg^+(v)$$

Rank

- Rank of an undirected graph is similar to rank of a matrix.
 - Demonstrate a graph into matrix.
 - Count number of linearly independent columns.

Path between two nodes is a string of vertices that are sequentially connected



$$p(v, v') = (v, v_1, v_2, \dots, v_k, v')$$

$$w(p) = \sum_{i=0}^k w(v_i, v_{i+1})$$

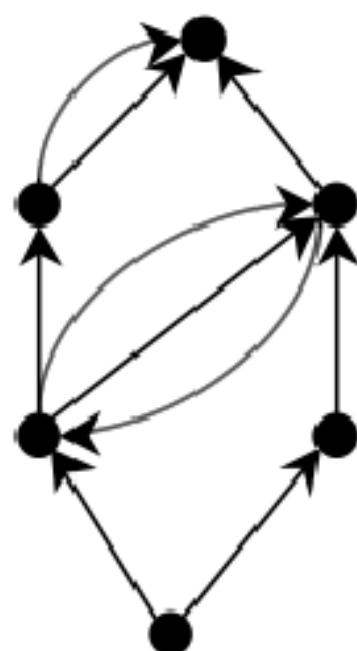
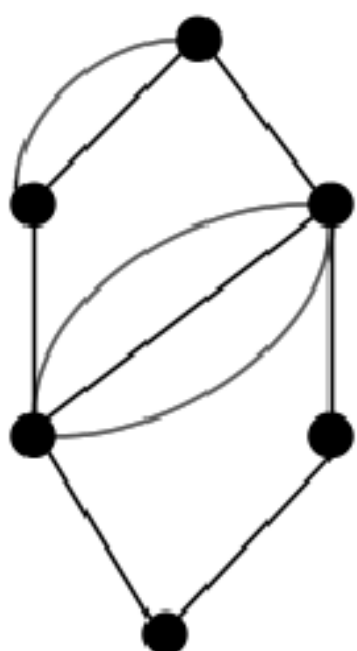
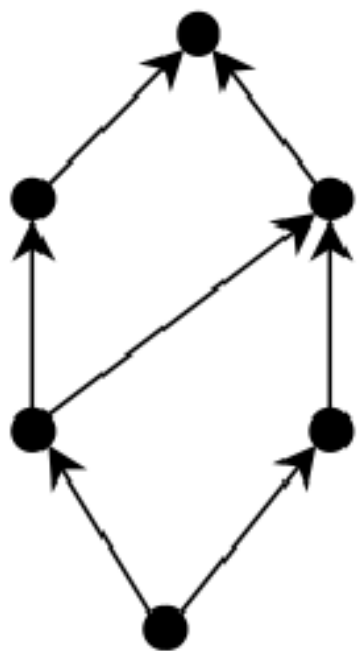
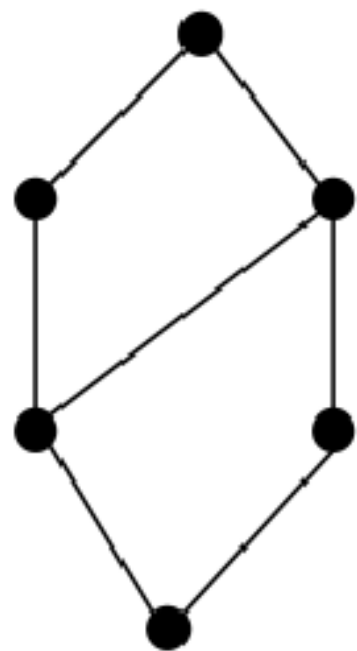
- **Shortest path:** The most efficient path between 2 nodes.

$$d_{w(v,v')} = \min w(p)$$

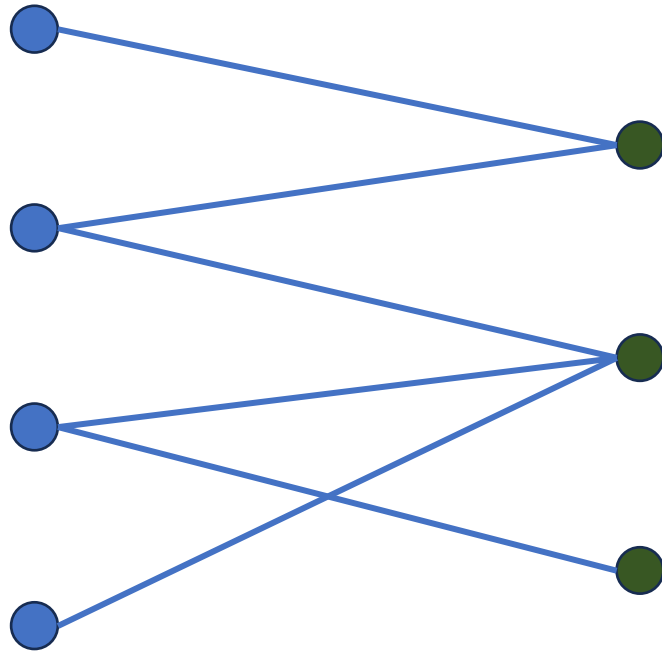
- There are quite many ways to find the shortest path.
- Keyword: *Shortest path problem*.

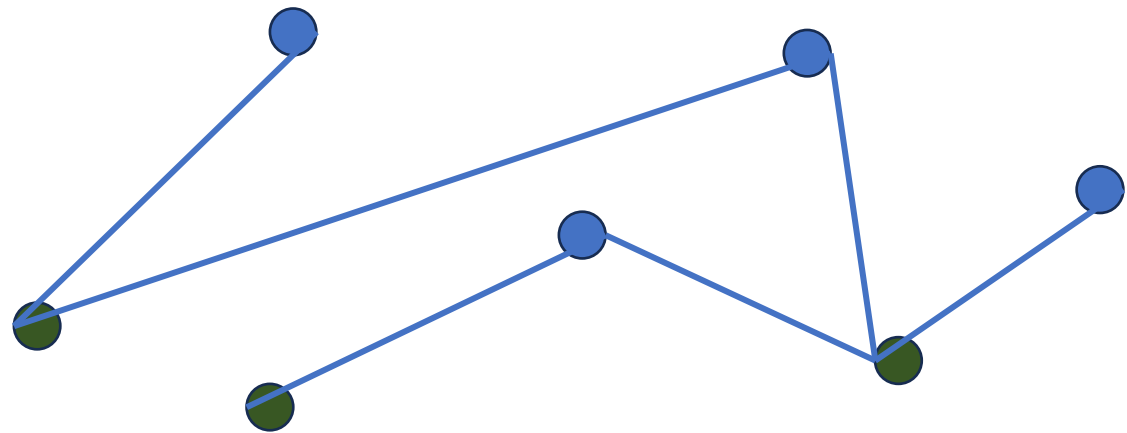
Some popular types

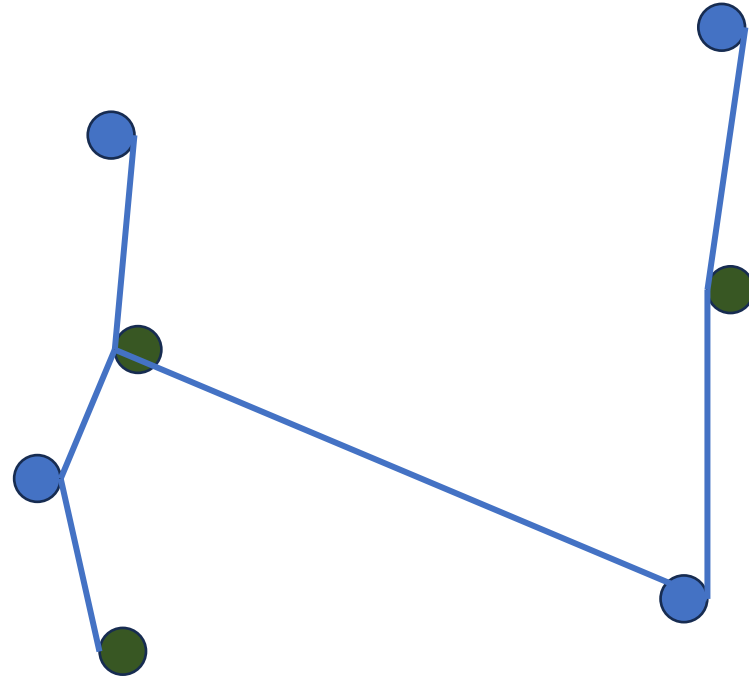
- Digraph (directed graph): Each edge has direction.
- Undirected graph: No direction for all edges.
- Simple graph: At most one edge between two nodes.
- Multigraph: Multiple edges between two nodes are allowed.

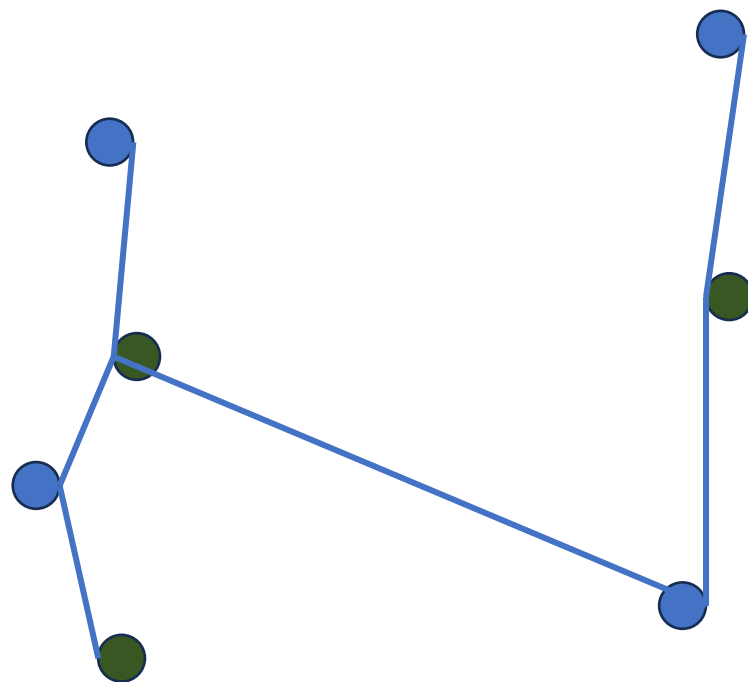
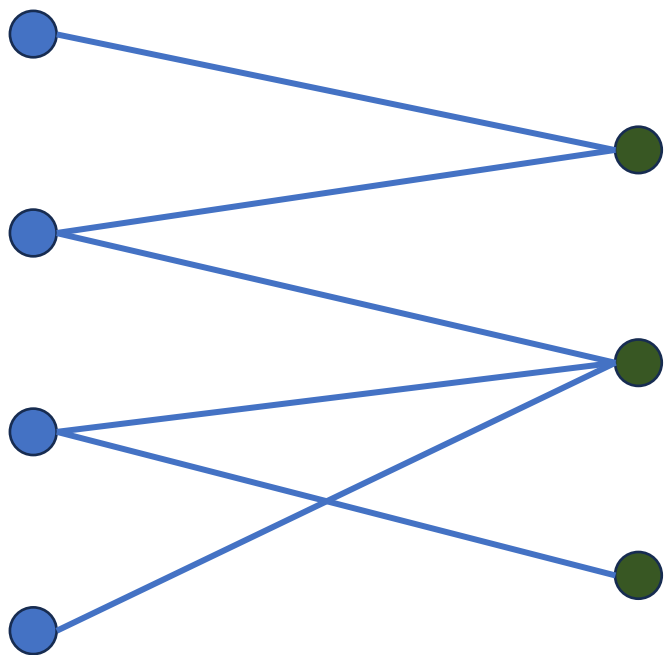


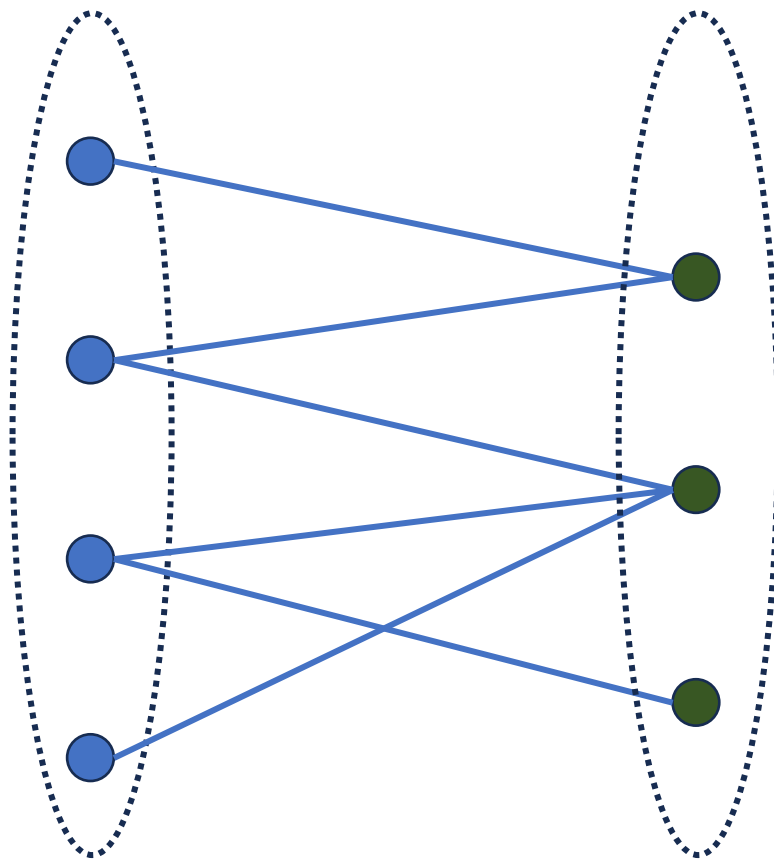
Bipartite graph







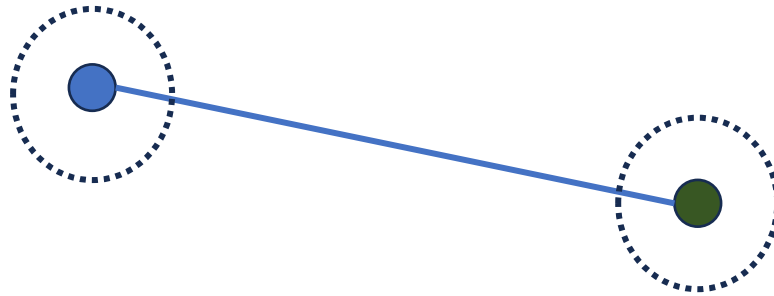




$K_{4,3}$

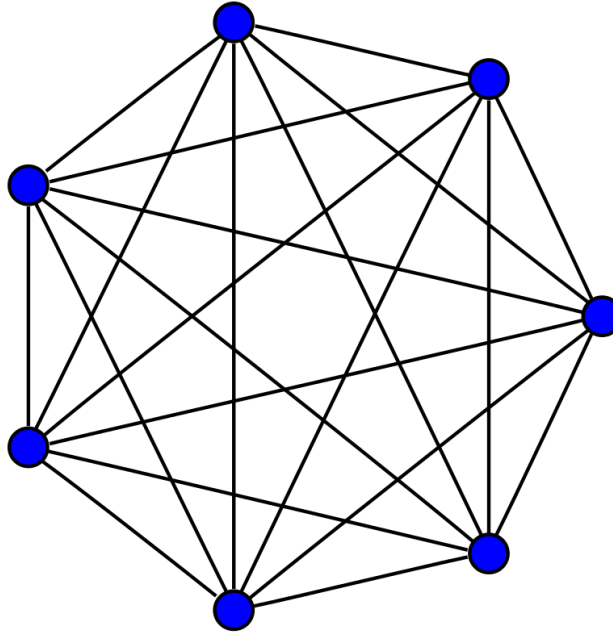
Bipartite graph: The set of vertices can be separated into two that any edge connects nodes of both.

A bipartite graph that has two equal size of set is called a **balanced bipartite graph**.



Complete graph

- An undirected simple graph
- There is one edge between any pair of nodes



A K_7 complete graph

Complete graph properties

- Number of vertices: n
- Number of edges: $C_n^2 = \frac{n(n-1)}{2}$
- All degrees are equal: $n - 1$
- Expression: K_n
- K_1 is also a complete graph.

Complete bipartite graph

- Given s, t as sizes of two sets, we get:

$$K_{s,t}$$

=> The expression is the same as bipartite graph, so we may need to clarify that the graph is also complete.

Exercises/homework

- 22/631 (individual)
- 1, 2, 3, 23/630, 631 (group)

Notes:

- Remember to determine the lacked property.
- Draw initial graph before finding the Hasse diagram.
- Divisibility: $\{(a, b)\}$ which b can be fully divided by a .