

1. Let $B = (F, P)$ be a basis for predicate logic with $F = \{0, 1, \dots, +, *, \dots\}$ and $P = \{<, \leq, \dots\}$. Furthermore, let I be the usual interpretation of the function and predicate symbols over the set Nat of natural numbers. Determine for each of the following formulas $w \in \text{WFF}_B$ the set of all assignments $\sigma: V \rightarrow \text{Nat}$ with $I(w)(\sigma) = \text{true}$:

- (a) $\neg \exists y. (x < y \wedge y < z)$
 (b) $\exists x. \forall x. y \leq x$
 (c) $x = 0 \rightarrow x = 1 \rightarrow x = 2$
 (d) $\exists y. x * x = x \rightarrow x \leq x$

Which of these formulas are valid in I ?

2. Let $B = (F, P)$ and I be as in Exercise 1. Give a formula $w \in \text{WFF}_B$ such that, for each assignment σ , $I(w)(\sigma) = \text{true}$ exactly when

- (a) $\sigma(z)$ divides $\sigma(x)$
 (b) $\sigma(x)$ and $\sigma(y)$ are relatively prime
 (c) $\sigma(z)$ is the gcd of $\sigma(x)$ and $\sigma(y)$.

3. For each of the following formulas give the set of all interpretations in which this formula is valid:

- (a) $x = y$ (b) $\forall x, y. x = y$
 (c) $\forall x. \exists y. x = y$ (d) $\exists y. \forall x. x = y$
 (e) $\exists x, y. \forall z (z = x \vee z = y)$ (f) $x = y \wedge y = z \rightarrow x = z$
 (g) $f(x) = f(y) \rightarrow x = y$ (h) $x = y \rightarrow f(x) = f(y)$

Which of these formulas are logically valid?

4. Show that for all formulas w, w_1, w_2 each of the following pairs of formulas are logically equivalent:

- (a) $w_1 \rightarrow w_2$ and $\neg w_1 \vee w_2$
 (b) $\forall x. w$ and $\neg \exists x. \neg w$
 (c) true and $w \vee \neg w$

5. Substitute the variable y by the term $y + z + x$ in the formula $\exists z. y * z = x$.
 6. Let V be the set of variables in predicate logic, and let $P(V)$ denote the power set of V , i.e. the set of all subsets of variables. Give an inductive definition of function

$$\text{free}: \text{WFF}_B \rightarrow P(V),$$

that maps every predicate logic formula $w \in \text{WFF}_B$ to the set $\text{free}(w)$ of all variables occurring free in w .

7. Let $B = (F, P)$ be a basis for predicate logic with $F = \{0, 1, +\}$ and $P = \emptyset$. Suppose $W \subseteq \text{WFF}_B$ consists of the following formulas:

- $\forall x, y, z. (x + (y + z) = (x + y) + z)$
 $\forall x. (x + 0 = x \wedge 0 + x = x)$
 $\forall x. \exists y. (x + y = 0 \wedge y + x = 0)$

Give a model of W such that the following formula is also valid over this model:

$$\forall x, y. \exists z, z'. (x + z = y \wedge z' + x = y)$$

8. Let $Z = \{x + x = x, \forall y. y + 0 = y\}$. Which of the following formulas are the logical consequences of Z in predicate logic?

- (a) $x = 0$ (b) $1 + 0 = 1$ (c) $0 + 1 = 1$ (d) $1 + 1 = 1$

9. Give a deduction of the formula $\forall x. w \rightarrow w_x^t$ from the empty set in the predicate calculus.

10. Let $W \subseteq \text{WFF}_B$ be a set of formulas with at least one model. Show that $\text{Cn}(W)$ is a theory.